

Algebraic Geometry Codes

Group 5.239A
Martin Sig Nørbjerg

Institute of Mathematics
Aalborg University



AALBORG UNIVERSITY
DENMARK

Linear Error Correcting Codes

Martin Sig Nørbjerg

Coding Theory

1

Bounds on the Parameters
of Codes

Divisors

Goppa Codes

Definition 1.1 & 1.5. Let $\mathcal{C} \subseteq \mathbb{F}_q^n$ be a linear subspace of dimension k , then \mathcal{C} is called a $[n, k]_q$ code. Furthermore if \mathcal{C} has minimum distance d , then \mathcal{C} is called a $[n, k, d]_q$ code.

Linear Error Correcting Codes

Martin Sig Nørbjerg

Coding Theory

1

Bounds on the Parameters
of Codes

Divisors

Goppa Codes

Definition 1.1 & 1.5. Let $\mathcal{C} \subseteq \mathbb{F}_q^n$ be a linear subspace of dimension k , then \mathcal{C} is called a $[n, k]_q$ code. Furthermore if \mathcal{C} has minimum distance d , then \mathcal{C} is called a $[n, k, d]_q$ code.

- If \mathcal{C} has minimum distance d , then it can correct $\lfloor \frac{d-1}{2} \rfloor$ errors but detect $d - 1$ errors.

Bounds on the Parameters of Codes.

Martin Sig Nørbjerg

Coding Theory

Bounds on the Parameters
of Codes

2

Divisors

Goppa Codes

Corollary 1.18. Let \mathcal{C} be a $[n, k, d]_q$ code, then $d - 1 \leq n - k$.

Bounds on the Parameters of Codes.

Martin Sig Nørbjerg

Coding Theory

Bounds on the Parameters
of Codes

2

Divisors

Goppa Codes

Corollary 1.18. Let \mathcal{C} be a $[n, k, d]_q$ code, then $d - 1 \leq n - k$.
Proof: Let H be a parity check matrix of \mathcal{C} .

Bounds on the Parameters of Codes.

Martin Sig Nørbjerg

Coding Theory

Bounds on the Parameters
of Codes

2

Divisors

Goppa Codes

Corollary 1.18. Let \mathcal{C} be a $[n, k, d]_q$ code, then $d - 1 \leq n - k$.
Proof: Let H be a parity check matrix of \mathcal{C} .

$$\blacktriangleright \dim(\mathcal{C}) = \dim(\text{null}(H)) = k \implies \text{rank}(H) = n - k$$

Bounds on the Parameters of Codes.

Martin Sig Nørbjerg

Coding Theory

Bounds on the Parameters
of Codes

2

Divisors

Goppa Codes

Corollary 1.18. Let \mathcal{C} be a $[n, k, d]_q$ code, then $d - 1 \leq n - k$.

Proof: Let H be a parity check matrix of \mathcal{C} .

- ▶ $\dim(\mathcal{C}) = \dim(\text{null}(H)) = k \implies \text{rank}(H) = n - k$
- ▶ $\text{rank}(H) \geq d - 1$, by the Proposition 1.16.



Divisors

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

3

Definition 2.81 & 2.82. Let \mathcal{X} be an absolutely irreducible regular projective plane curve, over \mathbb{F}_q .

Definition 2.81 & 2.82. Let \mathcal{X} be an absolutely irreducible regular projective plane curve, over \mathbb{F}_q .

► A divisor D on \mathcal{X} is a formal sum:

$$D = \sum_{P \in \mathcal{X}} n_P P$$

where $n_P = 0$ for all but a finite number of points $P \in \mathcal{X}$.

Definition 2.81 & 2.82. Let \mathcal{X} be an absolutely irreducible regular projective plane curve, over \mathbb{F}_q .

- ▶ A divisor D on \mathcal{X} is a formal sum:

$$D = \sum_{P \in \mathcal{X}} n_P P$$

where $n_P = 0$ for all but a finite number of points $P \in \mathcal{X}$.

- ▶ Let $f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\}$, then $(f) := \sum_{P \in \mathcal{X}} v_P(f) P$ is called a principal divisor.

Definition 2.81 & 2.82. Let \mathcal{X} be an absolutely irreducible regular projective plane curve, over \mathbb{F}_q .

- ▶ A divisor D on \mathcal{X} is a formal sum:

$$D = \sum_{P \in \mathcal{X}} n_P P$$

where $n_P = 0$ for all but a finite number of points $P \in \mathcal{X}$.

- ▶ Let $f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\}$, then $(f) := \sum_{P \in \mathcal{X}} v_P(f) P$ is called a principal divisor.
 - ▶ $\forall f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\}$ we have $\deg((f)) = 0$ by Proposition 2.83.

The vector space $L(D)$

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

4

Definition 2.84. Let $D \in \text{Div}(\mathcal{X})$, then we define the vector space $L(D)$ as:

$$L(D) := \{f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\} \mid (f) + D \text{ is effective}\} \cup \{0\}$$

and let $\ell(D) := \dim_{\overline{\mathbb{F}}_q}(L(D))$.

9

The vector space $L(D)$

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

4

Definition 2.84. Let $D \in \text{Div}(\mathcal{X})$, then we define the vector space $L(D)$ as:

$$L(D) := \{f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\} \mid (f) + D \text{ is effective}\} \cup \{0\}$$

and let $\ell(D) := \dim_{\overline{\mathbb{F}}_q}(L(D))$.

Proposition 2.88. (i) Let $D \in \text{Div}(\mathcal{X})$, then $\deg(D) < 0$ implies that $\ell(D) = 0$.

9

The vector space $L(D)$

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

4

Definition 2.84. Let $D \in \text{Div}(\mathcal{X})$, then we define the vector space $L(D)$ as:

$$L(D) := \{f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\} \mid (f) + D \text{ is effective}\} \cup \{0\}$$

and let $\ell(D) := \dim_{\overline{\mathbb{F}}_q}(L(D))$.

Proposition 2.88. (i) Let $D \in \text{Div}(\mathcal{X})$, then $\deg(D) < 0$ implies that $\ell(D) = 0$.

Proof: For all $f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\}$ we have:

$$\deg((f) + D) = \deg((f)) + \deg(D) = \deg(D) < 0$$

since $\deg((f)) = 0$.

9

The vector space $L(D)$

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

4

Definition 2.84. Let $D \in \text{Div}(\mathcal{X})$, then we define the vector space $L(D)$ as:

$$L(D) := \{f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\} \mid (f) + D \text{ is effective}\} \cup \{0\}$$

and let $\ell(D) := \dim_{\overline{\mathbb{F}}_q}(L(D))$.

Proposition 2.88. (i) Let $D \in \text{Div}(\mathcal{X})$, then $\deg(D) < 0$ implies that $\ell(D) = 0$.

Proof: For all $f \in \overline{\mathbb{F}}_q(\mathcal{X}) \setminus \{0\}$ we have:

$$\deg((f) + D) = \deg((f)) + \deg(D) = \deg(D) < 0$$

since $\deg((f)) = 0$. Meaning $L(D) = \{0\}$.

9

The Riemann-Roch Theorem

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

5

Let \mathcal{X} be a regular projective plane curve of genus g .

Theorem 2.91. Let $D \in \text{Div}(\mathcal{X})$, then for all canonical $W \in \text{Div}(\mathcal{X})$ we have

$$\ell(D) - \ell(W - D) = \deg(D) - g + 1$$

9

The Riemann-Roch Theorem

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

5

Let \mathcal{X} be a regular projective plane curve of genus g .

Theorem 2.91. Let $D \in \text{Div}(\mathcal{X})$, then for all canonical $W \in \text{Div}(\mathcal{X})$ we have

$$\ell(D) - \ell(W - D) = \deg(D) - g + 1$$

Corollary 2.92. If $\deg(D) > 2g - 2$, then $\ell(D) = \deg(D) - g + 1$.

9

The Riemann-Roch Theorem

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

5

Let \mathcal{X} be a regular projective plane curve of genus g .

Theorem 2.91. Let $D \in \text{Div}(\mathcal{X})$, then for all canonical $W \in \text{Div}(\mathcal{X})$ we have

$$\ell(D) - \ell(W - D) = \deg(D) - g + 1$$

Corollary 2.92. If $\deg(D) > 2g - 2$, then $\ell(D) = \deg(D) - g + 1$.

Proof:

► $\deg(W - D) < 0$.

The Riemann-Roch Theorem

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

5

Let \mathcal{X} be a regular projective plane curve of genus g .

Theorem 2.91. Let $D \in \text{Div}(\mathcal{X})$, then for all canonical $W \in \text{Div}(\mathcal{X})$ we have

$$\ell(D) - \ell(W - D) = \deg(D) - g + 1$$

Corollary 2.92. If $\deg(D) > 2g - 2$, then $\ell(D) = \deg(D) - g + 1$.

Proof:

- ▶ $\deg(W - D) < 0$.
- ▶ $\ell(W - D) = 0$ by Proposition 2.88 (i), combining this with Theorem 2.91 yields the result.

9



Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

6

Definition 3.3. Let $P_1, P_2, \dots, P_n \in \mathcal{X}$ be n distinct rational points.

9

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

6

Definition 3.3. Let $P_1, P_2, \dots, P_n \in \mathcal{X}$ be n distinct rational points.

- Let $D = \sum_i^n P_i$ and $G \in \text{Div}(\mathcal{X})$ such that $\text{supp}(D) \cap \text{supp}(G) = \emptyset$.

9

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

6

Definition 3.3. Let $P_1, P_2, \dots, P_n \in \mathcal{X}$ be n distinct rational points.

- ▶ Let $D = \sum_i^n P_i$ and $G \in \text{Div}(\mathcal{X})$ such that $\text{supp}(D) \cap \text{supp}(G) = \emptyset$.
- ▶ If $\mathcal{P} := (P_1, P_2, \dots, P_n)$, then $\mathcal{C}_{D,G} := \text{Ev}_{\mathcal{P}}(L(G))$ is called a *Goppa Code*.

9



Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

9

Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D)$.

9

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D).$

(ii) $d \geq n - \deg(G).$

9

Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D).$

(ii) $d \geq n - \deg(G).$

Proof:

- (i) Follows from the fact that $E_{v_P} \upharpoonright_{L(G)}$ is a surjective linear map, from $L(G)$ to $\mathcal{C}_{G,D}$.

9

Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D).$

(ii) $d \geq n - \deg(G).$

Proof:

(i) Follows from the fact that $E v_{\mathcal{P}} \upharpoonright_{L(G)}$ is a surjective linear map, from $L(G)$ to $\mathcal{C}_{G,D}$.

$$k = \dim_{\mathbb{F}_q}(\text{image}(E v_{\mathcal{P}})) = \ell(G) - \dim_{\mathbb{F}_q}(\ker(E v_{\mathcal{P}}))$$

9

Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D).$

(ii) $d \geq n - \deg(G).$

Proof:

(i) Follows from the fact that $E_{v_P} \upharpoonright_{L(G)}$ is a surjective linear map, from $L(G)$ to $\mathcal{C}_{G,D}$.

$$k = \dim_{\mathbb{F}_q}(\text{image}(E_{v_P})) = \ell(G) - \dim_{\mathbb{F}_q}(\ker(E_{v_P}))$$

But $\ker(E_{v_P}) = L(G - D)$ since:

Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D).$

(ii) $d \geq n - \deg(G).$

Proof:

(i) Follows from the fact that $E_{v_P} \upharpoonright_{L(G)}$ is a surjective linear map, from $L(G)$ to $\mathcal{C}_{G,D}$.

$$k = \dim_{\mathbb{F}_q}(\text{image}(E_{v_P})) = \ell(G) - \dim_{\mathbb{F}_q}(\ker(E_{v_P}))$$

But $\ker(E_{v_P}) = L(G - D)$ since:

- $f \in L(G) \text{ \& } E_{v_P}(f) = 0 \implies v_{P_i}(f) \geq 1 \text{ and hence } f \in L(G - D).$
As $(f) + G - D$ is effective.

9

Parameters of Goppa Codes

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

7

Theorem 3.5. If $\mathcal{C}_{D,G}$ is a $[n, k, d]_q$ code. Then:

(i) $k = \ell(G) - \ell(G - D).$

(ii) $d \geq n - \deg(G).$

Proof:

- (i) Follows from the fact that $E_{v_P} \upharpoonright_{L(G)}$ is a surjective linear map, from $L(G)$ to $\mathcal{C}_{G,D}$.

$$k = \dim_{\mathbb{F}_q}(\text{image}(E_{v_P})) = \ell(G) - \dim_{\mathbb{F}_q}(\ker(E_{v_P}))$$

But $\ker(E_{v_P}) = L(G - D)$ since:

- ▶ $f \in L(G)$ & $E_{v_P}(f) = 0 \implies v_{P_i}(f) \geq 1$ and hence $f \in L(G - D)$.
As $(f) + G - D$ is effective.
- ▶ $f \in L(G - D) \implies f(P) = 0 \forall P \in \text{supp}(D)$ as
 $\text{supp}(G) \cap \text{supp}(D) = \emptyset$. Meaning $f \in \ker(E_{v_P})$.



Parameters of Goppa Codes (Continued.)

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

8

Corollary 3.6 (i). If $\deg(G) < n$, then $k = \ell(G)$.

9



Parameters of Goppa Codes (Continued.)

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

8

Corollary 3.6 (i). If $\deg(G) < n$, then $k = \ell(G)$.

Proof: $\deg(G - D) < 0$, the rest follows by Proposition 2.88 (i) and Theorem 3.5 (i).

9

Parameters of Goppa Codes (Continued.)

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

8

Corollary 3.6 (i). If $\deg(G) < n$, then $k = \ell(G)$.

Proof: $\deg(G - D) < 0$, the rest follows by Proposition 2.88 (i) and Theorem 3.5 (i).

Remark 3.7.

- ▶ $k = \ell(G) \geq \deg(G) - g + 1$ by the Riemann-Roch Theorem 2.91.

9

Parameters of Goppa Codes (Continued.)

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

8

Corollary 3.6 (i). If $\deg(G) < n$, then $k = \ell(G)$.

Proof: $\deg(G - D) < 0$, the rest follows by Proposition 2.88 (i) and Theorem 3.5 (i).

Remark 3.7.

- ▶ $k = \ell(G) \geq \deg(G) - g + 1$ by the Riemann-Roch Theorem 2.91.
- ▶ Combining this with Theorem 3.5 (ii), we see:

$$d + k \geq \underbrace{(n - \deg(G))}_{\geq d} + \underbrace{(\deg(G) - g + 1)}_{\geq k} = n - g + 1 \quad (1)$$

9

Parameters of Goppa Codes (Continued.)

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

8

Corollary 3.6 (i). If $\deg(G) < n$, then $k = \ell(G)$.

Proof: $\deg(G - D) < 0$, the rest follows by Proposition 2.88 (i) and Theorem 3.5 (i).

Remark 3.7.

- ▶ $k = \ell(G) \geq \deg(G) - g + 1$ by the Riemann-Roch Theorem 2.91.
- ▶ Combining this with Theorem 3.5 (ii), we see:

$$d + k \geq \underbrace{(n - \deg(G))}_{\geq d} + \underbrace{(\deg(G) - g + 1)}_{\geq k} = n - g + 1 \quad (1)$$

- ▶ Combining this with the Singleton Bound we see that $n + 1 \geq d + k \geq n - g + 1$, and that $g = 0$ implies that $\mathcal{C}_{G,D}$ is an MDS code.

9



End

Martin Sig Nørbjerg

Coding Theory

Divisors

Goppa Codes

9

Thank you for listening.