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Convex Hull

1. Partially functioning code in Appendix
2. Theoretical Analysis:

The algorithm finding the convex hull given n points consists of two main parts: a *convex_hull* function and a *merge* function. *convex_hull* calls *merge* and itself recursively.

Procedure *convex_hull* splits the list of sorted points in half and then calls itself on each of these halves. Therefore, each recursive call takes half as long as its calling stack, introducing an $O(\log n)$ complexity.

The main task of *merge* is to find an upper and lower tangent between two polygons and to join them together. The upper and lower tangent methods start from the inner points and work their ways outwards until slope decreases. In a worst-case scenario, these will take $O(n/2)$ time or $O(1)$ if the first guess is correct. Running both upper and lower tangent is equivalent to:
 $O(n/2) + O(n/2) = O(n)$.

Bringing this together with regards to the master theorem, we can say:

$a = 2$, as we are creating 2 sub-tasks with each recursive call

$b = 2$, as the time halves with each recursive call

$d = 1$, as it takes each recursive call is $O(n)$.

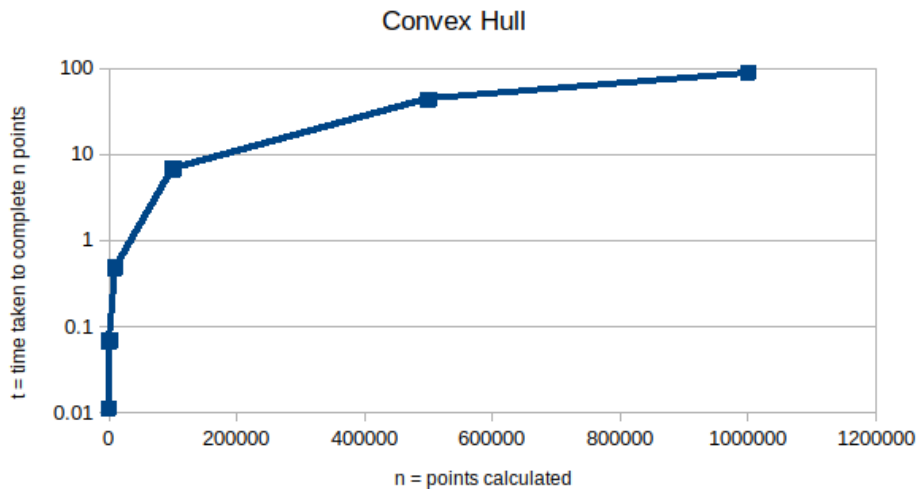
$T(n) = 2T(n/2) + O(n)$

$a/b^d = 2/2 = 1 \rightarrow O(n^d \log n) = O(n \log n)$

This intuitively makes sense, as we are running our $O(\log n)$ function n times.

3. Empirical Analysis:

| n | time(s) | mean time |
|---------|---------|-------------|
| 10 | 1E-24 | 1E+025 |
| 100 | 0.011 | 9090.909091 |
| 1000 | 0.068 | 14705.88235 |
| 10000 | 0.481 | 20790.02079 |
| 100000 | 6.718 | 14885.38255 |
| 500000 | 43.276 | 11553.74804 |
| 1000000 | 88.378 | 11315.03315 |



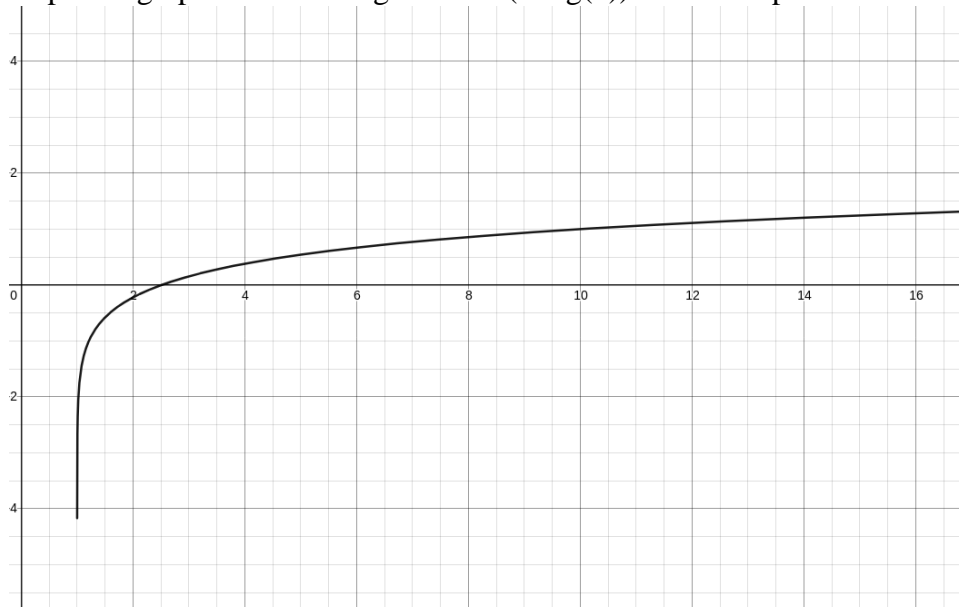
The order of growth which fits best is undoubtedly $O(n \log n)$. The reasoning behind the theoretical analysis would point to this being accurate.

By my estimate using empirical data, the constant of proportionality converged at about $8.838E-5$.

4. Observations

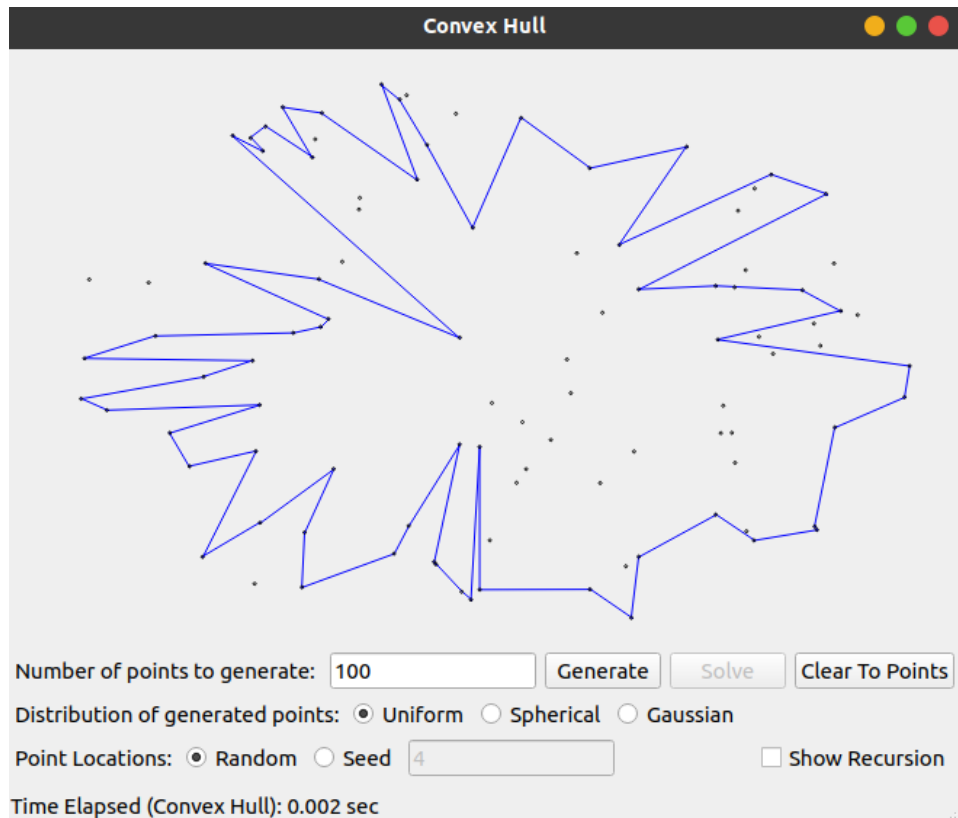
It is fascinating how closely we can predict the behavior and demands of an algorithm using asymptotic notation, the master theorem, among other tools. No major differences. The above

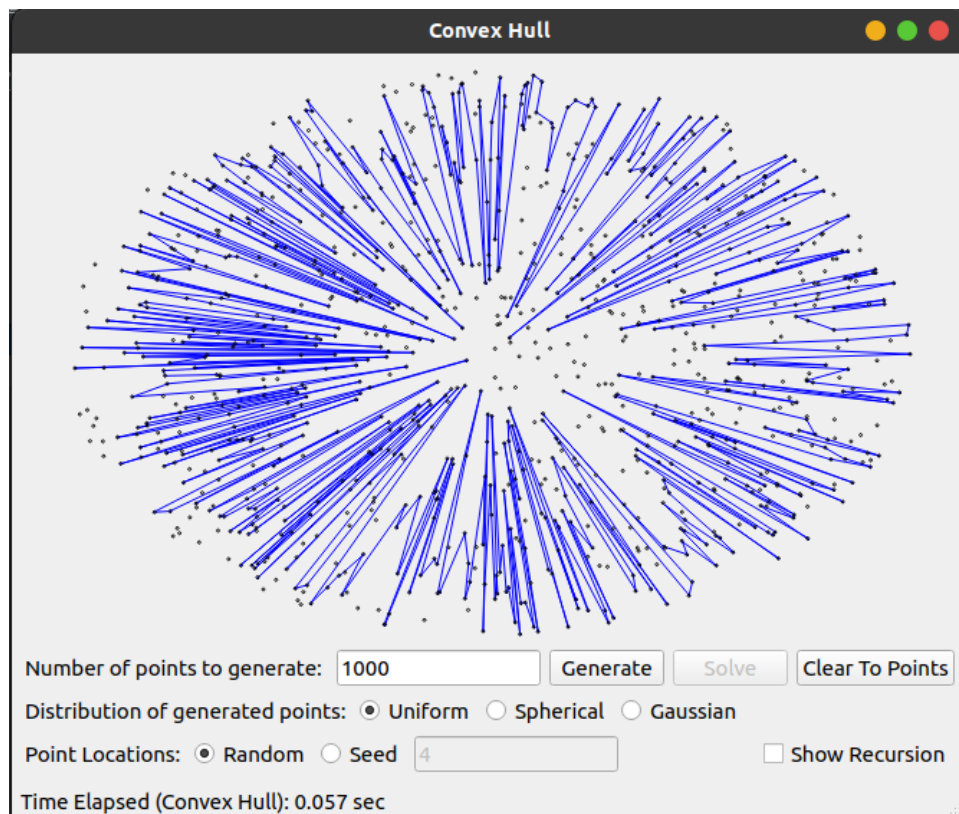
empirical graph looks like a generic $O(n \log(n))$ function I plotted in Desmos below:



5. Screenshots

These are obviously not completely functional, but hopefully the code demonstrates knowledge of divide-and-conquer concepts even if the implementation is poor.





Appendix A -- Python Code: