



PES UNIVERSITY, BENGALURU

Project Report

on

“16-Bit Booth’s Multiplier”

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

PES UNIVERSITY

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16-Bit Booth's Multiplier

Problem Description

Booth's multiplication algorithm is a multiplication algorithm that multiplies two signed binary numbers in two's complement notation.

Booth's algorithm examines adjacent pairs of bits of the 'N'-bit multiplier Y in signed two's complement representation, including an implicit bit below the least significant bit, $y_{-1} = 0$. For each bit y_i for i running from 0 to $N - 1$, the bits y_i and y_{i-1} are considered.

Where these two bits are equal, the product accumulator P is left unchanged.

Where $y_i = 0$ and $y_{i-1} = 1$, the multiplicand times 2^i is added to P; and where $y_i = 1$ and $y_{i-1} = 0$, the multiplicand times 2^i is subtracted from P. The final value of P is the signed product.

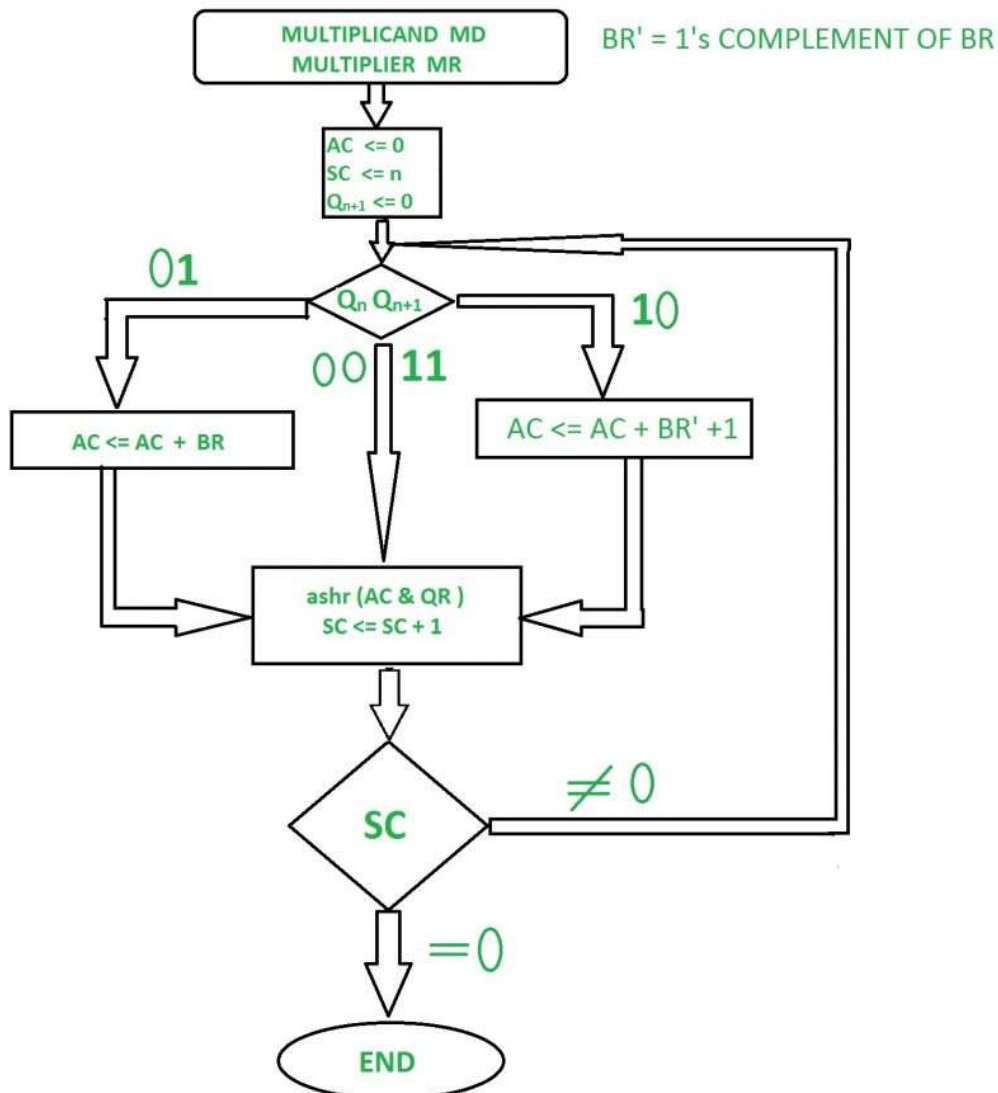
Implementation

Booth's algorithm can be implemented by repeatedly adding (with ordinary unsigned binary addition) one of two predetermined values A and S to a product P, then performing a rightward arithmetic shift on P. Let m and r be the multiplicand and multiplier, respectively; and let x and y represent the number of bits in m and r.

1. Determine the values of A and S, and the initial value of P. All of these numbers should have a length equal to $(x + y + 1)$.
 - a. A: Fill the most significant (leftmost) bits with the value of m. Fill the remaining $(y + 1)$ bits with zeros.
 - b. S: Fill the most significant bits with the value of $(-m)$ in two's complement notation. Fill the remaining $(y + 1)$ bits with zeros.
 - c. P: Fill the most significant x bits with zeros. To the right of this, append the value of r. Fill the least significant (rightmost) bit with a zero.

2. Determine the two least significant (rightmost) bits of P.
 - a. If they are 01, find the value of $P + A$. Ignore any overflow.
 - b. If they are 10, find the value of $P + S$. Ignore any overflow.
 - c. If they are 00, do nothing. Use P directly in the next step.
 - d. If they are 11, do nothing. Use P directly in the next step.
3. Arithmetically shift the value obtained in the 2nd step by a single place to the right. Let P now equal this new value.
4. Repeat steps 2 and 3 until they have been done y times.
5. Drop the least significant (rightmost) bit from P.
This is the product of m and r.

Flowchart Diagram





Example

Let A: 3 and B: 17

Multiplicand -	
Decimal:	3
Binary:	00000011
Multiplier -	
Decimal:	17
Binary:	00010001
Two's Complement:	11101111
Steps -	
Starting Out:	0000000000000011
Subtract:	1110111100000011
Shift:	1111011110000001
Shift:	1111101111000000
Add:	0000110011000000
Shift:	0000011001100000
Shift:	0000001100110000



Shift:	0000000110011000
Shift:	0000000011001100
Shift:	0000000001100110
Shift:	0000000000110011
Final Product (Binary):	0000000000110011
Final Product (Decimal):	51

Implementation of a Booth's algorithm uses various sub-modules, as described below.

Worst and Ideal Case

The **worst case of an implementation** using Booth's algorithm is when pairs of 01s or 10s occur very frequently in the multiplier.

Modules and Sub Modules

Implementation of a Booth's algorithm uses various sub-modules, as described below.

1. **boothmul()**: This module is the main module which uses the help of other sub-modules or counterparts to solve our problem.

This module takes in two **8-bit signed** inputs, which are our multiplicand and multiplier. It has one **16-bit signed** output. Inside the module, we have **eight** 8-bit signed wires hold the value of the changed bits after shifting so that we can manipulate them later.

2. **booth_substep()**: This sub-module does the main operation of either adding/subtracting or just shifting the bits according to the last two positions.

This module takes in an 8-bit signed **accumulator**, 8-bit signed **multiplier**, the last bit of the accumulator, 8-bit signed **multiplicand**, and the output consists of two 8-bit signed registers



containing first 8 and last 8 bits of the product, and cq0 is the changed q0 after the shift operation.

3. **Adder()**: This sub-module adds two 8-bit register values, and gives out their sum. This uses a library module of **fa** which is nothing but a simple full adder.
4. **Subtractor()**: This sub-module subtracts two 8-bit register values, and gives out their difference. This uses a library module of **invert** to invert each bit separately, and then uses **fa** which is nothing but a simple full adder as described above.
5. **Lib.v**:
 - a. **invert(output ib,input b);**
 - b. **and2 (input wire i0, i1, output wire o);**
 - c. **or2 (input wire i0, i1, output wire o);**
 - d. **xor2 (input wire i0, i1, output wire o);**
 - e. **nand2 (input wire i0, i1, output wire o);**
 - f. **nor2 (input wire i0, i1, output wire o);**
 - g. **xnor2 (input wire i0, i1, output wire o);**
 - h. **and3 (input wire i0, i1, i2, output wire o);**
 - i. **or3 (input wire i0, i1, i2, output wire o);**
 - j. **nor3 (input wire i0, i1, i2, output wire o);**
 - k. **nand3 (input wire i0, i1, i2, output wire o);**
 - l. **xor3 (input wire i0, i1, i2, output wire o);**
 - m. **xnor3 (input wire i0, i1, i2, output wire o);**
 - n. **fa (input wire i0, i1, cin, output wire sum, cout);**

Apart from using the above modules, we use a testbench to supply the initial values.

```
a = 8'b11110000;
```

```
b = 8'b11110000;
```

```
#10
```

```
a = 8'b10010101;
```

```
b = 8'b100000;
```

And so on and so forth.

Final Result on Screen

The results on the screen are printed like this:

VCD info: dumpfile tb_boothsalgo.vcd opened for output.

0	-16	X	-16	=	256
10	-107	X	32	=	-3424
20	7	X	0	=	0
30	1	X	1	=	1
40	60	X	5	=	300
50	-86	X	35	=	-3010
60	17	X	28	=	476
70	8	X	-65	=	-520
