

## Complexity Order

$\log n < \sqrt{n} < n < n \log n < n \sqrt{n} < n^2 < \dots$

For  $i \leq n$  iteration, TC  $\Rightarrow O(\sqrt{n})$

For  $i \leq n/2$  iteration, TC  $\Rightarrow O(\log n)$

## Range of Long-

Min:  $-2^{63} = -9 \times (10^{18})$

Max:  $2^{63}-1 = 9 \times (10^{18})$

## Arithmetic Progression

$\text{sum} = N/2 \times (2 \times a_0 + (n-1) \times d)$

## Geometric Progression

$\text{sum} = (a \times (r^n - 1)) / (r - 1)$

## Prefix Array

$\text{Pf}[i] = \text{pf}[R] - \text{pf}[L-1] \quad , L > 0$   
 $\quad = \text{pf}[R] \quad , L = 0$

## Total Subarrays

$n \times (n+1) / 2$


## Number of subarrays of size k

$(n - k + 1)$

## Total the sum of all subarrays

$\text{Sum} += (i+1) \times (n-i) \times \text{arr}[i]$

## For Matrix[N][N], boundary in clockwise:

N-1 

N-1 

N-1 

N-1 

## Rotate a matrix by 90° clockwise

Transpose the matrix + reverse of each row

## Sliding windows

Sum = sum -a[s-1]+a[e]

### Joseph Problem

If number  $2^n$ , then 1

Otherwise  $2^{n+1}$ ,  $n = \text{number} - \text{nearest power of 2 less than the number}$

## Bit Manipulation

### k-based to decimal/2 based number

$a[n-1]*k^0 + a[n-2]*k^1 + a[n-3]*k^2 + \dots + a[0]*k^{(n-1)}$

### Decimal to k-based number

Until  $n=0$

```
{ n = n//k  
  rem = n%k  
  str = rem+str }
```

## Properties

1>  $A \& 1 = 0$ , if A is even  
     $= 1$ , if A is odd

2>  $A \& 0 = 0$

3>  $A \& A = A$

4>  $A | 0 = A$

5>  $A | 1 = A$

6>  $A \wedge 0 = A$

7>  $A \wedge A = 0$

8>  $A \wedge 1 = A[0 \rightarrow n-2 \text{ bits}] + 0$  if odd  
     $A[0 \rightarrow n-2 \text{ bits}] + 1$  if even

9> Communicative properties

$A \& B = B \& A$

$A | B = B | A$

$A \wedge B = B \wedge A$

10> Associative Properties

$(A \& B) \& C = A \& (B \& C)$

$(A | B) | C = A | (B | C)$

$(A \wedge B) \wedge C = A \wedge (B \wedge C)$

11> 1 Left Shift  $\Rightarrow$  Multiplies by 2

12> 1 Right Shift  $\Rightarrow$  Division by 2

13>  $A \ll n = A \cdot (2^n)$  and  $A \gg n = A / (2^n)$

14>  $1 \ll n \Rightarrow 2^n$

15>  $N \mid 1 \ll i \Rightarrow N$  ,if ith bit set  
           $N + 1 \ll i$  ,if ith bit is unset

16>  $N \wedge 1 \ll i \Rightarrow$  Flips or toggle ith bit

17>  $N \& 1 \ll i \Rightarrow 1 \ll i$  ,if ith bit is set  
          0 ,if ith bit is unset

### Check if the ith bit is set or not

1>  $N \& (1 \ll i) == 1 \ll i$

2>  $N \mid (1 \ll i) == N$

3>  $N \wedge (1 \ll i) == N - (1 \ll i)$

4>  $(N \gg i) \& 1 == 1$

Then do  $\Rightarrow N \wedge (1 \ll i)$  OR  $N + (1 \ll i)$

### Negative no.s

Step1> 1's complement  $A = A'$

Step2> 2's complement  $(A' + 1)$

MSB(Most Significant Bit) - nth bit in 8-digit binary number

MSB - 1 then -Ve

0 then +Ve

## Modular Arithmetic

1>  $(a+b)\%p = (a\%p + b\%p)\%p$

2>  $(a*b)\%p = (a\%p * b\%p)\%p$

3>  $(a-b)\%p = ((a\%p - b\%p)\%p + p)\%p$

4>  $(a\%p)\%p = a\%p$

5>  $(a\%p * b) = (a*b)\%p$

**Fermat Theorem:**  $a^{(p-2)\%p} = a^{(-1)\%p}$

## **String**

'A' - 65

'Z' - 90

'a' - 97

'Z' - 122

'0' - 48

'9' - 57

### **Toggling:**

Caps->small: s+32

small->Caps: s-32