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**1.**

$$\mathbf{a.} \quad T(0) = 0$$

$$T(n+1) = T(n) + 5$$

$$T(n) = T(n-1) + 5$$

$$= T(n-1) + 5$$

$$= (T(n-2) + 5) + 5$$

$$= ((T(n-3) + 5) + 5) + 5$$

$$= T(n-k) + 5k$$

Let  $k = n$ :

$$T(n) = 5n + T(0) = 5n + 0 = 5n$$

$$\mathbf{T(n) = 5n}$$

$$\mathbf{b.} \quad T(0) = 0$$

$$T(n+1) = n + T(n)$$

$$T(n) = (n-1) + T(n-1)$$

$$= (n-1) + (n-2) + T(n-2)$$

$$= (n-1) + (n-2) + T(n-2)$$

$$= (n-1) + (n-2) + \dots + (n-k+1) + (n-k) + T(n-k)$$

Let  $k = n$ :

$$T(n) = (n-1) + (n-2) + \dots + (n-n+1) + (n-n) + T(n-n)$$

$$T(n) = (n-1) + (n-2) + \dots + 1 + 0 + T(0)$$

$$T(n) = \sum_{i=0}^{n-1} i + T(0)$$

$$T(n) = \frac{(n-1)(n-1+1)}{2}$$

$$\mathbf{T(n)} = \frac{(n-1)(n)}{2}$$

**2.**

$$\mathbf{a.} \quad T(0) = 1$$

$$T(n+1) = 2 \cdot T(n)$$

$$T(n) = 2 \cdot T(n-1)$$

$$= 2 \cdot 2 \cdot T(n-2)$$

$$= 2 \cdot 2 \cdot 2 \cdot T(n-3)$$

$$= 2^k \cdot T(n-k)$$

Let  $k = n$

$$T(n) = 2^n \cdot T(n-n)$$

$$T(n) = 2^n \cdot T(0)$$

$$T(n) = 2^n \cdot 1$$

$$\mathbf{T(n)} = 2^n$$

$$\mathbf{b.} \quad T(0) = 1$$

$$T(n+1) = 2^{n+1} + T(n)$$

$$T(n) = 2^n + T(n-1)$$

$$= 2^n + 2^{n-1} + T(n-2)$$

$$= 2^n + 2^{n-1} + 2^{n-2} + T(n-3)$$

$$= 2^n + 2^{n-1} + 2^{n-2} + \dots + 2^{n-k+1} + T(n-k)$$

Let  $k = n$

$$T(n) = 2^n + 2^{n-1} + 2^{n-2} + \dots + 2^{n-n+1} + T(n-n)$$

$$T(n) = 2^n + 2^{n-1} + 2^{n-2} + \dots + 2^1 + T(0)$$

$$T(n) = \sum_{i=0}^n 2^i + T(0) - 2^0$$

$$T(n) = (2^{n+1} - 1) + 1 - 1$$

$$\mathbf{T(n)} = 2^{n+1} - 1$$

**3.**

$$\mathbf{a.} \quad T(1) = 1$$

$$T(n) = n + T(n/2) \quad \text{Assume } n \text{ has the form } n = 2^m$$

$$\begin{aligned}
T(2^m) &= 2^m + T\left(\frac{2^m}{2}\right) \\
&= 2^m + T(2^{m-1}) \\
&= 2^m + 2^{m-1} + T\left(\frac{2^{m-1}}{2}\right) \\
&= 2^m + 2^{m-1} + T(2^{m-2}) \\
&= 2^m + 2^{m-1} + 2^{m-2} + T\left(\frac{2^{m-2}}{2}\right) \\
&= 2^m + 2^{m-1} + 2^{m-2} + T(2^{m-3}) \\
&= 2^m + 2^{m-1} + 2^{m-2} + \dots + 2^{m-k+1} + T(2^{m-k})
\end{aligned}$$

Let  $k = m$ :

$$\begin{aligned}
T(2^m) &= 2^m + 2^{m-1} + 2^{m-2} + \dots + 2^{m-m+1} + T(2^{m-m}) \\
&= 2^m + 2^{m-1} + 2^{m-2} + \dots + 2^1 + T(2^0) \\
&= \left(\sum_{i=0}^m 2^i\right) - 2^0 + T(0) \\
&= (2^{m+1} - 1) - 1 + 1 \\
&= 2^{m+1} - 1
\end{aligned}$$

$$\Rightarrow n = 2^m \Rightarrow \lg(n) = m$$

$$\Rightarrow 2^{m+1} \Rightarrow (2^m)2 \Rightarrow (2^{\lg(n)})2 \Rightarrow 2n$$

$$T(n) = 2n - 1$$

$$\mathbf{b.} \quad T(1) = 1$$

$$T(n) = 1 + T(n/3) \quad \text{Assume } n \text{ has the form } n =$$

$$3^m$$

$$T(3^m) = 1 + T\left(\frac{3^m}{3}\right)$$

$$= 1 + T(3^{m-1})$$

$$= 1 + 1 + T\left(\frac{3^{m-1}}{3}\right)$$

$$= 1 + 1 + T(3^{m-2})$$

$$= k + T(3^{m-k})$$

Let  $k = m$ :

$$T(3^m) = m + T(3^{m-m})$$

$$= m + T(3^0)$$

$$= m + T(1)$$

$$= m + 1$$

$$\Rightarrow n = 3^m$$

$$\Rightarrow \log_3(n) = m$$

$$T(n) = \log_3(n) + 1$$

**6.**

**f.**

**i.**

*Theorem: For any  $n \in \mathbb{N}$ ,  $T_F(n) = F_{n+1} - 1$*

*Proof: By the strong form of induction*

*Observe that when  $n = 0$  we have:*

$$\begin{aligned} T_F(0) &= 0 = F_{0+1} - 1 \\ &= 1 - 1 \\ &= 0 \end{aligned}$$

*Observe that when  $n = 1$  we have:*

$$\begin{aligned} T_F(1) &= 0 = F_2 - 1 \\ &= 1 - 1 \\ &= 0 \end{aligned}$$

Assume  $T_F(k) = F_{k+1} - 1$  if  $0 \leq k < n$

$$\begin{aligned} T_F(k+1) &= 1 + T_F(k) + T_F(k-1) \\ &= 1 + F_{k+1} - 1 + F_k - 1 \\ &= F_{k+1} + F_k - 1 \end{aligned}$$

$$T_F(k+1) = F_{k+2} - 1$$

**ii.**

$$= \lim_{n \rightarrow \infty} \frac{\frac{1}{\sqrt{5}}(\varphi^{n+1} - \widehat{\varphi}^{n+1}) - 1}{\varphi^n}$$

$$\begin{aligned}
&= \lim_{n \rightarrow \infty} \frac{\frac{1}{\sqrt{5}}(\varphi^{n+1} - \widehat{\varphi}^{n+1} - \sqrt{5})}{\varphi^n} \\
&= \frac{1}{\sqrt{5}} \lim_{n \rightarrow \infty} \frac{(\varphi^{n+1} - \widehat{\varphi}^{n+1} - \sqrt{5})}{\varphi^n} \\
&= \frac{1}{\sqrt{5}} \left[ \left( \lim_{n \rightarrow \infty} \frac{\varphi^{n+1}}{\varphi^n} \right) - \left( \lim_{n \rightarrow \infty} \frac{\widehat{\varphi}^{n+1}}{\varphi^n} \right) - \left( \lim_{n \rightarrow \infty} \frac{\sqrt{5}}{\varphi^n} \right) \right] \\
&= \frac{1}{\sqrt{5}} (\varphi - 0 - 0) \\
&= \frac{\varphi}{\sqrt{5}}
\end{aligned}$$

**Since**  $\lim_{n \rightarrow \infty} \frac{T_F(n)}{\varphi^n} = c$  and  $c \neq 0$ , therefore  $T_F(n) \in \Theta(\varphi^n)$

**g.**

$$T_f(0) = 0$$

$$T_f(1) = 0$$

$$T_f(n + 1) = 1 + T_f(n)$$

$$T_f(n) = 1 + T_f(n - 1)$$

$$= 1 + 1 + T_f(n - 2)$$

$$= 1 + 1 + 1 + T_f(n - 3)$$

$$= k + T_f(n - k)$$

Let  $k = n$ :

$$T_f(n) = n + T_f(n - n) = n + T_f(0) = n + 0 = n$$

$$T_f(n) = n$$

**h.**

Theorem: For any  $n \in \mathbb{N}$ ,  $L^n(a, b) = (f(n;a,b), f(n+1;a,b))$ .

*Proof: By weak induction*

Observe when  $n = 0$  we have

$$\begin{aligned} L^0(a, b) &= (a,b) \\ &= (f(0;a,b), f(1;a,b)) \end{aligned}$$

$$\text{Assume } L^k(a, b) = (f(k;a,b), f(k+1;a,b))$$

$$\begin{aligned} L^{k+1}(a, b) &= L(L^k(a, b)) \\ &= L(f(k; a, b), f(k + 1; a, b)) \\ &= (f(k + 1; a, b), f(k; a, b) + f(k + 1; a, b)) \\ &= (f(k + 1; a, b), f(k + 2; a, b)) \end{aligned}$$

**i.**

Consider the vector  $\vec{x} = (x_1, x_2)$  and  $\vec{y} = (y_1, y_2)$

**Prove**  $L(\vec{x} + \vec{y}) = L(\vec{x}) + L(\vec{y})$  :

$$L(\vec{x} + \vec{y}) = L(x_1 + y_1, x_2 + y_2) = (x_2 + y_2, x_1 + y_1 + x_2 + y_2)$$

$$\begin{aligned} L(\vec{x}) + L(\vec{y}) &= L(x_1, x_2) + L(y_1, y_2) = (x_2, x_1 + x_2) + (y_2, y_1 + y_2) \\ &= (x_2 + y_2, x_1 + y_1 + x_2 + y_2) \end{aligned}$$



Since  $L(\vec{x} + \vec{y}) = (x_2 + y_2, x_1 + y_1 + x_2 + y_2)$  and  $L(\vec{x}) + L(\vec{y}) = (x_2 + y_2, x_1 + y_1 + x_2 + y_2)$ , therefore,  $L(\vec{x} + \vec{y}) = L(\vec{x}) + L(\vec{y})$ .

**Prove  $L(c \cdot \vec{x}) = c \cdot L(\vec{x})$ :**

$$L(c \cdot \vec{x}) = L(c \cdot (x_1, x_2)) = L(cx_1, cx_2) = (cx_2, cx_1 + cx_2)$$

$$c \cdot L(\vec{x}) = c \cdot L(x_1, x_2) = c \cdot (x_2, x_1 + x_2) = (cx_2, cx_1 + cx_2)$$

Since  $L(c \cdot \vec{x}) = (cx_2, cx_1 + cx_2)$  and  $c \cdot L(\vec{x}) = (cx_2, cx_1 + cx_2)$ , therefore,

$$L(c \cdot \vec{x}) = c \cdot L(\vec{x}).$$

Since  $L(\vec{x} + \vec{y}) = L(\vec{x}) + L(\vec{y})$  and  $L(c \cdot \vec{x}) = c \cdot L(\vec{x})$ , the function **L** is linear.

**j.**

$$\begin{aligned} L(a, b) &= L(a \cdot \vec{e}_1 + b \cdot \vec{e}_2) \\ &= L(a \cdot \vec{e}_1) + L(b \cdot \vec{e}_2) \\ &= a \cdot L(\vec{e}_1) + b \cdot L(\vec{e}_2) \\ &= a \cdot (0, 1) + b \cdot (1, 1) \\ &= (0, a) + (b, b) \end{aligned}$$

$$= (b, a+b)$$

$$= M \cdot (a,b) \quad \text{for some } 2 \times 2 \text{ matrix } M$$

Since  $M \cdot (a,b) = (b, a+b)$ , we can represent the matrix  $M$  in the form of system of equations:

$$aM_{11} + bM_{12} = b$$

$$aM_{21} + bM_{22} = a + b$$

Only  $m_{21} = m_{22} = 1$  would satisfy the second equation.

Only  $m_{11} = 0$  and  $m_{12} = 1$  would satisfy the first equation.

Therefore the matrix representation of  $L$  would be the matrix:

$$\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$$

**n.**  $\text{fibPow} \in \Theta(\log(n))$ , since the function called upon  $\text{matPow}()$  which requires in the worst case  $\Theta(\log(n))$  recursive calls to compute a result.

**o.**

**a.**  $\text{fib}$  is not a pseudo polynomial time because  $\Theta(\varphi^n)$  can not be bounded by a polynomial function since it is exponential.

- b.** fibIt is a pseudo polynomial time since  $\Theta(n)$  can be bounded by another polynomial function.
- c.** fibPow is a pseudo polynomial time since  $\Theta(\log(n))$  can be bounded by another polynomial function.