

Lambda calculus

Functional models of computation

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History

- 1928 - Hilbert's *Entscheidungsproblem*¹
 - Is there an *algorithm* for deciding whether a proposition in first-order logic is true or false?
- Replacement for set theory as foundation of mathematics
 - 1930 - Combinatory logic (*Curry, Schönfinkel*)
 - 1932 - λ -calculus (*Church*)
 - 1935 - Kleene-Rosser paradox
- Effective computability
 - 1935 - Untyped λ -calculus (*Church, Kleene, Rosser*)
 - 1936 - Turing machine
 - 1936 - Church-Turing thesis
- 1936 - Undecidability of first-order logic
 - Halting problem of Turing machine
 - Equivalence of λ -terms

¹German for “decision problem”

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David Hilbert

- [Haskell Curry](#)
- Wilhelm Ackermann
- John von Neumann
- Ernst Zermelo
- ...

Alonzo Church

- [Stephen Cole Kleene](#)
- [J. Barkley Rosser](#)
- Alan Turing
- Dana Scott
- Michael O. Rabin
- ...

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Grammar

$$term ::= \underbrace{var}_{\text{Variable}} \mid \underbrace{term\ term}_{\text{Application}} \mid \underbrace{\lambda var. term}_{\text{Abstraction}}$$

Examples

$$\lambda x. x \quad (\lambda x. xx)(\lambda y. yy) \quad \lambda f. \lambda x. f(fx)$$

Conventions

- Application is left associative
 $abc = (ab)c$
- Abstraction is right associative
 $\lambda x. \lambda y. x = \lambda x. (\lambda y. x)$
- Consecutive abstractions can be combined
 $\lambda x. \lambda y. x = \lambda xy. x$

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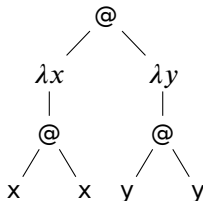
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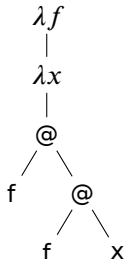
Tree representation

$\lambda x. x$
 λx
|
 x

$(\lambda x. xx)(\lambda y. yy)$



$\lambda f. \lambda x. f(fx)$



Free and bound variables

Substitution

α -equivalence

β -conversion

β -reduction

β -abstraction

Normal order reduction

First Church-Rosser theorem

Second Church-Rosser theorem

Normal order reduction

Fixed-point combinator

Curry's Y -combinator

$$Y = \lambda f. (\lambda x. f(xx)) (\lambda x. f(xx))$$

Turing's Θ -combinator

$$\Theta = (\lambda xy. x(xxy)) (\lambda xy. x(xxy))$$

Undecidability

Church numerals

Relation to folds

Algebraic data types

Predecessor

Q&A