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Weekday: Week 3- Day 13

Answer to the Q. No.1:

Algorithm initResult(G)  
S<-new Sequence

Algorithm preComponentVisit(G,v)  
S.insertLast(v)

Algorithm result(G)  
return S

Answer to the Q. No.2(a):

Algorithm BFS(G)  
Input graph G  
Output labeling of the edges and partition  
of the vertices of G

```
initResult( G )
for all u in G.vertices( )
    preInitVertex(u)
    setLabel(u, UNEXPLORED)
for all e in G.edges()
    preInitEdge(e)
    setLabel(e, UNEXPLORED)
for all v in G.vertices()
    if getLabel(v) = UNEXPLORED
        preComponentVisit(G, v)
        BFS(G, v)
        postComponentVisit(G,v)

result(G)
```

Algorithm BFS(G, s)

```
L <- new empty List
L.insertLast(s)
setLabel(s, VISITED)

while !L.isEmpty()
    v <- L.remove (L.first())
    vertexVisit(v)
    for all e in G.incidentEdges(v)
        if getLabel(e) = UNEXPLORED then
            w <- opposite(v,e)
            if getLabel(w) = UNEXPLORED then
                preDiscoveryTraversal(G, v, e, w)
                setLabel(e, DISCOVERY)
                setLabel(w, VISITED)
                L.insertLast(w)
                postDiscoveryTraversal(G, v, e, w)
            else
                setLabel(e, CROSS)
                crossTraversal(G,v,e,w)
```

finishVertexVisit(G,s)

Answer to the Q. No.2(b):

**Algorithm** findPathBFS(G,u,v)  
S<-new Sequence  
path<-null  
pathFound<-false  
z<-v  
for all n in G.vertices() do  
 setLabel(p,UNEXPLORED)  
for all l in G.edges() do

```

        setLabel(l, UNEXPLORED)
    BFS(G, u)

```

```

if !pathFound = false then
    return NO_SUCH_PATH
else
    return path

```

**Algorithm** vertexVisit(v)

```

if v=s then

```

```

    v.setPath(v) {path is a property of node}

```

**Algorithm** preDiscoveryTraversal(G, v, e, w)

```

if !pathFound then
    w.setPath(v.getPath()+e+w)

```

**Algorithm** postDiscoveryTraversal(G, v, e, w)

```

if z=w then
    pathFound=true
    path<-w.getPath()

```

Answer to the Q. No.2(c):

**Algorithm** vertexVisit(s)

```

if v=s then
    setParent(s, null, null)

```

**Algorithm** preDiscoveryTraversal(G, v, e, w)

```

if !cycleFound=false
    setParent(w, v, e) {set parent and related edge}

```

**Algorithm** crossTraversal(G, v, e, w)

```

cycleFound=true
S<-new Stack()
Q<-new Queue()
while getParent(v)!=getParent(w) then
    S.push(v)
    S.push(getParentConnectedEdge(v))
    Q.enqueue(w)
    Q.enqueue(getParentConnectedEdge(w))

```

```

    v<- getParent(v)
    w<-getParent(w)
    cyclePath<-new Sequence()
    cyclePath.insertLast(getParent(v))

```

```

while !S.isEmpty() then

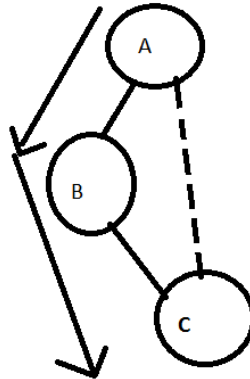
```

```

cyclePath.insertLast(s.pop())
cyclePath.insertLast(e)
while !Q.isEmpty() then
cyclePath.insertLast(Q.dequeue())

```

Answer to the Q. No.2(d):



If we see above example, where DFS is used to traverse a graph, To find path from node A to C,

It will traverse from node A to B, then from B to C. Even though shortest path is from A to C. So it's not guaranteed that template version or non-template version of DFS algorithm will find minimum edges path between two vertices.

Answer to the Q. No.3:

```

Algorithm findShortestPath(G,s,d)
z<-d {z is a subclass variable}
initResult( G )
DijkstraDistances(G, s)
result(G)

```

```

Algorithm DijkstraDistances(G, s)
Q <- new heap-based priority queue
for all v <- G.vertices()
    preInitVertex(u)
    if v = s
        setDistance(v, 0)
    else
        setDistance(v, INFINITY)
    Q.insertItem(getDistance(v), v)
while !Q.isEmpty()
    u <- Q.removeMin()
    vertexVisit(v)

```

```

for all  $e$  in  $G.incidentEdges(u)$ 
{ relax edge  $e$  }
     $z \leftarrow G.opposite(u,e)$ 
    preDiscoveryTraversal( $G, u, e, z$ )
     $r \leftarrow getDistance(u) + weight(e)$ 
    if  $r < getDistance(z)$ 
         $setDistance(z,r)$ 
        beforeDistanceChange( $G,u,e,z$ )
         $Q.replaceKey(z,r)$ 
        afterDistanceChange( $G,u,e,z$ )
    postDiscoveryTraversal( $G, u, e, z$ )

```

**Algorithm** *afterDistanceChange*( $G,u,e,z$ )  
 $setParent(z,u,e)$  {set parent and related edge}

**Algorithm** *result*( $G$ )  
 $S \leftarrow new\ Sequence()$   
while  $getParent(d) \neq s$  do  
 $S.insertLast(d)$   
 $e \leftarrow getParentConnectedEdge(d)$   
 $S.insertLast(e)$   
 $d \leftarrow getParent(d)$   
 $S.insertLast(d)$   
 $S.insertLast(getParentConnectedEdge(d))$   
 $S.insertLast(s)$

Answer to the Q. No.4:

**Algorithm** *initResult*(  $G$  )  
 $connectedComponent \leftarrow -1$  {  $connectedComponent$  is a subclass variable }

**Algorithm** *postComponentVisit*( $G,v$ )  
 $connectedComponent \leftarrow connectedComponent + 1$

**Algorithm** *startVertexVisit*( $s$ )  
 $setLabel(s, connectedComponent)$