CS 33: Introduction to Computer Organization

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Office Hours: Friday, 9:30-11:30AM

Outline

- Locking mechanism (quick recap)
- Exceptions
- Linking
- Virtual Memory
- Worksheet Problems

Locking Mechanism

```
What happens if
Thread A
                                      Thread B
                   threads are at
                 this point at the
lock (lock a);
                                      lock (lock b);
                    same time?
a += 5;
                                      b += 5;
lock (lock b);
                                      lock (lock a);
b += 7;
                                      a += 7;
a += b;
                                      a += b;
unlock (lock b);
                                      unlock (lock a);
                                      b += 11;
a += 11;
unlock (lock a);
                                      unlock (lock b);
```

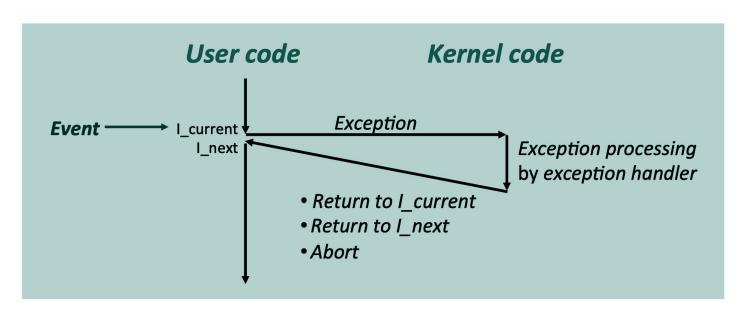
If there is an issue, how can we solve it?

Solution to Locking Problem

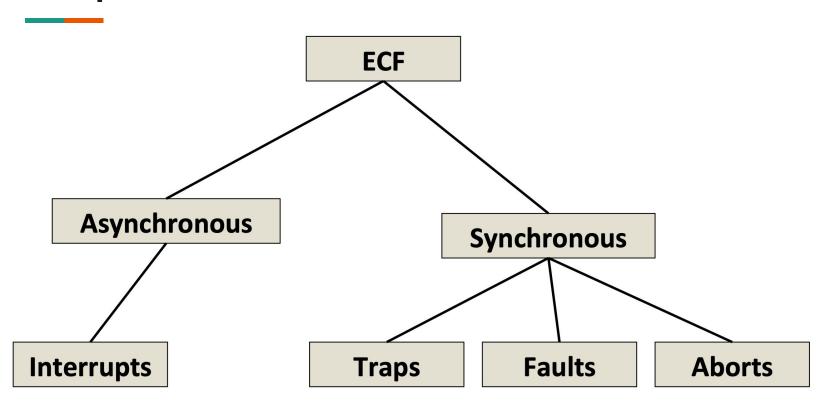
```
Threads must lock
                lock_a before lock_b
Thread A
                                       Thread B
lock (lock a);
                                       lock (lock a);
a += 5;
                                       lock (lock b);
lock (lock b);
                                       b += 5;
b += 7;
                                       a += 7;
a += b;
                                       a += b;
unlock (lock b);
                                       unlock (lock a);
a += 11;
                                       b += 11;
unlock (lock a);
                                       unlock (lock b);
```

Exception

An *exception* is a transfer of control to the OS *kernel* in response to some *event* (i.e., change in processor state)

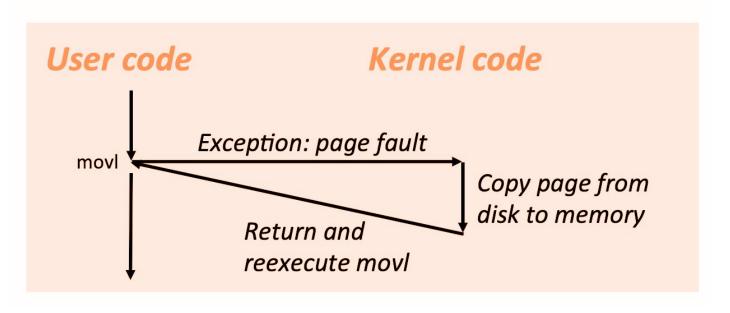


Exception

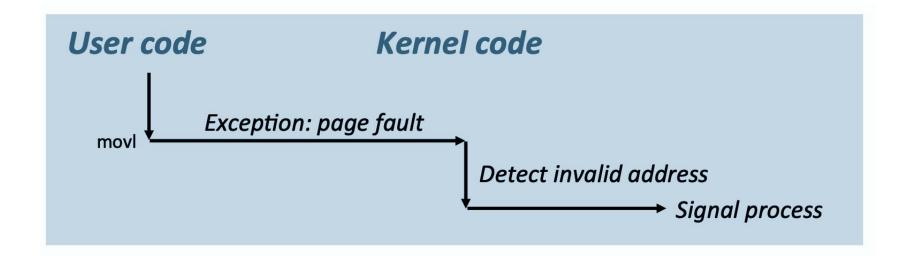


Fault Example: Page Fault

What if there was a page fault?



Fault Example: Invalid Memory Reference



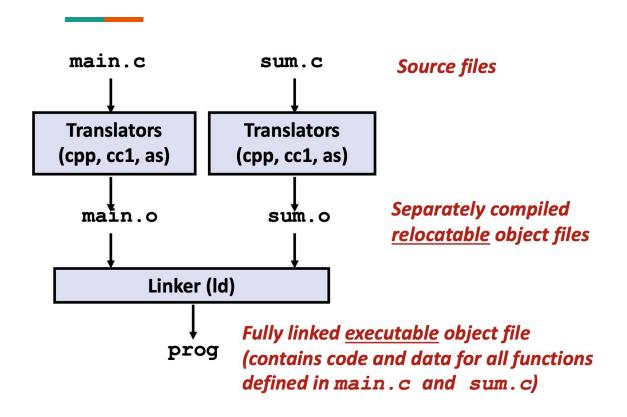
They can also lead to **process termination**

Linking

```
int sum(int *a, int n);
int array[2] = \{1, 2\};
int main(int argc, char** argv)
    int val = sum(array, 2);
    return val;
                       main.c
```

How do we find the **sum** function?

Linker



Why do we use **Linking**?

Linking

gcc -o prog main.c sum.c

gcc -o prog sum.c main.c

What do linkers do?

- Symbol resolution
- Relocation of code and data section

Example - Symbol Resolution

```
'sum' function is
                    That's defined in sum.c.
                                                            defined here
                                                                                      Local
                                                                                     variable
       int sum(int *a, int n);
                                                int sum (int *a, int n)
                                                     int i, s = 0;
       int array[2] = \{1, 2\};
                                                    for (i = 0; i < n; i++) {
       int main(int argc, char** argv)
Global
                                                         s += a[i];
definition
            int val = sum(array, 2);
            return val;
                                                    return s;
                                  main.c
                                                                           sum.c
        Local variable
                              Global reference
```

Check your understanding!

```
int time;
int foo(int a) {
  int b = a + 1;
  return b;
int main(int argc,
         char* argv[]) {
 printf("%d\n", foo(5));
  return 0;
```

What would be the symbols in **foo.o**?

Filename: foo.c

Check your understanding! - Solution

```
int time;
int foo(int a) {
  int b = a + 1;
  return b;
int main(int argc,
         char* argv[]) {
 printf("%d\n", foo(5));
  return 0;
```

What would be the symbols in **foo.o**?

- time
- foo
- main
- printf

Filename: foo.c

More Puzzles

```
Cannot initialize a function p1 twice
int x;
p1() {}
             p1() {}
int x;
             int x;
                                 'x' would refer to the same memory location
p1() {}
             p2() {}
             double x;
int x;
                                Write to 'x' in p2 file will override value in 'y' in p1 file
int y;
             p2() {}
p1() {}
int x=7;
             double x;
                                 Write to 'x' in p2 file will override value in 'y' in p1 file
int y=5;
             p2() {}
p1() {}
int x=7;
              int x;
                                 'x' might be initialized
p1() {}
             p2() {}
```

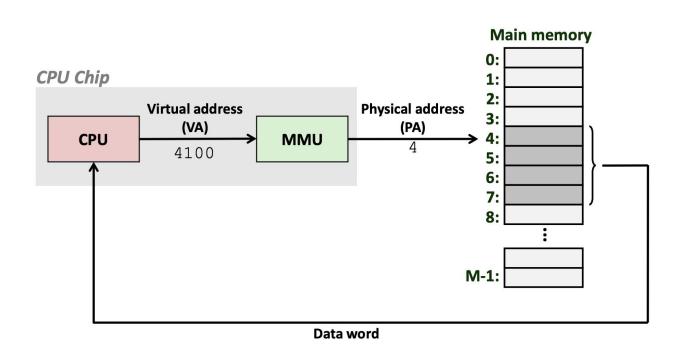
P1

X - mem

Y - mem + 4

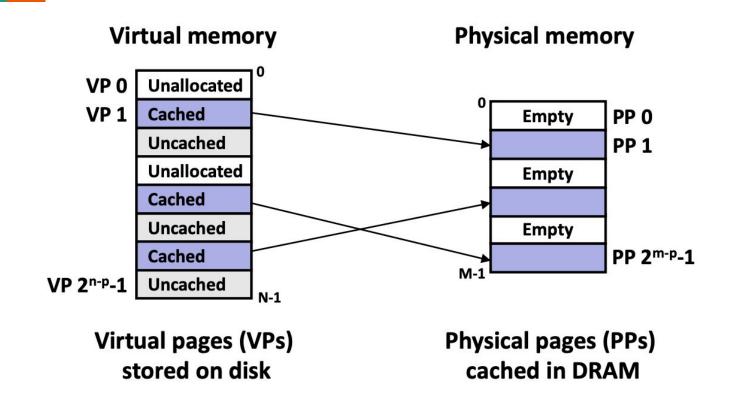
P2: X - mem

Virtual Memory

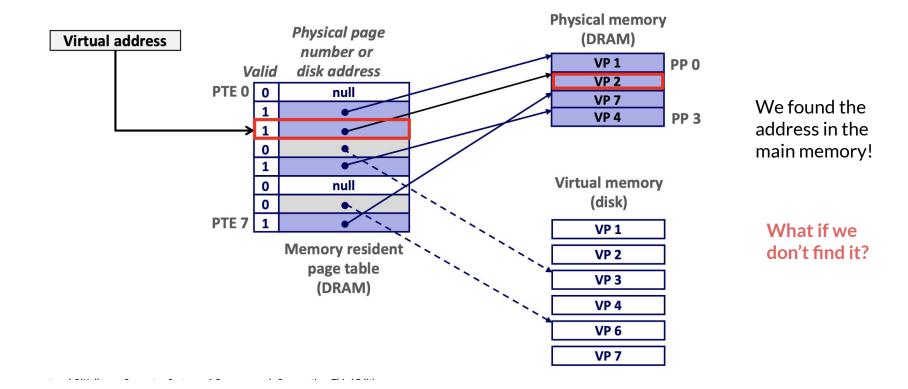


Virtual Memory Addressing

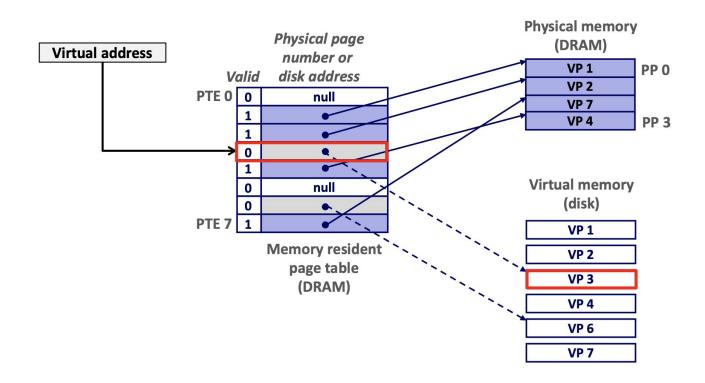
Virtual Memory - Caching



Virtual Memory - Page Hit

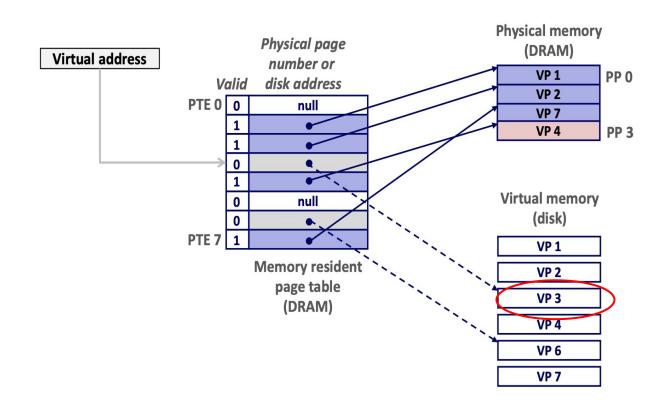


Virtual Memory - Page Fault



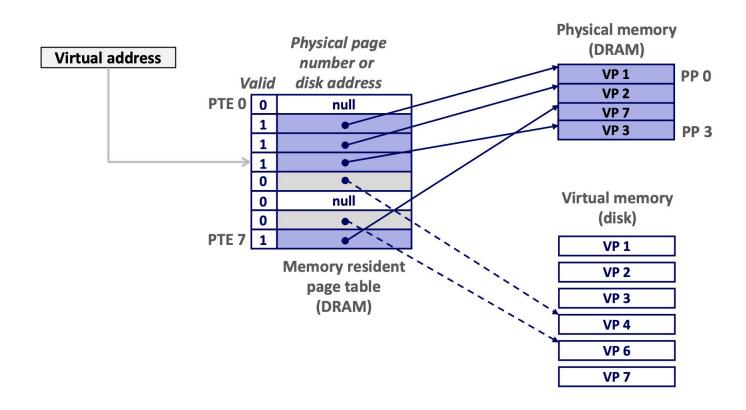
What should we do in this scenario?

Virtual Memory - Handling Page Fault

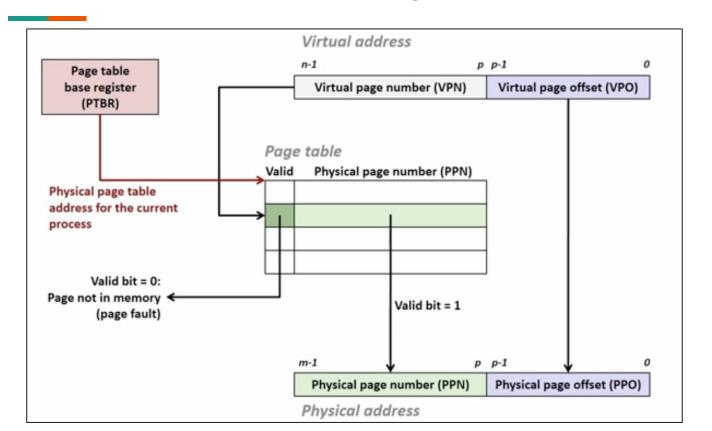


Virtual memory uses the idea of **Demand Paging**

Virtual Memory - Handling Page Fault



Address Translation - Page Table



Check your understanding!

A computer system has a 36-bit virtual address space with a page size of 8K, and 4 bytes per page table entry.

- 1. How many pages are in the virtual address space?
- 2. What is the maximum size of addressable physical memory in this system?

Check your understanding! - Solution

A computer system has a 36-bit virtual address space with a page size of 8K, and 4 bytes per page table entry.

1. How many pages are in the virtual address space?

$$=> 2^36 \text{ bytes} / (2^3*2^10 \text{ page-size}) = 2^36/2^13 = 2^23 \text{ pages}$$

2. What is the maximum size of addressable physical memory in this system?

=> 2^{4*8} = 2^32 pages. Page-size: 2^13 bytes. Size of the physical memory is $(2^32^2^13) = 2^45$ bytes

Worksheet

https://tinyurl.com/cs33-ucla-bday