



# CS 33: Introduction to Computer Organization

TA: Aalisha Dalal

LA: Jonathan Myong

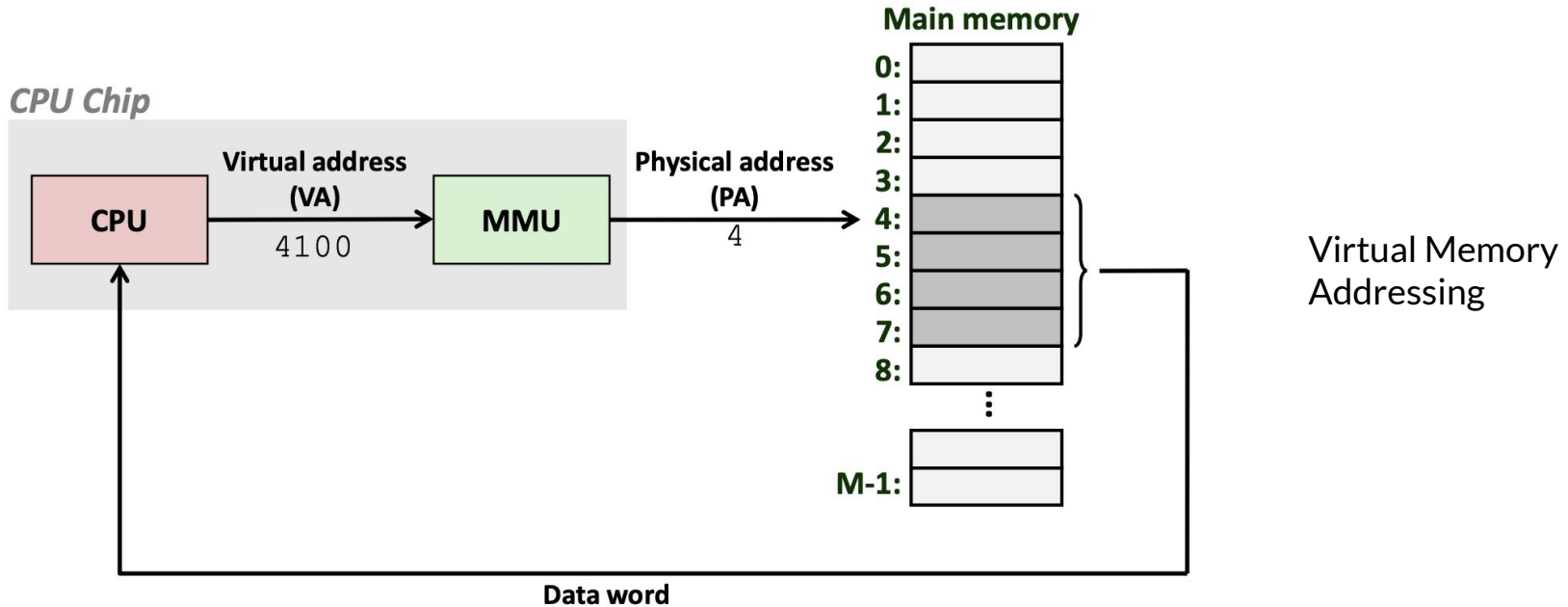
Office Hours: Friday, 9:30-11:30AM

# Outline

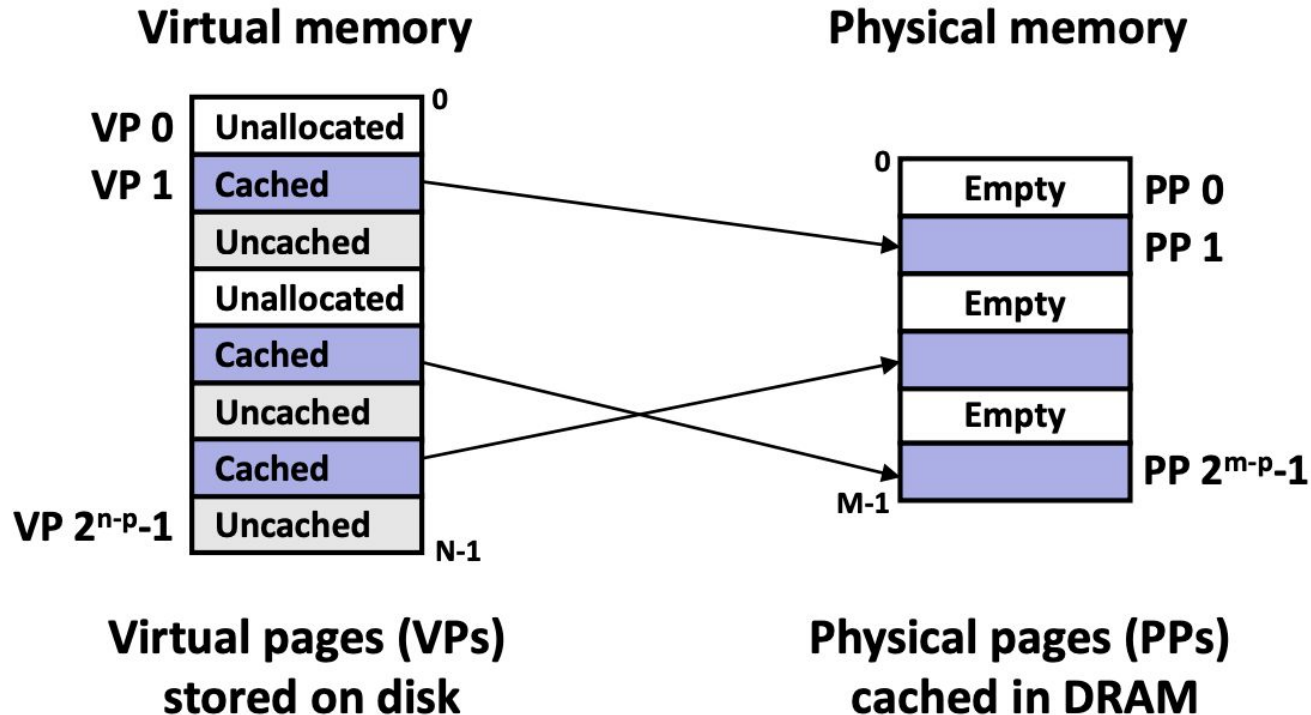


- **Virtual memory [Recap]**
- **MIPS**
- **Revision - Finals ( Part I )**
- **Worksheet Problems**

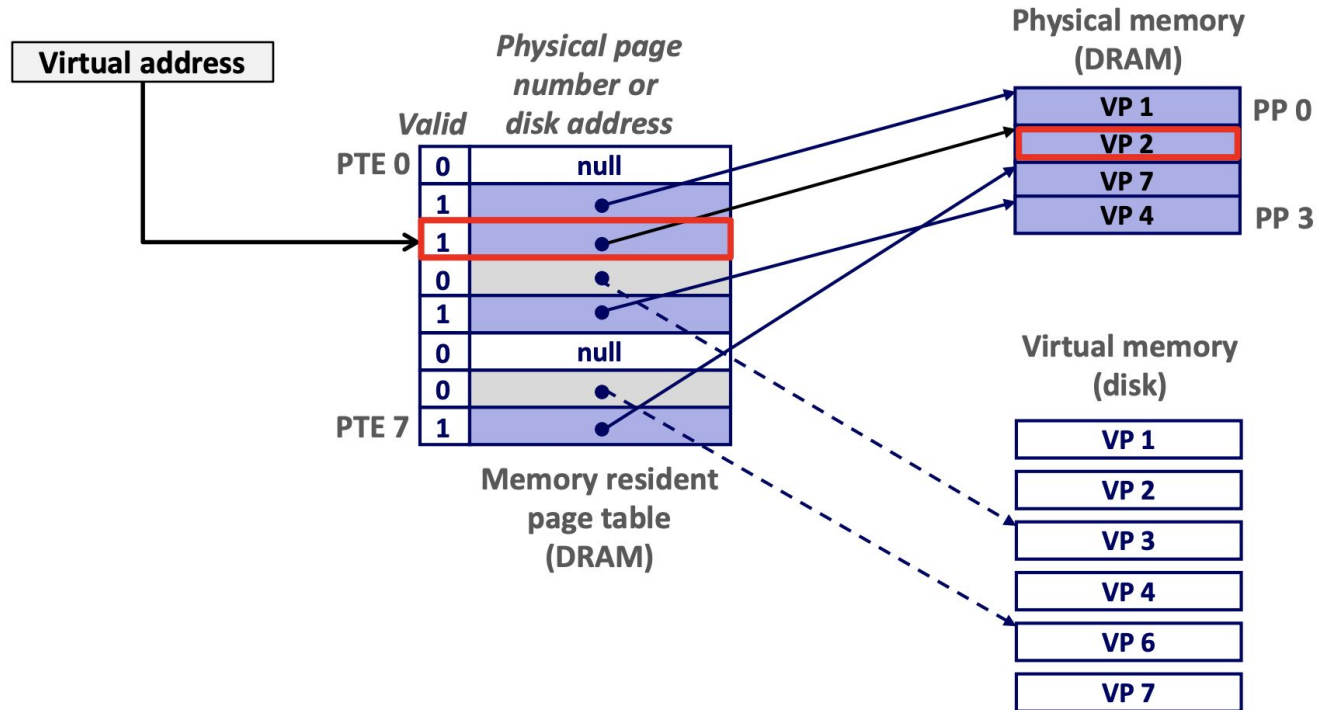
# Virtual Memory



# Virtual Memory - Caching



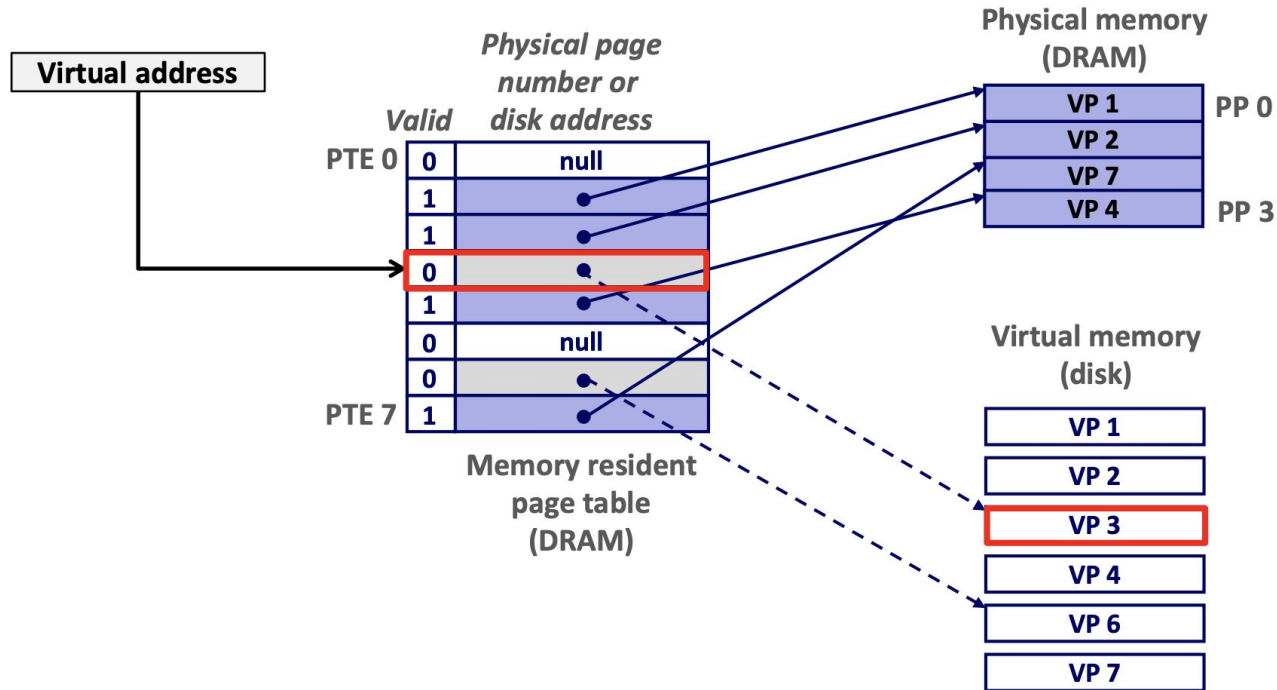
# Virtual Memory - Page Hit



We found the address in the main memory!

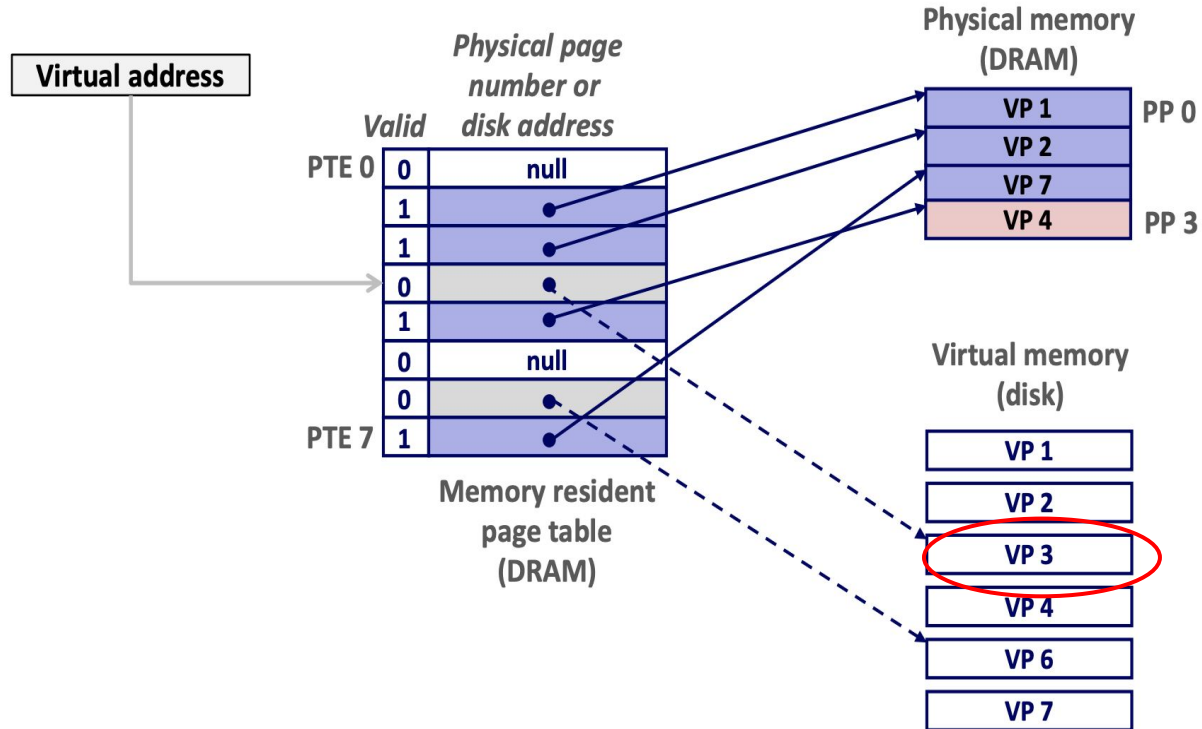
What if we don't find it?

# Virtual Memory - Page Fault



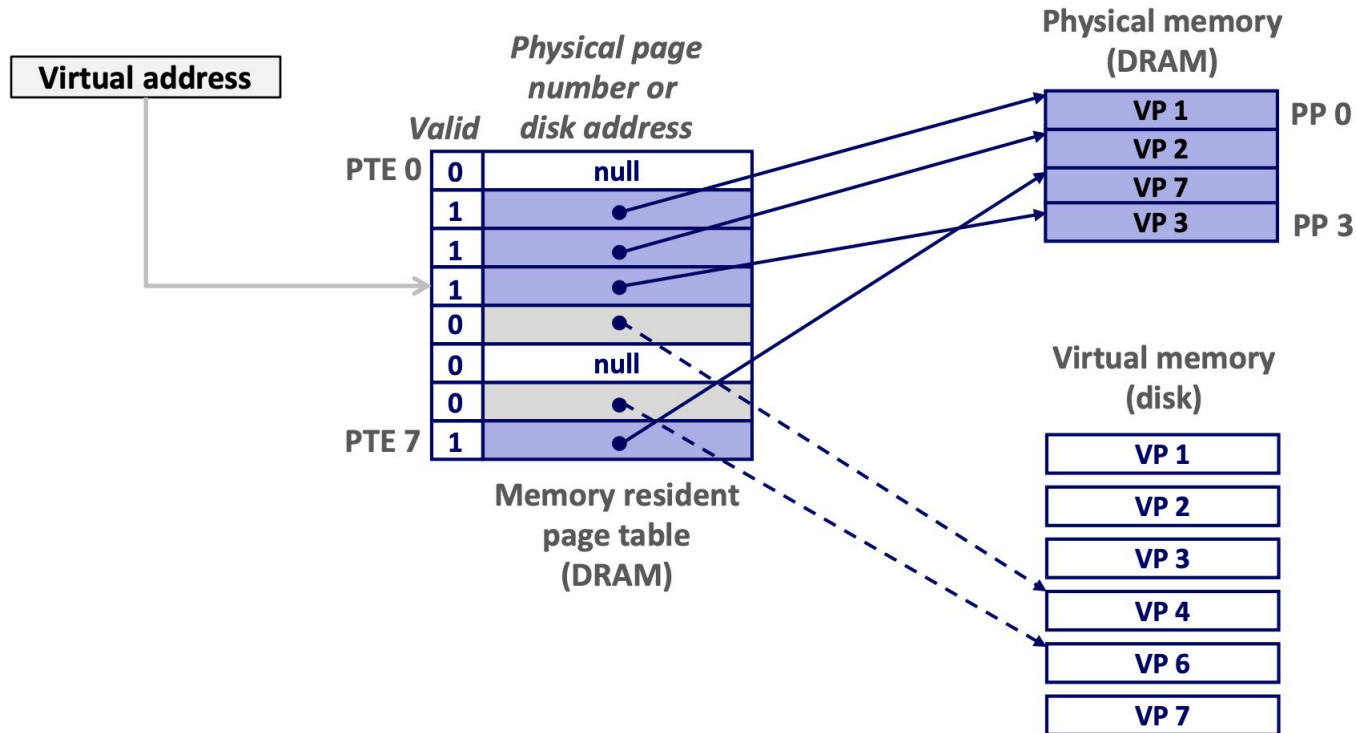
What should we do in this scenario?

# Virtual Memory - Handling Page Fault



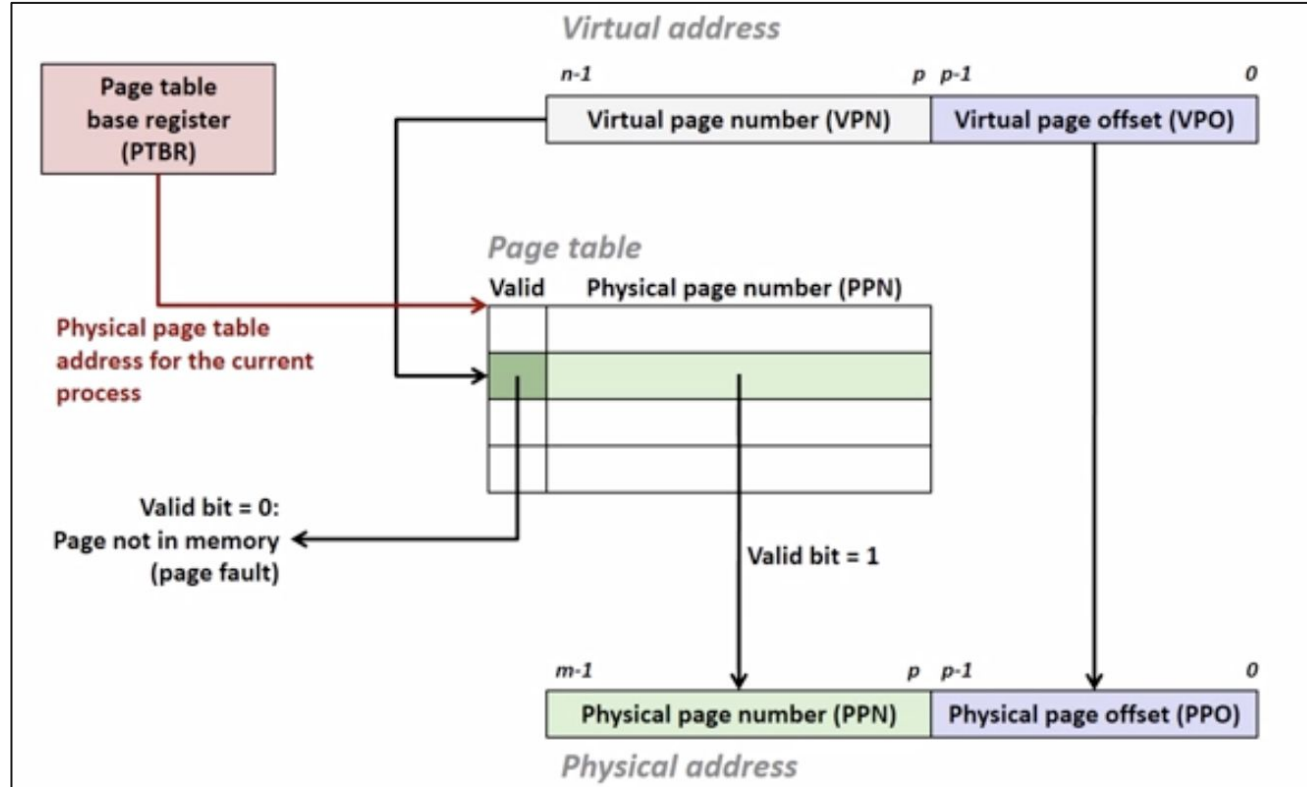
Virtual memory uses the idea of **Demand Paging**

# Virtual Memory - Handling Page Fault





# Address Translation - Page Table



# Check your understanding!



**A computer system has a 36-bit virtual address space with a page size of 8K, and 4 bytes per page table entry.**

1. How many pages are in the virtual address space?
2. What is the maximum size of addressable physical memory in this system?

# Check your understanding! - Solution



**A computer system has a 36-bit virtual address space with a page size of 8K, and 4 bytes per page table entry.**

1. How many pages are in the virtual address space?

=>  $2^{36} \text{ bytes} / (2^3 * 2^{10} \text{ page-size}) = 2^{36} / 2^{13} = 2^{23} \text{ pages}$

2. What is the maximum size of addressable physical memory in this system?

=>  $2^{(4*8)} = 2^{32} \text{ pages}$ . Page-size:  $2^{13} \text{ bytes}$ . Size of the physical memory is  $(2^{32} * 2^{13}) = 2^{45} \text{ bytes}$

# RISC vs. CISC

## RISC vs. CISC Machines

| Feature             | RISC                  | CISC             |
|---------------------|-----------------------|------------------|
| Registers           | ~32                   | 6, 8, 16         |
| Register Classes    | One                   | Some             |
| Arithmetic Operands | Registers             | Memory+Registers |
| Instructions        | 3-addr                | 2-addr           |
| Addressing Modes    | $r$<br>$M[r+c]$ (l,s) | several          |
| Instruction Length  | 32 bits               | Variable         |

## Example - CISC

Complex instruction such as 'mov' in which operands can be both memory address/ registers

```
//Copy %eax register val to %ebx  
mov %eax, %ebx
```

```
//Copy *(%esp+4) to %ebx  
mov 4(%esp), %ebx
```

```
//Copy %ebx register val to *(%esp+4)  
mov %ebx, 4(%esp)
```

# MIPS - RISC Load/Store Architecture

---

```
a = b + c;  
d = a + b;
```



```
Load b into register Rx  
Load c into register Ry  
Rz <- Rx + Ry  
Store Rz into a  
Rz <- Rz + Rx  
Store Rz into d
```

## Code Example



### Hello World

```
.text          # text segment

.global __start

__start:       # execution starts here

    la $a0,str  # put string address into a0

    li $v0,4    #

    syscall     # print

    li v0, 10   #

    syscall     # au revoir...

.data          # data segment

str: .asciiz "hello world\n"
```



# Revision



# Loop Unrolling



## Why is it beneficial?

- Allows faster execution as lesser 'jump' and 'branch' conditions have to be evaluated
- Increase program efficiency while reducing loop overhead

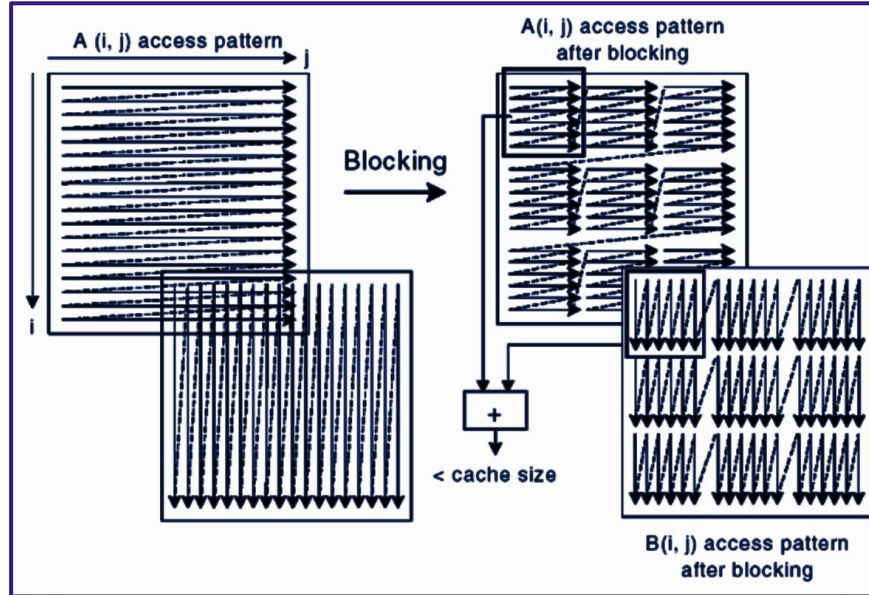
Cons: Code readability reduces

# Loop Unrolling

| Normal loop   | After loop unrolling   |
|---|--|
| <pre data-bbox="189 436 888 780">int x;<br/>for (x = 0; x &lt; 100; x++)<br/>{<br/>    delete(x);<br/>}</pre> | <pre data-bbox="927 436 1723 988">int x;<br/>for (x = 0; x &lt; 100; x += 5 )<br/>{<br/>    delete(x);<br/>    delete(x + 1);<br/>    delete(x + 2);<br/>    delete(x + 3);<br/>    delete(x + 4);<br/>}</pre> |

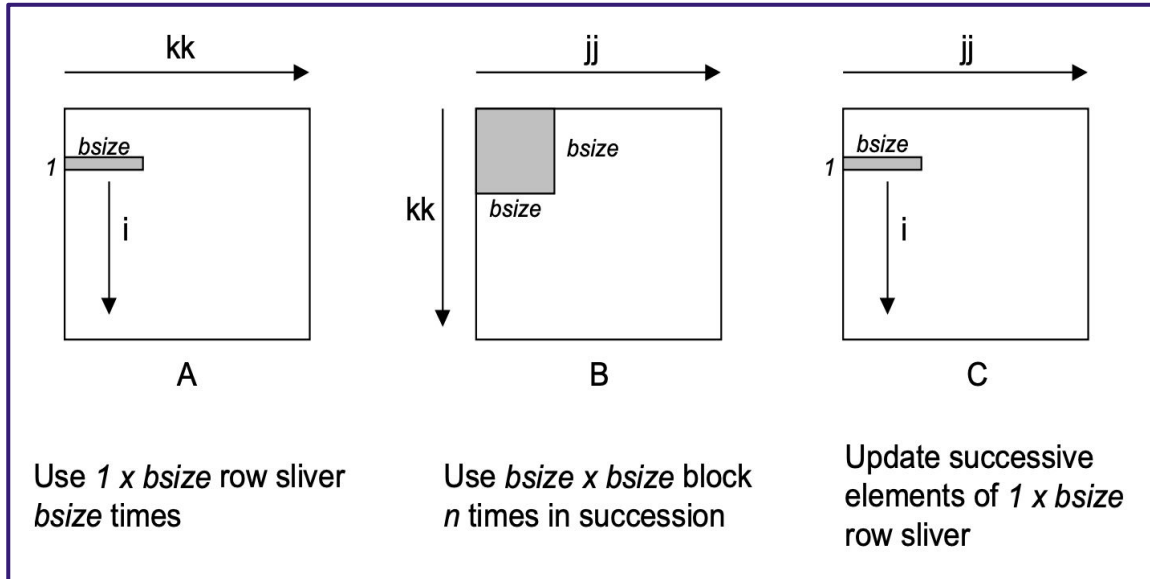
# Loop Blocking/Tiling

*“This technique uses the idea of temporal locality in cache”*



Example for Loop blocking

# Loop Blocking/Tiling



# Code Snippet Example - Matrix Multiplication

```
1 void bijk(array A, array B, array C, int n, int bsize)
2 {
3     int i, j, k, kk, jj;
4     double sum;
5     int en = bsize * (n/bsize); /* Amount that fits evenly into blocks */
6
7     for (i = 0; i < n; i++)
8         for (j = 0; j < n; j++)
9             C[i][j] = 0.0;
10
11     for (kk = 0; kk < en; kk += bsize) {
12         for (jj = 0; jj < en; jj += bsize) {
13             for (i = 0; i < n; i++) {
14                 for (j = jj; j < jj + bsize; j++) {
15                     sum = C[i][j];
16                     for (k = kk; k < kk + bsize; k++) {
17                         sum += A[i][k]*B[k][j];
18                     }
19                     C[i][j] = sum;
20                 }
21             }
22         }
23     }
24 }
```

# Floating Point Conversion



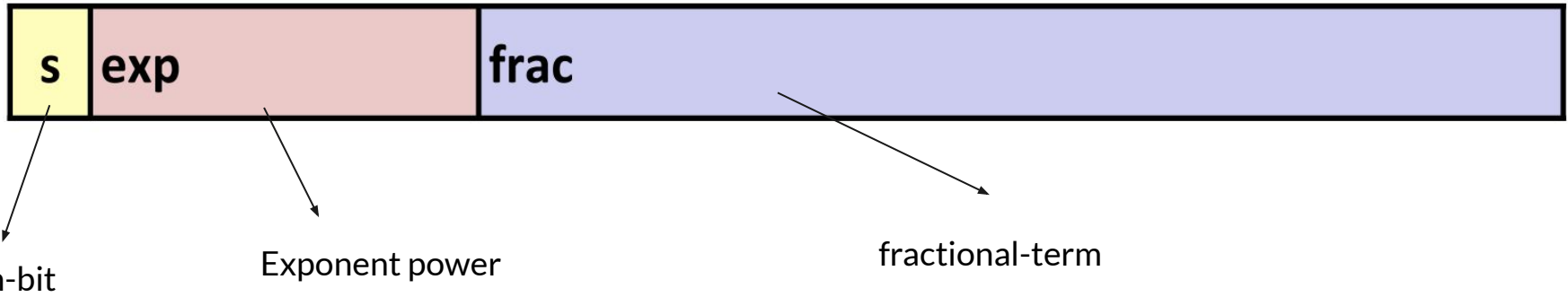
What is  $1011.101_2$ ?

# Revision - Floating Point



What is  $1011.101_2$ ?  $\longrightarrow$   $11 \frac{5}{8}$

## Floating point representation

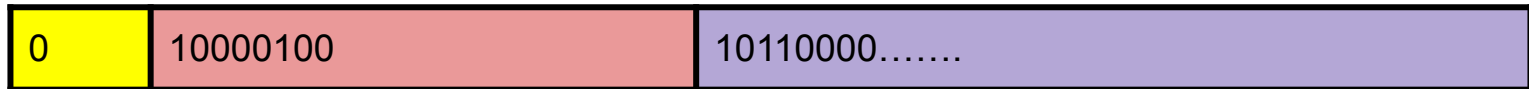


# How would we represent 27?



$$27_{10} = 11011_2$$

$$= 1.1011 * 2^4$$



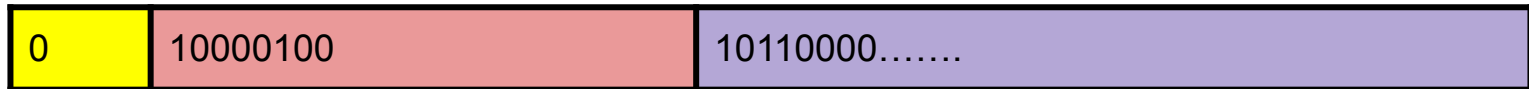


# How would we represent 27?



$$27_{10} = 11011_2$$

$$= 1.1011 * 2^4$$



## Another example!



**float: 0xC0A00000**

## Another example!

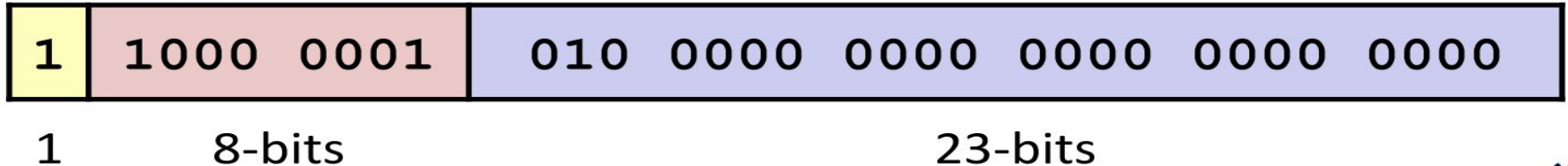
**float: 0xC0A00000**

**binary:** 1100 0000 1010 0000 0000 0000 0000 0000

## Another example!

**float: 0xC0A00000**

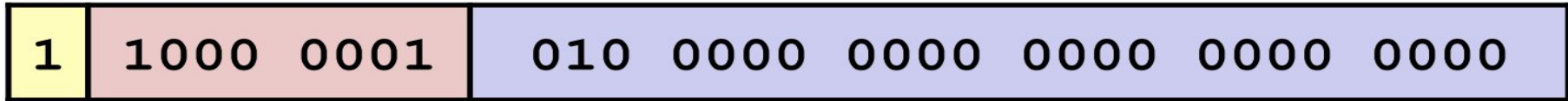
**binary:** 1100 0000 1010 0000 0000 0000 0000 0000



## Another Example!

**float: 0xC0A00000**

binary: 1100 0000 1010 0000 0000 0000 0000 0000



1

8-bits

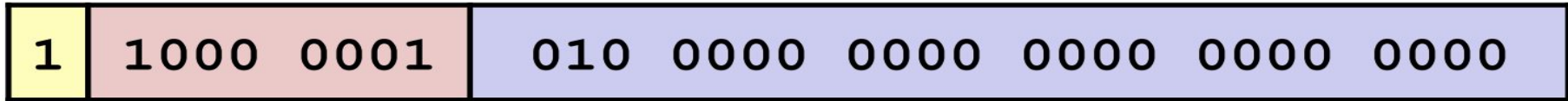
23-bits

**E = 129 -> Exp = 129 - 127 = 2** (decimal)

## Another Example!

**float: 0xC0A00000**

binary: 1100 0000 1010 0000 0000 0000 0000 0000



1

8-bits

23-bits

**M** = 1.010 0000 0000 0000 0000 0000

**E** = 129 -> **Exp** = 129 - 127 = **2** (decimal)



# Worksheet

**<https://tinyurl.com/cs33-2nd-to-last>**