EXPERIMENT 2

CPU SCHEDULING

1. OBJECTIVES

After completing this lab, you will be able to:

- Understand how to schedule processes/threads using round-robin strategy with a fixed time quantum.
- Understand how to create and reset timer interrupts to implement the fixed time quantum.

2. LABORATORY

Software Lab 1 (N4-01a-02) Software Lab 2 (N4-01c-06) SPL (N4-B1b-11)

3. EQUIPMENT

Pentium IV PC with Nachos 3.4.

4. MODE OF WORKING

You should be working alone. No group effort.

5. PROBLEM STATEMENT

In the previous experiment you observed that the short-term CPU scheduler in Nachos is very simple. It is non-preemptive and scheduler.cc in the threads directory implements a FIFO scheduling algorithm for the readyList ready queue. In this experiment, our aim is to implement a **preemptive round-robin CPU scheduling algorithm** for Nachos. For this purpose, we need to implement the following features:

- a) Initialize a periodic timer interrupt with a fixed quantum size of 40 time ticks in Nachos.
- b) Invoke the CPU scheduler to perform a context switch whenever the timer interrupt expires. The current running thread should be moved to the back of the readyList and the thread at the head of readyList should be scheduled for execution next.
- c) Adjust the timer interrupt whenever the current running thread completes (i.e., executes Thread > Finish();) in the middle of a time quantum. Let us understand this using an example. If a thread completes after 10 time ticks in the current time quantum, then the next timer interrupt is currently scheduled to expire 30 time ticks in the future (total time quantum is 40). This means that the new thread that will be scheduled for execution, will only get a time quantum of 30 time ticks. We therefore must "reset" the timer interrupt so that it expires 40 time ticks after the new thread is scheduled.

6. THREADS

The thread is a light-weight process in Nachos. In the first experiment, we have learnt the basic feature of thread execution in Nachos. In this experiment, we need to use/modify the following thread operations to finish the exercises. Please read the

source code under the directory threads (particularly for thread.h and thread.cc).

A thread gives up the CPU ownership and allows another thread to execute. Its state shifts from running to ready, and it is put to the back of the ready queue (readyList).

A thread executes this cleanup operation whenever it completes executing its function/procedure. Among other operations, it sets itself for destruction (threadToBeDestroyed = currentThread;) and then invokes the CPU scheduler for a context switch by calling Sleep();.

7. TIMERS

Timer can be used to trigger an interrupt (i.e., after a fixed number of time ticks). Please read the source code under the directories threads and machine (particularly system.cc, timer.cc and timer.h).

This function in <code>system.cc</code> executes whenever the timer expires and the interrupt is triggered. The timer itself is initialized in <code>system.cc</code> using the constructor for class <code>Timer</code> which is defined in <code>timer.cc</code>.

This function in timer.cc executes whenever the timer expires. It inturn invokes the interrupt handler which is defined in system.cc as described above.

This function returns the next time tick when the interrupt should be scheduled for the current <code>Timer</code> object. This function can be used to make the timer periodic as required for round-robin scheduling.

8. <u>INTERRUPTS</u>

The timer uses several functions from the Interrupt class. Pending timer interrupts in the system are maintained in a list called pending, comprising objects of the class PendingInterrupt. This list is sorted in increasing order of the time tick when the interrupt will be triggered. Please read the source code under the directory machine (particularly interrupt.h and interrupt.cc)

This function schedules a new interrupt and inserts it into the pending list in sorted order. It is also used in the Timer class constructor to initialize a timer interrupt.

This function processes a single time tick. Once the tick variable is updated (stats \rightarrow totalTicks), it calls CheckIfDue() to process any pending interrupts that would be triggered in the current tick. Thereafter, if the yieldOnReturn variable is set to true, then it triggers a context switch by invoking Thread \rightarrow Yield().

This function checks if the pendingInterrupt at the head of pending list should be triggered at the current time tick. If yes, the corresponding interrupt handler is invoked.

This function sets the variable <code>yieldOnReturn</code> to true so that a context switch is triggered when the interrupt processing is completed. It is used in the timer interrupt handler <code>TimerInterruptHandler()</code> defined in <code>system.cc</code>.

9. EXERCISE

We will implement the fixed time-quantum round-robin CPU scheduler in 3 steps. Look for code comments starting with /* Experiment 2*/ for guidance in the code.

- 1. Initialize the timer interrupt with a fixed time quantum of 40 time ticks.
 - a) Activate Timer in system.cc.
 - b) Initialize the timer with the fixed time quantum in timer.cc.
- 2. Make the timer interrupt periodic.
 - a) Modify function TimerExpired() in timer.cc to make the timer periodic. It should trigger an interrupt every 40 time ticks.
 - b) At this stage, you have an initial implementation of preemptive round-robin CPU scheduling. **QUESTION:** Why don't we have to handle the context switch from one thread to another? How does it happen?
 - c) Change your working directory to Experiment 2 by typing cd ~/nachos-exp1-2/exp2. Compile Nachos by typing make. If you see "In -sf arch/inte1-i386-linux/bin/nachos nachos" at the end of the compiling output, your compilation is successful. If you encounter any anomalies, type make clean to remove all object and executable files and then type make again for a clean compilation.
 - d) Trace a run of this Nachos test program by typing ./nachos -d > output_1.txt. Option -d is to display Nachos debugging messages.
 - e) Fill in **Table1.csv** (template is provided) whenever:
 - The timer interrupt is triggered, or
 - The current thread is changed, or
 - The current thread completes.

Tick	ready list	current thread	Timer Interrupt triggered	Thread completion	Context switch
40	child1, child2, child3	main	Timer interrupt scheduled at 80		main -> child1
50	child2, child3, main	child1		Thread 1 Completed	child1 -> child2

- 3. Reset the timer interrupt if a thread finishes in the middle of a time quantum.
 - a) When the current thread finishes, remove the pending timer interrupt from the pending list, and insert a new timer interrupt with the time quantum of 40 time ticks. **Note:** For this experiment, to keep things simple, we will assume that no other interrupts are pending in the list, except the timer interrupts created by us.
 - b) To accomplish the above task, you would need to modify files/functions Threads→Finish(), timer.cc, timer.h, interrupt.cc and interrupt.h.
 - c) Compile and execute Nachos as in Step 2 above (use filename output_2.txt to store your results), and fill Table2.csv (template is provided) whenever:
 - The timer interrupt is triggered, or
 - · The current thread is changed, or
 - The current thread completes.

Tick	ready list	current thread	Timer Interrupt triggered	Thread completion	Context switch
40	child1, child2, child3	main	Timer interrupt scheduled at 80		main -> child1
50	child2, child3, main	child1	Timer interrupt reset to 90	Thread 1 Completed	child1 -> child2

10. ASSESSMENT

- Assessment of your implementation. Please leave your code, the files output_1.txt and output_2.txt, as well as files Table1.csv and Table2.csv in the exp2 folder for TA/Supervisor to review. Deadline is 1 week after your lab session (e.g., if lab session is from 10AM-12PM on a Monday, then deadline is 9:59AM on the next Monday).
- Lab Quiz 1, which is an online multiple-choice quiz, will be administered through NTULearn during ecess week.