

Question 1 [25 Marks]

[CLO 3: Apply various security and risk management tools for achieving information security and privacy]

Fill in the blanks in the space provided and write working on the last page (extra work/calculations).

a) A captured encrypted message is suspected to be encrypted using Rail Fence 3 rails (depth=3):

~~WEGRLTEERDSOEEFEAOCAIVDEN~~

Fill in rail positions (just first few letters):

W E C R L  
T E R D S O E  
F K E A

[5 marks]

Rail 1: W \_ \_ L \_ \_ E \_ \_ S \_ \_

Rail 2: \_ E C \_ T E \_ R D \_ E \_

Rail 3: \_ \_ R \_ E \_ O \_ E \_ A \_

Recovered plaintext: ~~WEGR~~ ~~TE~~ W T E E E F

[5 marks]

b) Columnar Transposition Cipher (Decryption)

An employee receives the following ciphertext that was encrypted using Columnar Transposition.

The keyword used is LOCKED.

Keyword order (write numeric order under each letter):

L	O	C	K	E	D
5	6	1	4	3	2

A	
B	
C	-1
D	-2
E	-3
F	
G	
H	
I	
J	
K	
L	
M	
N	
O	
P	
Q	
R	
S	
T	
U	

Ciphertext: ØFRXUYAXLXXDOEXGCOMOKUS

Step 1 - Write ciphertext in the blanks given below to decrypt:

L	O	C	K	E	D
Ø	E	R	X	U	Y
A	X	L	R	X	X
Ø	O	E	X	G	C
O	M	Q	K	U	S
-	-	-	-	-	-
-	-	-	-	-	-
C	D	E	H	L	O

Step 2 - Write down the plaintext in the blank given below

Plaintext: RVUXOF LXXRAX ECGXDO OSUKOM

c) Frequency Analysis

[5 marks]

A cybersecurity analyst intercepts an encrypted message suspected to be encrypted using a monoalphabetic substitution cipher. The ciphertext is shown below:

NGXIQCTFGZSGGATRZITDTLLQUT

The analyst performs frequency analysis on the ciphertext and observes the following clues:

Cipher Letter	Frequency	Notes
T / L	High	Likely maps to E or T (most common English letters)
Z	Medium	Often maps to A or T
U	Medium	Possible O or A
N / G / X	Low	Likely consonants

Additional hint: The sentence appears to be a normal English statement.

#### Tasks

- Using the frequency hints and common English patterns, identify likely plaintext words in the sentence.
- Fill in the partially-decoded plaintext below (letters revealed as hints):  
NGX IQCT FGZ SGGATR ZT DTLLQUT  
EOR GAME BOT COOKED TIE JESSALE
- Write the final decrypted plaintext sentence below:

Plaintext: FOR GAME BOT COOKED TIE JESSALE

G - O

A  
B  
C  
D  
E  
F  
G  
H  
I  
J  
K  
L  
M  
N  
O  
P  
Q  
R  
S  
T

H  
J  
L  
N  
P

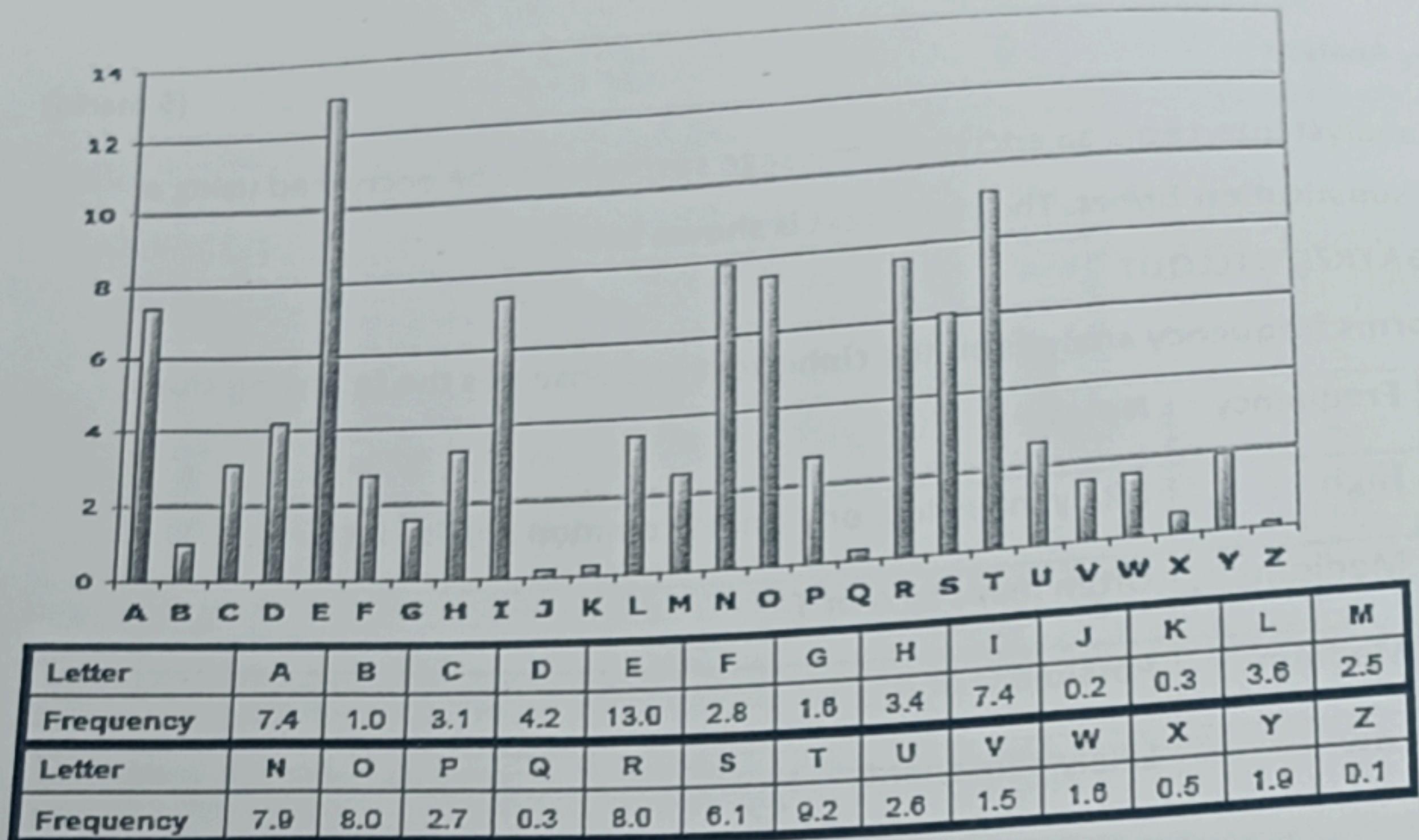
Q  
U

V  
W

X  
Y

Z

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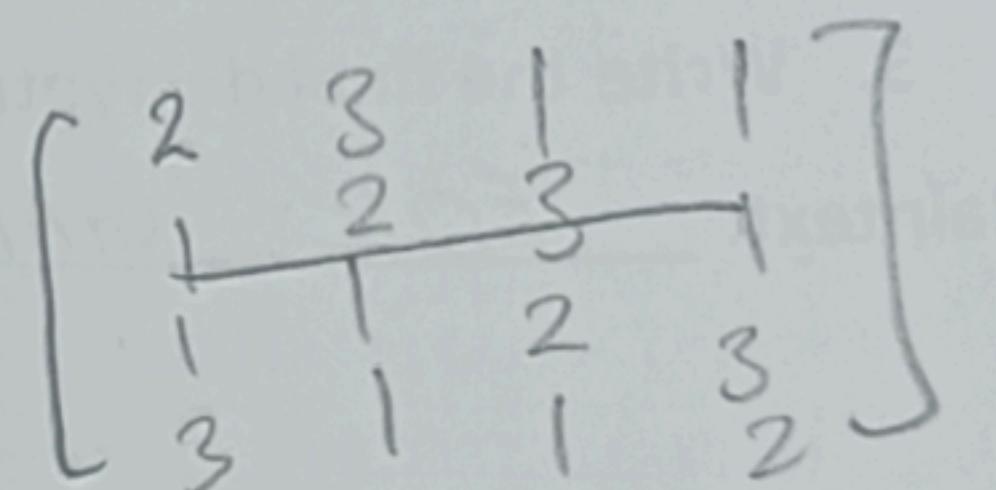


- d) In the AES-128 encryption algorithm, each round (except the last) includes a MixColumns transformation. You are given the following state matrix (in hexadecimal) before MixColumns:

$$\text{State} = \begin{bmatrix} 1 & 2 & 3 \\ D4 & E0 & B8 & 1E \\ BF & B4 & 41 & 27 \\ 5D & 52 & 11 & 98 \\ 30 & AE & F1 & E5 \end{bmatrix}$$

and the MixColumns transformation matrix:

$$M = \begin{bmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{bmatrix}$$



Compute only the 2nd byte of 3rd word after the MixColumns transformation.

Fill in the blanks and show your working on last page:

[5 marks]

$$b'_{(3,4)} = (\underline{1} \times \underline{B8}) \oplus (\underline{2} \times \underline{41}) \oplus (\underline{3} \times \underline{11}) \oplus (\underline{1} \times \underline{E1})$$

$$= \underline{B8} \oplus \underline{82} \oplus \underline{33} \oplus \underline{F1}$$

$$= \underline{\underline{F8}}$$

- e) You are given the following AES key schedule words (in hexadecimal):

[5 marks]

$$\text{Key Words } (W_0-W_7) = \begin{bmatrix} W_0 & W_1 & W_2 & W_3 & W_4 & W_5 & W_6 & W_7 \\ 2B & 28 & AB & 09 & A0 & 88 & 23 & 2A \\ 7E & AE & F7 & CF & FA & 54 & A3 & 6C \\ 15 & D2 & 97 & 4F & FE & 2C & 39 & 76 \\ 16 & A6 & 75 & 3C & 17 & B1 & 39 & 05 \end{bmatrix}$$

Rcon Matrix (Round Constants)

$$\text{Rcon} = \begin{bmatrix} 01 & 02 & 04 & 08 & 10 & 20 & 40 & 80 & 1B & 36 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \end{bmatrix}$$

Using the AES-128 key expansion process, compute the next two words:  
~~W<sub>8</sub> and W<sub>9</sub>.~~

$$W_8 = \underline{\underline{\hspace{1cm}}}$$

$$W_9 = \underline{\underline{\hspace{1cm}}}$$

	y															
x	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	63	7C	77	7B	F2	6B	6F	C5	30	01	67	2B	FE	D7	AB	76
1	CA	82	C9	7D	FA	59	47	F0	AD	D4	A2	AF	9C	A4	72	C0
2	B7	FD	93	26	36	3F	F7	CC	34	A5	E5	F1	71	D8	31	15
3	04	C7	23	C3	18	96	05	9A	07	12	80	E2	EB	27	B2	75
4	09	83	2C	1A	1B	6E	5A	A0	52	3B	D6	B3	29	E3	2F	84
5	53	D1	00	ED	20	FC	B1	5B	6A	CB	BE	39	4A	4C	58	CF
6	D0	EF	AA	FB	43	4D	33	85	45	F9	02	7F	50	3C	9F	A8
7	51	A3	40	8F	92	9D	38	F5	BC	B6	DA	21	10	FF	F3	D2
8	CD	0C	13	EC	5F	97	44	17	C4	A7	7E	3D	64	5D	19	73
9	60	81	4F	DC	22	2A	90	88	46	EE	B8	14	DE	5E	0B	DB
A	E0	32	3A	0A	49	06	24	5C	C2	D3	AC	62	91	95	E4	79
B	E7	C8	37	6D	8D	D5	4E	A9	6C	56	F4	EA	65	7A	AE	08
C	BA	78	25	2E	1C	A6	B4	C6	E8	DD	74	1F	4B	BD	8B	8A
D	70	3E	B5	66	48	03	F6	0E	61	35	57	B9	86	C1	1D	9E
E	E1	F8	98	11	69	D9	8E	94	9B	1E	87	E9	CE	55	28	DF
F	8C	A1	89	0D	BF	E6	42	68	41	99	2D	0F	BD	54	BB	16

(a) S-box

**Question 2 [ 25 Marks]**

To be answered on the last page

[CLO 1: Explain key concepts of information security such as design principles, cryptography, risk management]

1. In DES, during the Feistel function (f-function), the 32-bit right half is expanded to 48 bits using: [1 mark]

- a. Initial Permutation (IP)
- b. Key Schedule (PC-1)
- c. Expansion D-box (E)
- d. Substitution Boxes (S-boxes)

2. Given  $p = 61$ ,  $q = 53$ , and  $e = 17$ . Find the modulus  $n$  and Euler's totient  $\varphi(n)$ . [2 marks]

a.  $n = 3233$ ,  $\varphi(n) = 3120$

$$\begin{aligned} n &= p \times q \\ &= 61 \times 53 \\ &= 3233 \end{aligned}$$

$$\begin{aligned} \varphi(n) &= (p-1) \times (q-1) \\ &= 60 \times 52 \\ &= \end{aligned}$$

b.  $n = 3120$ ,  $\varphi(n) = 3233$

c.  $n = 3233$ ,  $\varphi(n) = 3180$

- d. None of above

3. For the same RSA setup, find the private exponent  $d$  such that  $17 \cdot d \equiv 1 \pmod{3120}$ . [2 marks]

a. 157

b. 2753

c. 187

- d. None of above

$p = 61 \quad q = 53 \quad c = 17$

$\varphi(n) = 3120 \quad n = 3233$

$17 \cdot d \pmod{3120} = 1$

$$\Rightarrow \frac{3120 \times d + 1}{17} = d$$

4. After XORing the expanded right half with the round key, which of the following steps reduces the data back to 32 bits? [1 mark]

- a. P-box permutation

$6.599 \times 10^{30}$

46800

- b. Key mixing

$6.599 \times 10^{30}$

$2 \times 10^{27}$

- c. S-box substitution

- d. Left circular shift

$2 \times 10^{27}$

5. Encrypt  $m = 65$  using the public key ( $n = 3233$ ,  $e = 17$ ). Find ciphertext  $c \equiv m^e \pmod{n}$ . [2 marks]

a. 2790

$65^{17} \pmod{3233}$  C7

b. 2805

- c. 2753
  - d. None of above
6. Given  $n = 2773$ ,  $e = 17$ ,  $d = 157$ , and ciphertext  $c = 1924$ , find the plaintext message  $m$ . [2 marks]
- a. 65
  - b. 157
  - c. 123
  - d. None of above

$$n = 2773 \quad e = 17 \quad d = 157 \quad c = 1924$$

$$1924^d \bmod n \quad C = m^e \bmod n$$

$$1924^{157} \bmod 2773 \quad 65^{17} \bmod 2773$$

7. Which of the following correctly describes one round of DES encryption? [2 marks]
- a.  $L_t = R_{t-1}$ ,  $R_t = L_{t-1} \oplus f(R_{t-1}, K_t)$
  - b.  $R_t = L_{t-1}$ ,  $L_t = R_{t-1} \oplus f(L_{t-1}, K_t)$
  - c.  $L_t = f(R_{t-1}, K_t)$ ,  $R_t = L_{t-1}$
  - d.  $L_t = L_{t-1}$ ,  $R_t = f(R_{t-1}, K_t)$

8. The DES key schedule starts with a 64-bit key, but only 56 bits are used because: [1 mark]
- a. The remaining 8 bits are for parity checking.
  - b. The remaining 8 bits are discarded for security reasons.
  - c. The 8 bits are added to the IV (Initialization Vector).
  - d. They are used in the expansion step.

8. Let  $p = 23$ ,  $g = 5$ ,  $a = 6$ ,  $b = 15$ . Find the shared secret  $s$  after both parties exchange their public keys. [2 marks]
- a. 8
  - b. 5
  - c. 19
  - d. None of above

$$P = 23 \quad g = 5 \quad a = 6 \quad b = 17$$

$$g^5 \bmod 23$$

$a$  &  $b$  are private keys or what?  
or use generally accepted notation.

9. In the key generation process, PC-2 (Permutation Choice 2) selects bits from: [1 mark]
- a. The output of the Expansion D-box
  - b. The output of PC-1 after left shifts
  - c. The plaintext halves (L and R)
  - d. The ciphertext feedback loop

National University of Computer and Emerging Sciences  
Islamabad Campus

[CLO 4: Identify appropriate techniques to tackle and solve problems of real life in the discipline of information security]

Analyze the following Cryptographic Flow / Functional Block Diagrams to determine which of the five security properties are provided: Confidentiality (C), Integrity (I), Authentication (A), Non-repudiation (NR), and Freshness (F) (replay safety). For each system:

- Understand what remains [Private] and what is exchanged on [Public] channel/ out-of-band [OOB].
- Evaluate resistance to MITM and Replay Attack. Then choose the most accurate option.

**C:** Information is accessible only to authorized entities (control of keys/credentials).

**I:** Unauthorized modification/insertion/deletion is detectable by crypto-graphic checks.

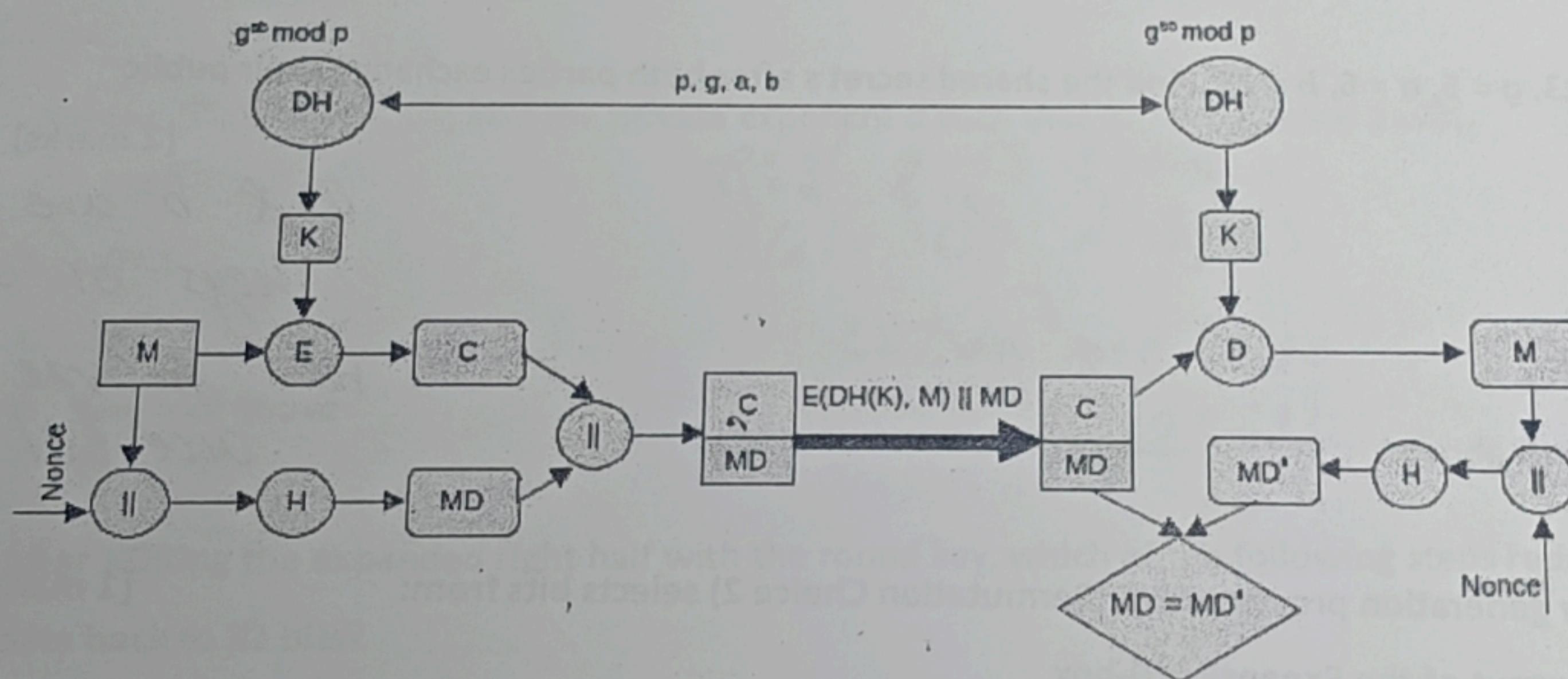
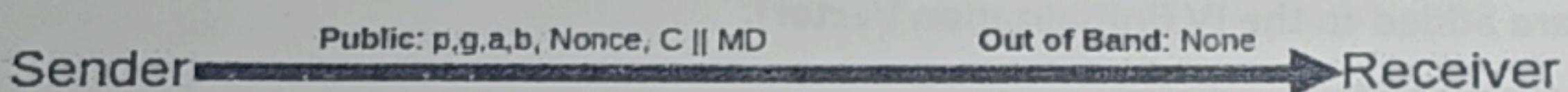
**A:** Sender or data source is verified via credentials/keys/signatures

**NR:** Cryptographic proof of origin/action prevents later denial.

**F:** Message is recent, i.e., not a replay of a previous valid exchange.

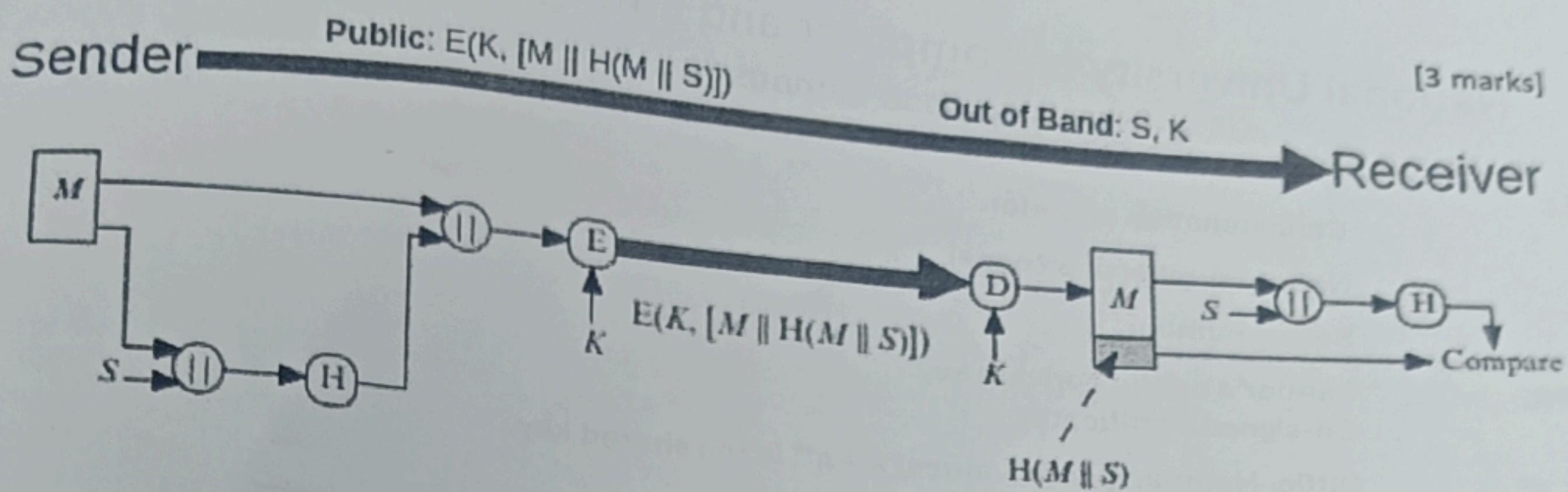
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[3 marks]



- C, I
- C, I, F
- C, I, A
- C, I, A, NR
- C, I, A, NR, F
- None of the Above

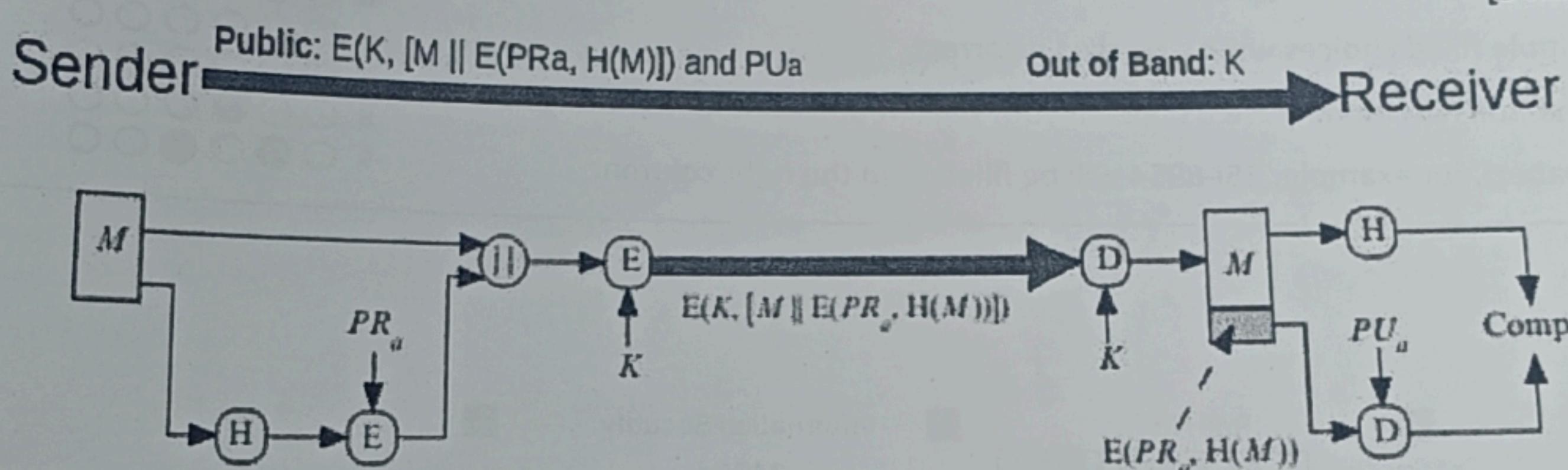
11.



- a) C, I
- b) C, I, F
- c) C, I, A
- d) C, I, A, NR
- e) C, I, A, NR, F
- f) None of the Above

12.

[3 marks]



- a) I, F
- b) I, A, F
- c) C, I, A
- d) C, I, A, NR
- e) C, I, A, NR, F
- f) None of the Above

#### Glossary (symbols)

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$M$	Plaintext message.
$H$	Cryptographic hash.
$MD$	Message Digest generated by hash function.
$E/D$	Encrypt / Decrypt primitives.

d) Third Word =  $\begin{bmatrix} B8 \\ 41 \\ 11 \\ F1 \end{bmatrix}$ ; For 2nd Byte [01 02 03 01]

$$\Rightarrow (01 \times B8) \oplus (02 \times 41) \oplus (03 \times 11) \oplus (01 \times F1)$$

Here,

$$01 \times B8 = B8 ;$$

$$01 \times F1 = F1$$

but,

$02 \times 41 \Rightarrow$  First convert  $41_{16}$  into binary

$$41_{16} = 0100\ 0001_2$$

$\because$  MSB = 0, we can simply do left shift by 1

$$\therefore 02 \times 41 = 1000\ 0010_2 = 82_{16}$$

$$8, 03 \times 11 = (02 \times 11) \oplus (01 \times 11)$$

$$\Rightarrow 02 \times 11 \Rightarrow 0001\ 0001_2 \quad \because \text{MSB} = 0 \text{ we can simply do left shift}$$

$$\Rightarrow 0010\ 0010_2$$

$$= 22_{16}$$

$$\Rightarrow (03 \times 11) = (22) \oplus (11) \quad \because 01 \times 11 = 11$$

$$= 33_{16}$$

$$\Rightarrow (B8) \oplus (82) \oplus (33) \oplus F1 = F8$$

#	Option	marks
1	c	1
2	A	2
3	b,d	2
4	c	1
5	a,d	2
6	c,d	2
7	a	2
8	a	1
8	d	2
9	b	1
10	f	3
11	c	3
12	c	3