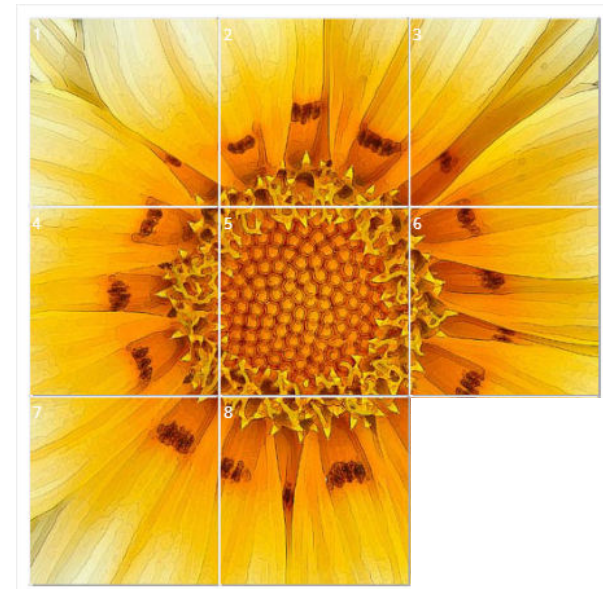
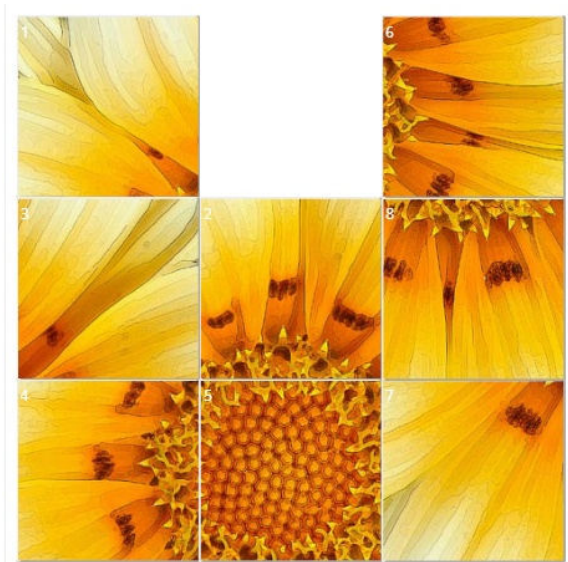


# The $A^*$ Algorithm

# 8 puzzle



# 8 puzzle

8	5	4
6	2	7
1	3	

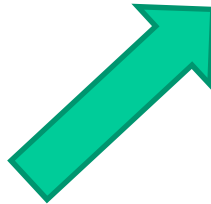


1	2	3
4	5	6
7	8	

- The goal is to go from the **initial state** to the **goal state**

# 8 puzzle

8	5	4
6	2	7
1	3	



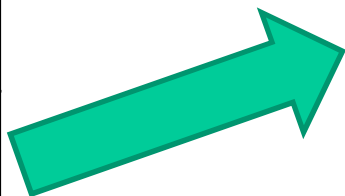
8	5	4
6	2	
1	3	7



8	5	4
6	2	7
1		3

# 8 puzzle

8	5	4
6	2	
1	3	7



8	5	
6	2	4
1	3	7



8	5	4
6	2	7
1	3	



8	5	4
6		2
1	3	7



# 8 puzzle

8	5	4
6	2	
1	3	7



8		4
6	5	2
1	3	7

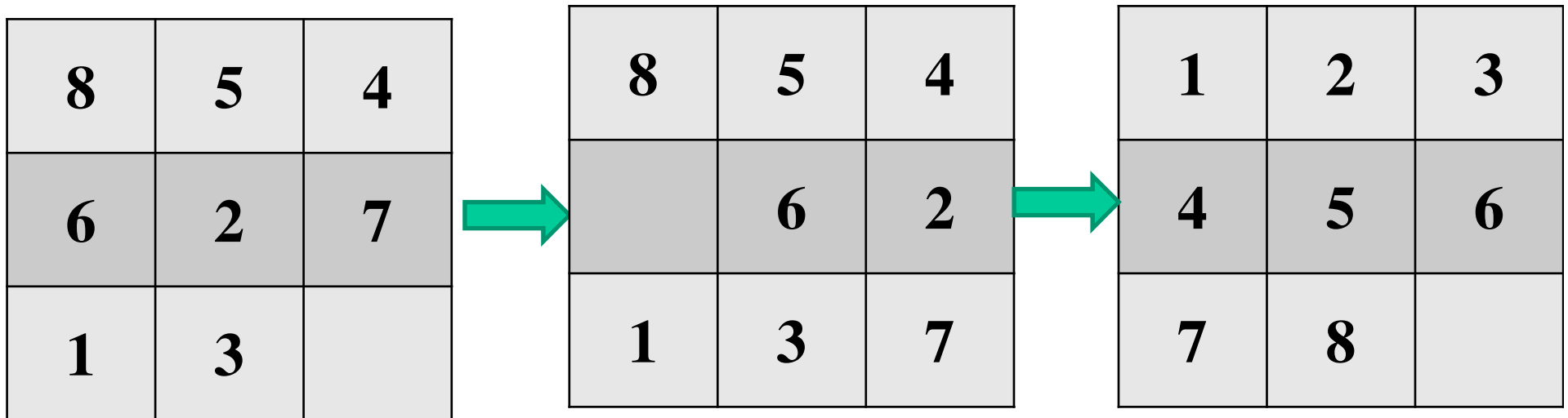
8	5	4
6	3	2
1		7



8	5	4
6		2
1	3	7

8	5	4
	6	2
1	3	7

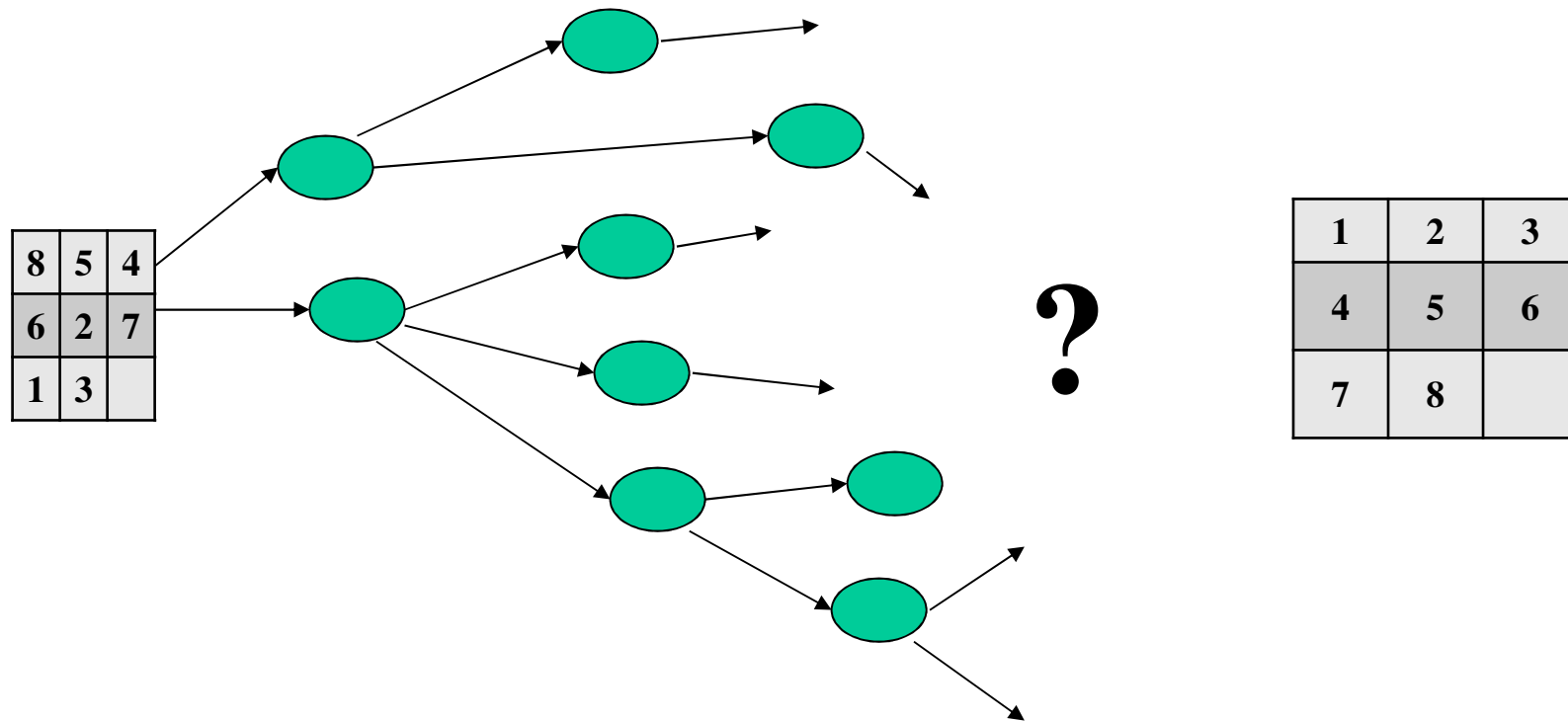
# 8 puzzle



- The goal is to go from the **initial state** to the **goal state** via other states in **minimum number of moves**

# Can be considered as a graph

Starting from a **node n** find the shortest path to a goal **node g**

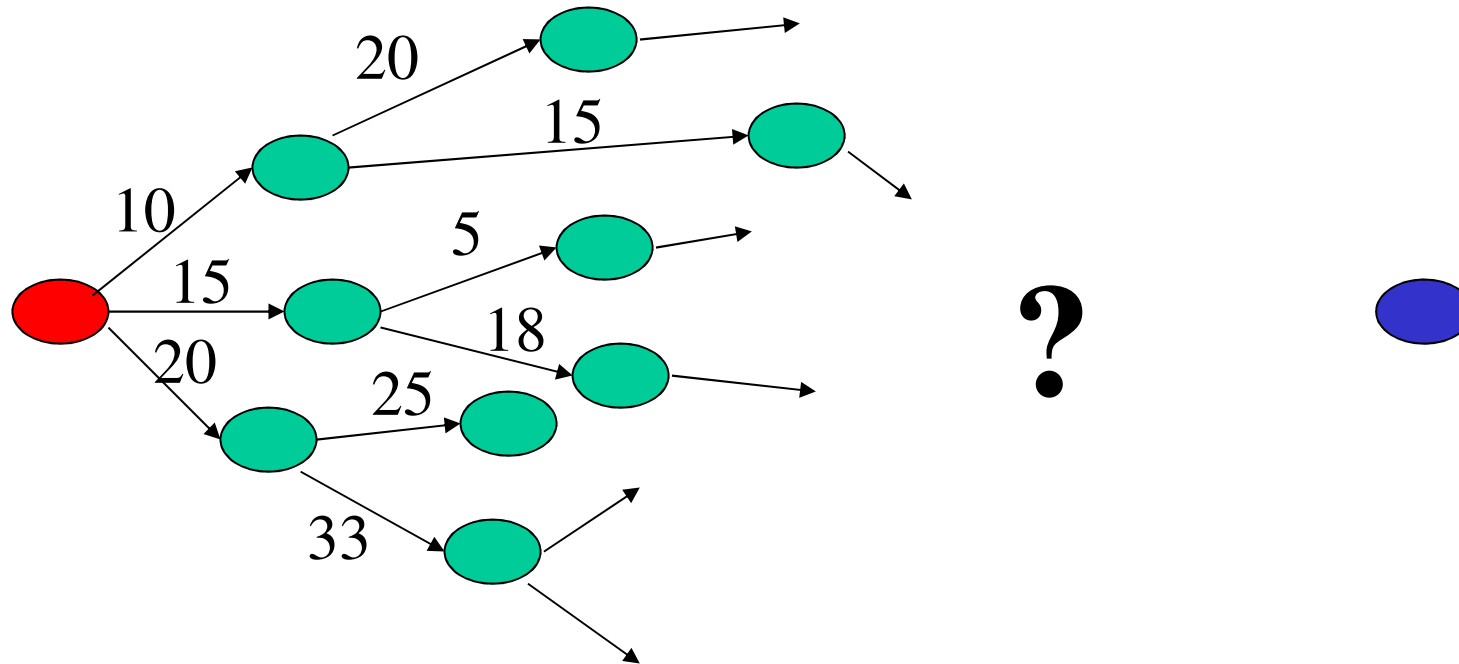


**State Space Search**



# The Generalized Search Problem

Starting from a **node n** find the shortest path to a goal **node g**



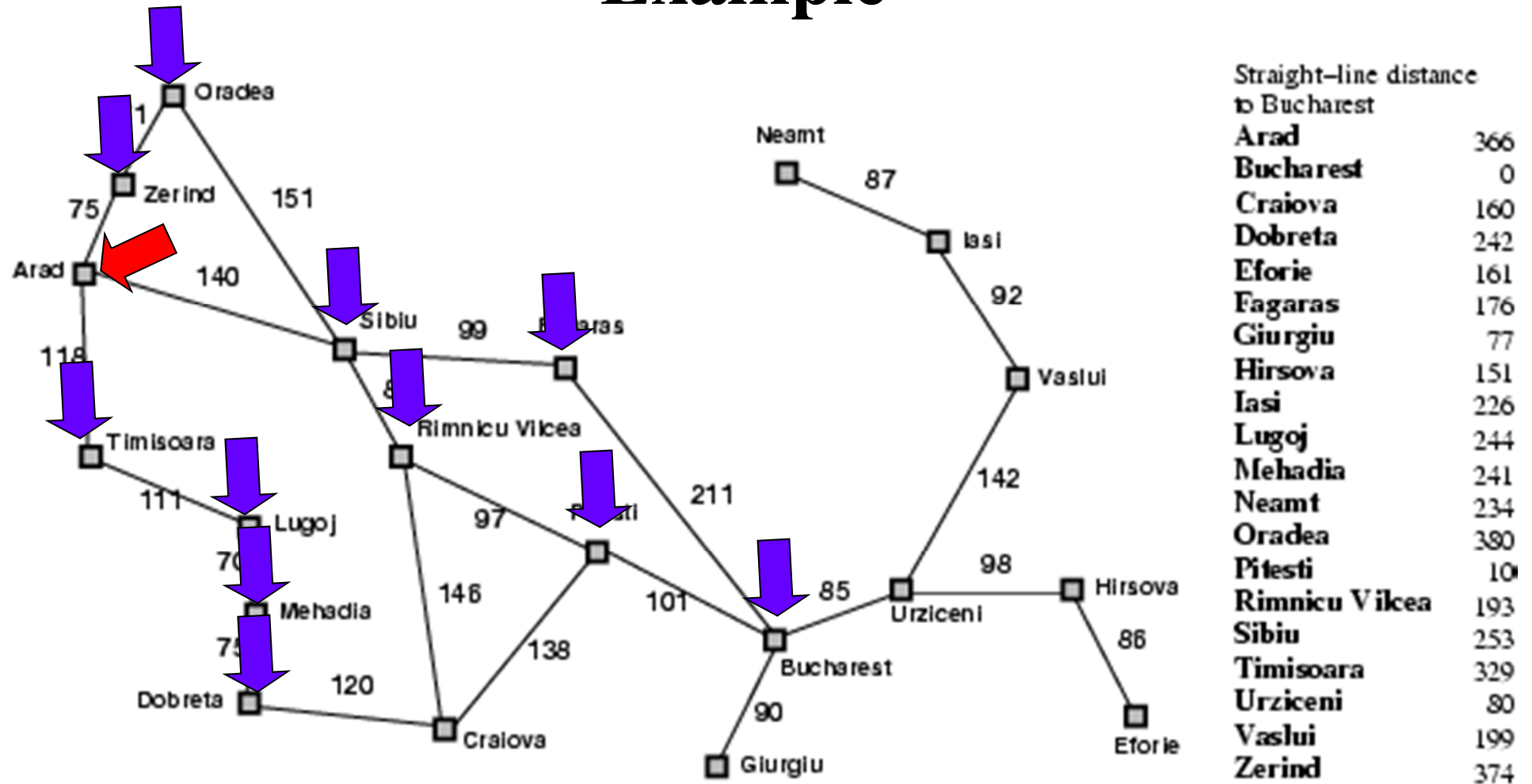
**We want:** A path  $\text{red node} \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow \text{blue node}$

Such that  $g(\text{blue node}) = \text{cost}(\text{red node} \rightarrow v_1) + \text{cost}(v_1 \rightarrow v_2) + \dots$   
 $+ \text{cost}(v_{n-1} \rightarrow \text{blue node})$  is minimum

How to solve this problem ??

Dijkstra Algorithm ???

# Example

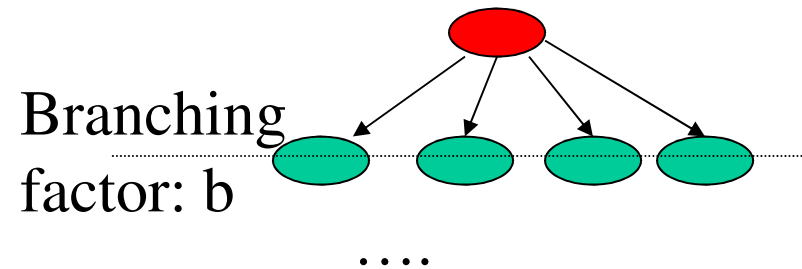


What does Dijkstra's algorithm will do? (minimizing  $g(n)$ )

**Problem:** Visit too many nodes, some *clearly* out of the question

# Complexity

- Actual complexity is  $O(|E|\log_2 |E|)$
- Is this good?  
Actually it is bad for very large graphs!



# nodes =  $b^{(\text{\# levels})}$

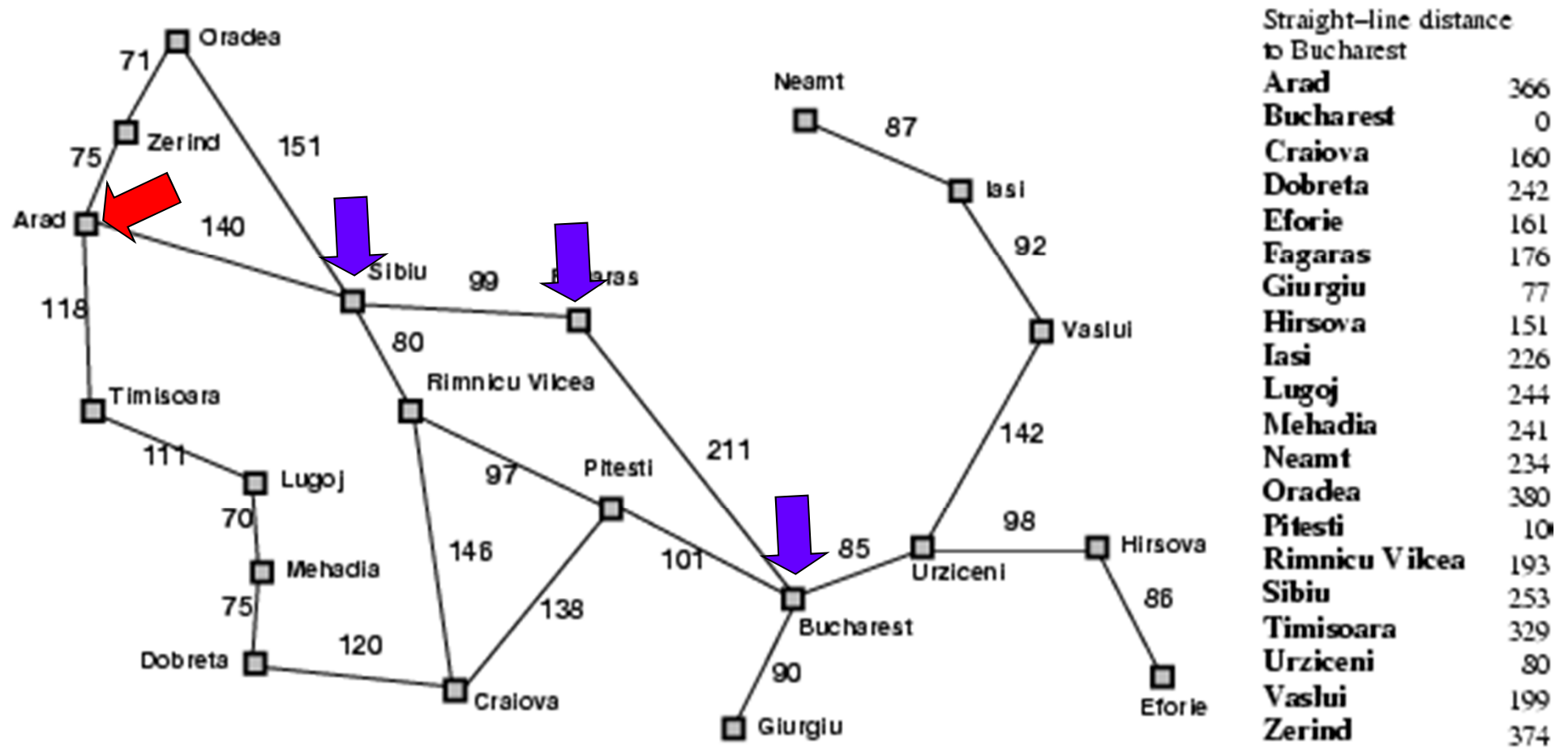
```
graph LR; A(( )) --- B(( )) --- C(( )) --- D(( )) --- E(( ))
```

**Another Example:** think of the search space in chess

# Better Solution: Make a “hunch”!

- Use *heuristics* to guide the search
  - **Heuristic**: estimation or “hunch” of how to search for a solution
- We define a heuristic function:  
 $h(n)$  = “estimate of the cost of the cheapest path from the **initial state** to the **goal state**”

# Lets Try A Heuristic



**Heuristic:** minimize  $h(n)$  = “Euclidean distance to destination”

**Problem:** not optimal (through Rimmici Viacea and Pitesti is shorter)

# The A\* Search

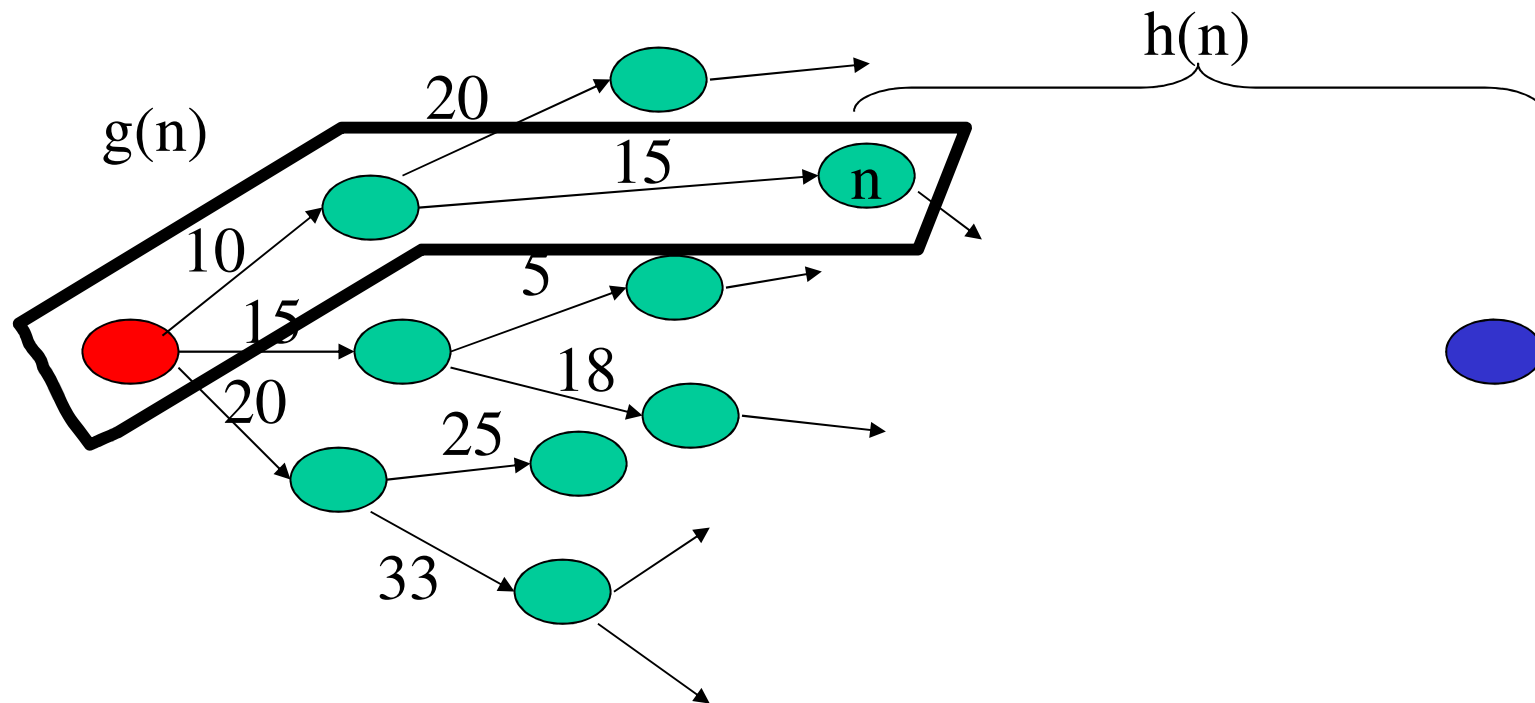
- **Difficulty:** we want to still be able to generate the path with minimum cost
- A\* is an algorithm that:
  - Uses heuristic to guide search
  - While ensuring that it will compute a path with minimum cost

- A\* computes the function  $f(n) = g(n) + h(n)$ 
  - “estimated cost”
  - “actual cost”

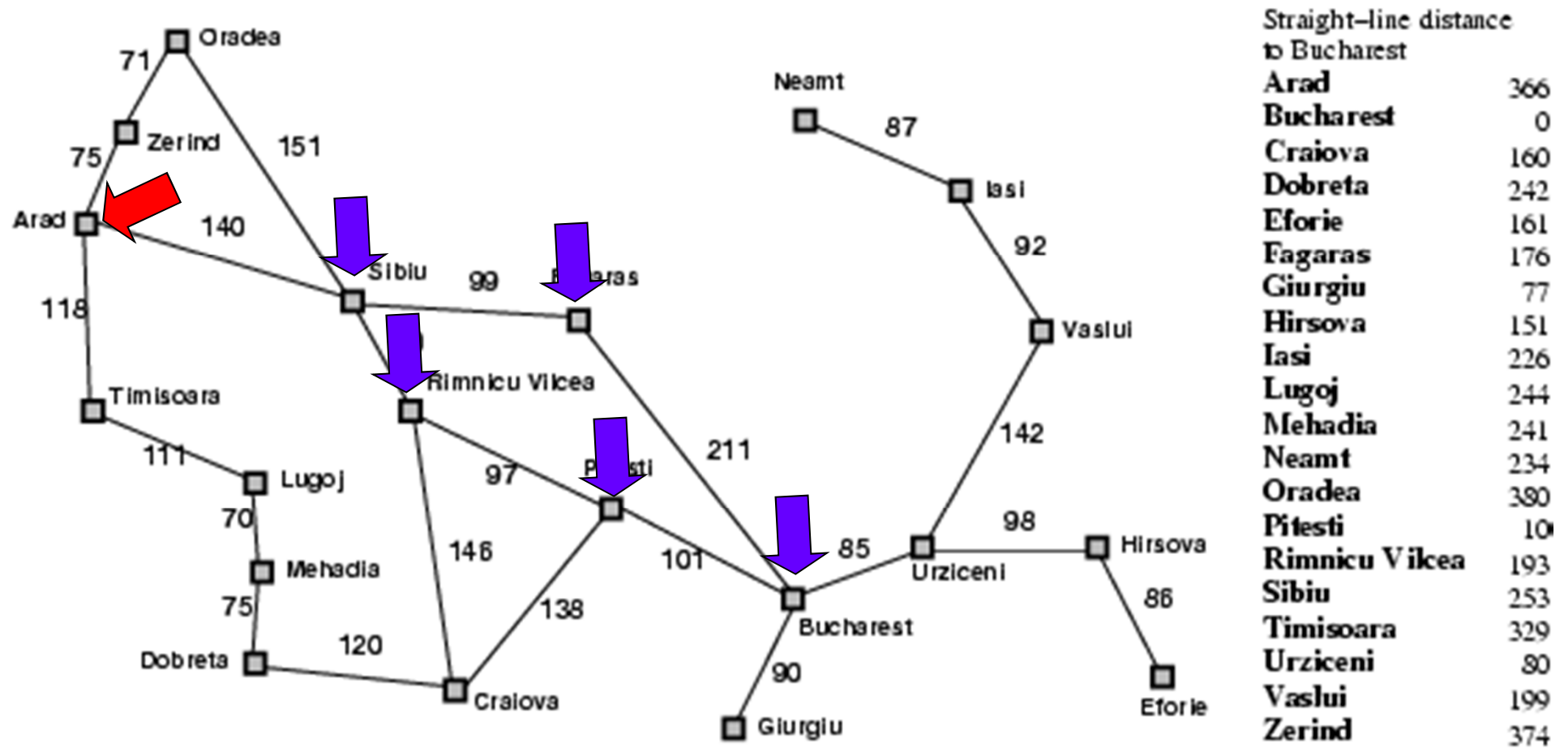


# A\*

- $f(n) = g(n) + h(n)$ 
  - $g(n)$  = “cost from **the starting node** to reach  $n$ ”
  - $h(n)$  = “estimate of the cost of the cheapest path from  $n$  to the **goal node**”



# Example

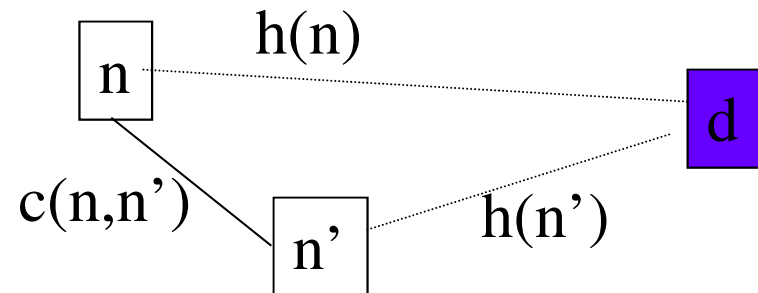


$A^*$ : minimize  $f(n) = g(n) + h(n)$

# Properties of A\*

- A\* generates an optimal solution if  $h(n)$  is an admissible heuristic and the search space is a tree:
  - $h(n)$  is **admissible** if it never overestimates the cost to reach the destination node
- A\* generates an optimal solution if  $h(n)$  is a consistent heuristic and the search space is a graph:
  - $h(n)$  is **consistent** if for every node  $n$  and for every successor node  $n'$  of  $n$ :

$$h(n) \leq c(n, n') + h(n')$$



- If  $h(n)$  is consistent then  $h(n)$  is admissible
- Frequently when  $h(n)$  is admissible, it is also consistent

# Admissible Heuristics

- A heuristic is admissible if it is optimistic, estimating the cost to be smaller than it actually is.
- Example:

In the road map domain,

$$h(n) = \text{“Euclidean distance to destination”}$$

is admissible as normally cities are not connected by roads that make straight lines

What if the heuristic is **too**  
optimistic??

# How to Create Admissible Heuristics

- Relax the conditions of the problem
  - This will result in admissible heuristics!
- Lets look at an 8-puzzle game:

/

- Possible heuristics?

# 8 puzzle

8	5	4
6	2	7
1	3	



1	2	3
4	5	6
7	8	

- The goal is to go from the **initial state** to the **goal state**

# Example: Admissible Heuristics in 8-Puzzle Game

- Heuristic: a tile A can be moved to any tile B
  - $H1(n)$  = “number of misplaced tiles in board n”
- Heuristic: a tile A can be moved to a tile B if B is adjacent to A
  - $H2(n)$  = “sum of distances of misplaced tiles to goal positions in board n”
- Some experimental results reported in Russell & Norvig (2002):
  - $A^*$  with  $h2$  performs up to 10 times better than  $A^*$  with  $h1$
  - $A^*$  with  $h2$  performs up to 36,000 times better than a classical uninformed search algorithm (iterative deepening)



# A\* in Games

- Path finding
  - <http://aigamedev.com/open/interviews/mario-ai/>
- A\* can be used for planning moves computer-controlled player (e.g., chess)

