Real-Coded Genetic Algorithm Benchmarked on Noiseless Black-Box Optimization Testbed

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ABSTRACT

Genetic algorithms, a class of stochastic population-based optimization algorithms, are widely realized as effective tools to solve optimization problems arising from diverse application domains. In this paper, a real-coded genetic algorithm (RCGA), which employs an adaptive-range variant of the well-known non-uniform mutation, is furnished with an independent restarts mechanism to benchmark the noise-free black-box optimization testbed. The maximum number of function evaluations for each run is set to 50000 times the search space dimension. For low search space dimensions the algorithm shows very good results on many functions. Although the algorithm is unable to solve all the functions to the highest required accuracy, for each type of functions, some of them can be solved, especially to lower precision, with dimensions up to 40.

Categories and Subject Descriptors

G.1.6 [Numerical Analysis]: Optimization—global optimization, unconstrained optimization; F.2.1 [Analysis of Algorithms and Problem Complexity]: Numerical Algorithms and Problems

General Terms

Algorithms

Keywords

Benchmarking, Black-box optimization, Evolutionary computation, Real-coded genetic algorithm

1. INTRODUCTION

Genetic algorithms (GAs) are general-purpose stochastic search algorithms belonging to the earliest and most fundamental branch of evolutionary computation. Soon after John Holland introduced the original form in 1975 [5], GAs have rapidly grown in popularity due to their simplicity and

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GECCO'10, July 7–11, 2010, Portland, Oregon, USA. Copyright 2010 ACM 978-1-4503-0073-5/10/07 ...\$10.00. ease of implementation. As a result, a multitude of variations to the original Holland GA have been developed based on various aspects of the basic framework, i.e. representation schemes, selection, crossover, mutation, and elitism operators.

Among representation schemes being available in GAs, floating-point technique, in which each candidate solution (or chromosome vector) is encoded as a vector of floating-point valued numbers of the same length as the dimension of the search space, is the most widely used. Among several selection strategies, tournament selection is a popular form and commonly used owing to its ease of implementation.

The role of selection phase is to create an intermediate population for the subsequent application of other operators. Its main principle is "the better is an individual, the higher is its chance of being parent" [9]. Once an intermediate population is established, crossover will be applied to produce offspring. Offspring may be created through various manners. Notably, arithmetical crossover [7] has become a favorite operator to recombine information of parent solutions so as to yield an offspring solution that benefits from the advantageous information of both parents. Besides, non-uniform mutation [6] has also gained in popularity as one of the most important operators to introduce new exploratory information into the intermediate population while ensuring exploitation in order to produce high precision solutions.

In this work, a GA using float-point representation together with tournament selection, arithmetical crossover and adaptive-range variation of non-uniform mutation, which is commonly referred to as a real-coded genetic algorithm (RCGA), is implemented and benchmarked on the noiseless functions of the BBOB-2010 testbed. More details about the RCGA being in use is presented in the following sections.

2. REAL-CODED GENETIC ALGORITHM

Over the years, several variations have been developed within GAs, however they still conform to the general structure as illustrated in Algorithm 1

2.1 Selection Operator

Genetic algorithms, and evolutionary algorithms in general, search for near-optimal solutions by evolving population of candidate solutions gradually over generations, in consequence, what individuals of the current population selected for crossover may have significant influence on the next population. The underlying idea for selection methods is that better individuals, i.e. higher fitness, have higher chance of being parents in recombination process (crossover

Algorithm 1 General structure of a Genetic Algorithm

```
Input: Fitness function f; parameters; stopping_condition
Output: Best solution, \mathbf{x}_b, f(\mathbf{x}_b)
 1: Initialize \mathcal{P}(0), \mathcal{P}(0) = \{ \mathbf{x}_i(0), 1 \le i \le N \}
 2: \{f(\mathbf{x}_i(0)), 1 \leq i \leq N\} \leftarrow \text{Evaluate}(\mathcal{P}(0))
 3: t \leftarrow 1 // generation counter
 4: while not stopping_condition do
            \mathcal{P}'(t) \leftarrow \text{Selection}(\mathcal{P}(t)) \ // \ \mathcal{P}'(t) = \{\mathbf{x}_i'(t)\}
 5:
            \mathcal{P}'(t) \leftarrow \text{Crossover}(\mathcal{P}'(t))
 6:
            \mathcal{P}'(t) \leftarrow \operatorname{Mutation}(\mathcal{P}'(t))
 7:
            \{f(\mathbf{x}_i'(t))\} \leftarrow \text{Evaluate}(\mathcal{P}'(t))
 8:
            \mathcal{P}(t+1) \leftarrow \text{Replace}(\mathcal{P}(t), \mathcal{P}'(t))
 9:
10:
            t \leftarrow t + 1 // advance to next generation
11: end while
12: return \mathbf{x}_b, f(\mathbf{x}_b)
```

and mutation). The rationale for this principle is that, in a population, highly fit members possess good properties that, if recombined in the right way, could lead to even better solutions. The impact of a selection phase can be regulated by a selection pressure which may drive the population to better members. On the contrary, worst members should not be discarded and should have some chance to be selected since they may lead to useful genetic material. In generational genetic algorithms, however, not every individual makes it into the intermediate population. This issue is concerned in all selection schemes and known as the loss of diversity.

A common selection strategy is fitness proportional selection, also known as "roulette-wheel selection". As the name suggests, individuals get drawn with probabilities directly proportional to their fitness values. This is in contrast to using the rank of an individual to assign these probabilities. This selection strategy, however, may introduce a bias towards an exceptionally fit individual at the beginning of the search which will probably dominate the intermediate population, leading to a loss of diversity and premature convergence. For a more suitable selection strategy, tournament selection is used in this work.

Tournament selection selects a group of n_{ts} individuals uniformly at random from the population, where $n_{ts} < N$, N is population size and n_{ts} is tournament size. The performance of n_{ts} selected individuals is then compared in terms of fitness and the best individual from this group is selected and returned by the operator. To select N individuals, the tournament procedure is repeated N times. Pseudo-code of this strategy is presented in Algorithm 2.

Algorithm 2 Tournament selection

```
Input: \mathcal{P}(t) = \{\mathbf{x}_i(t)\}, \, \mathcal{F}(t) = \{f(\mathbf{x}_i(t))\}, \, N, \, n_{ts}
Output: \mathcal{P}'(t) = \{\mathbf{x}_i'(t)\}, \, \mathcal{F}'(t) = \{f(\mathbf{x}_i'(t))\}
1: for i = 1 to N do
2: Tournament\_Group \leftarrow \text{Pickup}(\mathcal{P}(t), \mathcal{F}(t), n_{ts})
3: (\mathbf{x}_{win}, f(\mathbf{x}_{win})) \leftarrow \text{ChooseBest}(Tournament\_Group)
4: (\mathbf{x}_i', f(\mathbf{x}_i')) \leftarrow (\mathbf{x}_{win}, f(\mathbf{x}_{win}))
5: end for
6: return \mathcal{P}'(t) = \{\mathbf{x}_i'(t)\}, \, \mathcal{F}'(t) = \{f(\mathbf{x}_i'(t))\}
```

2.2 Crossover Operator

The role of crossover operators is to inherit some characteristics of the two parents to generate offspring. For RC-

GAs, several crossover operators have been developed. In this version of RCGA, we employ the arithmetical crossover. The implementation of this operator is described as follows. Suppose $\mathbf{x}_1(t) = [x_{1,1}^t, x_{1,2}^t, \cdots, x_{1,D}^t]$ and $\mathbf{x}_2(t) = [x_{2,1}^t, x_{2,2}^t, \cdots, x_{2,D}^t]$ are two parent individuals, chosen uniformly at random from the population, subject to be crossed at generation t, $\mathbf{x}_1'(t) = [x_{1,1}^{tt}, x_{1,2}^{tt}, \cdots, x_{1,D}^{tt}]$ and $\mathbf{x}_2'(t) = [x_{2,1}^{tt}, x_{2,2}^{tt}, \cdots, x_{2,D}^{tt}]$ are the corresponding offspring pair generated by crossover with probability $p_c \in (0,1]$. Then each element of the produced offspring is a combination of two corresponding elements from the two parents:

$$x_{1,j}^{\prime t} = \lambda_j x_{1,j}^t + (1 - \lambda_j) x_{2,j}^t$$
$$x_{2,j}^{\prime t} = (1 - \lambda_j) x_{1,j}^t + \lambda_j x_{2,j}^t$$

where λ_j is a random number in (0,1) which is uniformly drawn anew for each element index j.

2.3 Mutation Operator

Several mutation schemes have been suggested to enhance RCGAs. Among remarkable schemes is the non-uniform mutation introduced by Janikow and Michalewicz [6]. It is designed for giving a fine-tuning capabilities aimed at achieving high precision. For a given parent individual $\mathbf{x}_i(t) = [x_{i,1}^t, \cdots, x_{i,j}^t, \cdots, x_{i,D}^t]$ at generation t, each element $x_{i,j}^t$ $(1 \leq j \leq D)$ of it has exactly equal chance of undergoing the mutative process. Assuming that the element $x_{i,k}^t$ is, with a mutation probability p_m , selected for mutation, then the resulting offspring is $\mathbf{x}_i'(t) = [x_{i,1}^t, \cdots, x_{i,k}^{tt}, \cdots, x_{i,D}^t]$. The mutation-generated element $x_{i,k}^t$ is produced by adding to or subtracting from the original element $x_{i,k}^t$ a perturbation amount as follows:

$$x_{i,k}^{\prime t} = \begin{cases} x_{i,k}^t + \Delta(t, ub_j - x_{i,k}^t) \text{ with prob. } q\\ x_{i,k}^t - \Delta(t, x_{i,k}^t - lb_j) \text{ with prob. } 1 - q \end{cases}$$
 (1)

where $\Delta(t,y)$ is the perturbation function, dependent on generation t and position y of the original value relative to the search boundaries. This function returns a value in the range (0,y) such that the returned value approaches 0 as t increases. This property causes the operator to search the space uniformly at early stages when t is small and very locally at later stages. The function $\Delta(t,y)$ is given as follows:

$$\Delta(t,y) = y \cdot \left(1 - r^{\gamma(t)}\right) \tag{2}$$

with r is a random number uniformly distributed in (0,1), and $\gamma(t)$ provides the fine-tuning capability according to

$$\gamma = \left(1 - \frac{t}{T}\right)^{\beta} \tag{3}$$

Here T is the maximum number of generation and β the exogenous strategy parameter which determines the degree of non-uniformity across generations.

Originally, the value of q is proposed to be 0.5 [7] which means that mutation to the left and right of the original value is equally likely. Neubauer [8] demonstrated that non-uniform muation is not a zero-mean deviation operator. By forming an expression for the expected value of mutation and using the above perturbation, mutation was shown to concentrate the search between the parent value and the center of the search range. As a modification, an adaptive range variation was introduced [1]. The modified operator establishes a mutation range $x \pm \Delta(t,y)$ based on the generation number t and a fixed preset value y, then randomly

selects a point within this range. The mutation range is redefined as follows

$$\begin{cases}
\sigma_L = \max\{lb_i, x_{i,k}^t - \Delta(t, y)\} \\
\sigma_U = \min\{ub_i, x_{i,k}^t + \Delta(t, y)\}
\end{cases}$$
(4)

while $y = ub_i - lb_i$. The mutation returns a random value within the range $[\sigma_L, \sigma_U]$ with the assurance of symmetry about the parent value $x_{i,k}^t$

$$x_{i,k}^{\prime t} = \begin{cases} x_{i,k}^t - (1 - 2p)(x_{i,k}^t - \sigma_L) & \text{if } q \le 0.5\\ x_{i,k}^t + (2p - 1)(\sigma_U - x_{i,k}^t) & \text{otherwise} \end{cases}$$
 (5)

2.4 Replacement Strategy

The replacement phase concerns the survivor selection of both the parent and the offspring populatios. Since the size of the population is constant, it allows to withdraw individuals according to a give strategy. There are two main strategies in GAs, namely generational and steady-state. Generational replacement will concern the whole population of size N. The offspring population will replace systematically the parent population. On the contrary, in steady-state replacement, at each generation, only one offspring is generated. For instance, it replaced the worst individual of the parent population. In this implementation of RCGA, the generational replacement strategy is ultilized. Besides, an elitist strategy is also used to ensure the best individual of the current population survive to the next generation.

3. INDEPENDENT RESTARTS SCHEME

An independent restarts version of the RCGA is implemented in study. For each start, the initial solutions $\{\mathbf{x}_i(0)\}$ are uniformly sampled at random within the search bound, i.e $[-5,5]^D$. Whenever the stopping conditions are met, the algorithm will be reinitialized and restarted without inheriting any information about the last run. This process is iterated until the mission is accomplished, i.e. objective function value is less than the target function value, or the total number of function evaluations surpasses 5×10^4 .

There are two stopping conditions for the independent restarts scheme. The first one is the maximum number of iterations in each run of the algorithm. This value is computed according to the search space dimensionality, viz floor $\left(100+3800D\sqrt{D}\right)$. The second condition is that the best objective function values obtained so far and during the last 50+25D generations do not vary more than 10^{-12} . It means there is no significant improvement in the population after evolving over several consecutive generations. Whenever each of these two conditions is met, the algorithm will be terminated and start from scratch.

4. PARAMETER SETTINGS

In order to implement the described RCGA, we used the population size N = 100; crossover rates $p_c = 0.95$, tournament size $n_{ts} = 2$; mutation rate $p_m = \{0.05, 0.1, 0.2\}$ which was in turn assigned to each start, and the non-uniformity degree for mutation was $\beta = 5$.

No parameter tuning has been conducted. The final parameter setting were identical on all functions and therefore the crafting effort [3] computes to CrE = 0.

5. CPU TIMING EXPERIMENTS

CPU timing experiments were conducted using the same independent restarts RCGA which was particularly run on f_8 with a maximum of $10^5 \times D$ function evaluations and restarted until at least 30 seconds had passed (according to Figure 2 in [3]). These experiments have been implemented on an Intel Core 2 Duo CPU E8400 running at 3.00 GHz under Windows XP SP2 with Matlab 7.8.0 (R2009a). The results are given in Table 3. These results show that the independent restarts implementation of the RCGA takes on approximately 30, 20, 12, 6.0, 3.0, 3.5, and 4.7 times 10^{-5} seconds per function evaluation for 2-, 3-, 5-, 10-, 20-, 40-, and 80-dimensional search space, respectively. The implemented RCGA is a CPU-inexpensive algorithm, i.e. increase in dimension, up to 40-D, has negligible influence on the CPU time. Moreover up to 80-D, a dependency of CPU time on the search space dimensionality is trivial.

6. EXPERIMENTAL RESULTS

Results from experiments according to [3] on the benchmark functions given in [2, 4] are presented in Figures 1, 2, 3 and in Tables 1, 2.

The obtained results from Figures 1 show that RCGA has encouraging performance on all types of test functions, there is at least one representative function within each type that can be solved in dimensions of 5 or even 10. Besides, it can be observed that the expected number of function evaluations to reach a given target function value often scales quadratically with the dimension on almost all functions (Figures 1). To low precision, such as 10^{-3} , the scaling is maintained or even better than quadratical on functions f_5 , f_6 , f_{14} , f_{17} , f_{21} .

Some functions that turn out to be laborious for this RCGA implementation are Skew Rastrigin-Bueche (f_4) , Rosenbrock (f_8, f_9) , Ellipsoid (f_{10}) , Discuss (f_{11}) , Lunacek bi-Rastrigin (f_{24}) . Apart from these functions, which fall into different types, performance of the algorithm with a not-so-high resolution requirement is acceptable within the given computational budget, fruitful results are accomplished on functions Sphere (f_1) , Linear slope (f_5) , Attractive sector (f_6) , Sum of different powers (f_{14}) , Schaffer F7 with condition $10 (f_{17})$, Gallagher 101 peaks (f_{21}) . No matter whether these reported results are competitive or not, the overall outcome from all experiments demonstrate that the implemented version of RCGA is unbiased towards any certain problem types.

7. CONCLUSIONS

The real-coded genetic algorithm, which is made up of the a tournament selection, an arithmetical crossover and an adaptive-range variant of non-uniform mutation, incorporated with an independent restarts mechanism has presented encouraging results when benchmarked on noiseless black-box optimization testbed. Without producing impressive results, the algorithm is unable to be considered as a competitive algorithm against other state-of-the-art evolutionary computation approaches, such as some modifications of covariance matrix adaptation evolution strategy (CMA-ES); the implemented RCGA, nevertheless, not only display potential scaling property but also exhibit the robustness on various types of problems with relatively small number of variables (less than 20). On the other hand, the multiple independent restarts induce comparatively effective search

Table 2: ERT loss ratio (see Figure 3) compared to the respective best result from BBOB-2009 for budgets given in the first column. The last row $\mathrm{RL_{US}}/\mathrm{D}$ gives the number of function evaluations in unsuccessful runs divided by dimension. Shown are the smallest, 10%-tile, 25%-tile, 50%-tile, 75%-tile and 90%-tile values (smaller values are better).

	f	1 ⁻ <i>f</i> 24	in 5-D	$_{ m maxFF}$	E/D=50	000
#FEs/D	$_{ m best}$	10%	25%	\mathbf{med}	75%	90%
2	1.2	1.6	2.0	2.6	3.7	8.3
10	1.6	3.1	3.5	5.1	8.3	50
100	4.2	6.7	8.7	17	24	96
1e3	5.0	11	29	61	89	4.2e2
1e4	22	31	81	2.9e2	4.7e2	2.3e3
1e5	23	61	1.3e2	6.7e2	2.3e3	3.6e3
RL_{US}/D	5e4	5e4	5e4	5e4	5e4	5e4
	•					
	f_1	− <i>f</i> 24 i	n 20-D	, maxFI	E/D=50	000
#FEs/D	f_1	- f 24 i	n 20-D 25%	, maxFI med	E/D=50 75%	000 90%
#FEs/D	-	-				
	best	10%	25%	med	75%	90% 40
2	best 1.0	10% 5.1	25% 11	med 31	75% 40	90% 40
2 10	best 1.0 4.8	10% 5.1 6.9	25% 11 11	med 31 51	75% 40 2.0e2	90% 40 2.0e2
2 10 100	best 1.0 4.8 6.4	10% 5.1 6.9 6.8	25% 11 11 14	med 31 51 26	75% 40 2.0e2 37	90% 40 2.0e2 2.9e2
2 10 100 1e3	best 1.0 4.8 6.4 14	10% 5.1 6.9 6.8 18	25% 11 11 14 25	med 31 51 26 60	75% 40 2.0e2 37 1.8e2	90% 40 2.0e2 2.9e2 4.7e2
2 10 100 1e3 1e4	best 1.0 4.8 6.4 14 6.6	10% 5.1 6.9 6.8 18 32	25% 11 11 14 25 62	med 31 51 26 60 2.9e2	75% 40 2.0e2 37 1.8e2 5.9e2	90% 40 2.0e2 2.9e2 4.7e2 4.7e3

Table 3: Results for CPU timing experiments. Shown are dimension, D, number of times the timing procedure is repeated within 30 seconds, n, and CPU-time per function evaluation, $t (\times 10^{-5} \text{ second})$.

D	2	3	5	10	20	40	80
\overline{n}	125574	119864	108579	78598	6378	1	1
t	30	20	12	6.0	3.0	3.5	4.7

on unstructured multi-modal landscapes. Further modifications so as to adaptively control parameters, such as crossover rate, mutation rate, non-uniform degree and/or restart strategy during the course of evolution may improve the performance and efficiency of the algorithm.

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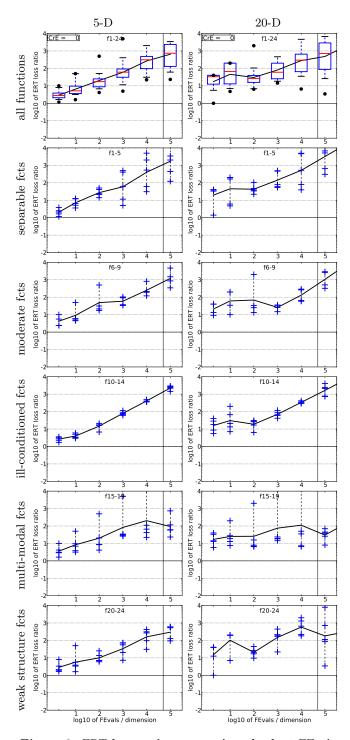


Figure 3: ERT loss ratio versus given budget FEvals. The target value $f_{\rm t}$ for ERT (see Figure 1) is the smallest (best) recorded function value such that ERT($f_{\rm t}$) \leq FEvals for the presented algorithm. Shown is FEvals divided by the respective best ERT($f_{\rm t}$) from BBOB-2009 for functions f_1-f_{24} in 5-D and 20-D. Each ERT is multiplied by $\exp({\rm CrE})$ correcting for the parameter crafting effort. Line: geometric mean. Box-Whisker error bar: 25-75%-tile with median (box), 10-90%-tile (caps), and minimum and maximum ERT loss ratio (points). The vertical line gives the maximal number of function evaluations in this function subset.

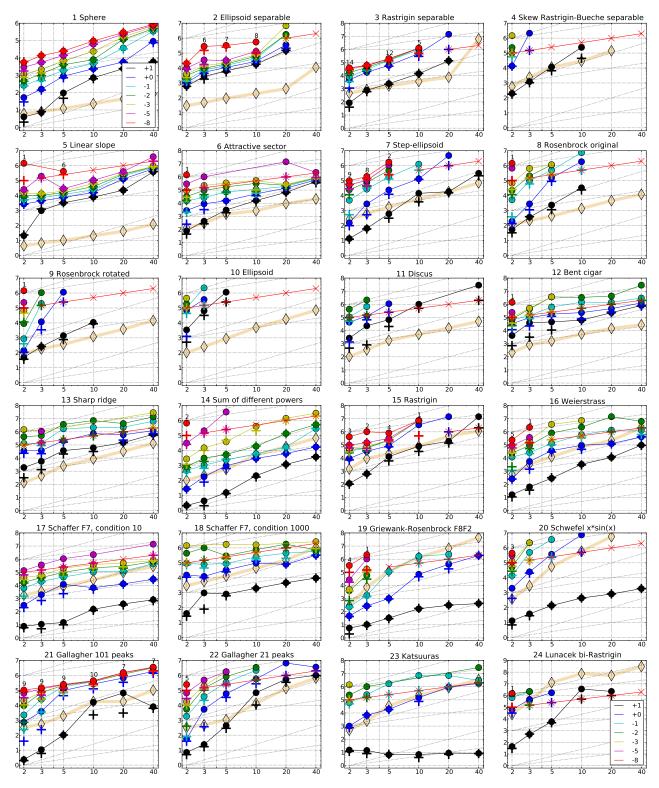


Figure 1: Expected Running Time (ERT, ullet) to reach $f_{\mathrm{opt}}+\Delta f$ and median number of f-evaluations from successful trials (+), for $\Delta f=10^{\{+1,0,-1,-2,-3,-5,-8\}}$ (the exponent is given in the legend of f_1 and f_{24}) versus dimension in log-log presentation. For each function and dimension, $\mathrm{ERT}(\Delta f)$ equals to $\#\mathrm{FEs}(\Delta f)$ divided by the number of successful trials, where a trial is successful if $f_{\mathrm{opt}}+\Delta f$ was surpassed. The $\#\mathrm{FEs}(\Delta f)$ are the total number (sum) of f-evaluations while $f_{\mathrm{opt}}+\Delta f$ was not surpassed in the trial, from all (successful and unsuccessful) trials, and f_{opt} is the optimal function value. Crosses (×) indicate the total number of f-evaluations, $\#\mathrm{FEs}(-\infty)$, divided by the number of trials. Numbers above ERT-symbols indicate the number of successful trials. Y-axis annotations are decimal logarithms. The thick light line with diamonds shows the single best results from BBOB-2009 for $\Delta f=10^{-8}$. Additional grid lines show linear and quadratic scaling.

Δf f1 in 5-D, N=15, mFE=29500 # ERT 10% 90% RT _{succ}	f1 in 20-D, N=15, mFE=292100 # ERT 10% 90% RT _{succ}	Δf f2 in 5-D, N=15, mFE=250000 Δf # ERT 10% 90% RT _{succ}	f2 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}
10 15 9.5e1 1.0e1 2.7e2 9.5e1 1 15 8.8e2 5.2e2 1.6e3 8.8e2	15 2.4e3 1.8e3 3.1e3 2.4e3 15 5.7e3 4.7e3 6.8e3 5.7e3	10 15 6.5e3 3.4e3 1.0e4 6.5e3	15 1.5e5 1.3e5 1.6e5 1.5e5 13 3.6e5 1.7e5 1.2e6 2.0e5
1e-1 15 1.9e3 1.1e3 2.8e3 1.9e3 1e-3 15 6.8e3 5.9e3 8.8e3 6.8e3	15 3.7e4 1.2e4 6.8e4 3.7e4 15 1.6e5 1.5e5 1.8e5 1.6e5	1e-1 15 1.5e4 1.2e4 1.9e4 1.5e4 $1e-3$ 15 2.3e4 1.7e4 2.6e4 2.3e4	6 1.7e6 2.1e5 4.7e6 2.2e5 2 6.8e6 7.8e5 1.3e7 2.8e5
le-5 15 1.4e4 1.2e4 1.6e4 1.4e4	$15 \ \ 2.3 \mathrm{e}5 \ \ 2.3 \mathrm{e}5 \ \ 2.4 \mathrm{e}5 \qquad 2.3 \mathrm{e}5$	1e-5 15 3.2e4 2.9e4 3.5e4 3.2e4	0 51e-2 47e-5 16e-1 3.4e5
1e-8 15 2.5e4 2.3e4 2.9e4 2.5e4 f3 in 5-D, N=15, mFE=250000		1e-8 7 3.3e5 4.0e4 7.9e5 4.0e4 f4 in 5-D , N=15, mFE=250000	f4 in 20-D, N=15, mFE=1.00e6
Δf # ERT 10% 90% RT _{SUCC} 10 15 2.5e3 1.2e3 3.3e3 2.5e3	# ERT 10% 90% RT _{succ} 15 1.4e5 1.3e5 1.6e5 1.4e5	Δf # ERT 10% 90% RT _{SUCC} 10 15 1.2e4 4.6e3 3.1e4 1.2e4	# ERT 10% 90% RT _{SUCC} 0 22e+0 18e+0 26e+0 2.0e5
1 15 6.0 e4 4.0 e4 1.0 e5 6.0 e4	1 1.4e7 2.2e6 3.2e7 2.4e5	1 0 $30e-1$ $20e-1$ $40e-1$ $1.1e5$	
1e-1 13 1.5e5 4.7e4 3.1e5 1.1e5 $1e-3$ 12 1.9e5 5.8e4 4.4e5 1.3e5	0 41e-1 10e-1 60e-1 3.3e5	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	
1e-5 12 1.9e5 6.5e4 3.9e5 1.3e5 1e-8 12 2.0e5 7.2e4 4.0e5 1.4e5		1e-5	
f5 in 5-D, N=15, mFE=250000	f5 in 20-D, N=15, mFE=1.00e6	f6 in 5-D, N=15, mFE=250000	f6 in 20-D, N=15, mFE=1.00e6
Δf # ERT 10% 90% RT _{SUCC} 10 15 3.3e3 2.9e3 3.8e3 3.3e3	# ERT 10% 90% RT _{SUCC} 15 2.3e4 2.0e4 2.5e4 2.3e4		# ERT 10% 90% RT _{SUCC} 15 6.7e4 5.3e4 8.0e4 6.7e4
1 15 7.5e3 6.5e3 8.4e3 7.5e3 1e-1 15 1.1e4 1.0e4 1.3e4 1.1e4	15 1.3e5 1.2e5 1.4e5 1.3e5 15 2.1e5 2.0e5 2.2e5 2.1e5	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	15 1.4e5 1.3e5 1.5e5 1.4e5 15 1.9e5 1.9e5 2.0e5 1.9e5
1e-3 15 2.0e4 1.8e4 2.2e4 2.0e4	15 2.9e5 2.9e5 3.0e5 2.9e5	1e-3 7 $3.5e5$ $3.5e4$ $8.3e5$ $6.1e4$	$15\ \ 2.7e5\ \ 2.6e5\ \ 2.8e5\ \ \ \ 2.7e5$
1e-5 15 2.9e4 2.6e4 3.1e4 2.9e4 1e-8 6 4.1e5 3.8e4 6.6e5 3.9e4	14 4.0e5 3.3e5 1.3e6 3.3e5 0 64e-7 33e-7 97e-7 3.4e5	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	1 1.4e7 1.3e6 3.2e7 3.2e5 0 63e-6 16e-6 28e-5 3.4e5
Δf in 5-D, N=15, mFE=250000 Δf # ERT 10% 90% RT _{succ}	f7 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}	Δf # ERT 10% 90% RT _{succ}	f8 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}
10 15 6.0 e2 2.0 e2 8.8 e2 6.0 e2	15 1.8e4 1.2e4 3.4e4 1.8e4	10 15 2.3e3 7.7e2 7.1e3 2.3e3	$\begin{array}{cccccccccccccccccccccccccccccccccccc$
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	3 4.7e6 8.7e5 1.1e7 7.5e5 0 13e-1 81e-2 19e-1 4.9e5	1 11 1.5e5 1.0e4 3.9e5 6.2e4 1e-1 6 4.9e5 1.8e4 1.1e6 1.1e5	
1e-3 3 1.1e6 8.3e4 2.7e6 8.5e4 1e-5 3 1.1e6 7.5e4 2.8e6 8.5e4		$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	
1e-8 2 1.7e6 1.0e5 4.6e6 9.0e4		1e-8	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	f9 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}	Δf in 5-D, N=15, mFE=250000 Δf # ERT 10% 90% RT _{succ}	f10 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}
10 15 1.4e3 6.3e2 3.1e3 1.4e3 1 3 1.1e6 1.6e5 2.8e6 1.4e5	0 17e+0 16e+0 18e+0 3.4e5	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0 48e+2 25e+2 82e+2 3.4e5
1e-1 0 17e-1 20e-2 24e-1 1.5e5 1e-3		1e-1	
1e-5		1e-5	
1e-8	f ₁₁ in 20-D, N=15, mFE=1.00 e6	1e-8	
Δf # ERT 10% 90% RT _{SUCC} 10 13 6.7e4 3.4e2 2.5e5 2.9e4	# ERT 10% 90% RT _{SUCC} 0 51e+0 13e+0 64e+0 4.4e5	Δf # ERT 10% 90% RT _{SUCC} 10 14 4.6e4 8.7e3 1.5e5 2.8e4	# ERT 10% 90% RT _{SUCC} 15 2.3e5 2.2e5 2.4e5 2.3e5
1 3 1.1e6 1.3e5 2.2e6 1.0e5		1 11 1.8e5 1.3e4 5.1e5 8.7e4	13 4.7e5 2.5e5 1.0e6 3.1e5
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$		1e-1 5 6.2e5 8.3e4 1.3e6 1.2e5 1e-3 1 3.7e6 4.3e5 9.9e6 1.8e5	7 1.4e6 2.7e5 3.3e6 2.8e5 0 12e-2 45e-4 13e-1 3.4e5
1e-5		1e-5 0 29e-2 10e-3 94e-1 1.7e5 1e-8	
Δf # ERT 10% 90% RT _{succ}	f13 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}	Δf # ERT 10% 90% RT _{succ}	f14 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}
10 15 3.7e4 7.3e3 9.5e4 3.7e4 1 11 2.1e5 2.9e4 4.4e5 1.2e5	15 1.5e5 1.4e5 1.6e5 1.5e5 11 6.0e5 2.2e5 1.2e6 2.4e5	10 15 1.5e1 4.0e0 2.6e1 1.5e1 1 15 9.6e2 4.1e2 2.2e3 9.6e2	15 1.2e3 8.0e2 1.5e3 1.2e3 15 6.1e3 5.3e3 6.7e3 6.1e3
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	5 2.3e6 2.7e5 6.3e6 2.7e5 0 28e-2 37e-4 22e-1 3.4e5	1e-1 15 2.3e3 7.9e2 3.8e3 2.3e3 1e-3 15 3.8e4 7.2e3 7.5e4 3.8e4	15 1.7e4 1.0e4 3.3e4 1.7e4 7 1.4e6 2.5e5 3.3e6 2.6e5
1e-5		1e-5 1 $3.7e6$ $4.9e5$ $9.9e6$ $2.4e5$	0 11e-4 69e-5 15e-4 3.4e5
1e-8	f ₁₅ in 20-D, N=15, mFE=1.00e6	1e-8 0 90e-6 16e-6 18e-5 1.9e5 f16 in 5-D, N=15, mFE=250000	f16 in 20-D, N=15, mFE=1.00e6
Δf # ERT 10% 90% RT _{SUCC} 10 15 1.3 e4 1.2 e3 4.6 e4 1.3 e4	# ERT 10% 90% RT _{SUCC} 15 1.9e5 1.2e5 1.6e5 1.9e5	Δf # ERT 10% 90% RT _{SUCC} 10 15 2.9e2 3.4e1 6.2e2 2.9e2	# ERT 10% 90% RT _{SUCC} 15 1.1e4 6.7e3 1.5e4 1.1e4
1 15 8.8e4 5.4e4 1.5e5 8.8e4	1 1.4e7 1.2e6 3.0e7 2.4e5	1 15 3.7e4 1.1e4 5.4e4 3.7e4	15 1.3e5 8.6e4 1.7e5 1.3e5
1e-1 11 2.2e5 8.3e4 4.4e5 1.3e5 1e-3 11 2.3e5 6.6e4 4.6e5 1.4e5	0 50e-1 30e-1 77e-1 3.4e5 	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	7 1.4e6 2.1e5 3.2e6 2.2e5 0 12e-2 18e-3 43e-2 3.4e5
1e-5 9 3.2e5 7.6e4 6.2e5 1.5e5 1e-8 4 8.1e5 8.3e4 2.0e6 1.2e5		1e-5 0 $16e-3$ $15e-4$ $27e-2$ $1.1e5$ $1e-8$	
Δf f17 in 5-D, N=15, mFE=250000 Δf # ERT 10% 90% RT _{succ}	f17 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}	Δf # ERT 10% 90% RT _{succ}	f18 in 20-D, N=15, mFE=1.00e6 # ERT 10% 90% RT _{succ}
10 15 1.5e1 3.0e0 3.9e1 1.5e1	15 3.9e2 1.4e2 8.4e2 3.9e2	10 15 8.1e2 2.3e2 1.8e3 8.1e2	15 4.7e3 2.9e3 7.2e3 4.7e3
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	15 1.3e4 8.9e3 2.0e4 1.3e4 15 1.2e5 9.6e4 1.4e5 1.2e5	1 15 3.3e4 3.8e3 6.7e4 3.3e4 1e-1 15 9.0e4 5.1e4 1.4e5 9.0e4	15 7.9e4 3.4e4 1.2e5 7.9e4 12 4.5e5 1.7e5 1.2e6 2.0e5
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	12 5.1e5 2.5e5 1.3e6 2.6e5 0 12e-5 16e-6 78e-4 3.4e5	1e-3 2 1.7e6 1.2e5 3.4e6 1.1e5 $1e-5$ 0 $53e-4$ $77e-5$ $38e-3$ 1.1e5	0 16e-3 34e-4 18e-2 3.4e5
1e-8 0 30e-5 32e-7 38e-4 8.5 e4 f19 in 5-D, N=15, mFE=250000	f19 in 20-D, N=15, mFE=1.00e6	1e-8	f20 in 20-D, N=15, mFE=1.00e6
Δf # ERT 10% 90% RT _{succ}	# ERT 10% 90% RT _{succ}	Δf # ERT 10% 90% RT _{succ}	# ERT 10% 90% RT _{succ}
10 15 3.3e1 5.0e0 5.1e1 3.3e1 1 15 1.0e3 4.1e2 1.7e3 1.0e3	15 3.1e2 2.1e2 3.5e2 3.1e2 15 4.1e5 5.5e4 9.9e5 4.1e5	10	15 8.2 e2 5.3 e2 1.1 e3 8.2 e2 0 22e-1 16e-1 28e-1 3.7 e5
1e-1 15 1.1e5 2.5e4 2.1e5 1.1e5 1e-3 0 63e-3 22e-3 96e-3 1.5e5	5 2.4e6 1.9e5 5.1e6 3.7e5 0 23e-2 65e-3 77e-2 3.4e5	$ \begin{array}{ccccccccccccccccccccccccccccccccc$	
1e-5		1e-5	
1e-8	f21 in 20-D, N=15, mFE=1.00e6	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	f22 in 20-D, N=15, mFE=1.00e6
Δf # ERT 10% 90% RT _{succ} 10 15 1.0 e2 2.0 e1 1.8 e2 1.0 e2	# ERT 10% 90% RT _{Succ} 15 6.7e4 2.3e3 1.1e4 6.7e4	Δf # ERT 10% 90% RT _{Succ} 10 15 4.4e2 3.8e1 6.5e2 4.4e2	# ERT 10% 90% RT _{succ} 11 5.5e5 2.9e3 1.3e6 1.9e5
1 12 9.7 e4 2.1 e3 3.1 e5 3.5 e4	10 6.7e5 5.8e3 2.0e6 1.7e5	1 14 6.0 e4 6.4 e2 1.9 e5 4.2 e4	2 6.7e6 7.0e5 1.5e7 2.0e5
1e-3 9 2.2e5 1.7e4 5.8e5 5.4e4	8 1.2e6 1.2e4 3.2e6 3.3e5 8 1.3e6 4.9e4 2.9e6 3.8e5	1e-3 2 1.7e6 2.0e4 4.3e6 8.5e4	0 73e-1 69e-2 11e+0 3.2e5
1e-5 9 2.3e5 2.3e4 5.5e5 6.7e4 1e-8 9 2.5e5 6.0e4 5.7e5 8.6e4	7 1.5e6 1.2e5 3.7e6 3.6e5 7 1.6e6 2.0e5 3.8e6 4.1e5	1e-5 2 1.8e6 1.7e5 3.7e6 1.5e5 1e-8 0 90e-3 79e-7 73e-2 1.5e5	
f23 in 5-D, N=15, mFE=250000	f23 in 20-D, N=15, mFE=1.00e6	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	f_{24} in 20-D, N=15, mFE=1.00 e6 # ERT 10% 90% RT _{Succ}
10 15 6.5e0 1.0e0 1.3e1 6.5e0	15 8.7e0 3.0e0 2.4e1 8.7e0	10 15 5.7e3 2.2e3 9.6e3 5.7e3	6 2.3e6 6.8e5 5.4e6 7.6e5
1 15 1.9e4 5.6e3 4.6e4 1.9e4 1e-1 2 1.7e6 7.5e4 3.9e6 7.3e4	15 9.2e5 9.6e5 9.9e5 9.2e5 2 7.1e6 9.9e5 1.6e7 6.2e5	le-1 0 55e-1 82e-2 68e-1 6.9e4	0 25e+0 38e-1 60e+0 1.0e6
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	0 26e-2 94e-3 46e-2 1.0e6 	1e-3	
1e-8		1e-8	1

Table 1: Shown are, for a given target difference to the optimal function value Δf : the number of successful trials (#); the expected running time to surpass $f_{\rm opt}+\Delta f$ (ERT, see Figure 1); the 10%-tile and 90%-tile of the bootstrap distribution of ERT; the average number of function evaluations in successful trials or, if none was successful, as last entry the median number of function evaluations to reach the best function value (RT_{succ}). If $f_{\rm opt}+\Delta f$ was never reached, figures in *italics* denote the best achieved Δf -value of the median trial and the 10% and 90%-tile trial. Furthermore, N denotes the number of trials, and mFE denotes the maximum of number of function evaluations executed in one trial. See Figure 1 for the names of functions.

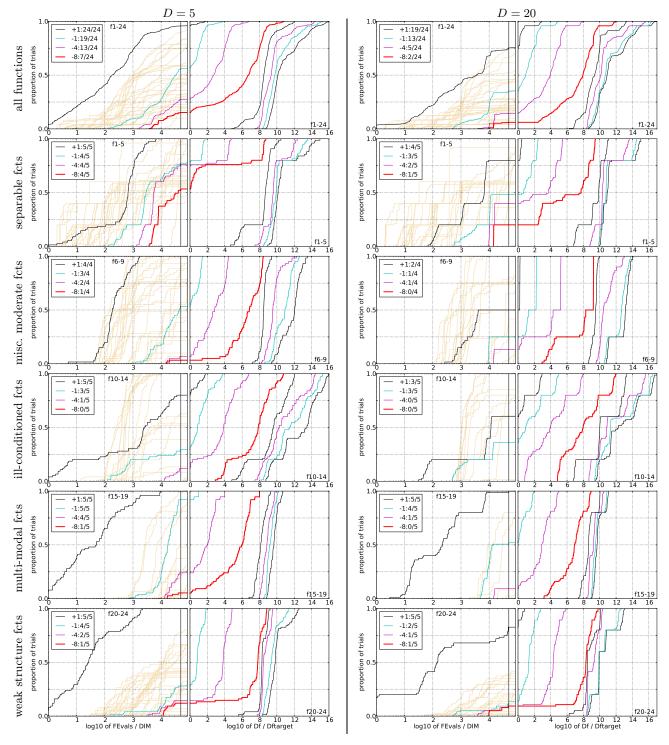


Figure 2: Empirical cumulative distribution functions (ECDFs), plotting the fraction of trials versus running time (left subplots) or versus Δf (right subplots). The thick red line represents the best achieved results. Left subplots: ECDF of the running time (number of function evaluations), divided by search space dimension D, to fall below $f_{\rm opt} + \Delta f$ with $\Delta f = 10^k$, where k is the first value in the legend. Right subplots: ECDF of the best achieved Δf divided by 10^k (upper left lines in continuation of the left subplot), and best achieved Δf divided by 10^{-8} for running times of $D, 10\,D, 100\,D\dots$ function evaluations (from right to left cycling black-cyan-magenta). The legends indicate the number of functions that were solved in at least one trial. FEvals denotes number of function evaluations, D and DIM denote search space dimension, and Δf and Df denote the difference to the optimal function value. Light brown lines in the background show ECDFs for target value 10^{-8} of all algorithms benchmarked during BBOB-2009.

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