

What is induction? It refers to two different definitions. One is about the composition of the set of natural numbers \mathbb{N} . The other is a method of proof of mathematical statements. Lets start by the first one.

Set of Natural Numbers

Definition: Set of Natural Numbers

The set of Natural Numbers is the smallest set such that,

- $0 \in \mathbb{N}$
- $n \in \mathbb{N} \rightarrow n + 1 \in \mathbb{N}$

This definition contains 3 of the 5 ideas behind Peano axioms. They are,

- $0 \in \mathbb{N}$
- $n \in \mathbb{N} \rightarrow n + 1 \in \mathbb{N}$
- $n + 1 = 0 \rightarrow n \notin \mathbb{N}$
- $n + 1 = m + 1 \rightarrow n = m$
- $(0 \in A \wedge n \in A \rightarrow n + 1 \in A) \rightarrow A \subseteq \mathbb{N}$

Proof by Induction

To prove a family of statements $\forall n \in \mathbb{N}: P(n)$, we can use proof by induction.

Definition: Principle of Mathematical Induction

By proving the *initial case* $P(0)$ and the *induction step* $P(k) \rightarrow P(k + 1)$ we conclude $\forall n \in \mathbb{N}: P(n)$. Formally,

$$(P(0) \wedge P(k) \rightarrow P(k + 1)) \rightarrow \forall n \in \mathbb{N}: P(n)$$

Example 1 Intro to Uni Math Sheet 1: Upper bounded sum and the lower bounded product

Let $x_1 + x_2 + \dots + x_n \leq \frac{1}{3}$. Show that $(1 - x_1)(1 - x_2) \dots (1 - x_n) \geq \frac{2}{3}$

Doubt

What is the trick here?

Example : Sum of first n natural numbers

We want to prove the following,

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

Initial Case: $n=1$

$$\sum_{i=1}^1 i = 1 = \frac{1(2)}{2}$$

□

Induction step:

$$\left(\sum_{i=1}^n i = \frac{n(n+1)}{2}\right) \rightarrow \left(\sum_{i=1}^{n+1} i = \frac{(n+1)(n+2)}{2}\right)$$

To prove the induction step,

$$\begin{aligned}\sum_{i=1}^{n+1} i &= (n+1) + \sum_{i=1}^n i \\ &= (n+1) + \frac{n(n+1)}{2} \\ &= \frac{2n+2}{2} + \frac{n(n+1)}{2} \\ &= \frac{2n+2+n^2+n}{2} \\ &= \frac{n^2+3n+2}{2} \\ &= \frac{(n+1)(n+2)}{2}\end{aligned}$$

□

Using the induction hypothesis we get

We can also reduce the scope of the statement to a subset of natural numbers. Instead of the initial case being $P(0)$, we can start with the property $P(k)$ of a given natural number k .

Principle: Advice on proving a statement by Induction

Always try to decompose conclusion into induction hypothesis and other term.

Strong Induction

The induction hypothesis is now,

$$(P(0) \wedge P(1) \wedge \dots \wedge P(k)) \rightarrow P(k+1)$$

Example University of Illinois: Proof Recurrence Relation by Strong Induction

Let a_n be a sequence where $a_1 = 1$ and $a_2 = 8$ and $a_n = a_{n-1} + 2a_{n-2}$. We want to prove that,

$$a_n = 3 \cdot 2^{n-1} + 2(-1)^n$$

We prove by induction on n . Initial case: $n = 3$

$$a_3 = 10$$

□

Inductive Step: $3 \leq n \leq k \rightarrow a_n = 3 \cdot 2^{n-1} + 2(-1)^n$

$$\begin{aligned}
 a_{n+1} &= a_n + 2a_{n-1} \\
 &= 3 \cdot 2^{n-1} + 2(-1)^n + 2(3 \cdot 2^{n-2} + 2(-1)^{n-1}) \\
 &= 2(3 \cdot 2^{n-1}) + 2(-1)^n + 2^2(-1)^{n-1} \\
 &= 3 \cdot 2^n + 2(-1)^{n-1}(-1 + 2) \\
 &= 3 \cdot 2^n + 2(-1)^{n-1} \\
 &= 3 \cdot 2^n + 2(-1)^{n+1} \\
 &\square
 \end{aligned}$$

We can use the induction hypothesis twice

Principle: Strong Induction

We notice that weak induction is not enough when we need more than a single hypothesis in the induction step.

Forward-Backward Induction

We have the same 2 steps in a regular proof by induction, the initial case and the inductive step, but the inductive step is composed by two different steps.

- Initial case: $P(0)$
- Inductive step:
 - Forward step: $P(k) \rightarrow P(f(k))$ where f is an increasing function.
 - Backward step: $P(k) \rightarrow P(k-1)$

Quest: More patterns in forward backward induction

Find an application of a proof that requires more than one forward and/or backward step. Example:

- $f_1(k) = 2^k$
- $f_2(k) = 2^k - 1$
- $b(k) = k - 2$

Example brilliant.org: AM-GM Inequality

Initial Case: $\frac{a+b}{2} \geq \sqrt{ab}$

$$\begin{aligned}
 (a-b)^2 &\geq 0 \leftrightarrow \\
 a^2 - 2ab + b^2 &\geq 0 \leftrightarrow \\
 a^2 + 2ab + b^2 - 4ab &\geq 0 \leftrightarrow \\
 a^2 + 2ab + b^2 &\geq 4ab \leftrightarrow \\
 a^2 + 2ab + b^2 &\geq 4ab \leftrightarrow \\
 (a+b)^2 &\geq 4ab \leftrightarrow \\
 |a+b| &\geq 2\sqrt{ab} \leftrightarrow \\
 \rightarrow a+b &\geq 2\sqrt{ab} \\
 \rightarrow \frac{a+b}{2} &\geq \sqrt{ab}
 \end{aligned}$$

Since both a and b are positive so is their addition

Inductive Step:

Induction Hypothesis:

$$\frac{\sum_{i=1}^k a_i}{k} \geq \sqrt[k]{\prod_{i=1}^k a_i}$$

Forward pass: We want to show,

$$\frac{\sum_{i=1}^{2k} a_i}{2k} \geq \sqrt[2k]{\prod_{i=1}^{2k} a_i}$$

We can start by splitting the summation in the left hand side,

$$\begin{aligned} a_1 + a_2 + \cdots + a_{2k} &= \frac{a_1 + a_2 + \cdots + a_k}{k} + \frac{a_{k+1} + a_{k+2} + \cdots + a_{2k}}{k} \\ &\geq \frac{\sqrt[k]{a_1 a_2 \cdots a_k} + \sqrt[k]{a_{k+1} a_{k+2} \cdots a_{2k}}}{2} \\ &\geq \sqrt{\sqrt[k]{a_1 a_2 \cdots a_k} \sqrt[k]{a_{k+1} a_{k+2} \cdots a_{2k}}} \\ &= \sqrt[2k]{a_1 a_2 \cdots a_{2k}} \end{aligned}$$

□

This completes the proof for the forward pass.

Backward pass: We want to show,

$$\frac{\sum_{i=1}^{k-1} a_i}{k-1} \geq \sqrt[k-1]{\prod_{i=1}^{k-1} a_i}$$

We start by using the inductive hypothesis,

$$\begin{aligned} \frac{a_1 + a_2 + \cdots + a_k}{k} &\geq \sqrt[k]{a_1 a_2 \cdots a_k} \\ \frac{a_1 + a_2 + \cdots + a_{k-1} + \frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1}}{k} &\geq \sqrt[k]{a_1 a_2 \cdots a_{k-1} \cdot \frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1}} \\ \frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1} &\geq \sqrt[k]{a_1 a_2 \cdots a_{k-1} \cdot \frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1}} \\ \left(\frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1} \right)^k &\geq a_1 a_2 \cdots a_{k-1} \cdot \frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1} \\ \left(\frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1} \right)^{k-1} &\geq a_1 a_2 \cdots a_{k-1} \\ \frac{a_1 + a_2 + \cdots + a_{k-1}}{k-1} &\geq \sqrt[k-1]{a_1 a_2 \cdots a_{k-1}} \end{aligned}$$

□

This completes the backward pass and the proof of the AM-GM Inequality.

Double Induction

We start by defining the addition of two natural numbers recursively,

Definition: Recursive Definition of Addition of Two Natural Numbers

Let m and n be natural numbers.

- $m+0=m$
- $m+(n+1)=(m+n)+1$

Example Intro to Uni Math ex.2: Commutativity of Addition of Natural Numbers

We want to show that the addition is commutative using its recursive definition,

$$m + n = n + m$$

Let $n = 0$. We want to show that $m + 0 = 0 + m$. We show by induction on m . The base case is skipped. We assume $k + 0 = 0 + k$. We want to prove $0 + (k + 1) = (k + 1) + 0$

$0 + (k + 1) = (0 + k) + 1$	Recursive Part of Definition
$= (k + 0) + 1$	Inductive Hypothesis
$= k + 1$	Base Case for Recursive Definition LR
$= (k + 1) + 0$	Base Case for Recursive Definition RL

Let $n = 1$. We want to show that $m + 1 = 1 + m$. We show by induction on m . Base case $0 + 1 = 1 + 0$ was already proven when $n = 0 \wedge m = 1$. We assume $k + 1 = 1 + k$. We want to prove $1 + (k + 1) = (k + 1) + 1$

$(k + 1) + 1 = (1 + k) + 1$	Inductive Hypothesis
$= 1 + (k + 1)$	Recursive Part of Definition RL
□	

We want to show that $m + n = n + m$. We prove by induction on n . The base case was already shown. We assume $m + k = k + m$. We want to prove $m + (k + 1) = (k + 1) + m$

$(m + k) + 1 = (k + m) + 1$	Inductive Hypothesis
$= 1 + (k + m)$	Using (ii)
$= (1 + k) + m$	Recursive Part of Definition
$= (k + 1) + m$	Using (ii)
□	

Infinite Descent

Example Wikipedia: Proof by Infinite Descent

The square root of a non-integer is always irrational, formally,

$$k \notin \mathbb{N} \rightarrow \sqrt{k} \notin \mathbb{R} - \mathbb{Q}$$

Let \sqrt{k} be a rational number and q the last integer before \sqrt{k} ,

$$\begin{aligned} \sqrt{k} \in \mathbb{Q} &\leftrightarrow \sqrt{k} = \frac{m}{n}, \quad m, n \in \mathbb{N} \\ q \in \mathbb{N} \wedge q &< \sqrt{k} \wedge q + 1 > \sqrt{k} \end{aligned}$$

We start by describing k ,

$$\begin{aligned}\sqrt{k} &= \frac{m}{n} \\ &= \frac{m(\sqrt{k} - q)}{n(\sqrt{k} - q)} \\ &= \frac{m\sqrt{k} - mq}{n\sqrt{k} - nq} \\ &= \frac{n\sqrt{k}\sqrt{k} - mq}{n\frac{m}{n} - nq} \\ &= \frac{nk - mq}{m - nq}\end{aligned}$$

□

Since there is an irreducible fraction for every rational number then the last expression is a contraction.

Structural Induction

The power set of a set A can be defined recursively as follows

- $\emptyset \in \mathbb{P}(A)$
- $B \in \mathbb{P}(A) \wedge x \in A \leftrightarrow B \cup \{x\} \in \mathbb{P}(A)$

Doubt: Recursive Definition of Power Set

Does the previous definition qualify as a recursive one?

Tricks

TODO