

The Davis-Hodge Isomorphism

Functorial Equivalence of Cohomology and Translator Cycles on Information Manifolds

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Abstract

We prove a conditional theorem: **The Hodge Conjecture holds for Davis Manifolds.** Every cohomology class (topological feature) on an Information Manifold M admits a realization as a rational linear combination of Translator Cycles via the functor $G_S : \text{SamGeom}^0(S) \rightarrow \text{SamTrans}^0(S)$.

The classical Hodge Conjecture asks whether every harmonic form on a projective algebraic variety is a rational combination of algebraic cycles. We translate this into the Davis Framework: algebraic cycles become **Translator Realizations**—functional, compositional structures that patch the manifold together.

The key insight: if a topological feature exists in M , it must be detectable by the Davis-Wilson Map $\Gamma(A)$. By the Error Budget Transfer Theorem, any detectable feature in M has a corresponding realization in the Translator System $T = G_S(M)$. Therefore, every cohomology class corresponds to a Translator Cycle.

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1 Introduction

1.1 The Classical Hodge Conjecture

The Hodge Conjecture, one of the seven Millennium Prize Problems, asks:

Conjecture 1.1 (Hodge, 1950). *Let X be a non-singular complex projective algebraic variety. Then every Hodge class on X is a rational linear combination of classes of algebraic cycles.*

In plain terms: every “hole” (topological feature) in a smooth algebraic shape can be described by polynomial equations.

1.2 The Davis Translation

We propose a geometric reformulation using the Davis Framework:

- **Projective variety $X \rightarrow$ Davis Manifold M** (Information Manifold with semantic structure)
- **Hodge class \rightarrow Cohomology class** (topological feature, e.g., winding number r)
- **Algebraic cycle \rightarrow Translator Cycle** (functional realization via chart-transition maps)

The question becomes: Can every topological feature of an Information Manifold be realized functorially?

1.3 Main Result

Hypothesis Block (Davis Framework Assumptions):

- (DF1) **Smooth Chartability:** The manifold M admits a smooth atlas compatible with the stitching functor F_S
- (DF2) **Functorial Equivalence:** The functors F_S, G_S are quasi-inverses establishing ε -equivalence between $\text{SamTrans}^0(S)$ and $\text{SamGeom}^0(S)$
- (DF3) **Error Budget Transfer:** Detectable features in M are preserved in $T = G_S(M)$ up to controlled slack
- (DF4) **Davis-Wilson Detectability:** Non-trivial cohomology classes are witnessed by the Davis-Wilson Map

Theorem 1.2 (Conditional: Davis-Hodge Isomorphism). ***Assume (DF1)–(DF4).** Let M be a Davis Manifold satisfying the smooth chartability condition. Let $H^{p,p}(M, \mathbb{Q})$ denote the rational Hodge classes. Then:*

$$H^{p,p}(M, \mathbb{Q}) \cong \text{TransCyc}^p(G_S(M), \mathbb{Q}) \quad (1)$$

where TransCyc^p denotes degree- p Translator Cycles in the unfolded Translator Graph $T = G_S(M)$.

Remark 1.3. This is **not** a derivation of the classical Hodge Conjecture from known theorems of algebraic geometry. It is a **conditional theorem**: if the Davis Framework axioms hold (smooth chartability, functorial equivalence, error budget transfer), then the Hodge Conjecture holds for Davis Manifolds. The result depends on the foundational program of *The Geometry of Sameness*, not on standard Hodge theory.

2 The Davis Framework

We recall the key structures from *The Geometry of Sameness*.

2.1 Semantic Sameness Structure

Definition 2.1 (Semantic Sameness Structure). A **semantic sameness structure** S is a hidden space of entities and their relationships—the “globe” that underlies all observations.

2.2 Two Categories of Realizations

Definition 2.2 (Translation-Based Realization). $\text{SamTrans}(S)$ is the category whose objects are systems on heterogeneous observation spaces with bounded translator drift.

Definition 2.3 (Manifold-Based Realization). $\text{SamGeom}(S)$ is the category whose objects are Riemannian manifolds with path families and configuration margins.

2.3 The Functors

Definition 2.4 (The Stitching Functor F_S). $F_S : \text{SamTrans}^0(S) \rightarrow \text{SamGeom}^0(S)$ “stitches the maps into a globe”—it takes a translator system and constructs the underlying manifold.

Definition 2.5 (The Unfolding Functor G_S). $G_S : \text{SamGeom}^0(S) \rightarrow \text{SamTrans}^0(S)$ “unfolds the globe back into maps”—it takes a manifold and produces the translator graph.

2.4 The Key Construction

Proposition 2.6 (Chart-Transition Translators). For a manifold M with atlas $\{(U_i, \psi_i)\}$, the Translator Graph $T^M = G_S(M)$ has edges:

$$T_{ij}^M = \chi_j^M \circ \psi_i^M \quad (2)$$

These are the chart-transition maps—functional (algebraic) objects that patch the manifold together.

3 Topological Features and Detectability

3.1 The Davis-Wilson Map

Definition 3.1 (Davis-Wilson Map). For a configuration space \mathcal{A}/\mathcal{G} , the Davis-Wilson Map is:

$$\Gamma(A) = (\Phi(A), r(A)) \in \mathbb{R}^{d_\Phi} \times \mathbb{Z} \quad (3)$$

where Φ encodes local curvature (Wilson loop traces) and r encodes global topology (winding number).

3.2 Detectability Principle

Proposition 3.2 (Topological Detectability). *If a topological feature exists in M (e.g., a non-trivial cohomology class $[c] \in H^p(M)$), it must be detectable by the Davis-Wilson Map:*

$$[c] \neq 0 \implies \exists A, A' \text{ such that } \Gamma(A) \neq \Gamma(A') \text{ witness } [c] \quad (4)$$

Remark 3.3. *This is analogous to: “If there’s a hole in the manifold, you can detect it by measuring how loops fail to shrink.”*

4 The Error Budget Transfer Theorem

The key tool is the Error Budget Transfer Theorem from *The Geometry of Sameness*.

Theorem 4.1 (Error Budget Transfer). *For $T \in \text{SamTrans}^0(S)$ and $M = F_S(T)$:*

$$|E_{\text{tot}}^{\text{TSP}}(T) - E_{\text{tot}}^{\text{MSP}}(M)| \leq C_F(\varepsilon_{\text{trans}}, \varepsilon_{\text{dist}}, \delta_{\text{chart}}) \quad (5)$$

Any feature detectable in one realization is detectable in the other, up to controlled slack.

Corollary 4.2 (Feature Preservation). *If $[c] \in H^p(M)$ is detected by the Davis-Wilson Map on M , then $G_S([c])$ defines a corresponding cycle in the Translator Graph $T = G_S(M)$.*

5 Proof of the Davis-Hodge Isomorphism

Proof of Theorem 1.2 (Conditional on DF1–DF4). We construct the isomorphism in both directions, **assuming the Davis Framework hypotheses**.

Step 1: Cohomology \rightarrow Translator Cycles.

Let $[c] \in H^{p,p}(M, \mathbb{Q})$ be a rational Hodge class. By Proposition 3.2, $[c]$ is detected by the Davis-Wilson Map—there exist configurations whose cache values witness the non-triviality of $[c]$.

By the Error Budget Transfer Theorem 4.1, any detectable feature in M has a corresponding realization in $T = G_S(M)$ up to ε -slack.

The functor G_S explicitly constructs this realization: the cohomology class $[c]$ is supported by a specific configuration of chart-transition translators $\{T_{ij}^M\}$. These form a **Translator Cycle** $\gamma_c \in \text{TransCyc}^p(T, \mathbb{Q})$.

Step 2: Translator Cycles \rightarrow Cohomology.

Conversely, let $\gamma \in \text{TransCyc}^p(T, \mathbb{Q})$ be a Translator Cycle. By the stitching functor F_S , the Translator Graph assembles into a manifold $M = F_S(T)$.

The cycle γ corresponds to a closed chain of chart-transition maps. On the stitched manifold, this defines a homology class $[\gamma] \in H_p(M, \mathbb{Q})$, and by Poincaré duality, a cohomology class in $H^{n-p}(M, \mathbb{Q})$.

For $p = n/2$ (middle dimension), this gives a Hodge class in $H^{p,p}(M, \mathbb{Q})$.

Step 3: Functoriality (assuming DF2).

The maps $[c] \mapsto \gamma_c$ and $\gamma \mapsto [c_\gamma]$ are inverses up to ε -equivalence. This follows from hypothesis (DF2): the functors F_S, G_S are quasi-inverses. On well-behaved subcategories satisfying (DF1), they become exact inverses.

Note: Each step invokes results from *The Geometry of Sameness*. This proof does not derive Hodge from scratch; it shows that **within the Davis Framework**, the Hodge property is automatic. \square

6 Plain English: The Library Analogy

Remark 6.1 (The Library Analogy). *The Hodge Conjecture in library terms:*

The manifold M is a library. Cohomology classes are “holes”—topological features like “this library has a reading room with a courtyard in the middle.”

The question: Can every architectural feature be described by the floor plan—the way rooms (charts) connect to each other via doorways (translators)?

Our answer: Yes. The functor G_S unfolds the library into a graph of rooms and doorways. Every hole in the library corresponds to a cycle of doorways—a path that goes through a sequence of rooms and returns to the start, encircling the hole.

The Davis-Wilson Map is the surveyor’s instrument: it detects holes by measuring how loops fail to shrink. If a hole exists, the surveyor will find it, and then the functor G_S will express it as a doorway-cycle.

That’s the Hodge Conjecture for Information Manifolds.

7 What Remains

The Davis-Hodge Isomorphism is conditional on:

1. **Smooth Chartability:** The manifold M must admit a smooth atlas compatible with the functor F_S . This is the “globe exists” condition.
2. **Rationality:** The theorem applies to rational Hodge classes. Extension to integral classes requires additional control on torsion.
3. **Projective Structure:** Classical Hodge requires M to be projective algebraic. For general Davis Manifolds, we need to specify the analogous structure (Kähler, symplectic, etc.).
4. **Dimension Matching:** The isomorphism in Step 2 uses Poincaré duality, which requires compact, oriented manifolds.

Remark 7.1 (Experimental Validation). *The Davis-Wilson Map is computable on lattice simulations. This provides an **experimental test**: compute cohomology via topological charge distribution, then verify that every observed class admits a Translator Cycle representation via G_S .*

8 Conclusion

The Hodge Conjecture asks whether topology (holes) can be captured by algebra (polynomial cycles).

The Davis Framework provides a **conditional answer**: for Information Manifolds satisfying the Davis axioms (DF1–DF4), every cohomology class corresponds to a Translator Cycle. The functor G_S provides the explicit bridge.

The Davis-Hodge Isomorphism (Conditional):

Assuming smooth chartability and functorial ε -equivalence:

$$H^{p,p}(M, \mathbb{Q}) \cong \text{TransCyc}^p(G_S(M), \mathbb{Q})$$

Remark 8.1. *This result does **not** resolve the classical Hodge Conjecture for arbitrary projective varieties. It shows that **within the Davis Framework ecosystem**, the Hodge property follows from the foundational axioms. Whether this framework captures all projective varieties—and thus implies classical Hodge—remains an open question.*