The Field Equations of Semantic Coherence

A Geometric Theory of Meaning, Curvature, and Reasoning
in Transformer Architectures

Complete Conjecture Reference
89 Mathematical Results with Dependency Structure

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November 30th, 2025

Contents

Note on Mathematical Status

The results in this document are stated as **conjectures** establishing a geometric research program for semantic coherence in transformer architectures. Formal proofs are under development. Scaling relationships use $\approx$ or $\propto$ to indicate functional form rather than exact equality. Optimality claims represent design hypotheses to be validated. Topological results assume benign manifold conditions (compact, finite-type, bounded curvature) unless otherwise specified.
This release establishes priority for the theoretical framework.

Part I
Overview and Master Results

1 The Master Equation

The Davis Law

The fundamental equation governing inference from incomplete information:

$$\boxed{C = \frac{\tau}{K}} \tag{1}$$

Where:

- C = **Inference Capacity** (Completion) — A measure of the degree to which unobserved states are uniquely determined by observed constraints on a geometric manifold
- τ = **Tolerance Budget** — The acceptable error threshold; the slack in the system
- K = **Curvature** — The geometric complexity of the space where information lives

1.1 The Three Master Equations

1. **Static Form (Existence):**
$$C = \frac{\tau}{K} \tag{2}$$
How much can be completed from incomplete information.
2. **Variational Form (Selection):**
$$\delta \oint \mathrm{Hol} = 0 \tag{3}$$
Which completion is chosen among possibilities — the Principle of Stationary Holonomy.
3. **Dynamic Form (Evolution):**
$$C_{\mathrm{dyn}} = \frac{\tau}{K + \eta T} \tag{4}$$

How completion capacity changes during learning, where η is the learning rate and T is time.

2 The Extended Master Trichotomy

Geometric Trichotomy

Every completion problem falls into exactly one regime, determined by the parameter:

$$\Gamma = \frac{m \cdot \tau_{\text{budget}}}{\hat{K}_{\max} \cdot \log |S|} \tag{5}$$

Static Regime ($\partial M/\partial t = 0$):

- $\Gamma > 1$: **DETERMINED** — Unique completion, stable, fast consensus
- $\Gamma = 1$: **CRITICAL** — Phase transition, power-law behavior, slow dynamics
- $\Gamma < 1$: **UNDERDETERMINED** — Multiple completions, unstable, no consensus

Dynamic Regime ($\partial M/\partial t \neq 0$):

$$\Gamma_{\text{eff}} = \Gamma \cdot \left(1 - \frac{\eta T}{\tau}\right) \tag{6}$$

Learning pushes the system toward critical/underdetermined. Three dynamic phases:

- $\eta < \eta_{\text{safe}}$: **STABLE LEARNING** — Guarantees preserved
- $\eta = \eta_{\text{safe}}$: **CRITICAL LEARNING** — Guarantees marginal
- $\eta > \eta_{\text{safe}}$: **UNSTABLE LEARNING** — Cache invalidation

3 Summary of Results

Order	Results	Cumulative
Foundation (T1–T5, E0, Corollary)	7	7
First Order	8	15
Second Order	24	39
Third Order	33	72
Fourth Order	13	85
Fifth Order: Energy (E1–E5)	5	90
Fifth Order: Dynamics (D1–D6)	6	96
Total Derived Results	**89**	

Note: The count of 89 excludes the foundational axioms and counts only derived results.

Part II
Foundational Theorems

4 The Five Foundational Theorems

T1: Geometric Completion Uniqueness

Conjecture 4.1 (Geometric Completion Uniqueness). *Given a partial world state W_0 on a Davis manifold (M, g) with curvature bound $\hat{K}_{\text{loc}} < \hat{K}_{\text{max}}$, and a set of observed constraints C, if the holonomy around all constraint-bounded regions satisfies*

$$\|\text{Hol} - I\| < \tau_{\text{budget}} \tag{7}$$

then there exists at most one completion W^* consistent with C up to ε-equivalence, where ε is determined by the Davis distortion radius.

Interpretation: This is the Sudoku theorem — sufficient constraints plus bounded curvature implies a unique solution.

T2: Harmonization Preserves Completion

Conjecture 4.2 (Harmonization Preserves Completion). *Let H be a harmonization operator that forces path-independence on observable operations O. For any non-deterministic completion process A, the harmonized completion $H(A(W_0, C))$ is observationally equivalent to the deterministic completion $D(W_0, C)$ on all operations in O.*

Interpretation: This imports the BRIDGE result into the reasoning setting — harmonization makes non-deterministic processes behave deterministically on observables.

T3: Gap-Filling Complexity Reduction

Conjecture 4.3 (Gap-Filling Complexity Reduction). *For a world model with n unobserved variables and m geometric constraints satisfying the benign curvature condition, the effective search space for valid completions is bounded by:*

$$|S_{\text{valid}}| \leq |S_{\text{unconstrained}}| \cdot \exp\left(-\frac{m \cdot \tau_{\text{budget}}}{\hat{K}_{\text{max}}}\right) \tag{8}$$

In the limit of tight curvature bounds ($\hat{K}_{\text{max}} \to 0$), valid completions converge to a unique solution.

Interpretation: This quantifies how geometry compresses the hypothesis space — constraints exponentially shrink the space of valid completions.

T4: Reasoning Fidelity Under Partial Observation

Conjecture 4.4 (Reasoning Fidelity Under Partial Observation)**.** *Let γ be a reasoning path from premises P to conclusion Q on (M, g), with k intermediate steps unobserved. If the observed steps satisfy the holonomy budget and the manifold has bounded curvature, then any valid completion of γ produces conclusions Q' satisfying:*

$$d_g(Q, Q') \leq k \cdot \ell_c \cdot \sqrt{\hat{K}_{\max}} + \varepsilon_{\text{disc}} \tag{9}$$

The reasoning error grows at most linearly in gap size, not exponentially.

Interpretation: This is the anti-hallucination guarantee — geometry prevents drift even when you can't observe every step.

T5: Davis Cache Sufficiency

Conjecture 4.5 (Davis Cache Sufficiency)**.** *For gap-filling on benign paths, the state (Φ_t, r_t) — continuous potential plus topological residue — is sufficient to determine valid completions. The cache size remains $O(1)$ in the number of gaps, provided total path length stays within the holonomy horizon $s_{\max}$.*

Interpretation: You don't need to store the whole world; the geometric summary is enough to constrain completions.

5 The Sudoku Principle Corollary

Corollary: The Sudoku Principle Test

Corollary 5.1 (The Sudoku Principle)**.** *A world model is "sudoku-complete" if its geometric constraints uniquely determine all unobserved states. The Davis framework provides a constructive test: compute holonomy around gap boundaries; if $\|\mathrm{Hol} - I\| < \tau$ for all such loops, the completion is unique.*

6 The Energy Principle

E0: Principle of Least Holonomy

Conjecture 6.1 (Principle of Least Holonomy)**.** *Among all paths connecting premises to conclusions, the realized path minimizes total holonomy. Define the Davis Energy Functional:*

$$E[\gamma] = \int_0^L \left(\lambda_1 + \lambda_2 \hat{K}_{\text{loc}}(s) + \lambda_3 \|\mathrm{Hol}_{\gamma_s} - I\|\right) ds \tag{10}$$

Optimal paths satisfy $\delta E / \delta\gamma = 0$.

Interpretation: This is the variational principle for reasoning — nature chooses the path of least resistance (lowest curvature and holonomy).

Part III
First-Order Derivations

These 8 results follow directly from the foundational theorems.

7 (1) Constraint Saturation Threshold

Conjecture 7.1 (Constraint Saturation Threshold m^*). *There exists a critical threshold m^* where $|S_{\text{valid}}| \approx 1$:*

$$m^* = \frac{\hat{K}_{\max} \cdot \log |S_{\text{unconstrained}}|}{\tau_{\text{budget}}} \tag{11}$$

At $m = m^*$, the completion becomes unique. Below m^*, multiplicity remains. Above m^*, constraints are redundant or inconsistent.

Derived from: T1 (Uniqueness) + T3 (Complexity Reduction)
Interpretation: This is the "sudoku moment" — the phase transition in the system.

8 (2) Compositional Holonomy

Conjecture 8.1 (Compositional Holonomy). *Holonomy composes additively to first order when individual holonomies are small:*

$$\|\text{Hol}_{\gamma_A \cup \gamma_B} - I\| \leq \|\text{Hol}_{\gamma_A} - I\| + \|\text{Hol}_{\gamma_B} - I\| + \mathcal{O}(\|\text{Hol}_{\gamma_A} - I\| \cdot \|\text{Hol}_{\gamma_B} - I\|) \tag{12}$$

Derived from: T1 (Uniqueness) + Corollary (Sudoku Principle)
Interpretation: This justifies doing completions incrementally — the locality property.

9 (3) Constraint Consistency Test

Lemma 9.1 (Constraint Consistency Test). *Constraints C are geometrically consistent iff there exists a path γ passing through $\bigcap_c R_c$ with holonomy below budget. Equivalently: if the holonomy around the boundary of $\bigcap_c R_c$ exceeds τ_{budget}, no valid completion exists.*

Derived from: T1 (Uniqueness)
Interpretation: The unsolvable sudoku detector.

10 (4) Completion Stability Under Perturbation

Conjecture 10.1 (Completion Stability). *If W_0' satisfies $d_g(W_0, W_0') < \delta$ and $\hat{K}_{\text{loc}} < \hat{K}_{\max}$ throughout, then the completions satisfy:*

$$d_g(W^*, W'^*) \leq \delta \cdot \exp\left(\sqrt{\hat{K}_{\max}} \cdot L\right) \tag{13}$$

where L is the path length through the gap region.

Derived from: T1 (Uniqueness) + T4 (Fidelity)
Interpretation: Small input perturbations don't cause large completion changes in the benign regime.

11 (5) Optimal Constraint Ordering

Lemma 11.1 (Optimal Constraint Ordering). *Given constraints $\{c_1, \ldots, c_m\}$, the optimal projection order minimizes cumulative holonomy. Greedily: at each step, project to the constraint whose region R_c is closest to the current path in geodesic distance.*

Derived from: T3 (Complexity) + F2 (Compositional Holonomy)
Interpretation: The greedy constraint scheduling lemma.

12 (6) Cache Compression Bound

Conjecture 12.1 (Cache Compression Bound). *The Davis Cache (Φ_t, r_t) contains at most:*

$$I[(\Phi_t, r_t)] \leq d_\Phi \cdot \log(1/\varepsilon) + \log K + \log B + W \cdot \log 3 \tag{14}$$

bits, where d_Φ is potential dimension, K charts, B basins, W winding components. This is independent of path length L and gap count k.

Derived from: T5 (Cache Sufficiency)
Interpretation: The cache has bounded entropy regardless of problem size.

13 (7) Maximum Gap Size from Holonomy Horizon

Lemma 13.1 (Maximum Gap Size). *A gap of geodesic length g can be uniquely completed iff:*

$$g < \frac{\tau_{\text{budget}}}{\sqrt{\hat{K}_{\max}}} \tag{15}$$

Gaps larger than this threshold admit multiple completions even with perfect boundary constraints.

Derived from: T4 (Fidelity) + T1 (Uniqueness)
Interpretation: The holonomy horizon *is* the maximum gap size.

14 (8) Multi-Agent Consensus

Conjecture 14.1 (Multi-Agent Consensus). *If agents A and B share anchor topology and both receive (Φ_t, r_t) with holonomy below budget, their completions W_A and W_B satisfy:*

$$d_g(W_A, W_B) \leq 2\varepsilon_{\text{disc}} \tag{16}$$

where $\varepsilon_{\text{disc}}$ is the discretization slack.

Derived from: T5 (Cache Sufficiency) + T1 (Uniqueness)
Interpretation: The consensus theorem for teleportation — agents converge.

Part IV
Second-Order Derivations

These 24 results follow from the first-order derivations.

15 From (1) Saturation Threshold

15.1 (1a) Constraint Redundancy Detection

Conjecture 15.1 (Constraint Redundancy Detection)**.** *If $m > m^*$, then at least $(m - m^*)$ constraints are redundant — they can be removed without changing the unique completion. Constructively: a constraint c_i is redundant iff removing it does not increase $|S_{\text{valid}}|$ above 1.*

Derived from: F1 (Saturation)

15.2 (1b) Constraint Value Ordering

Conjecture 15.2 (Constraint Value Ordering)**.** *Not all constraints contribute equally to shrinking S_{valid}. Define the information value of constraint c as:*

$$V(c) = \log |S_{\text{valid}}^{-c}| - \log |S_{\text{valid}}| \tag{17}$$

where S_{valid}^{-c} is the valid set without c. Constraints with higher $V(c)$ are more "informative" geometrically.

Derived from: F1 (Saturation)
Interpretation: Geometric entropy measure on constraints.

15.3 (1c) Phase Transition Sharpness

Corollary 15.3 (Phase Transition Sharpness)**.** *Near m^*, the transition from multiple completions to unique completion is sharp:*

$$\left.\frac{d|S_{\text{valid}}|}{dm}\right|_{m=m^*} = -\frac{\tau_{\text{budget}}}{\hat{K}_{\max}} |S_{\text{valid}}| \tag{18}$$

The transition sharpens as $\hat{K}_{\max} \to 0$ (flatter geometry = sharper phase transition).

Derived from: F1 (Saturation)
Interpretation: The manifold has a critical exponent.

16 From (2) Compositional Holonomy

16.1 (2a) Holonomy Algebra

Conjecture 16.1 (Holonomy Algebra)**.** *The set of holonomy operators $\{\mathrm{Hol}_\gamma\}$ under composition forms a group $H \subset GL(d)$. In the benign regime (all $\|\mathrm{Hol} - I\| < \tau$), H is approximately abelian:*

$$[\mathrm{Hol}_{\gamma_1}, \mathrm{Hol}_{\gamma_2}] = \mathcal{O}(\tau^2) \tag{19}$$

Completion order doesn't matter to first order.

Derived from: F2 (Compositional Holonomy)
Interpretation: The holonomy group is nearly commutative in the benign regime.

16.2 (2b) Holonomy Decomposition

Lemma 16.2 (Holonomy Decomposition). *Any holonomy* Hol_γ *for a complex loop* γ *can be decomposed into elementary loops around single gaps:*

$$\mathrm{Hol}_\gamma = \prod_i \mathrm{Hol}_{\gamma_i} + \mathcal{O}(\tau^2) \tag{20}$$

where γ_i *are simple loops around individual gaps.*

Derived from: F2 (Compositional Holonomy)
Interpretation: Holonomy factorization — complex reasoning decomposes into local completions.

16.3 (2c) Parallel Gap-Filling

Corollary 16.3 (Parallel Gap-Filling). *If gaps* $\{g_1, \ldots, g_k\}$ *have non-overlapping loop boundaries, their completions can be computed in parallel with combined error:*

$$\varepsilon_{\text{parallel}} = \varepsilon_{\text{sequential}} + \mathcal{O}(\tau^2) \tag{21}$$

Derived from: F2 (Compositional Holonomy)
Interpretation: Parallelization with bounded additional error.

17 From (3) Consistency Test

17.1 (3a) Inconsistency Localization

Conjecture 17.1 (Inconsistency Localization). *If constraints* C *are inconsistent, there exists a minimal inconsistent subset* $C' \subseteq C$ *such that:*

- $|C'| \leq d + 1$ *where* $d = \dim(M)$
- *The holonomy around* $\bigcap_{c \in C'} R_c$ *exceeds* τ_{budget}

Inconsistency is always localizable to at most $(d + 1)$ *constraints.*

Derived from: F3 (Consistency Test)
Interpretation: This is a geometric Helly theorem.

17.2 (3b) Inconsistency Resolution

Algorithm 17.2 (Inconsistency Resolution). Given inconsistent C, iteratively:

1. Find minimal inconsistent C' (at most $d + 1$ constraints)
2. Compute holonomy excess: $\Delta\tau = \|\mathrm{Hol}\| - \tau_{\text{budget}}$
3. Relax the weakest constraint in C' by $\Delta\tau$

Terminates in at most $|C|$ iterations with a consistent relaxed constraint set.

Derived from: F3 (Consistency Test)
Interpretation: How to fix an unsolvable sudoku by minimal relaxation.

18 From (4) Stability

18.1 (4a) Stability Radius

Conjecture 18.1 (Stability Radius). *For a completion W^* under constraints C, define the stability radius:*

$$r_{\text{stable}}(W^*) = \sup\{\delta : d_g(W_0, W_0') < \delta \Rightarrow d_g(W^*, W'^*) < \varepsilon\} \tag{22}$$

Then:

$$r_{\text{stable}} \geq \frac{\varepsilon}{\exp(\sqrt{\hat{K}_{\max}} \cdot L)} \tag{23}$$

Flatter geometry (smaller $\hat{K}_{\max}$) yields larger stability radius.

Derived from: F4 (Stability)
Interpretation: Quantifies robustness to input noise.

18.2 (4b) Adversarial Perturbation Bound

Corollary 18.2 (Adversarial Perturbation Bound). *To change the completion from W^* to a different W'^* with $d_g(W^*, W'^*) > \varepsilon$, an adversary must perturb inputs by at least:*

$$\delta_{\text{adv}} \geq \varepsilon \cdot \exp(-\sqrt{\hat{K}_{\max}} \cdot L) \tag{24}$$

Derived from: F4 (Stability)
Interpretation: Geometric adversarial robustness certificate.

19 From (5) Optimal Ordering

19.1 (5a) Ordering Regret Bound

Conjecture 19.1 (Ordering Regret Bound). *Let σ be any constraint ordering and σ^* the optimal ordering. The excess holonomy from suboptimal ordering is bounded:*

$$\sum_i \|\text{Hol}_{\sigma(i)}\| - \sum_i \|\text{Hol}_{\sigma^*(i)}\| \leq m \cdot \hat{K}_{\max} \cdot D^2 \tag{25}$$

where D is the diameter of $\bigcap_c R_c$.

Derived from: F5 (Ordering)
Interpretation: Greedy is near-optimal.

19.2 (5b) Constraint Ordering is Submodular

Corollary 19.2 (Submodularity). *The cumulative holonomy reduction from adding constraints is submodular:*

$$\Delta\text{Hol}(c|C_1) \geq \Delta\text{Hol}(c|C_2) \quad \textit{whenever } C_1 \subseteq C_2 \tag{26}$$

Adding a constraint helps more when you have fewer constraints.

Derived from: F5 (Ordering)
Interpretation: Greedy achieves $(1 - 1/e)$ optimal by standard submodular optimization.

20 From (6) Cache Compression

20.1 (6a) Cache Rate-Distortion

Conjecture 20.1 (Cache Rate-Distortion). *For any compression of the reasoning state to fewer than $I[(\Phi_t, r_t)]$ bits, there exist completions with error exceeding ε. The Davis Cache is conjectured to be rate-distortion optimal.*

Derived from: F6 (Compression)

20.2 (6b) Incompressibility of Winding Code

Corollary 20.2 (Incompressibility of Winding). *The winding code component r_t cannot be compressed below $W \cdot \log 3$ bits without losing topological distinguishability of paths.*

Derived from: F6 (Compression)
Interpretation: Topology is incompressible.

20.3 (6c) Cache Sufficiency is Tight

Conjecture 20.3 (Cache Tightness). *There exists a family of manifolds and path pairs (γ_1, γ_2) with identical Φ_t but different r_t that yield completions differing by $\Omega(\varepsilon)$. Both components are necessary.*

Derived from: F6 (Compression)
Interpretation: You can't drop either Φ or r.

21 From (7) Maximum Gap

21.1 (7a) Gap Additivity

Conjecture 21.1 (Gap Additivity). *Multiple gaps of sizes $g_1, \ldots, g_k$ can all be uniquely completed iff:*

$$\sum_i g_i \cdot \sqrt{\hat{K}_{\text{loc}}(g_i)} < \tau_{\text{budget}} \tag{27}$$

Gaps compete for the shared holonomy budget.

Derived from: F7 (Max Gap)
Interpretation: Gap budget allocation.

21.2 (7b) Optimal Gap Distribution

Corollary 21.2 (Optimal Gap Distribution). *Given total unknown length $G = \sum g_i$, uniqueness is maximized when gaps are distributed to minimize $\sum g_i \cdot \sqrt{\hat{K}_{\text{loc}}(g_i)}$. In uniform curvature: equal-sized gaps are optimal.*

Derived from: F7 (Max Gap)
Interpretation: Spread your ignorance evenly.

21.3 (7c) Critical Gap Ratio

Lemma 21.3 (Critical Gap Ratio). *Define the critical gap ratio:*

$$\rho^* = \frac{g_{\max}}{L_{\text{total}}} \tag{28}$$

where $g_{\max}$ is the largest gap and L_{total} is total path length. Unique completion requires:

$$\rho^* < \frac{\tau_{\text{budget}}}{\sqrt{\hat{K}_{\max}} \cdot L_{\text{total}}} \tag{29}$$

Derived from: F7 (Max Gap)
Interpretation: Gap fraction determines completion success.

22 From (8) Consensus

22.1 (8a) Consensus Convergence Rate

Conjecture 22.1 (Consensus Convergence). *If k agents iteratively share (Φ_t, r_t) and re-complete, their completions converge:*

$$d_g(W_i^{(n)}, W_j^{(n)}) \leq 2\varepsilon_{\text{disc}} \cdot \lambda^n \tag{30}$$

where $\lambda < 1$ depends on anchor alignment quality.

Derived from: F8 (Consensus)
Interpretation: Exponential convergence through cache exchange.

22.2 (8b) Byzantine Fault Tolerance

Corollary 22.2 (Byzantine Fault Tolerance). *If $f < k/3$ agents send corrupted (Φ_t, r_t), the remaining agents can still reach consensus on the correct completion by majority filtering on r_t (discrete) and median filtering on Φ_t (continuous).*

Derived from: F8 (Consensus)
Interpretation: Protocol is BFT for free.

22.3 (8c) Anchor Misalignment Tolerance

Conjecture 22.3 (Anchor Misalignment Tolerance). *If agents share anchors with alignment error δ_A (i.e., $d_g(a_i^A, a_i^B) < \delta_A$ for all anchors), consensus completions satisfy:*

$$d_g(W_A, W_B) \leq 2\varepsilon_{\text{disc}} + K \cdot \delta_A \tag{31}$$

where K is the number of charts traversed.

Derived from: F8 (Consensus)
Interpretation: Approximate anchor alignment is enough.

Part V
Third-Order Derivations

These 33 results follow from the second-order derivations. For brevity, we state them with minimal commentary.

23 From (1a) Constraint Redundancy

Conjecture 23.1 (Minimal Constraint Basis). *Every sudoku-complete constraint set C contains a minimal basis $B \subseteq C$ of exactly m^* constraints such that:*

- *Removing any $b \in B$ breaks uniqueness*
- *All $c \in C \setminus B$ are expressible as geometric combinations of B*

The basis is unique up to holonomy-preserving equivalence. This is the ***geometric matroid*** *underlying the constraint system.*

Algorithm 23.2 (Basis Extraction). Given C with $|C| > m^*$:

1. Order constraints by information value $V(c)$
2. Greedily add c to B if it reduces $|S_{\text{valid}}|$
3. Stop when $|S_{\text{valid}}| = 1$

Outputs minimal basis in $O(m \cdot |C|)$ holonomy computations.

Corollary 23.3 (Constraint Dimension). *The dimension of the constraint space is:*

$$\dim(C) = m^* = \frac{\hat{K}_{\max} \cdot \log |S_{\text{unconstrained}}|}{\tau_{\text{budget}}} \tag{32}$$

This is invariant under constraint reparameterization. The system has a ***geometric rank****.*

24 From (1b) Constraint Value

Conjecture 24.1 (Information Monotonicity). *Constraint information value satisfies:*

$$V(c|C_1) \geq V(c|C_2) \quad \text{when } C_1 \subseteq C_2 \tag{33}$$

Later constraints are always less informative than earlier ones (diminishing returns).

Lemma 24.2 (Curvature-Information Duality). *For a constraint c with region R_c:*

$$V(c) = \int_{R_c} \hat{K}_{\text{loc}}(x)\, dV_g(x) + \mathcal{O}(\tau^2) \tag{34}$$

Information value equals integrated curvature over the constraint region. ***Curvature is information.***

Corollary 24.3 (Optimal Observation Strategy). *To maximally reduce $|S_{\text{valid}}|$ with k observations, sample constraints from regions of highest integrated curvature. This is* ***geometric active learning****.*

25 From (1c) Phase Transition

Conjecture 25.1 (Critical Exponent). *Near the phase transition m^*, define order parameter $\phi = |S_{\text{valid}}| - 1$. Then:*

$$\phi \sim (m^* - m)^\beta, \quad \beta = \frac{\hat{K}_{\max}}{\tau_{\text{budget}}} \tag{35}$$

The critical exponent β is determined by the curvature-to-budget ratio. The manifold has ***universality class****.*

Corollary 25.2 (Finite-Size Scaling). *For finite systems (bounded M), the transition smooths:*

$$\Delta m_{\text{transition}} \sim \frac{1}{\sqrt{\text{Vol}(M)}} \tag{36}$$

Larger manifolds have sharper transitions.

26 From (2a) Holonomy Algebra

Conjecture 26.1 (Holonomy Lie Algebra). *In the infinitesimal limit ($\tau \to 0$), the holonomy operators generate a Lie algebra $\mathfrak{h}$ with bracket:*

$$[A_{\gamma_1}, A_{\gamma_2}] = \oint_{\gamma_1 \cap \gamma_2} R \tag{37}$$

where R is the curvature 2-form and $A_\gamma = \text{Hol}_\gamma - I$. ***The holonomy algebra is the curvature.***

Corollary 26.2 (Holonomy Dimension Bound).

$$\dim(\mathfrak{h}) \leq \frac{d(d-1)}{2} \tag{38}$$

where $d = \dim(M)$. Equality holds iff curvature spans all antisymmetric matrices.

Lemma 26.3 (Abelianization Error). *The error from treating H as abelian is:*

$$\|\text{Hol}_{\gamma_1}\text{Hol}_{\gamma_2} - \text{Hol}_{\gamma_2}\text{Hol}_{\gamma_1}\| \leq \|R\|_{L^\infty} \cdot A(\gamma_1 \cap \gamma_2) \tag{39}$$

where A is the area of loop intersection. Non-commutativity is ***localized to intersections****.*

27 From (2b) Holonomy Decomposition

Conjecture 27.1 (Prime Loop Decomposition). *Under suitable topological conditions, every loop γ on M decomposes into prime loops $\{\pi_i\}$ (not decomposable into smaller loops) such that:*

$$\text{Hol}_\gamma = \prod_i \text{Hol}_{\pi_i} \tag{40}$$

The prime loops generate the fundamental group $\pi_1(M)$.

Corollary 27.2 (Holonomy Basis). *The number of independent holonomy operators is bounded by a function of the topology, related to $\beta_1(M)$ (the first Betti number) and the holonomy group dimension.* ***Topology bounds holonomy complexity.***

Algorithm 27.3 (Loop Factorization). Given complex loop γ:

1. Compute homology class $[\gamma] \in H_1(M)$
2. Express $[\gamma] = \sum n_i[\pi_i]$ in prime basis
3. Approximate $\mathrm{Hol}_\gamma \approx \prod(\mathrm{Hol}_{\pi_i})^{n_i}$

Reduces holonomy computation from $O(\mathrm{length}(\gamma))$ to $O(\beta_1)$.

28 From (2c) Parallel Gap-Filling

Conjecture 28.1 (Parallelization Overhead). *The overhead from parallel vs. sequential gap-filling is exactly:*

$$\varepsilon_{\mathrm{overhead}} = \sum_{i<j} \|[\mathrm{Hol}_{\gamma_i}, \mathrm{Hol}_{\gamma_j}]\| \tag{41}$$

This is computable before execution.

Corollary 28.2 (Optimal Parallelization Partition). *Given gaps $G = \{g_1, \ldots, g_k\}$, the optimal partition into parallel batches minimizes:*

$$\sum_{\mathrm{batches}} \sum_{i<j\in\mathrm{batch}} A(\gamma_i \cap \gamma_j) \tag{42}$$

Gaps with non-intersecting boundaries should be parallelized.

Lemma 28.3 (Amdahl's Law for Geometric Completion). *Maximum speedup from parallelization:*

$$S_{\max} = \frac{1}{f_{\mathrm{sequential}} + \frac{\varepsilon_{\mathrm{overhead}}}{\varepsilon_{\mathrm{total}}}} \tag{43}$$

where $f_{\mathrm{sequential}}$ is the fraction of inherently sequential holonomy.

29 From (3a) Inconsistency Localization

Conjecture 29.1 (Geometric Helly Number). *The Helly number of constraint consistency on a d-dimensional Davis manifold is at most $d+1$ under geodesic convexity assumptions. That is: C is consistent iff every subset of size $\leq d+1$ is consistent.*

Corollary 29.2 (Inconsistency Detection Complexity). *Checking consistency of m constraints requires at most:*

$$\binom{m}{d+1} = O(m^{d+1}) \tag{44}$$

holonomy computations. For fixed d, this is polynomial in m.

Algorithm 29.3 (Fast Inconsistency Detection). Using geometric hashing:

1. Hash each constraint by its boundary holonomy signature
2. Constraints with incompatible signatures cannot be jointly consistent
3. Only check $(d+1)$-tuples with compatible signatures

Expected complexity $O(m^2)$ for randomly distributed constraints.

30 From (3b) Inconsistency Resolution

Conjecture 30.1 (Minimal Relaxation). *Among all relaxations of inconsistent C to consistent C', the minimal relaxation $C^{*'}$ satisfies:*

$$\sum_{c \in C} d_{\text{relax}}(c, c') \geq \oint_{\partial(\cap R_c)} \|\text{Hol} - I\| - \tau_{\text{budget}} \tag{45}$$

with equality for $C^{*'}$. ***Minimum edit distance to consistency equals holonomy excess.***

Corollary 30.2 (Relaxation is Unique). *If the holonomy excess is distributed among constraints proportionally to their boundary curvature contribution, the relaxation is unique.*

Lemma 30.3 (Relaxation Preserves Structure). *Minimal relaxation preserves constraint topology:*

$$\pi_1\left(\bigcap_{c \in C'} R_{c'}\right) \cong \pi_1\left(\bigcap_{c \in C} R_c\right) \tag{46}$$

*when relaxation is below the injectivity radius. You don't **tear** the constraint space, just stretch it.*

31 From (4a) Stability Radius

Conjecture 31.1 (Stability is Curvature-Determined). *The stability radius satisfies:*

$$r_{\text{stable}} = \frac{\tau_{\text{budget}}}{\sqrt{\hat{K}_{\max}}} \cdot \frac{1}{L} \tag{47}$$

Stability degrades linearly with path length and with square root of curvature.

Corollary 31.2 (Condition Number of Completion). *Define the geometric condition number:*

$$\kappa_g = \frac{L \cdot \sqrt{\hat{K}_{\max}}}{\tau_{\text{budget}}} \tag{48}$$

*Completions with $\kappa_g > 1$ are ill-conditioned. This is **numerical stability** for geometric inference.*

Lemma 31.3 (Stability Under Constraint Perturbation). *If constraint c is perturbed to c' with $d_g(R_c, R_{c'}) < \delta_c$, then:*

$$d_g(W^*, W'^*) \leq \delta_c \cdot V(c) \tag{49}$$

High-information constraints are more sensitive to perturbation.

32 From (4b) Adversarial Bound

Conjecture 32.1 (Adversarial Budget). *To force completion to a target W_{adv} with $d_g(W^*, W_{\text{adv}}) = \Delta$, an adversary must spend budget:*

$$B_{\text{adv}} \geq \Delta \cdot \exp(-\kappa_g) \tag{50}$$

Attacks are expensive when condition number is low.

Corollary 32.2 (Certified Radius). *No perturbation of size* $\delta < r_{\text{stable}}$ *can change the completion. This is a* ***geometric certificate****.*

Algorithm 32.3 (Adversarial Detection via Stability). Given input W_0 and completion W^*:

1. Compute r_{stable}
2. Sample perturbations of size $r_{\text{stable}}/2$
3. If completions vary by more than ε, flag as adversarial

Detects attacks that reduce stability radius.

33 From (5a) Regret Bound

Conjecture 33.1 (Regret Decomposition). *Total regret decomposes:*

$$\text{Regret}(\sigma) = \sum_i \sum_{j>i} \|[\text{Hol}_{\sigma(i)}, \text{Hol}_{\sigma(j)}]\| \tag{51}$$

Regret is sum of commutator norms over ordering inversions.

Corollary 33.2 (Optimal Order is Curvature-Sorted). *When constraints have nested regions ($R_{c_1} \supset R_{c_2} \supset \cdots$), optimal order is by decreasing integrated curvature. Large, curved constraints first.*

34 From (5b) Submodularity

Conjecture 34.1 (Greedy Approximation Ratio). *Greedy constraint ordering achieves holonomy reduction within factor $(1 - 1/e)$ of optimal:*

$$\text{Hol}_{\text{greedy}} \leq (1 - 1/e) \cdot \text{Hol}_{\text{optimal}} + \tau_{\text{budget}} \tag{52}$$

Corollary 34.2 (Online Constraint Selection). *Constraints arriving online can be greedily accepted/rejected with competitive ratio $(1 - 1/e)$ against offline optimal.*

35 From (6a) Rate-Distortion

Conjecture 35.1 (Cache is Sufficient Statistic). *(Φ_t, r_t) is a minimal sufficient statistic for completion:*

$$I(W^*; W_0, \gamma) = I(W^*; \Phi_t, r_t) \tag{53}$$

All completion-relevant information is captured.

Corollary 35.2 (No Better Cache Exists). *Any cache with fewer than $I[(\Phi_t, r_t)]$ bits must either lose completions or introduce errors exceeding ε.*

36 From (6b) Winding Incompressibility

Conjecture 36.1 (Winding Code is Homological). *The winding code r_t encodes the homology class of the path:*

$$r_t \cong [\gamma_{0:t}] \in H_1(M;\mathbb{Z}) \tag{54}$$

Winding counts crossings of homology generators.

Corollary 36.2 (Winding Dimension Equals Betti Number).

$$|r_t| = \beta_1(M) \cdot \log 3 \tag{55}$$

bits (for winding in $\{-1, 0, +1\}$ per generator).

37 From (6c) Cache Tightness

Lemma 37.1 (Φ-Distinguishability). *Paths with different Φ_t are geometrically separated:*

$$\Phi_t \neq \Phi'_t \Rightarrow d_g(\gamma(t), \gamma'(t)) > \varepsilon \tag{56}$$

Lemma 37.2 (r-Distinguishability). *Paths with different r_t are topologically separated:*

$$r_t \neq r'_t \Rightarrow [\gamma] \neq [\gamma'] \in \pi_1(M) \tag{57}$$

Conjecture 37.3 (Cache Separates All Paths). *Two paths yield the same completion iff they have identical (Φ_t, r_t):*

$$W^*_\gamma = W^*_{\gamma'} \Leftrightarrow (\Phi_t, r_t) = (\Phi'_t, r'_t) \tag{58}$$

Complete invariant.

38 From (7a) Gap Additivity

Conjecture 38.1 (Gap Budget Allocation). *Given total gap length G and holonomy budget τ, optimal allocation minimizes:*

$$\min_{\{g_i\}} \sum_i g_i \sqrt{\hat{K}_{\mathrm{loc}}(g_i)} \quad s.t. \quad \sum_i g_i = G \tag{59}$$

Solution: allocate proportionally to $1/\hat{K}_{\mathrm{loc}}$. ***Put gaps where curvature is low.***

Corollary 38.2 (Gap Capacity). *Maximum total gap length achievable:*

$$G_{\max} = \frac{\tau_{\mathrm{budget}}}{\min_x \sqrt{\hat{K}_{\mathrm{loc}}(x)}} \tag{60}$$

Capacity is determined by the flattest region.

Lemma 38.3 (Gap Interference). *Gaps g_i and g_j interfere iff their loop boundaries share edges. Interference cost:*

$$I(g_i, g_j) = \|\mathrm{Hol}_{\gamma_i \cap \gamma_j}\| \tag{61}$$

39 From (7b) Optimal Distribution

Conjecture 39.1 (Uniform Distribution Optimality). *In constant curvature ($\hat{K}_{\text{loc}} = \hat{K}_{\max}$ everywhere), equal-sized gaps minimize total holonomy:*

$$g_i = G/k \quad \forall i \tag{62}$$

Corollary 39.2 (Gap Fragmentation Principle). *Many small gaps are better than few large gaps:*

$$\sum_i \sqrt{g_i} \leq \sqrt{k} \cdot \sqrt{G/k} = \sqrt{G} \tag{63}$$

with equality for uniform distribution. ***Divide your ignorance.***

40 From (7c) Critical Ratio

Conjecture 40.1 (Observability Threshold). *A world model is observable (admits unique completion) iff:*

$$\rho_{\text{observed}} = \frac{L_{\text{observed}}}{L_{\text{total}}} > 1 - \frac{\tau_{\text{budget}}}{\sqrt{\hat{K}_{\max}} \cdot L_{\text{total}}} \tag{64}$$

Must observe at least this fraction.

Corollary 40.2 (Minimum Observation Density). *The minimum observation density for unique completion:*

$$\rho_{\min} = 1 - \frac{\tau_{\text{budget}}}{\sqrt{\hat{K}_{\max}} \cdot L_{\text{total}}} \tag{65}$$

In flat geometry ($\hat{K}_{\max} \to 0$), even sparse observations suffice.

41 From (8a) Convergence

Conjecture 41.1 (Convergence Rate).

$$\lambda = 1 - \frac{\tau_{\text{budget}}}{\hat{K}_{\max} \cdot D^2} \tag{66}$$

where D is manifold diameter. Flatter geometry = faster consensus.

Corollary 41.2 (Mixing Time). *Agents reach ε-consensus in:*

$$T_{\text{mix}} = \frac{\log(1/\varepsilon)}{\log(1/\lambda)} = O\left(\frac{\hat{K}_{\max} \cdot D^2}{\tau_{\text{budget}}} \log(1/\varepsilon)\right) \tag{67}$$

42 From (8b) BFT

Conjecture 42.1 (Geometric BFT Threshold). *With f Byzantine agents among k total:*

$$f < \frac{k}{3} \cdot \frac{\tau_{\text{budget}}}{\hat{K}_{\max}} \tag{68}$$

Curvature reduces fault tolerance.

Corollary 42.2 (Flat Geometry Maximizes Fault Tolerance). *As $\hat{K}_{\max} \to 0$, Byzantine threshold approaches $k/3$ (classical optimal).*

43 From (8c) Misalignment

Conjecture 43.1 (Alignment-Consensus Tradeoff).

$$d_g(W_A, W_B) \leq 2\varepsilon_{\text{disc}} + K \cdot \delta_A + \frac{\delta_A^2}{\tau_{\text{budget}}} \tag{69}$$

Quadratic penalty for large misalignment.

Corollary 43.2 (Maximum Tolerable Misalignment).

$$\delta_A^{\max} = \sqrt{\tau_{\text{budget}} \cdot \varepsilon} \tag{70}$$

Beyond this, consensus degrades rapidly.

Algorithm 43.3 (Anchor Alignment Protocol). Given agents A, B with potentially misaligned anchors:

1. Exchange cache states (Φ_t^A, r_t^A), (Φ_t^B, r_t^B)
2. Compute alignment error: $\delta_A \approx \|\Phi_t^A - \Phi_t^B\|$ (on shared test paths)
3. If $\delta_A > \delta_A^{\max}$, run anchor recalibration
4. Else proceed with completion

Part VI
Fourth-Order Derivations (Synthesis)

These 13 results synthesize across third-order branches.

44 Basis-Cache Duality

Conjecture 44.1 (Basis-Cache Duality). *The minimal constraint basis B and the Davis cache (Φ_t, r_t) are dual representations:*

$$|B| = m^* = \frac{I[(\Phi_t, r_t)]}{\log(1/\varepsilon)} \tag{71}$$

*Constraints and cache carry the same information, differently encoded. This is **duality between observations and state**.*

Derived from: T1a-i (Minimal Basis) + T6a-i (Sufficient Statistic)

45 Information-Curvature Conservation

Conjecture 45.1 (Information-Curvature Conservation). *Total information required for completion equals total curvature over gaps:*

$$I_{\text{required}} = \int_{\text{gaps}} \hat{K}_{\text{loc}}(x)\, dV_g(x) \tag{72}$$

Information and curvature are conserved quantities.

Derived from: T1b-ii (Curvature-Info Duality) + T7a-i (Gap Allocation)

Corollary 45.2 (Observation-Gap Complementarity).

$$\int_{\text{observed}} \hat{K}_{\text{loc}}\, dV + \int_{\text{gaps}} \hat{K}_{\text{loc}}\, dV = \int_M \hat{K}_{\text{loc}}\, dV = \text{const} \tag{73}$$

What you observe and what you infer sum to the total manifold curvature.

46 Structure Theorem for Davis Cache

Conjecture 46.1 (Structure Theorem). *The Davis cache decomposes as:*

$$(\Phi_t, r_t) \cong \mathfrak{h}^* \times H_1(M; \mathbb{Z}) \tag{74}$$

Continuous part lives in dual of holonomy algebra; discrete part lives in first homology.

Derived from: T2a-i (Lie Algebra) + T6b-i (Winding Homological)

Corollary 46.2 (Cache Dimension Formula).

$$\dim(\Phi_t) = \dim(\mathfrak{h}) \leq \frac{d(d-1)}{2} \tag{75}$$

$$|r_t| = \beta_1(M) \tag{76}$$

Both determined by manifold topology.

47 Constraint-Loop Duality

Conjecture 47.1 (Constraint-Loop Duality). *The minimal constraint basis B has size:*

$$|B| = \beta_1(M) + d + 1 \tag{77}$$

First Betti number (topological) plus Helly number (geometric). Constraints split into ***topological and geometric components****.*

Derived from: T2b-ii (Holonomy Basis) + T3a-i (Helly Number)

48 Condition-Convergence Relationship

Conjecture 48.1 (Condition-Convergence).

$$T_{\text{mix}} = O(\kappa_g^2 \cdot \log(1/\varepsilon)) \tag{78}$$

Consensus time scales with square of condition number. ***Ill-conditioned problems have slow consensus.***

Derived from: T4a-ii (Condition Number) + T8a-i (Convergence Rate)

49 Greedy Gap-Filling Near-Optimal

Conjecture 49.1 (Greedy Gap-Filling). *Greedily filling smallest gaps first achieves:*

$$\text{Hol}_{\text{greedy}} \leq (1 + 1/e) \cdot \text{Hol}_{\text{optimal}} \tag{79}$$

Combined with constraint ordering, total approximation ratio is $(1 - 1/e^2)$.

Derived from: T5b-i (Greedy Ratio) + T7b-ii (Fragmentation)

50 Phase Diagram of Completion

Conjecture 50.1 (Phase Diagram). *The $(\rho_{\text{observed}}, \hat{K}_{\max})$ plane divides into:*

1. ***Unique completion region****: $\rho > \rho_{\min}(\hat{K})$*
2. ***Multiple completion region****: $\rho < \rho_{\min}(\hat{K})$*
3. ***Critical line****: $\rho = \rho_{\min}(\hat{K})$ with phase transition*

The critical line is:

$$\rho_{\text{crit}}(\hat{K}) = 1 - \frac{\tau_{\text{budget}}}{\sqrt{\hat{K}} \cdot L} \tag{80}$$

Derived from: T1c-i (Critical Exponent) + T7c-i (Observability)

51 Fault-Constraint Duality

Conjecture 51.1 (Fault-Constraint Duality). *Byzantine agents act as inconsistent constraints. Maximum tolerable:*

$$f_{\max} = \frac{k-(d+1)}{3} \tag{81}$$

Must have enough honest agents to satisfy Helly bound.

Derived from: T3a-i (Helly) + T8b-i (BFT Threshold)

52 Universal Cache Protocol

Conjecture 52.1 (Universal Cache Protocol). *For any two agents A, B with bounded misalignment $\delta_A < \delta_A^{\max}$:*

1. *Cache exchange: $O(d_\Phi + \beta_1)$ bits*
2. *Alignment check: $O(1)$ paths*
3. *Consensus: $O(T_{\text{mix}})$ rounds*

Total communication: $O(d_\Phi + \beta_1 + T_{\text{mix}})$. ***Optimal*** *for the given geometric constraints.*

Derived from: T6c-iii (Cache Separates) + T8c-iii (Alignment Protocol)

53 Representation Theorem

Conjecture 53.1 (Representation Theorem). *The following are equal:*

1. *Minimal constraint basis size $|B| = m^*$*
2. *Cache information $I[(\Phi_t, r_t)]/\log(1/\varepsilon)$*
3. *Holonomy algebra dimension + Betti number + 1*
4. *Geometric degrees of freedom for completion*

Fundamental invariant *of the completion problem.*

Derived from: T1a-iii (Constraint Dim) + T6a-ii (No Better Cache)

54 Parallel-Sequential Equivalence

Conjecture 54.1 (Parallel-Sequential Equivalence). *Total holonomy is invariant under execution strategy:*

$$\text{Hol}_{\text{sequential}} + \text{Regret}(\sigma) = \text{Hol}_{\text{parallel}} + \varepsilon_{\text{overhead}} \tag{82}$$

for any ordering σ and any parallel partition. ***Conservation law for holonomy.***

Derived from: T2c-i (Overhead) + T5a-i (Regret Decomposition)

55 Attack Surface Geometry

Conjecture 55.1 (Attack Surface Geometry)**.** *The adversarial attack surface has measure:*

$$\mu(\text{Attack}) = G_{\max} \cdot \exp(-\kappa_g) \tag{83}$$

Large gap capacity but high condition number = small attack surface.

Derived from: T4b-i (Adversarial Budget) + T7a-ii (Gap Capacity)

Part VII
Fifth-Order Derivations

56 Energy Derivations (E1–E5)

From E0: Principle of Least Holonomy

Conjecture 56.1 (Geodesic Completion). *In regions where $\hat{K}_{\text{loc}} < \hat{K}_{\max}$ uniformly, energy-minimizing paths are geodesics:*

$$\nabla_{\dot{\gamma}}\dot{\gamma} = 0 \tag{84}$$

Optimal reasoning follows straight lines in flat regions.

Conjecture 56.2 (Energy-Holonomy Equivalence). *For benign paths, total energy equals integrated holonomy:*

$$E[\gamma] = \lambda_3 \oint_{\partial\gamma} \|\text{Hol} - I\| + O(\tau^2) \tag{85}$$

***Energy is holonomy** (to first order).*

Conjecture 56.3 (Noether's Theorem for Reasoning). *If the manifold has a symmetry (isometry group G), then there exists a conserved quantity along optimal paths:*

$$J_G = \langle \dot{\gamma}, \xi_G \rangle = \text{const} \tag{86}$$

where ξ_G is the Killing field of G. ***Symmetries produce conservation laws.***

Conjecture 56.4 (Hamilton-Jacobi Equation). *The optimal completion value function $V(W_0, W_{\text{goal}})$ satisfies:*

$$\|\nabla V\|^2 = \lambda_1 + \lambda_2 \hat{K}_{\text{loc}} + \lambda_3 \|\text{Hol} - I\| \tag{87}$$

This PDE characterizes all optimal completions simultaneously.

Corollary 56.5 (Principle of Stationary Holonomy). *Among all completions connecting W_0 to W^*, the realized completion is the one for which holonomy is stationary:*

$$\delta \oint \text{Hol} = 0 \tag{88}$$

The variational principle for reasoning.

57 Dynamics Derivations (D1–D6)

Manifold Evolution: $\partial M/\partial t$

Conjecture 57.1 (Cache Invalidation Condition). *The Davis Cache (Φ_t, r_t) remains valid under manifold update $\partial M/\partial t$ iff:*

$$\left\| \frac{\partial g}{\partial t} \right\|_{L^\infty} < \frac{\tau_{\text{budget}}}{T \cdot L} \tag{89}$$

where T is time since cache computation and L is path length. Beyond this rate, cache must be recomputed.

Conjecture 57.2 (Learning-Completion Tradeoff). *During active learning with rate η, the completion capacity degrades:*

$$C_{\text{dyn}} = \frac{\tau}{K + \eta \cdot T} \tag{90}$$

Fast learning temporarily reduces completion capacity.

Conjecture 57.3 (Manifold Stability Under Updates). *A manifold remains in the benign regime during learning iff:*

$$\frac{\partial \hat{K}_{\max}}{\partial t} < \frac{\tau_{\text{budget}}}{s_{\max}^2} \tag{91}$$

Curvature cannot grow faster than the holonomy horizon can accommodate.

Conjecture 57.4 (Basin Drift Bound). *If basin B has center c_B at time t, then under learning:*

$$\|c_B(t + \Delta t) - c_B(t)\| \leq \eta \cdot \Delta t \cdot \sqrt{\hat{K}_{\text{loc}}(c_B)} \tag{92}$$

Basins in high-curvature regions drift faster.

Conjecture 57.5 (Anchor Persistence). *An anchor set A remains valid (preserves cache structure) iff no anchor crosses a basin boundary:*

$$d_g(a_i(t), \partial B_j) > \eta \cdot T \cdot \sqrt{\hat{K}_{\max}} \quad \forall i, j \tag{93}$$

Anchors must stay away from basin boundaries by a margin proportional to learning rate.

Corollary 57.6 (Safe Learning Rate). *The maximum learning rate that preserves all completion guarantees:*

$$\eta_{\text{safe}} = \min\left(\frac{\tau}{T \cdot L \cdot K}, \frac{d_{\min}(A, \partial B)}{T \cdot \sqrt{\hat{K}_{\max}}} \right) \tag{94}$$

where $d_{\min}$ is minimum anchor-to-boundary distance. ***This is the speed limit for training.***

Part VIII
The Davis Manifold Relaxation Algorithm

58 Problem Formulation

Given:

- Partial world W_0, goal W_{goal}, constraints C
- Curvature bound $\hat{K}_{\max}$, holonomy budget τ_{budget}

Goal: Compute a benign, low-curvature path γ^* from W_0 to W_{goal} whose completions obey the Sudoku Principle.

59 Algorithm Steps

1. **Initialization (Anchor Step)**
   - Map W_0, W_{goal} to anchors z_{start}, z_{end}
   - Initialize γ as a geodesic between anchors
2. **Constraint Projection (Sudoku Step)**
   - For each constraint $c \in C$, define valid region $R_c \subset M$
   - Project γ toward R_c; treat violations as forces pulling γ back into $\bigcap_c R_c$
3. **Manifold Relaxation (Energy Minimization)**
   - Define Davis Energy Functional:
   $$E[\gamma] = \lambda_1 \int_0^L ds + \lambda_2 \int_0^L \hat{K}_{\text{loc}}(s)\, ds + \lambda_3 \int_0^L \|\text{Hol}_{\gamma_s} - I\|\, ds \tag{95}$$
   - Gradient descent or geodesic shooting to minimize functional subject to boundary conditions
4. **Discretization & Harmonization (BRIDGE Step)**
   - Discretize γ^* via speed-of-thought CFL: $\Delta s \approx \ell_c / \sqrt{1 + \hat{K}_{\text{loc}}}$
   - Execute reasoning via BRIDGE-like engine
   - Use Davis cache to store (Φ_t, r_t)
   - Apply harmonization H on observable operations
   - On failure (e.g., exit code 126 or holonomy spike), add new constraint at collision point and re-run relaxation
5. **Output**
   - Verified deterministic path γ^* with:
     - Bounded curvature and holonomy
     - Unique completion guarantees
     - Zero hallucinated observables under H

60 Correctness Links to Theorems

- **Theorem 1**: Uniqueness of γ^*'s induced world completion
- **Theorem 2**: Harmonization ensures deterministic replay
- **Theorem 3**: Effective search-space reduction via energy functional
- **Theorem 4**: Linear error bounds for any remaining gaps
- **Theorem 5**: Use of Davis cache as sufficient state
- **E0**: Energy minimization selects optimal path

Part IX
Applications

61 Curvature-Aware Training Data Filtering

Use $\hat{K}_{\text{loc}}$ as a pretrain filter:

- Estimate curvature induced by documents/tasks
- Exclude high-curvature samples to learn flatter world manifolds
- Extends the Sudoku regime

62 Teleporting Reasoning State Between Agents

Protocol:

- Each agent maintains (Φ_t, r_t)
- On handoff, send only this cache and relevant constraints
- Receiver reconstructs local view and continues Davis Manifold Relaxation
- Communication cost: $O(d_\Phi + \beta_1)$ bits

63 Adversarial Detection via Holonomy Spikes

Prompt/observation injection shatters geometry:

- Sudden increase in $\hat{K}_{\text{loc}}$ or $\|\text{Hol} - I\|$ across loops that used to be benign
- Detect and block completions when holonomy exceeds τ_{budget}
- Detection before generation via curvature monitoring

64 Mental Stack Trace Debugger

Log $(\Phi_t, r_t, \hat{K}_{\text{loc}}, \text{Hol})$ along reasoning:

- On failure, locate step where:
  - Curvature exceeded $\hat{K}_{\text{max}}$, or
  - Holonomy budget violated
- Provide human-readable "where it lost the Sudoku solution" stack trace
- Geometric debugging: "model lost coherence at step 43 because curvature exceeded 1.0"

Appendices

A Notation Reference

Symbol	Definition
(M, g)	Davis manifold with metric g
$\hat{K}_{\text{loc}}$	Local normalized curvature
$\hat{K}_{\max}$	Maximum curvature bound
τ_{budget}	Holonomy/tolerance budget
Hol_γ	Holonomy operator around loop γ
$s_{\max}$	Holonomy horizon (maximum coherent path length)
ℓ_c	Characteristic semantic length
C	Constraint set
R_c	Valid region for constraint c
W, W_0, W^*	World state, partial world, completion
Φ_t	Continuous potential (location on manifold)
r_t	Topological residue (chart, basin, winding code)
(Φ_t, r_t)	Davis cache
m^*	Constraint saturation threshold
S_{valid}	Set of valid completions
$V(c)$	Information value of constraint c
κ_g	Geometric condition number
$\beta_1(M)$	First Betti number of M
$\mathfrak{h}$	Holonomy Lie algebra
$\varepsilon_{\text{disc}}$	Discretization slack
η	Learning rate
η_{safe}	Maximum safe learning rate
C (in $C = \tau/K$)	Inference capacity / completion
Γ	Trichotomy parameter
$E[\gamma]$	Davis energy functional

B Dependency Summary

Foundation → First Order:

- T1 + T3 → F1 (Saturation)
- T1 + Cor → F2 (Compositional)
- T1 → F3 (Consistency)
- T1 + T4 → F4 (Stability)
- T3 + F2 → F5 (Ordering)
- T5 → F6 (Compression)

- T4 + T1 → F7 (Max Gap)
- T5 + T1 → F8 (Consensus)

First → Second: Each Fi generates multiple Six results (see main text).
Second → Third: Each Six generates multiple Tix-y results (see main text).
Third → Fourth: Cross-synthesis produces FO1–FO14.
Foundation → Fifth: E0 generates E1–E5 (Energy). T1, T5, FO, S6, S8 generate D1–D6 (Dynamics).
Fourth + Fifth → Master: FO9, FO12, E5, D6 → Extended Trichotomy → Master Equations.

"The amount you can know from incomplete information is inversely proportional to the curvature of the space where that information lives."

The Davis Law: $C = \tau/K$
