

Attack Lab

CSE251, Spring 2019

Recitation 1: Wed, March 13th, 2019

Changmin Yi

ulistar93@unist.ac.kr

*Reference : CMU 15-213: Intro to Computer Systems Fall 2015
Recitation 3 - Dhruven Shah, Ben Spinelli*

Agenda

- Stack review
- Attack lab overview
 - Phases 1-3: Buffer overflow attacks
 - Phases 4-5: ROP attacks

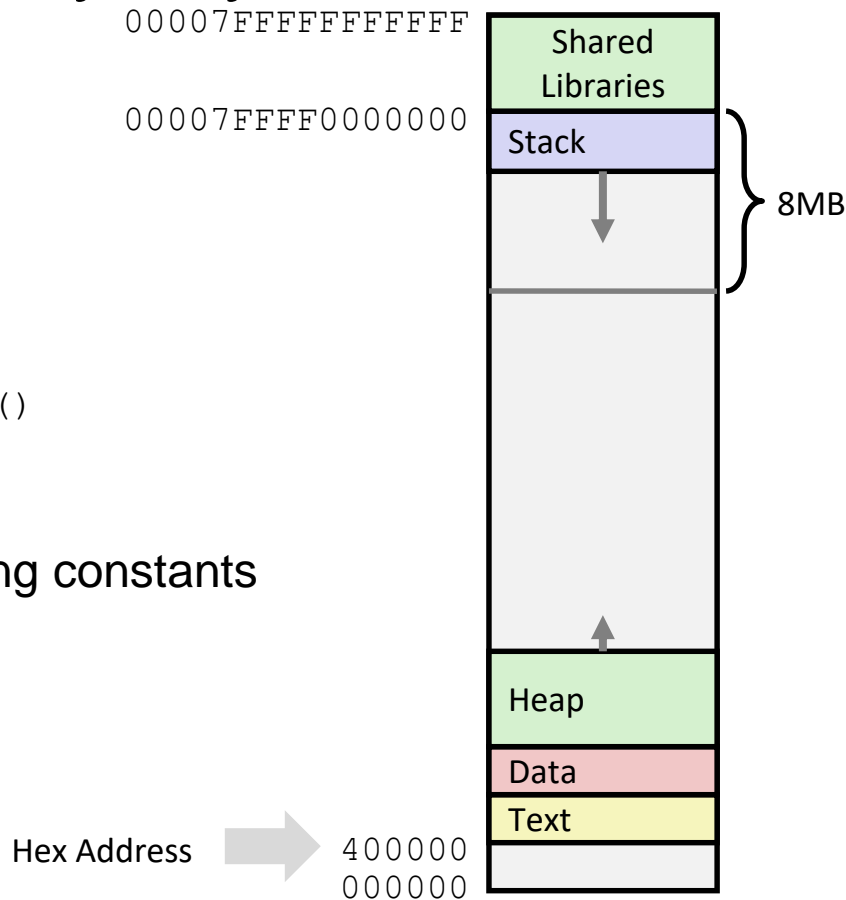
Notices

- Typo in pdf,
Due date is **Apr 8. Monday.**

x86-64 Linux Memory Layout

not drawn to scale

- Stack
 - Runtime stack (8MB limit)
 - E. g., local variables
- Heap
 - Dynamically allocated as needed
 - When call `malloc()`, `calloc()`, `new()`
- Data
 - Statically allocated data
 - E.g., global vars, `static` vars, string constants
- Text / Shared Libraries
 - Executable machine instructions
 - Read-only

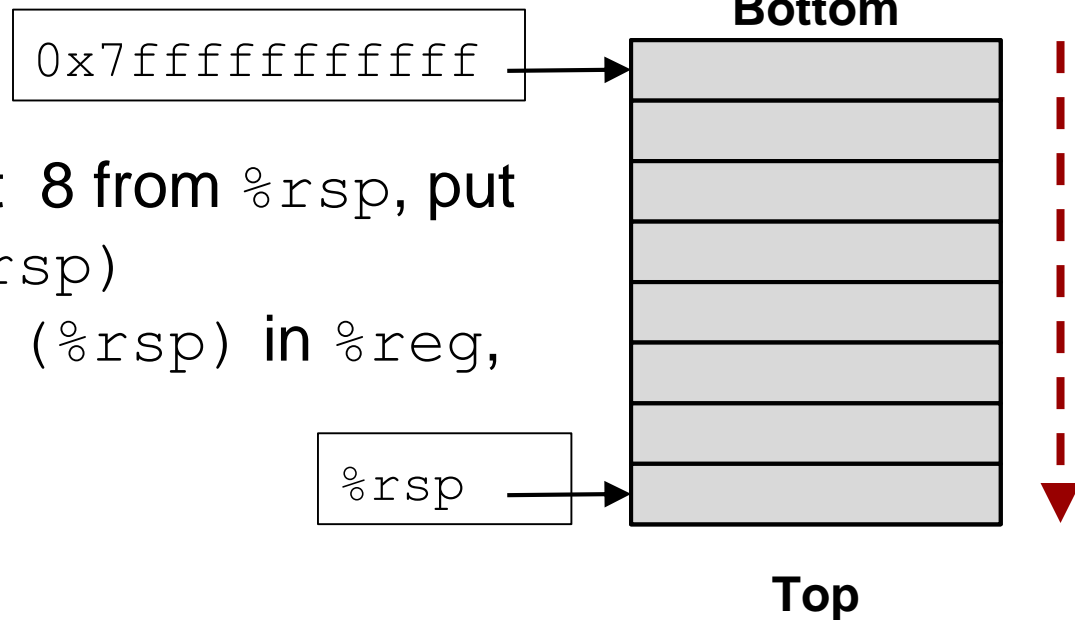


x86-64: Register Conventions

- Arguments passed in registers:
`%rdi, %rsi, %rdx, %rcx, %r8, %r9`
- Return value: `%rax`
- Callee-saved: `%rbx, %r12, %r13, %r14, %rbp, %rsp`
- Caller-saved: `%rdi, %rsi, %rdx, %rcx, %r8, %r9, %rax, %r10, %r11`
- Stack pointer: `%rsp`
- Instruction pointer: `%rip`

x86-64: The Stack

- Grows **downward** towards **lower** memory addresses
- `%rsp` points to **top** of stack



- `push %reg`: subtract 8 from `%rsp`, put `val` in `%reg` at `(%rsp)`
- `pop %reg`: put `val` at `(%rsp)` in `%reg`, add 8 to `%rsp`

Buffer Overflow Stack Example #1

After call to gets

Stack Frame for call_echo			
00	00	00	00
00	40	06	f6
00	32	31	30
39	38	37	36
35	34	33	32
31	30	29	28
27	26	25	24
23	22	21	20

```
void echo()
{
    char buf[4];
    gets(buf);
    . . .
}
```

```
echo:
    subq    $24, %rsp
    movq    %rsp, %rdi
    call    gets
    . . .
```

call_echo:

```
. . .
4006f1:  callq    4006cf <echo>
4006f6:  add      $0x8,%rsp
. . .
```

buf ← %rsp

```
unix> ./bufdemo-nsp
Type a string: 01234567890123456789012
01234567890123456789012
```

```
"01234567890123456789012\0"
```

Overflowed buffer, but did not corrupt state

Buffer Overflow Stack Example #2

After call to gets

Stack Frame for call_echo			
00	00	00	00
00	40	00	34
33	32	31	30
39	38	37	36
35	34	33	32
31	30	29	28
27	26	25	24
23	22	21	20

```
void echo()
{
    char buf[4];
    gets(buf);
    . . .
}
```

```
echo:
    subq    $24, %rsp
    movq    %rsp, %rdi
    call    gets
    . . .
```

call_echo:

```
. . .
4006f1:  callq    4006cf <echo>
4006f6:  add      $0x8,%rsp
. . .
```

buf ← %rsp

```
unix> ./bufdemo-nsp
Type a string: 0123456789012345678901234
Segmentation fault
```


Buffer Overflow Stack Example #3

After call to gets

Stack Frame for call_echo			
00	00	00	00
00	40	06	00
33	32	31	30
39	38	37	36
35	34	33	32
31	30	39	38
37	36	35	34
33	32	31	30

```
void echo()
{
    char buf[4];
    gets(buf);
    . . .
}
```

```
echo:
    subq    $24, %rsp
    movq    %rsp, %rdi
    call    gets
    . . .
```

call_echo:

```
. . .
4006f1:  callq    4006cf <echo>
4006f6:  add      $0x8,%rsp
. . .
```

buf ← %rsp

```
unix> ./bufdemo-nsp
Type a string: 012345678901234567890123
012345678901234567890123
```

Buffer Overflow Stack Example #3 Explained

After call to gets

Stack Frame for <code>call_echo</code>			
00	00	00	00
00	40	06	00
33	32	31	30
39	38	37	36
35	34	33	32
31	30	39	38
37	36	35	34
33	32	31	30

register_tm_clones:

```

. . .
400600:  mov    %rsp,%rbp
400603:  mov    %rax,%rdx
400606:  shr    $0x3f,%rdx
40060a:  add    %rdx,%rax
40060d:  sar    %rax
400610:  jne    400614
400612:  pop    %rbp
400613:  retq

```

`buf ← %rsp`

“Returns” to unrelated code

Lots of things happen, without modifying critical state

Eventually executes `retq` back to main

Attack Lab Overview: Phases 1-3

Overview

- Exploit x86-64 by overwriting the stack
- Overflow a buffer, overwrite return address
- Execute injected code

Key Advice

- Brush up on your x86-64 conventions!
- **Use objdump -d** to determine relevant offsets
- **Use GDB** to determine stack addresses

objdump/GDB – example1 / reci 2

- **Purpose:** To predict outputs of sample programs by analyzing assembly codes.

```
[cs20111100@uni06 gdb_ex]$ ls
ex1  ex2  ex3
[cs20111100@uni06 gdb_ex]$ objdump -d ex1

ex1:      file format elf64-x86-64

Disassembly of section .init:

0000000000400418 <_init>:
  400418:    48 83 ec 08      sub    $0x8,%rsp
  40041c:    48 8b 05 d5 0b 20 00  mov    0x200bd5(%rip),%rax
```

- There are 3 unknown binary files.
- “\$ objdump –d ex1” command generates an assembly code for ex1.
- You can store it as a file by redirection.
ex) \$ objdump –d ex1 > ex1_assembly

Attack Lab Overview: Phases 4-5

Overview

- Utilize return-oriented programming to execute arbitrary code
 - Useful when stack is **non-executable** or **randomized**
- Find **gadgets**, string together to form injected code

Key Advice

- Use mixture of pop & mov instructions + constants to perform specific task

ROP Example: Solution

Gadgets:

Address 1: `mov %rbx, %rax; ret`

Address 2: `pop %rbx; ret`

```
void foo(char *input){  
    char buf[32];  
    ...  
    strcpy (buf, input);  
    return;  
}
```

Old Return
address

buf

Next address in ROP chain....

Address 1

0xBBBBBBBB

Address 2

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF

0xFFFFFFFF (filler.....)

Gadget example

■ In rtarget

```
void setval_210(unsigned *p)
{
    *p = 3347663060U;
}
```

```
0000000000400f15 <setval_210>:
400f15:    c7 07 d4 48 89 c7    movl    $0xc78948d4, (%rdi)
400f1b:    c3                  retq
```

↑
400f18

A. Encodings of movq instructions

movq *S*, *D*

Source <i>S</i>	Destination <i>D</i>							
	%rax	%rcx	%rdx	%rbx	%rsp	%rbp	%rsi	%rdi
%rax	48 89 c0	48 89 c1	48 89 c2	48 89 c3	48 89 c4	48 89 c5	48 89 c6	48 89 c7
%rcx	48 89 c8	48 89 c9	48 89 ca	48 89 cb	48 89 cc	48 89 cd	48 89 ce	48 89 cf
%rdx	48 89 d0	48 89 d1	48 89 d2	48 89 d3	48 89 d4	48 89 d5	48 89 d6	48 89 d7
%rbx	48 89 d8	48 89 d9	48 89 da	48 89 db	48 89 dc	48 89 dd	48 89 de	48 89 df
%rsp	48 89 e0	48 89 e1	48 89 e2	48 89 e3	48 89 e4	48 89 e5	48 89 e6	48 89 e7
%rbp	48 89 e8	48 89 e9	48 89 ea	48 89 eb	48 89 ec	48 89 ed	48 89 ee	48 89 ef
%rsi	48 89 f0	48 89 f1	48 89 f2	48 89 f3	48 89 f4	48 89 f5	48 89 f6	48 89 f7
%rdi	48 89 f8	48 89 f9	48 89 fa	48 89 fb	48 89 fc	48 89 fd	48 89 fe	48 89 ff

Tools

- **objdump -d**
 - View byte code and assembly instructions, determine stack offsets
- **./hex2raw**
 - Pass raw ASCII strings to targets
- **gdb**
 - Step through execution, determine stack addresses
- **gcc -c**
 - Generate object file from assembly language file

More Tips

- Draw stack diagrams
- Memo what data is assigned register
- Be careful of byte ordering (little endian)

Questions?