

# CSE232 Assignment 1

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1. (a)  $A = \{x^2 | x \in N \wedge x \leq 5\}$   
(b)  $B = \{(x, y) | x, y \in N \wedge (x + y) < 3\}$
2. (a)  $A = \{1\}, B = \{1, 1\}$   
It is true that  $A \subseteq B$  and  $B \subseteq A$ . Thus,  $A = B$ .  
(b)  $A = \{1, 2\}, B = \{2, 1\}$   
It is true that  $A \subseteq B$  and  $B \subseteq A$ . Thus,  $A = B$ .  
(c)  $A = \{1, 2\}, B = (1, 2)$   
Obviously  $A \neq B$  because  $A$  is a set. Whereas,  $B$  is not.
3. Prove that  $\mathcal{P}(A \cup B) = \mathcal{P}(A) \cup \mathcal{P}(B)$

Suppose  $A = \{x | x \in A \wedge x \notin B\}$  and the size of it is  $n$ . Moreover, let's say that the size of the set  $B$  is  $m$ . Then it is obvious that the size of the set  $A \cup B$  should equal  $n + m$ .

Now, as the size of powerset of a set is always 2 to the power the size of that set, we can see that  $\mathcal{P}(A \cup B)$  has  $2^{(n+m)}$  elements. While,  $\mathcal{P}(A) \cup \mathcal{P}(B)$  has at most  $2^n + 2^m$  elements. Hence, the equation is not always true.

4. (a) Please refer to the handwritten sample attached to the paper.  
(b)  $(A - B) \cup (B - A) = (A \cup B) \cap (\overline{A} \cup \overline{B})$   
 $(A - B) \cup (B - A) = (A - B) \cup (B \cap \overline{A})$   
 $= (A \cap \overline{B}) \cup (B \cap \overline{A})$   
 $= ((A \cap \overline{B}) \cup B) \cap ((A \cap \overline{B}) \cup \overline{A})$   
 $= (A \cup B) \cap (\overline{B} \cup B) \cap ((A \cup \overline{A}) \cap (\overline{A} \cup \overline{B})) = (A \cup B) \cap (\overline{A} \cup \overline{B})$   
q.e.d.

$p$	$q$	$p \wedge \neg q$	$\neg p \wedge q$	$p \vee q$	$\neg p \vee \neg q$	$(p \wedge \neg q) \vee (\neg p \wedge q)$	$(p \vee q) \wedge (\neg p \wedge \neg q)$
$T$	$T$	$F$	$F$	$T$	$F$	$F$	$F$
$T$	$F$	$T$	$F$	$T$	$T$	$T$	$T$
$F$	$T$	$F$	$T$	$T$	$T$	$T$	$T$
$F$	$F$	$F$	$F$	$F$	$T$	$F$	$F$

5.  $A \times B = B \times A$  is true if one of the following is satisfied.

(a)  $A = B$

(b)  $A = \emptyset$  or  $B = \emptyset$

6.  $(p \wedge (q \oplus r)) \vee (q \wedge (p \oplus r)) \vee (r \wedge (q \oplus p))$

7. (a) Please refer to the table at the top of this page.

(b)  $(p \wedge \neg q) \vee (\neg p \wedge q) \equiv (p \vee q) \wedge (\neg p \vee \neg q)$

Distributive [1]:  $(p \wedge \neg q) \vee (\neg p \wedge q) \equiv (\neg p \vee (p \wedge \neg q)) \wedge (q \vee (p \wedge \neg q)) \equiv$

Distributive [2]:  $((\neg p \vee p) \wedge (\neg p \vee \neg q)) \wedge ((q \vee p) \wedge (q \vee \neg q)) \equiv$

$(\neg p \vee \neg q) \wedge (p \vee q)$

q.e.d.

8.  $\forall x \forall y \exists z (P(x, z) \wedge P(y, z) \vee P(x, y))$

as  $\forall x P(x, x) \equiv \text{True}$ , thus  $\forall x \forall y \exists z (P(x, z) \wedge P(y, z))$  is sufficient for this problem.

9.  $\neg \forall x \forall y (P(x, y) \vee Q(x, y)) \equiv \exists x \exists y \neg (P(x, y) \vee Q(x, y))$

Via De Morgan's law it's  $\equiv \exists x \exists y (\neg P(x, y) \wedge \neg Q(x, y))$