Bits, Bytes, and Integers

CSE251: System Programming 2nd and 3rd Lectures, Feb 27 and Mar 6

Instructor:

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Administrivia

Lab 0

Will be out today, due on Mar 5

REC 1 on Mar 4.

- Instructor will be out of town.
- We may meet at different class room please check your UNIST email and the class Gitlab.

Gitlab accounts

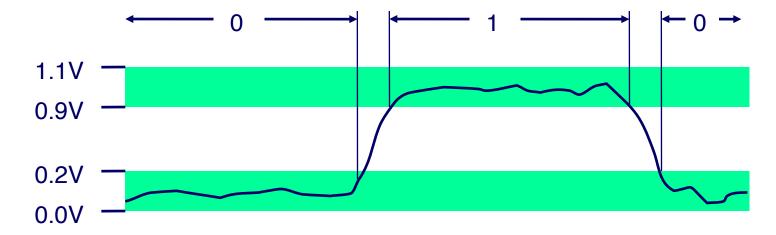
- All students registered by Monday night should have one.
- Will create more Friday night (once the course change is finalized).
- You are supposed to read the book also.

Today: Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- Integers
 - Representation: unsigned and signed
 - Conversion, casting
 - Expanding, truncating
 - Addition, negation, multiplication, shifting
 - Summary
- Representations in memory, pointers, strings

Everything is bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
 - Computers determine what to do (instructions)
 - ... and represent and manipulate numbers, sets, strings, etc...
- Why bits? Electronic Implementation
 - Easy to store with bistable elements
 - Reliably transmitted on noisy and inaccurate wires



For example, can count in binary

Base 2 Number Representation

- Represent 15213₁₀ as 11101101101101₂
- Represent 1.20₁₀ as 1.0011001100110011[0011]...₂
- Represent 1.5213 X 10⁴ as 1.1101101101101₂ X 2¹³

Encoding Byte Values

- Byte = 8 bits
 - Binary 000000002 to 111111112
 - Decimal: 010 to 25510
 - Hexadecimal 00₁₆ to FF₁₆
 - Base 16 number representation
 - Use characters '0' to '9' and 'A' to 'F'
 - Write FA1D37B₁₆ in C as
 - 0xFA1D37B
 - 0xfa1d37b

Hex Decimal

0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
В	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Example Data Representations

C Data Type	Typical 32-bit	Typical 64-bit	x86-64
char	1	1	1
short	2	2	2
int	4	4	4
long	4	8	8
float	4	4	4
double	8	8	8
long double	-	-	10/16
pointer	4	8	8

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Boolean Algebra

Developed by George Boole in 19th Century

- Algebraic representation of logic
 - Encode "True" as 1 and "False" as 0

And

Or

■ A&B = 1 when both A=1 and B=1

■ A | B = 1 when either A=1 or B=1

ı	0	1
0	0	1
1	1	1

Not

Exclusive-Or (Xor)

■ ~A = 1 when A=0

■ A^B = 1 when either A=1 or B=1, but not both

٨	0	1
0	0	1
1	1	0

General Boolean Algebras

- Operate on Bit Vectors
 - Operations applied bitwise

```
01101001 01101001 01101001

& 01010101 | 01010101 ^ 01010101 ~ 01010101

01000001 01111101 00111100 1010101
```

All of the Properties of Boolean Algebra Apply

Example: Representing & Manipulating Sets

Representation

- Width w bit vector represents subsets of {0, ..., w−1}
- $a_i = 1 \text{ if } j \in A$
 - 01101001 { 0, 3, 5, 6 }
 - **76543210**
 - 01010101 { 0, 2, 4, 6 }
 - **76543210**

Operations

&	Intersection	01000001	{ 0, 6 }
• 1	Union	01111101	{ 0, 2, 3, 4, 5, 6 }
^	Symmetric difference	00111100	{ 2, 3, 4, 5 }
■ ~	Complement	10101010	{ 1, 3, 5, 7 }

Bit-Level Operations in C

- Operations &, |, ~, ^ Available in C
 - Apply to any "integral" data type
 - long, int, short, char, unsigned
 - View arguments as bit vectors
 - Arguments applied bit-wise
- Examples (Char data type)
 - $\sim 0x41 \rightarrow 0xBE$
 - $\sim 010000012 \rightarrow 101111102$
 - ~0x00 → 0xFF
 - $\sim 000000002 \rightarrow 1111111112$
 - $0x69 \& 0x55 \rightarrow 0x41$
 - $011010012 & 010101012 \rightarrow 010000012$
 - $0x69 \mid 0x55 \rightarrow 0x7D$
 - $011010012 \mid 010101012 \rightarrow 011111012$

Contrast: Logic Operations in C

- Contrast to Logical Operators
 - **&**&, ||, !
 - View 0 as "Fall
 - Anything nonzo
 - Alway
 - Early
- Example
 - !0x41
 - !0x00
 - !!0x41

Watch out for && vs. & (and || vs. |)... one of the more common oopsies in C programmingz

- $0x69 \&\& 0x55 \rightarrow 0x01$
- $0x69 | 1 | 0x55 \rightarrow 0x01$
- p && *p (avoids null pointer access)

Shift Operations

- Left Shift: x << y
 - Shift bit-vector x left y positions
 - Throw away extra bits on left
 - Fill with 0's on right
- Right Shift: x >> y
 - Shift bit-vector x right y positions
 - Throw away extra bits on right
 - Logical shift
 - Fill with 0's on left
 - Arithmetic shift
 - Replicate most significant bit on left

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J	HU			ч	שכ	ı ıa	VIU	, .

Shift amount < 0 or ≥ word size</p>

Argument x	01100010
<< 3	00010 <i>000</i>
Log. >> 2	00011000
Arith. >> 2	00011000

Argument x	10100010
<< 3	00010 <i>000</i>
Log. >> 2	<i>00</i> 101000
Arith. >> 2	<i>11</i> 101000

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Encoding Integers

Unsigned

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

Two's Complement

$$B2T(X) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$$

Sign Bit

C short 2 bytes long

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
У	-15213	C4 93	11000100 10010011

Sign Bit

- For 2's complement, most significant bit indicates sign
 - 0 for nonnegative
 - 1 for negative

Two's complement Encoding Example (Cont.)

x = 15213: 00111011 01101101y = -15213: 11000100 10010011

Weight	152	13	-152	213
1	1	1	1	1
2	0	0	1	2
4	1	4	0	0
8	1	8	0	0
16	0	0	1	16
32	1	32	0	0
64	1	64	0	0
128	0	0	1	128
256	1	256	0	0
512	1	512	0	0
1024	0	0	1	1024
2048	1	2048	0	0
4096	1	4096	0	0
8192	1	8192	0	0
16384	0	0	1	16384
-32768	0	0	1	-32768

Numeric Ranges

Unsigned Values

$$UMax = 2^w - 1$$

$$111...1$$

■ Two's Complement Values

■
$$TMin = -2^{w-1}$$
100...0

■
$$TMax = 2^{w-1} - 1$$

011...1

Other Values

Minus 1111...1

Values for W = 16

	Decimal	Hex	Binary
UMax	65535	FF FF	11111111 11111111
TMax	32767	7F FF	01111111 11111111
TMin	-32768	80 00	10000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	00000000 00000000

Values for Different Word Sizes

		W				
	8	16	32	64		
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615		
TMax	127	32,767	2,147,483,647	9,223,372,036,854,775,807		
TMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808		

Observations

- \blacksquare | TMin | = TMax + 1
 - Asymmetric range
- UMax = 2 * TMax + 1

C Programming

- #include <limits.h>
- Declares constants, e.g.,
 - ULONG_MAX
 - LONG_MAX
 - LONG_MIN
- Values platform specific

Unsigned & Signed Numeric Values

X	B2U(<i>X</i>)	B2T(<i>X</i>)	
0000	0	0	
0001	1	1	
0010	2	2	
0011	3	3	
0100	4	4	
0101	5	5	
0110	6	6	
0111	7	7	
1000	8	-8	
1001	9	- 7	
1010	10	– 6	
1011	11	- 5	
1100	12	-4	
1101	13	- 3	
1110	14	-2	
1111	15	-1	

Equivalence

Same encodings for nonnegative values

Uniqueness

- Every bit pattern represents unique integer value
- Each representable integer has unique bit encoding

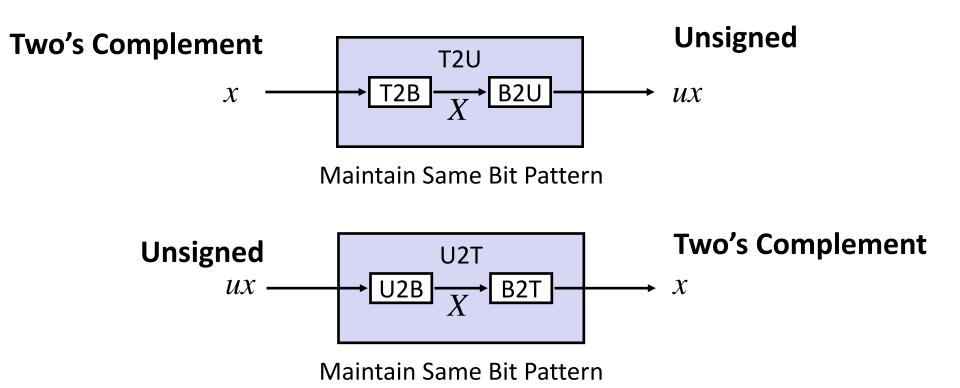
■ ⇒ Can Invert Mappings

- U2B(x) = B2U⁻¹(x)
 - Bit pattern for unsigned integer
- $T2B(x) = B2T^{-1}(x)$
 - Bit pattern for two's comp integer

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Mapping Between Signed & Unsigned

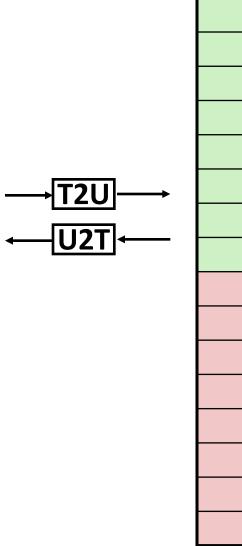


Mappings between unsigned and two's complement numbers:
 Keep bit representations and reinterpret

Mapping Signed ↔ Unsigned

Bits
0000
0001
0010
0011
0100
0101
0110
0111
1000
1001
1010
1011
1100
1101
1110
1111

Signed
0
1
2
3
4
5
6
7
-8
-7
-6
-5
-4
-3
-2
-1

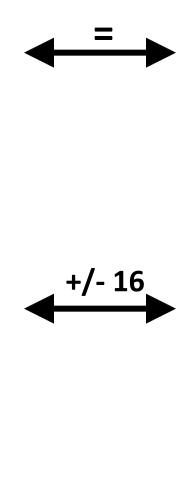


Unsigned
0
1
2
3
4
5
6
7
8
9
10
11
12
13
14
15

Mapping Signed ↔ Unsigned

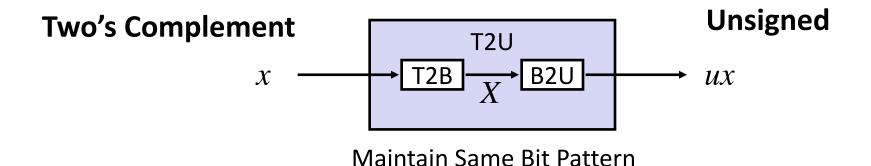
Bits
0000
0001
0010
0011
0100
0101
0110
0111
1000
1001
1010
1011
1100
1101
1110
1111

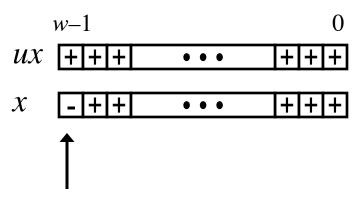
Signed	
0	
1	
2	
3	
4	
5	
6	
7	
-8	
-7	
-6	
-5	
-4	
-3	
-2	
-1	



Unsigned
0
1
2
3
4
5
6
7
8
9
10
11
12
13
14
15

Relation between Signed & Unsigned





Large negative weight becomes

Large positive weight

Conversion Visualized

2's Comp. \rightarrow Unsigned **UMax Ordering Inversion** UMax - 1Negative → Big Positive TMax + 1Unsigned TMax **TMax** Range 2's Complement Range

Signed vs. Unsigned in C

Constants

- By default are considered to be signed integers
- Unsigned if have "U" as suffixOU, 4294967259U

Casting

Explicit casting between signed & unsigned same as U2T and T2U

```
int tx, ty;
unsigned ux, uy;
tx = (int) ux;
uy = (unsigned) ty;
```

Implicit casting also occurs via assignments and procedure calls

```
tx = ux;

uy = ty;
```

Casting Surprises

Expression Evaluation

- If there is a mix of unsigned and signed in single expression, signed values implicitly cast to unsigned
- Including comparison operations <, >, ==, <=, >=
- **Examples for** W = 32: **TMIN = -2,147,483,648**, **TMAX = 2,147,483,647**

■ Constant ₁	Constant ₂	Relation	Evaluation
0	0U	==	unsigned
-1	0	<	signed
-1	0U	>	unsigned
2147483647	-2147483647-1	>	signed
2147483647U	-2147483647-1	<	unsigned
-1	-2	>	signed
(unsigned)-1	-2	>	unsigned
2147483647	2147483648U	<	unsigned
2147483647	(int) 2147483648U	>	signed

Summary Casting Signed ↔ Unsigned: Basic Rules

- Bit pattern is maintained
- But reinterpreted
- Can have unexpected effects: adding or subtracting 2^w
- Expression containing signed and unsigned int
 - int is cast to unsigned!!
 - You may want to give -Wall option to your compiler to recognize such implicit castings.

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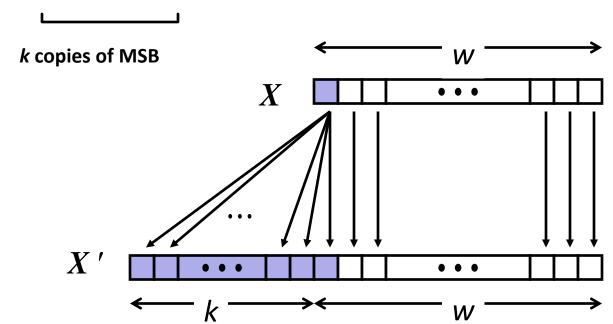
Sign Extension

Task:

- Given w-bit signed integer x
- Convert it to w+k-bit integer with same value

Rule:

- Make k copies of sign bit:
- $X' = X_{w-1}, ..., X_{w-1}, X_{w-1}, X_{w-2}, ..., X_0$



Sign Extension Example

```
short int x = 15213;
int        ix = (int) x;
short int y = -15213;
int        iy = (int) y;
```

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
ix	15213	00 00 3B 6D	00000000 00000000 00111011 01101101
У	-15213	C4 93	11000100 10010011
iy	-15213	FF FF C4 93	11111111 11111111 11000100 10010011

- Converting from smaller to larger integer data type
- C automatically performs sign extension

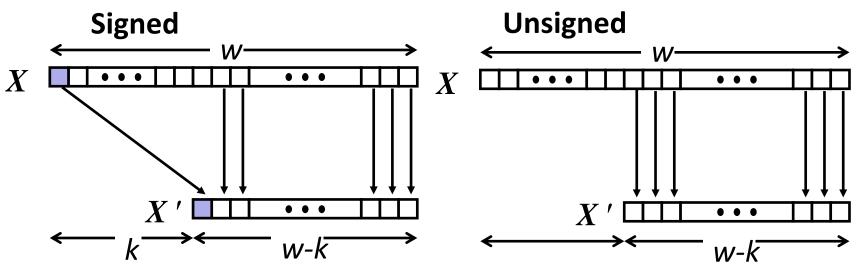
Truncation

Task:

- Given w-bit signed integer x
- Convert it to w-k-bit integer with same value

Rule:

- Preserve the sign bit for signed integers:
- $X' = x_{w-1}, x_{w-k-2}, x_{w-k-3}, ..., x_0$ If signed
- $X' = x_{w-k-1}, x_{w-k-2}, x_{w-k-3}, ..., x_0$ If unsigned



Summary: Expanding, Truncating: Basic Rules

- Expanding (e.g., short int to int)
 - Unsigned: zeros added
 - Signed: sign extension
 - Both yield expected result
- Truncating (e.g., unsigned to unsigned short)
 - Unsigned/signed: bits are truncated
 - Result reinterpreted
 - Unsigned: mod operation
 - Signed: similar to mod
 - For small numbers yields expected behavior