

## Part 1

### Task 1.1 Relation A - Employee

1) A superkey is any set of attributes that uniquely identifies a tuple. Based on the sample data, here are six examples:

- {EmpID, SSN, Email, Phone, Name, Department, Salary} (The entire relation)
- {EmpID}
- {SSN}
- {Email}
- {SSN, Phone}
- {EmpID, Email}
- {Email, Name} (Assuming Name is not unique, this may not be a superkey in a larger dataset. A safer superkey would be {SSN, Department})

2) A candidate key is a minimal superkey.

**EmpID:** Unique in the sample data.

**SSN:** Unique by definition (Social Security Number).

**Email:** Unique in the sample data (company email is typically unique per employee).

These are all single-attribute keys, so they are minimal.

3) I would choose **EmpID** as the primary key.

4) Based on the data shown, all phone numbers are unique. However, the sample size is very small. In a real-world scenario, it's entirely possible for two employees to share a phone number (e.g., a shared office line or a household with multiple employees). Therefore, **Phone cannot be considered a candidate key** based on this limited sample. The business rules would need to specify if this is allowed or not.

### Task 1.1 Relation B - Course Registration

1) StudentID, CourseCode, Section, Semester & Year

2)

**StudentID:** Identifies *which* student.

**CourseCode:** Identifies *which* course.

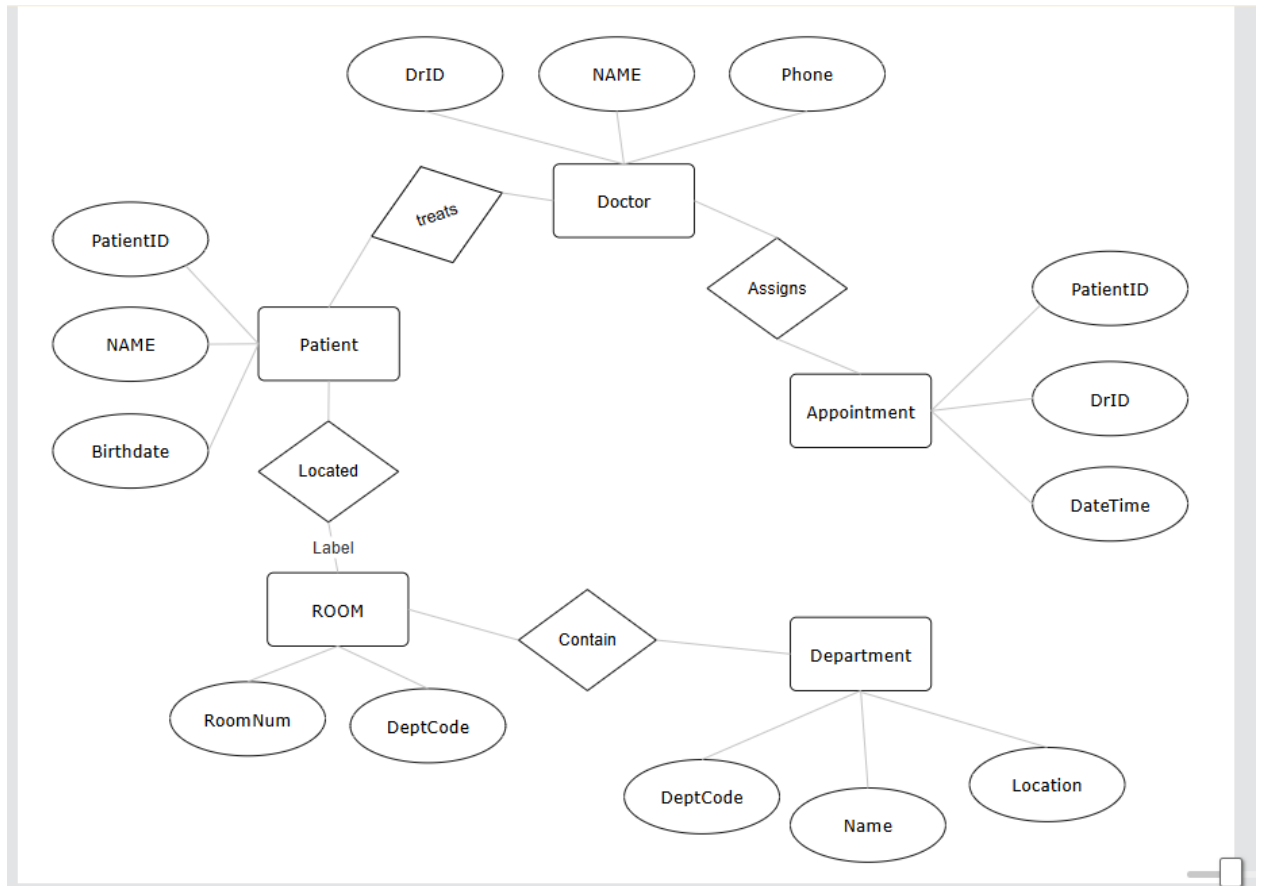
**Section:** Necessary because a course can have multiple sections (e.g., Lecture 1, Lab 3) in the same semester.

**Semester & Year:** Necessary because a student can re-take the same course (and section) in a future term. Without these, the combination of (StudentID, CourseCode, Section) would not be unique over time.

3) Another potential candidate key could be a synthetic key like RegistrationID. However, based on the given attributes, the composite key {**StudentID, CourseCode, Section, Semester, Year**} is the only natural candidate key.

## Part 2

### Task 2.1



**Patient (Strong)**

**Doctor (Strong)**

**Department (Strong)**

**Appointment (Weak)**

**Prescription (Weak)**

**Room (Strong)**

**Phone (Weak, Multi-valued)**

#### Attribute Classification:

**Composite:** Patient Address -> {Street, City, State, Zip}

**Multi-valued:** Doctor.Specialization. This would be modeled as a separate weak entity Specialization(DoctorID, Specialization).

**Derived:** (Possible) Patient.Age (derived from Birthdate).

#### Relationships & Cardinalities:

Patient *makes* Appointment (1:N) (*A patient can have many appointments, an appointment is for one patient*)

Doctor *has* Appointment (1:N) (*A doctor can have many appointments, an appointment is with one doctor*)

Doctor *works\_in* Department (N:1) (*A doctor works in one department, a department has many doctors*)

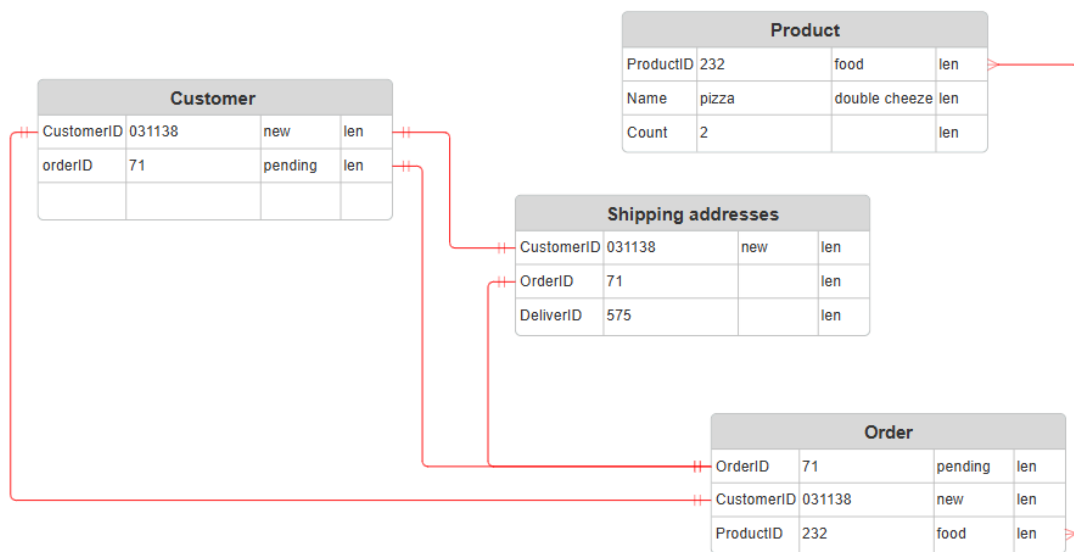
Doctor *prescribes* Prescription (1:N)

Patient *is\_prescribed* Prescription (1:N)

Room *is\_located\_in* Department (N:1) (*A room belongs to one department, a department has many rooms*)

Patient *has* Phone (1:N) (*Modeling the multi-valued attribute as an entity*)

## Task 2.2



**Weak Entity: OrderItem.** It is weak because its existence is dependent on the Order entity. An OrderItem cannot exist without an Order. Its primary key would be a composite of OrderID (from its owner, Order) and ProductID or a line item number.

**Many-to-Many with Attributes:** The relationship between Order and Product is M:N. This relationship itself has attributes Quantity and PriceAtTimeOfOrder. This is precisely why we create the associative entity OrderItem to hold these attributes.

## Part4 : Normalization Workshop

### Task 4.1

1. Functional Dependencies (FDs):

- StudentID  $\rightarrow$  StudentName, StudentMajor
- ProjectID  $\rightarrow$  ProjectTitle, ProjectType, SupervisorID
- SupervisorID  $\rightarrow$  SupervisorName, SupervisorDept

- (StudentID, ProjectID) → Role, HoursWorked, StartDate, EndDate
2. Redundancy and anomalies:
- Redundancy: StudentName and StudentMajor repeated for each project of a student; SupervisorName and SupervisorDept repeated for each project supervised by same supervisor.
  - Update anomaly example: Changing SupervisorName requires updating many rows.
  - Insert anomaly example: To add a new supervisor with no project yet, ProjectID required (depending on PK design).
  - Delete anomaly example: Deleting last project by a student may remove supervisor info if stored only in that row.
3. 1NF: No repeating groups apparent; ensure multi-valued attributes (if any) are moved to separate tables (e.g., multiple supervisors or roles). Assume table is in 1NF.
4. 2NF: Primary key = (StudentID, ProjectID) assuming each student may work on multiple projects. Partial dependencies: StudentID → StudentName, StudentMajor (depends only on StudentID) and ProjectID → ProjectTitle, ProjectType, SupervisorID (depends only on ProjectID).

2NF decomposition:

- Student(StudentID PK, StudentName, StudentMajor)
  - Project(ProjectID PK, ProjectTitle, ProjectType, SupervisorID)
  - StudentProject(StudentID FK, ProjectID FK, Role, HoursWorked, StartDate, EndDate)
5. 3NF: Transitive dependency: SupervisorID → SupervisorName, SupervisorDept in Project table; to remove transitive dependency, create Supervisor/Professor table:
- Supervisor(SupervisorID PK, SupervisorName, SupervisorDept)
- Final 3NF tables: Student, Supervisor, Project (with SupervisorID FK), StudentProject (associative).

## Task 4.2

1. Primary key: (StudentID, CourseID, TimeSlot) or more precisely (StudentID, CourseSectionID) if a CourseSectionID exists. Reason: student can take multiple course sections; a course section is uniquely defined by CourseID+TimeSlot+Room (or a SectionID).
2. Functional dependencies:
- StudentID → StudentMajor
  - CourseID → CourseName
  - InstructorID → InstructorName
  - Room → Building (rooms unique across campus)
  - (CourseID, TimeSlot, Room) → InstructorID (each course section taught by one instructor at one time in one room)

3. BCNF check: The table is not in BCNF because  $\text{StudentID} \rightarrow \text{StudentMajor}$  (non-key  $\rightarrow$  attribute) and  $\text{Room} \rightarrow \text{Building}$  are FDs violating BCNF if keys include  $\text{StudentID} + \text{CourseID} + \text{TimeSlot}$ .

4. BCNF decomposition:

- $\text{Student}(\text{StudentID PK}, \text{StudentMajor})$
- $\text{Course}(\text{CourseID PK}, \text{CourseName})$
- $\text{Instructor}(\text{InstructorID PK}, \text{InstructorName})$
- $\text{Room}(\text{Room PK}, \text{Building})$
- $\text{CourseSection}(\text{CourseSectionID PK}, \text{CourseID FK}, \text{InstructorID FK}, \text{TimeSlot}, \text{Room FK})$
- $\text{Enrollment}(\text{StudentID FK}, \text{CourseSectionID FK}, \text{PRIMARY KEY}(\text{StudentID}, \text{CourseSectionID}))$

5. Decomposition is lossless because  $\text{CourseSection}$  references original course+time+room combination;  $\text{Enrollment}$  links students to sections. No information loss if FKs are maintained; ensure data migration preserves  $\text{CourseSection}$  identity.