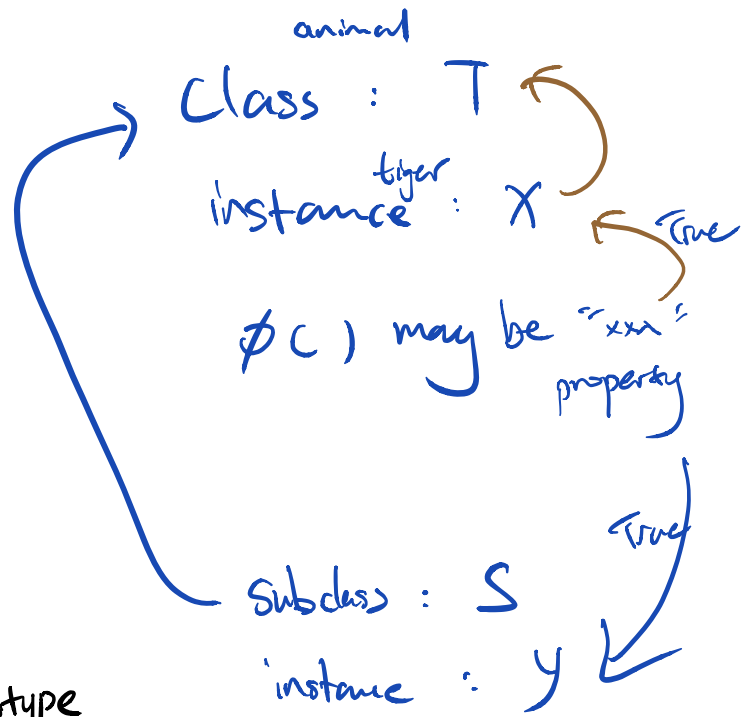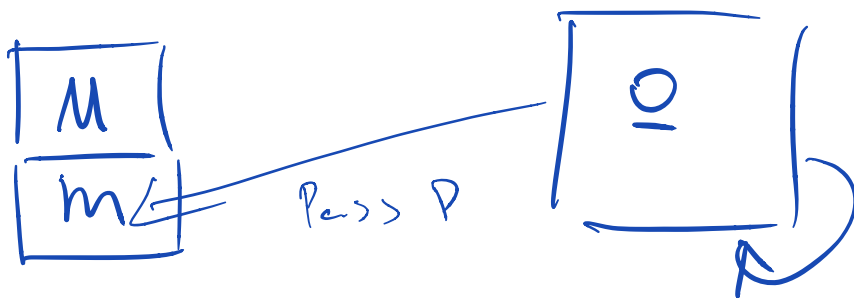# Liskov's Substitution Principle

invaendor

Let $\phi(x)$ be a property → instance
provable about objects x of
type T. Then $\phi(y)$ should
be true for objects y of
type S where S is a subtype
where S is a subtype of T

animal
Class : T
                tiger
instance : X
                        True
$\phi()$ may be "xxx"
                    property
                        True
Subclass : S
instance : y

$$B \xrightarrow{\ F\ } f_{oo}$$

M
m
Pass P
O

UI Class

- └→ Separated _ Line
- └→ Out_After_Add
- └→ Greeting
- └→ Out_After_Delete
- └→ print_List

Storage

- └→ Task_File
- └→ Loady_Task
- └→ Saviy_Task
- └→ Encoder
- └→ Decoder

TaskList

- └→ Task Class
- └→ Todo Class
- └→ Event Class
- └→ Deadline Class
- └→ Add Task
- └→ Delete Task
- └→ Update Task
- └→ Description ( )
- └→ Time ( )

Parser

- └→ Todo_Input_Check
- └→ E_D_Input_check
- └→ Month_Check
- └→ Time_check
- └→ Input_length_Check
- └→ IS Numeric
- └→ Done & Delete Number Check

Month_Enum

{ 1 ~ 12 Month }

Task_Enum

- └→ { Todo, Event, Deadline }
- └→ Not Case Sensitive.