

*Rms*

# Multiplexers/Demultiplexers

TEKNIK INFORMATIKA

# Overview

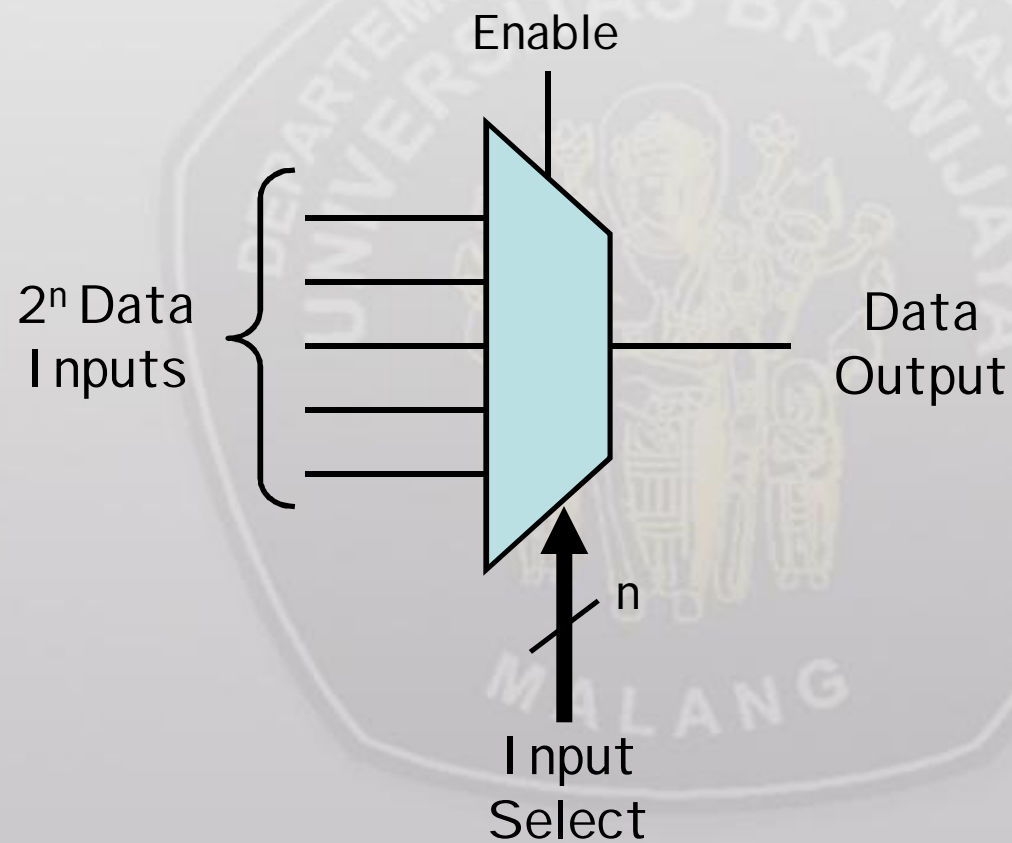
- Multiplexers (MUXs)
  - Functionality
  - Circuit implementation with MUXs
- Demultiplexers

# Multiplexer

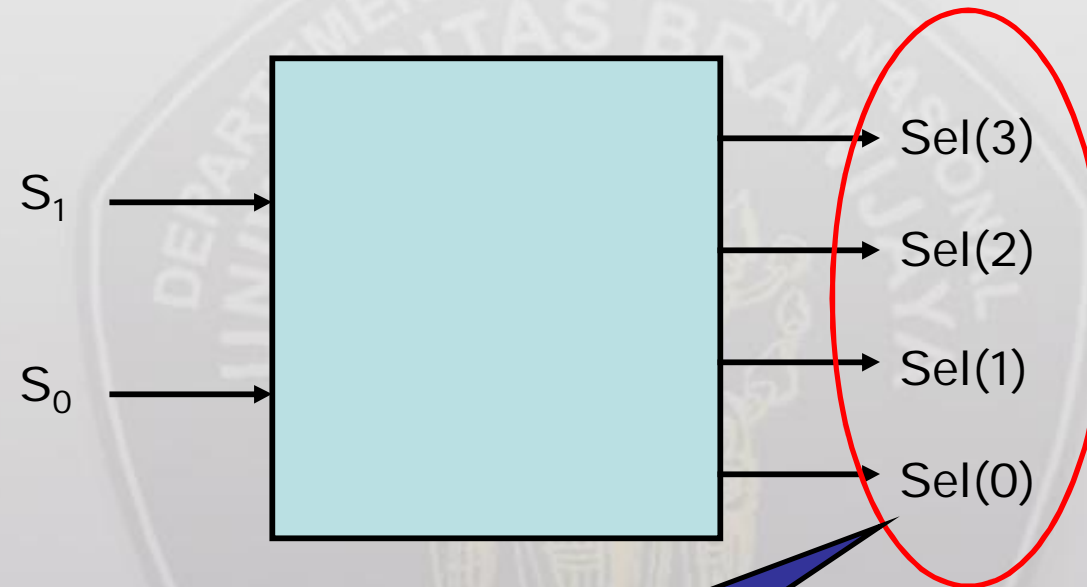


- “Selects” binary information from one of many input lines and directs it to a single output line.
- Also known as the “selector” circuit,
- Selection is controlled by a particular set of inputs lines whose # depends on the # of the data input lines.
- For a  $2^n$ -to-1 multiplexer, there are  $2^n$  data input lines and  $n$  selection lines whose bit combination determines which input is selected.

# MUX

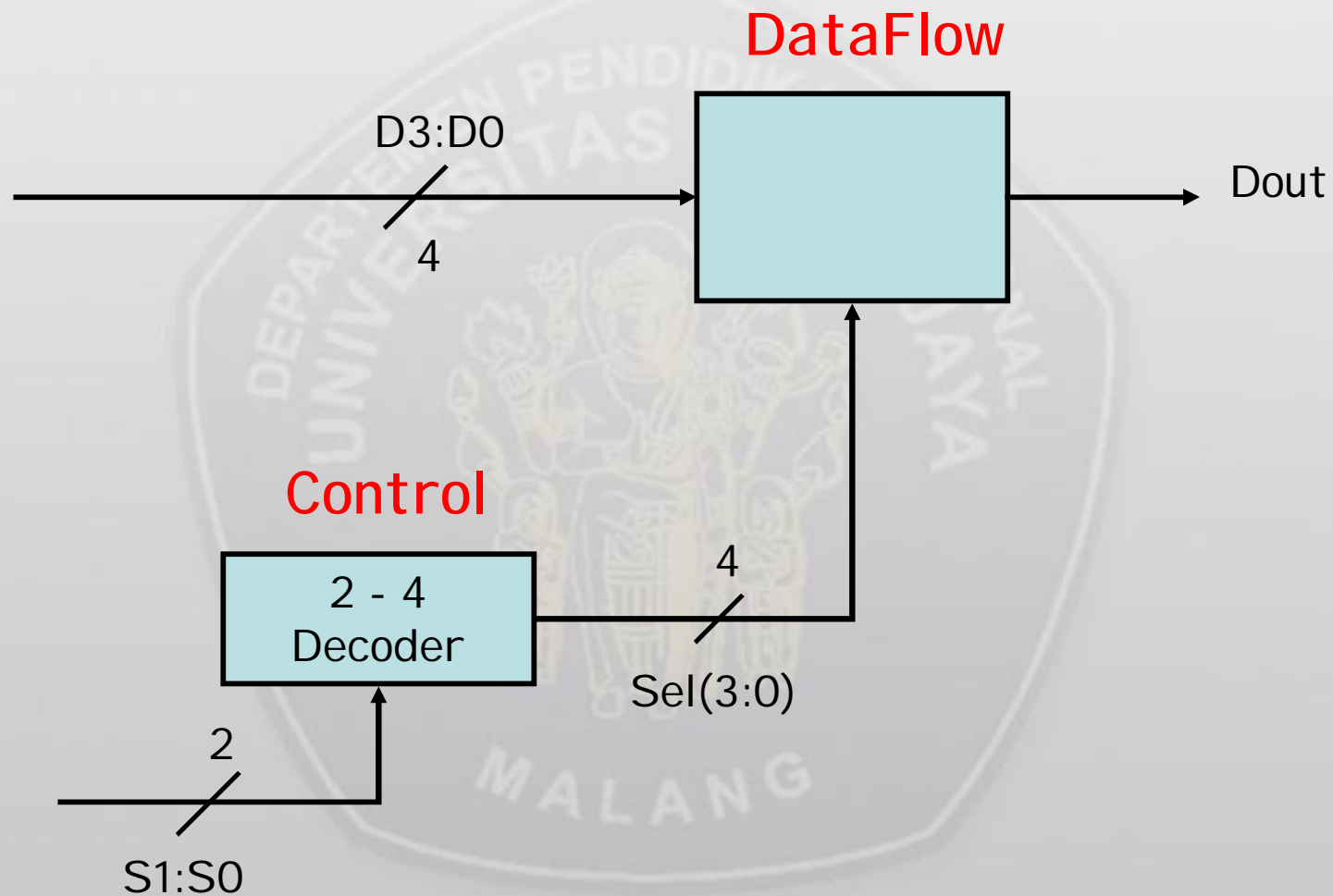


# Remember the 2 – 4 Decoder?

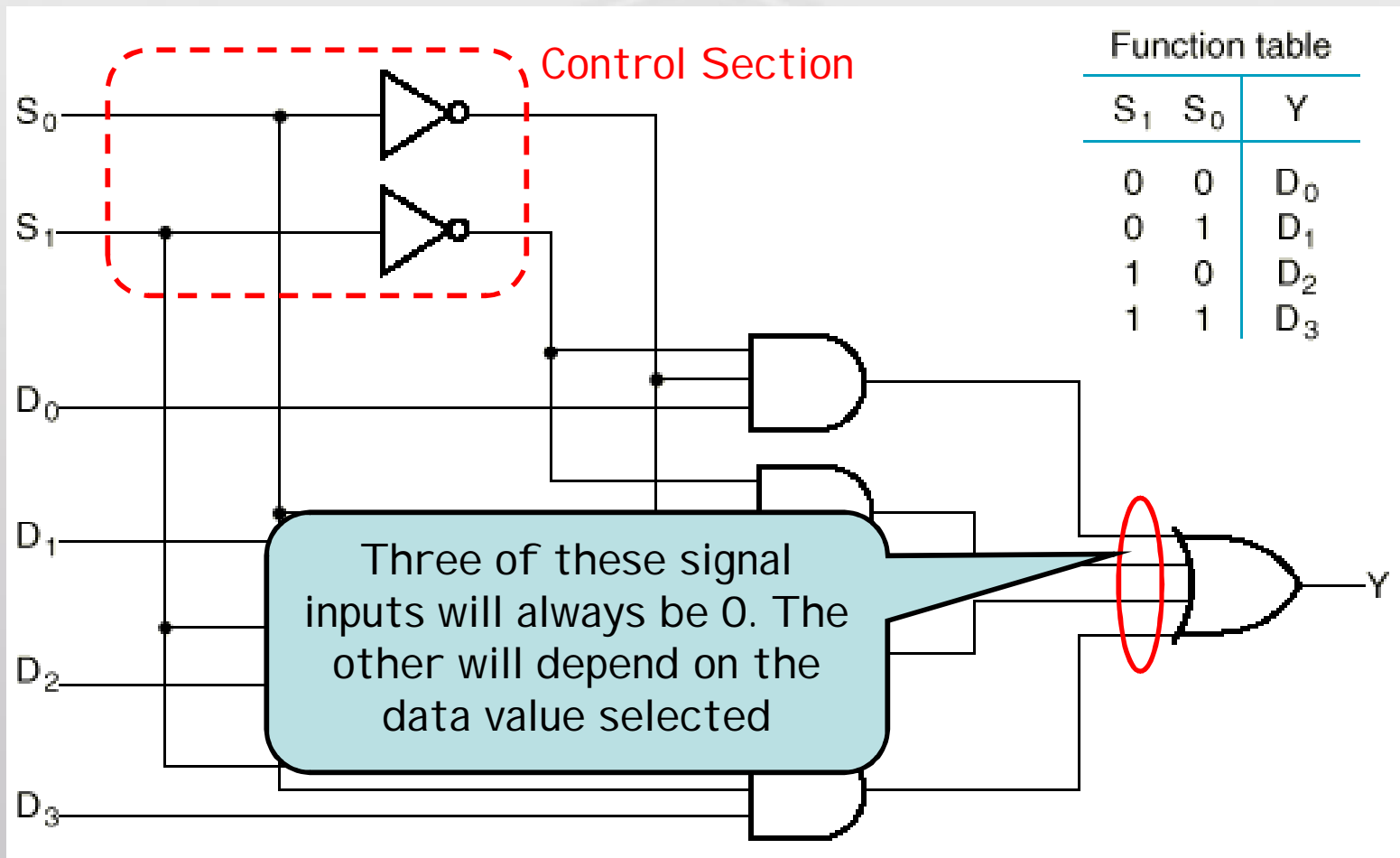


Mutually Exclusive  
(Only one O/P asserted at  
any time)

# 4 to 1 MUX



# 4-to-1 MUX (Gate level)



## Multiplexer (cont.)

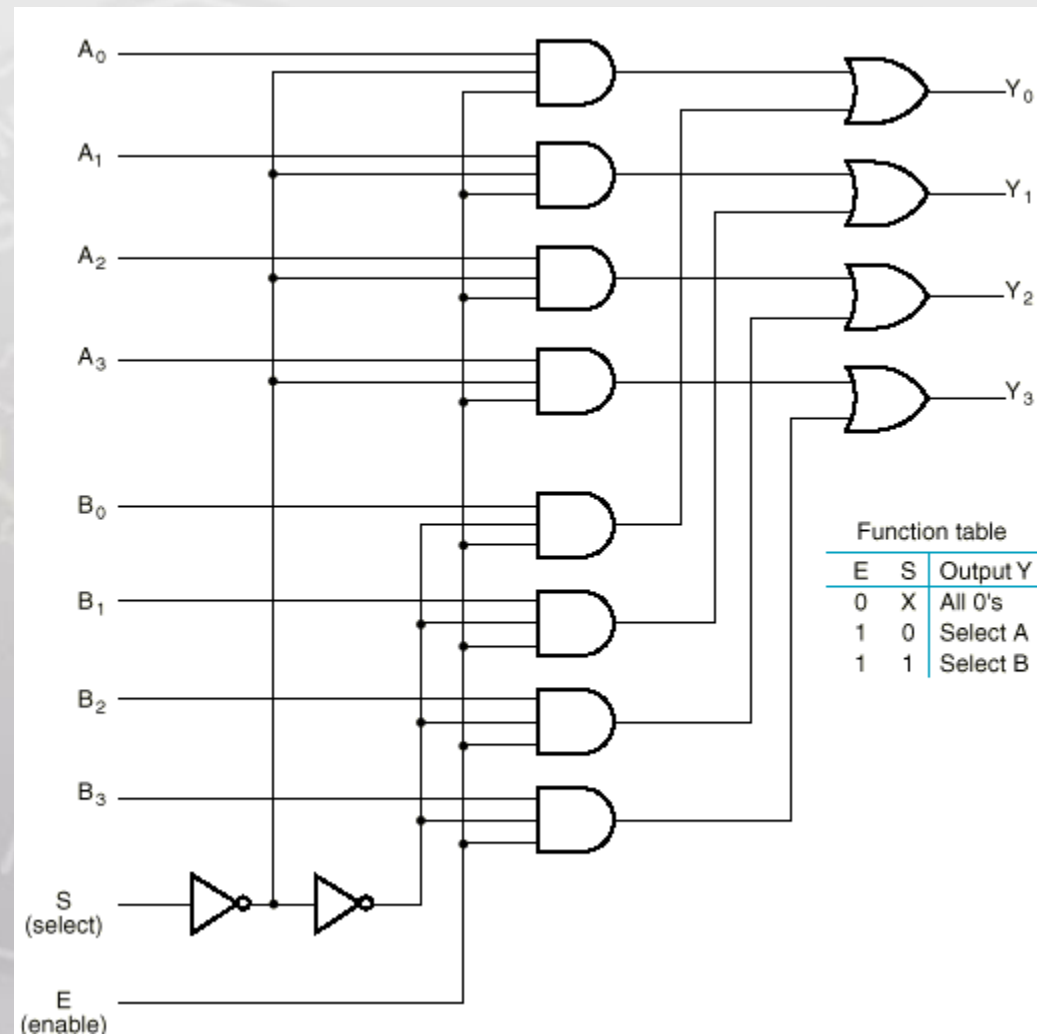
- Until now, we have examined single-bit data selected by a MUX. What if we want to select m-bit data/words?  
à Combine MUX blocks in parallel with common select and enable signals
- Example: Construct a logic circuit that selects between 2 sets of 4-bit inputs (see next slide for solution).



# Example: Quad 2-to-1 MUX



- Uses four 4-to-1 MUXs with common select (S) and enable (E).
- Select line chooses between  $A_i$ 's and  $B_i$ 's. The selected four-wire digital signal is sent to the  $Y_i$ 's
- Enable line turns MUX on and off ( $E=1$  is on).



# Implementing Boolean functions with Multiplexers



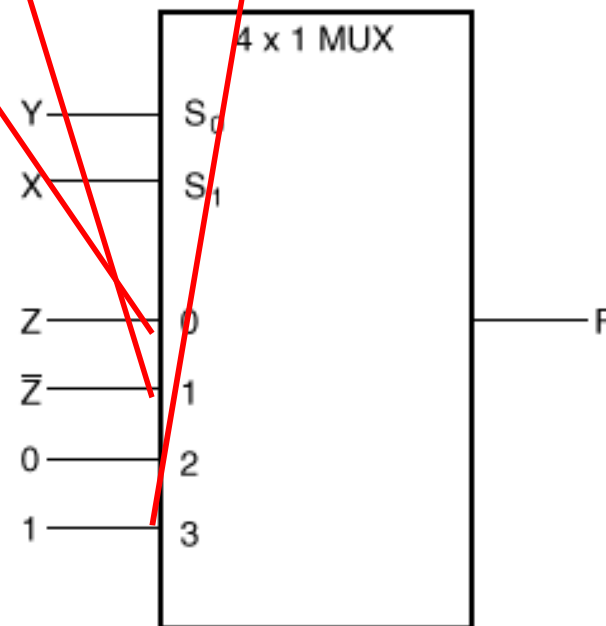
- Any Boolean function of  $n$  variables can be implemented using a  $2^{n-1}$ -to-1 multiplexer. A MUX is basically a decoder with outputs ORed together, hence this isn't surprising.
- The SELECT signals generate the minterms of the function.
- The data inputs identify which minterms are to be combined with an OR.

# Example

- $F(X,Y,Z) = X'Y'Z + X'YZ' + XYZ' + XYZ = \Sigma m(1,2,6,7)$
- There are  $n=3$  inputs, thus we need a  $2^2$ -to-1 MUX
- The first  $n-1$  ( $=2$ ) inputs serve as the selection lines

| X | Y | Z | F |               |
|---|---|---|---|---------------|
| 0 | 0 | 0 | 0 | $F = Z$       |
| 0 | 0 | 1 | 1 |               |
| 0 | 1 | 0 | 1 | $F = \bar{Z}$ |
| 0 | 1 | 1 | 0 |               |
| 1 | 0 | 0 | 0 | $F = 0$       |
| 1 | 0 | 1 | 0 |               |
| 1 | 1 | 0 | 1 | $F = 1$       |
| 1 | 1 | 1 | 1 |               |

(a) Truth table



(b) Multiplexer implementation

# Efficient Method for implementing Boolean functions



- For an  $n$ -variable function (e.g.,  $f(A,B,C,D)$ ):
  - Need a  $2^{n-1}$  line MUX with  $n-1$  select lines.
  - Enumerate function as a truth table with consistent ordering of variables (e.g.,  $A,B,C,D$ )
  - Attach the most significant  $n-1$  variables to the  $n-1$  select lines (e.g.,  $A,B,C$ )
  - Examine pairs of adjacent rows (only the least significant variable differs, e.g.,  $D=0$  and  $D=1$ ).
  - Determine whether the function output for the  $(A,B,C,0)$  and  $(A,B,C,1)$  combination is  $(0,0)$ ,  $(0,1)$ ,  $(1,0)$ , or  $(1,1)$ .
  - Attach 0,  $D$ ,  $D'$ , or 1 to the data input corresponding to  $(A,B,C)$  respectively.

# Another Example



- Consider  $F(A,B,C) = \sum m(1,3,5,6)$ . We can implement this function using a 4-to-1 MUX as follows.
- The index is ABC. Apply A and B to the  $S_1$  and  $S_0$  selection inputs of the MUX (A is most sig,  $S_1$  is most sig.)
- Enumerate function in a truth table.

# MUX Example (cont.)



When  $A=B=0$ ,  $F=C$  — [

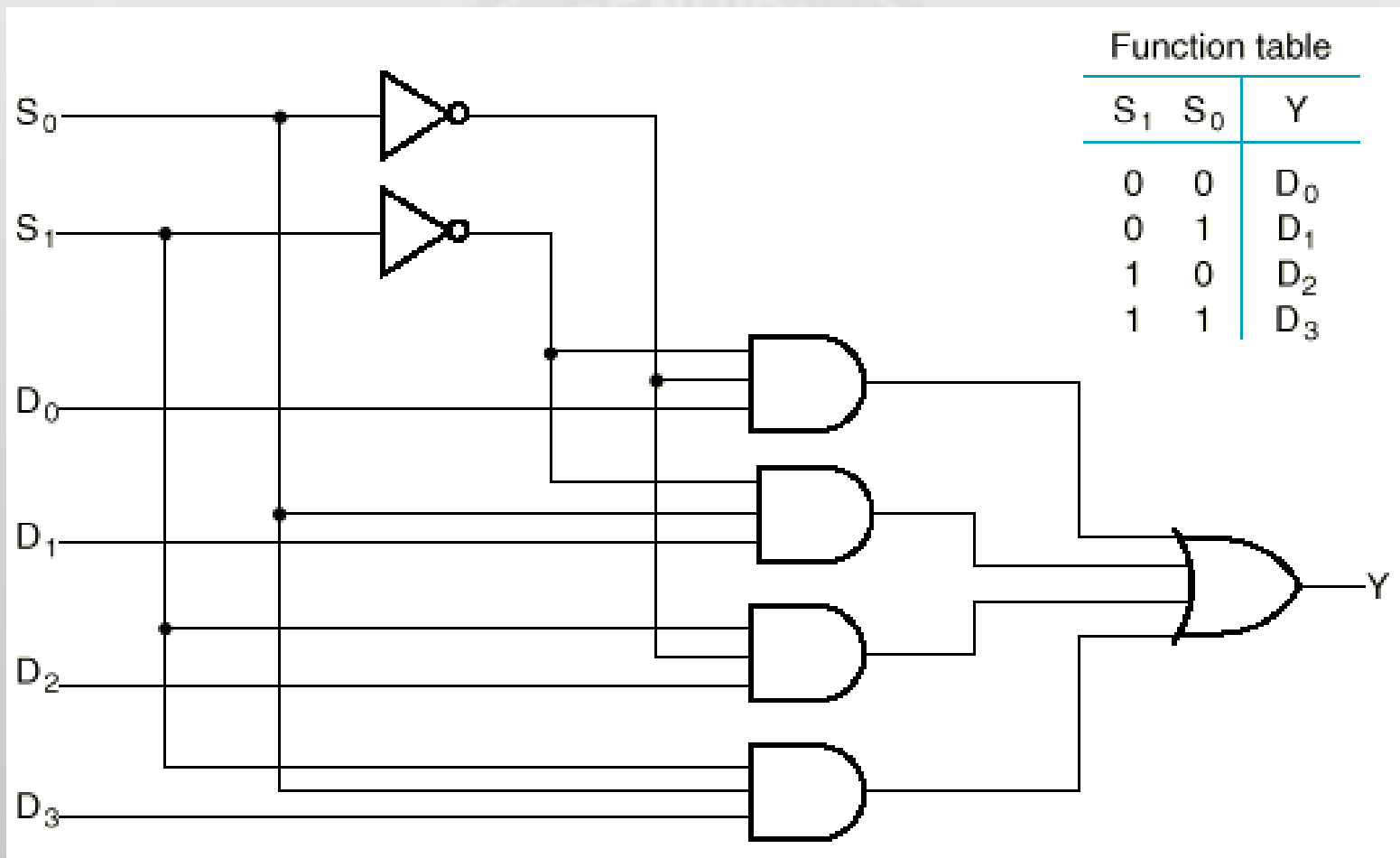
When  $A=0$ ,  $B=1$ ,  $F=C$  — [

When  $A=1$ ,  $B=0$ ,  $F=C$  — [

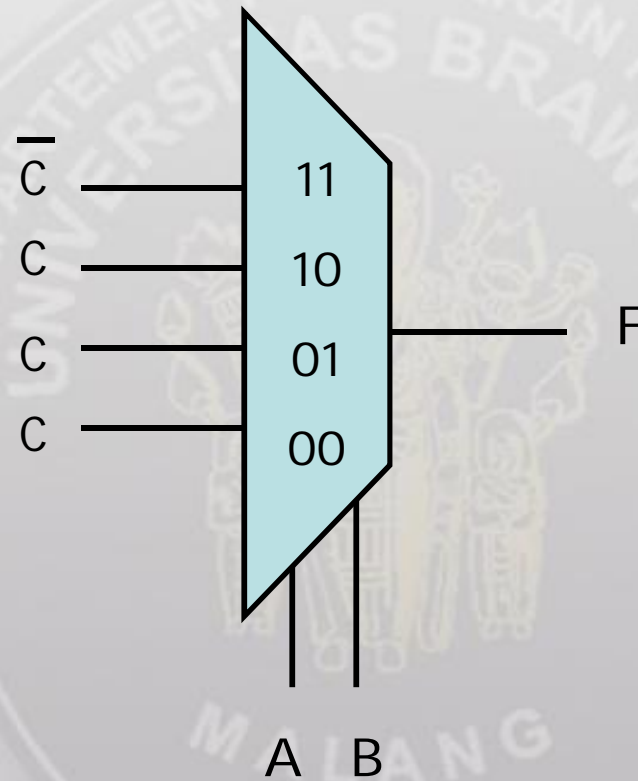
When  $A=B=1$ ,  $F=C'$  — [

| A | B | C |  | F |
|---|---|---|--|---|
| 0 | 0 | 0 |  | 0 |
| 0 | 0 | 1 |  | 1 |
| 0 | 1 | 0 |  | 0 |
| 0 | 1 | 1 |  | 1 |
| 1 | 0 | 0 |  | 0 |
| 1 | 0 | 1 |  | 1 |
| 1 | 1 | 0 |  | 1 |
| 1 | 1 | 1 |  | 0 |

# MUX implementation of $F(A,B,C) = \sum m(1,3,5,6)$



# Or Simply....

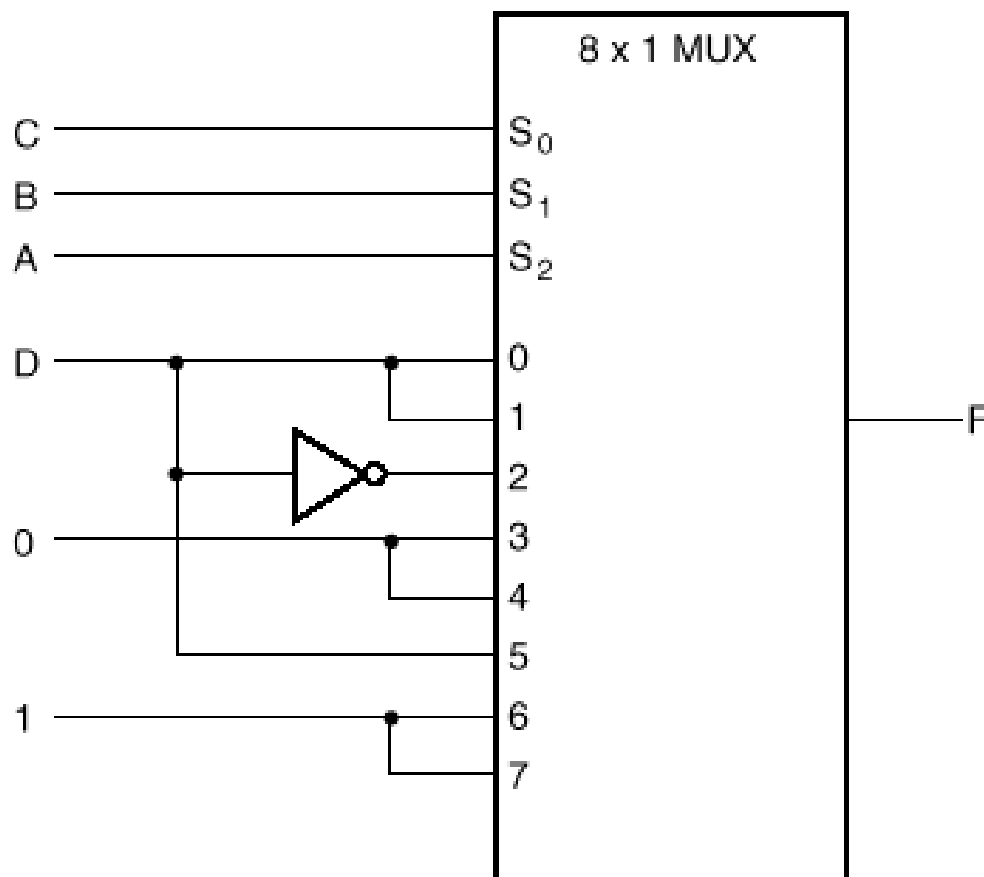




# A larger Example

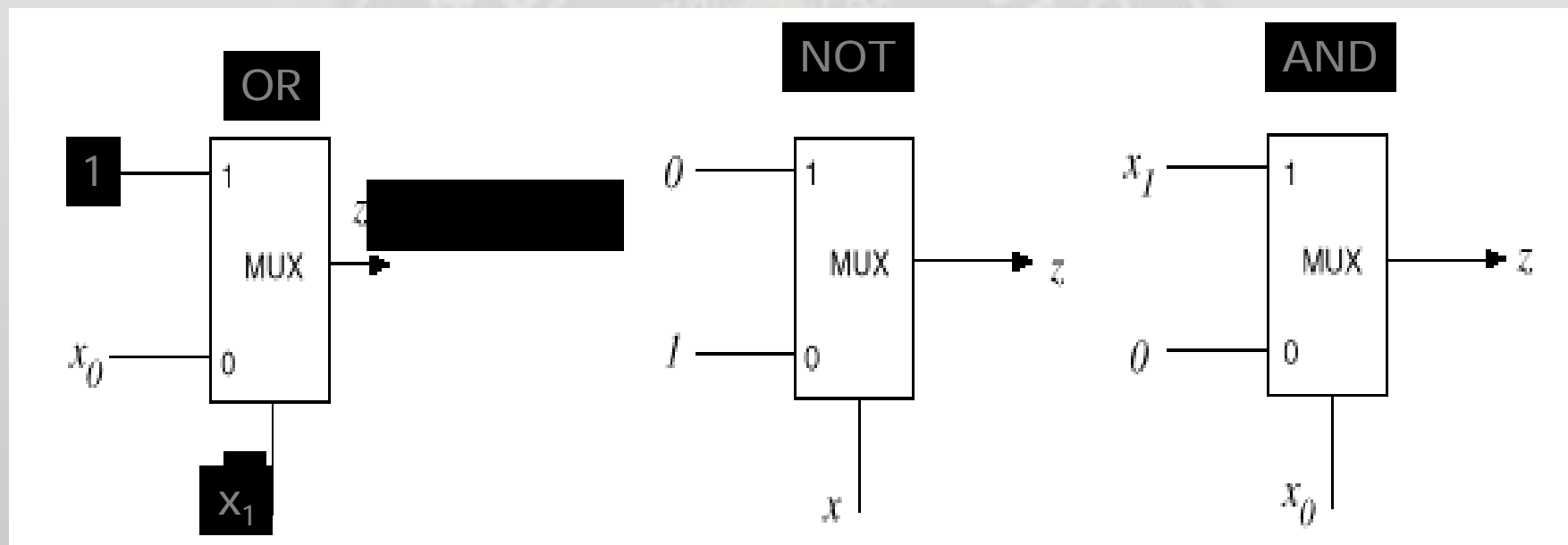


| A | B | C | D | F |               |
|---|---|---|---|---|---------------|
| 0 | 0 | 0 | 0 | 0 | $F = D$       |
| 0 | 0 | 0 | 1 | 1 |               |
| 0 | 0 | 1 | 0 | 0 | $F = D$       |
| 0 | 0 | 1 | 1 | 1 |               |
| 0 | 1 | 0 | 0 | 1 | $F = \bar{D}$ |
| 0 | 1 | 0 | 1 | 0 |               |
| 0 | 1 | 1 | 0 | 0 | $F = 0$       |
| 0 | 1 | 1 | 1 | 0 |               |
| 1 | 0 | 0 | 0 | 0 | $F = 0$       |
| 1 | 0 | 0 | 1 | 0 |               |
| 1 | 0 | 1 | 0 | 0 | $F = D$       |
| 1 | 0 | 1 | 1 | 1 |               |
| 1 | 1 | 0 | 0 | 1 | $F = 1$       |
| 1 | 1 | 0 | 1 | 1 |               |
| 1 | 1 | 1 | 0 | 1 | $F = 1$       |
| 1 | 1 | 1 | 1 | 1 |               |



# MUX as a Universal Gate

- We can construct OR, AND, and NOT gates using 2-to-1 MUXs. Thus, 2-to-1 MUX is a universal gate.



$$Z = X_1 + X_1'X_0$$

$$= X_1X_0' + X_1X_0 + X_1'X_0 = X_1 + X_0$$

$$Z = 0X + 1X' = X'$$

$$Z = X_1X_0 + 0X_0' = X_1X_0$$

# Demultiplexers (DMUX)



- Performs the inverse of a multiplexing operation:
  - Receives data from a single line
  - Transmit it to one of the  $2^n$  possible output lines
  - Selection of a specific output is controlled by the  $n$  select lines
  - Demultiplexers are basically decoders! For example, a 2-to-4 DMUX is a 2-to-4 decoder with enable input.