

Agenda

DAY **3**

8

**Object Oriented Programming (OOP)
– Inheritance**

9

Inter-Thread Communications



10

Functional Coverage



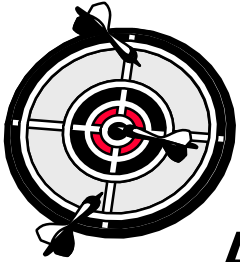
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SystemVerilog UVM Preview

CS

Customer Support

Unit Objectives



After completing this unit, you should be able to:

- **Establish order of execution using events**
- **Avoid resource collision with Semaphores**
- **Pass data between threads via Mailbox**

Inter-Thread Communications (ITC)

- Concurrent threads require communication to establish control for sequence of execution
- Three types are covered in this unit

Event based



Resource sharing



Data passing



Event Based ITC



- **Synchronize operation of concurrent threads via event variables:**
 - Thread **waits** for event to be triggered
 - An executing thread triggers the event to enable the waiting thread to be placed into the READY queue
- **Mainly used to control sequence of execution**

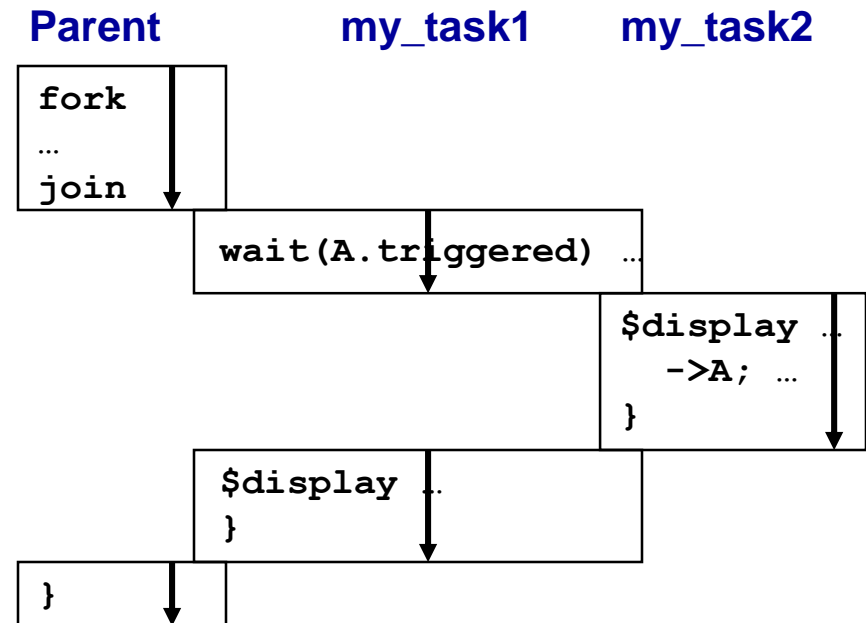
Event Based ITC Example

Controlling order of execution.

```
initial begin
    event A;
    fork
        my_task1();
        my_task2();
    join
end

task my_task1();
    wait(A.triggered);
    $display("Print 2nd");
endtask: my_task1

task my_task2();
    $display("Print 1st");
    ->A;
endtask: my_task2
```



Event Wait Syntax

■ Wait until event has been triggered (one shot):

```
@(event_var [, event2]);
```

- Will be satisfied by events which occur after the execution of this statement

■ Wait until event has been triggered (persistent):

```
wait(event_var.triggered);
```

- Similar to @(event_var), but will also be satisfied by events which happened earlier during the same simulation time
- Can be used to eliminate potential race condition

Trigger Syntax

`->eventN;`

- event variables are handles
- Assigning an event to another makes both event the same event

Example:

```
event a, b, c;  
a = b;  
-> c;  
-> a; // also triggers b  
-> b; // also triggers a  
a = c;  
b = a;  
-> a; // also triggers b and c  
-> b; // also triggers a and c  
-> c; // also triggers a and b
```

Controlling Termination of Simulation

```
program automatic test(router_io.TB rtr_io);  
    event DONE;  
    initial begin  
        fork  
            gen();  
            check();  
        join_none  
        for (int i=0; i<16; i++) begin  
            int j = i;  
            fork  
                send(j);  
                recv(j);  
            join_none  
        end  
        wait(DONE.triggered);  
    end  
    ...  
endprogram;
```

```
task check();  
    forever begin  
        ...  
        if ($get_coverage() == 100)  
            ->DONE;  
    end  
endtask: check;
```

Blocking statement to prevent termination of simulation until done

Trigger event when termination condition is detected

Resource Sharing ITC



■ Synchronize operation of concurrent threads via access to shared resources:

- Thread requests for a shared resource before executing a critical section of code
- Thread waits if the requested resource is unavailable
- Thread resumes execution when the requested resource becomes available

Semaphores

- **A Semaphore is a bucket in which keys can be deposited and removed:**
 - A thread tries to **acquire keys** from a semaphore bucket
 - Thread execution **waits** until keys requested are available in the semaphore bucket
 - When the requested keys becomes available, thread execution **resumes** and the requested keys removed from the semaphore bucket
- **Mainly used to prevent multiple threads from accessing the same hardware signal or using same software resource**

Semaphores

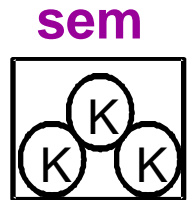
- Semaphores are supported via built-in semaphore class:

```
class semaphore;  
    function new(int keyCount = 0);  
    task put(int keyCount = 1);  
  
    // the following is blocking:  
    task get(int keyCount = 1);  
  
    // the following is non-blocking: 1 for success 0 for failure  
    function int try_get(int keyCount = 1);  
endclass
```

Using Semaphores

■ Semaphore buckets created using the constructor `new()`

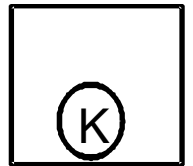
```
semaphore sem  
sem = new(3);           // bucket with 3 keys is created
```



`new(3)`

■ Acquiring Keys

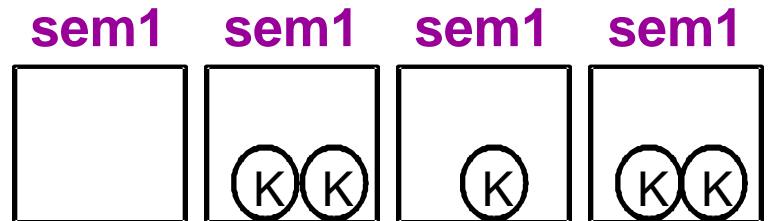
```
sem.get(2);              // bucket with 1 key left  
if (!sem.try_get(2))    // does not block  
    $display("Failed");  
sem.get(2);              //blocks until 2 keys available
```



`get(2)`

■ Returning/Creating Keys

```
semaphore sem1  
sem1 = new(); // 0 keys  
sem1.put(2);  // Create 2 keys  
sem1.get(1);  // Use 1 key  
sem1.put(1);  // Return 1 key
```

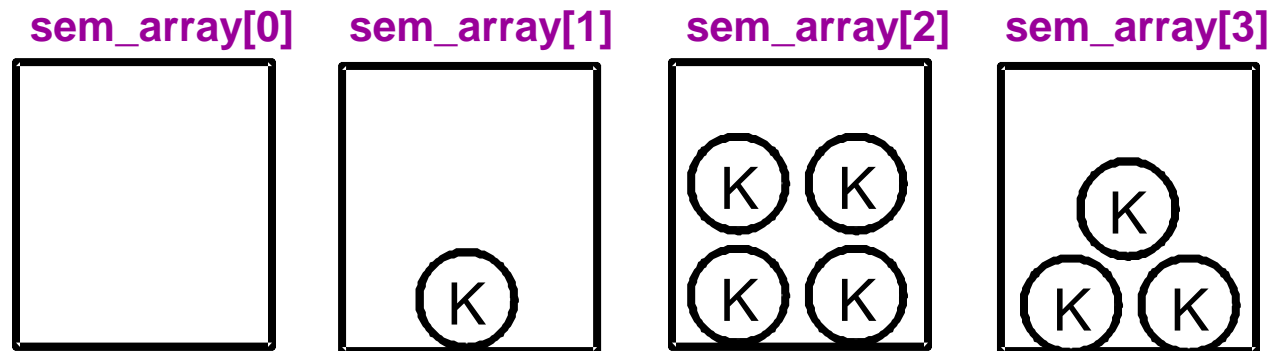


`new()` `put(2)` `get(1)` `put(1)`

Using Semaphore Arrays

- Semaphores in an array must be constructed before use

```
semaphore sem_array[];  
sem_array = new[4]; // array of 4 null semaphore handles created  
sem_array[2].put(2); // ERROR: semaphore bucket does not exist!  
foreach (sem_array[i])  
    sem_array[i] = new(i); // construct semaphore with i keys  
sem_array[2].put(2); // okay. sem_array[2] now has 4 keys
```



Arbitration Example

```
program automatic test(router_io.TB rtr_io);  
  semaphore sem[];  
  ...;  
  sem = new[16];  
  foreach(sem[i])  
    sem[i] = new(1);  
  ...;  
  task send();  
    sem[da].get(1);  
    send_addrs();  
    send_pad();  
    send_payload();  
    sem[da].put(1);  
  endtask: send  
  ...;  
endprogram: test
```

Create semaphore array to represent each output port

Construct each individual semaphore bucket

Block if others are driving the chosen output port

Re-deposit keys when done with driving port

Mailbox



■ Messages are passed between threads via mailbox:

- A thread **sends data** by putting messages a mailbox
- A thread **retrieves data** by getting messages from the mailbox
- If a message is not available, the thread can **wait** until there is a message to retrieve
- Thread **resumes** execution once the message becomes available

■ Mainly used for passing data between concurrent threads

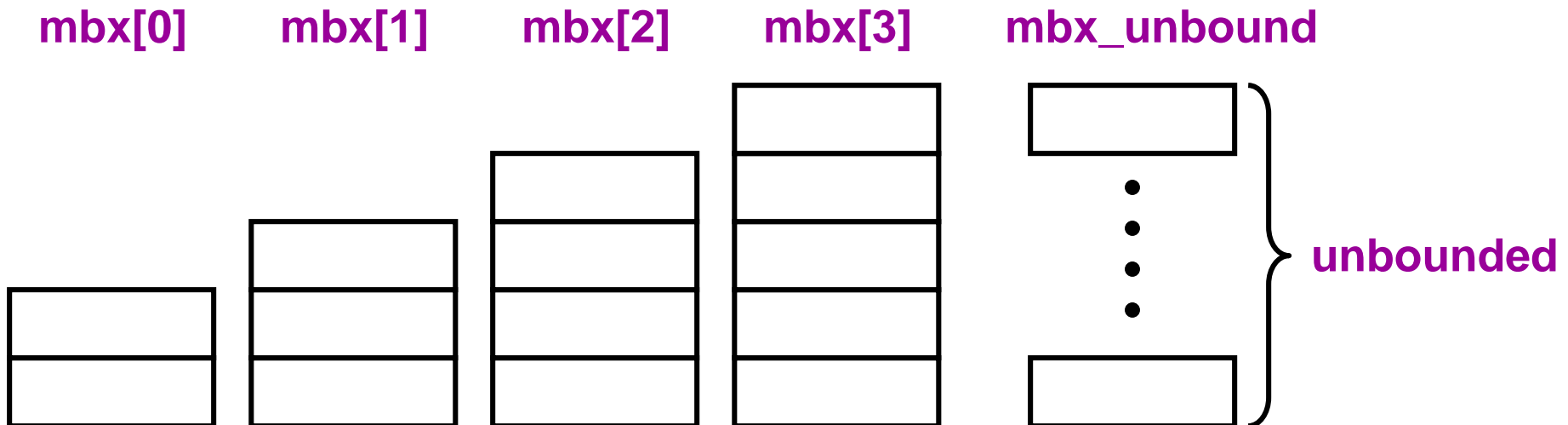
Mailbox Class

- Mailboxes are supported via built-in mailbox class:

```
class mailbox #(type T = dynamic_type) ;  
    function new(int bound = 0) ;  
    function int num() ;    // return # of messages  
    task put(T message) ;  // wait if mailbox full  
    task get(ref T message) ; // wait if no message  
    task peek(ref T message) ; // wait if no message  
    function int try_put(T message) ;  
    function int try_get(ref T message) ;  
    function int try_peek(ref T message) ;  
endclass
```


Creating Mailboxes

```
mailbox #(Packet) mbx[];           //Packets only mailbox array
mailbox #(instr_e) mbx_unbound;    //enum type instr_e
mbx_unbound = new();                // mailbox size is unbounded
mbx = new[4];                       // array of 4 null mailbox handles created
for (int i=0; i<mbx.size(); i++) begin
    mbx[i] = new(i+2);              // bound to max of i+2 messages
end
```



Putting Messages into Mailboxes

task **put** (message) ; // block if exceeds bound

function int **try_put** (message) ; // non-blocking

Example:

```
int status;  
typedef enum {ADD=1, SUB, MUL, DIV} instr_e;  
instr_e instr = ADD;  
mailbox #(instr_e) mbox = new(1);  
Packet pkt = new();  
mbox.put(instr);  
status = mbox.try_put(instr); //status = 0 - mailbox full  
mbox.put(pkt);
```

mbox

ADD

//incorrect type - compiler error

Retrieving Messages from Mailboxes (1/2)

`task get(ref message) ; // block if mailbox is empty`

`function int try_get(ref message) ; // non-blocking`
`// if successful, return number of messages in mailbox before try_get,`
`// 0 if no message`

Example:

```
int status;
typedef enum {ADD=1, SUB, MUL, DIV} instr_e;
instr_e instr1 = MUL, instr2; Packet pkt;
mailbox #(instr_e) mbox = new(1);
mbox.put(instr1); mbox.get(instr2);
$display("instr2 = %s", instr2.name()); // instr2 = MUL
status = mbox.try_get(instr2); // status = 0
status = mbox.try_put(instr1); // status = 1
mbox.get(instr2); // instr2 = MUL
mbox.get(instr2); // blocking – mailbox empty
mbox.get(pkt); // incorrect type - compiler error
```

mbox

MUL

Retrieving Messages from Mailboxes (2/2)

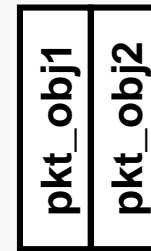
```
task peek (ref message) ; // block if mailbox is empty

function int try_peek (ref message) ; // non-blocking
// if successful, return number of message in mailbox before try_peek,
// 0 if no message
```

Example:

```
Packet pkt_obj1 = new() ;
Packet pkt_obj2 = new() , pkt2drv;
mailbox #(Packet) mbox = new() ;
mbox.put(pkt_obj1) ;
status = mbox.try_put(pkt_obj2) ;           // status = 1
$display("%0d messages" , mbox.num()) ; // 2 messages
mbox.peek(pkt2drv) ;                       // still 2 messages left
$display(mbox.try_peek(pkt2drv)) ;         // displays 2
$display(mbox.try_get(pkt2drv)) ;         // displays 2
```

mbox

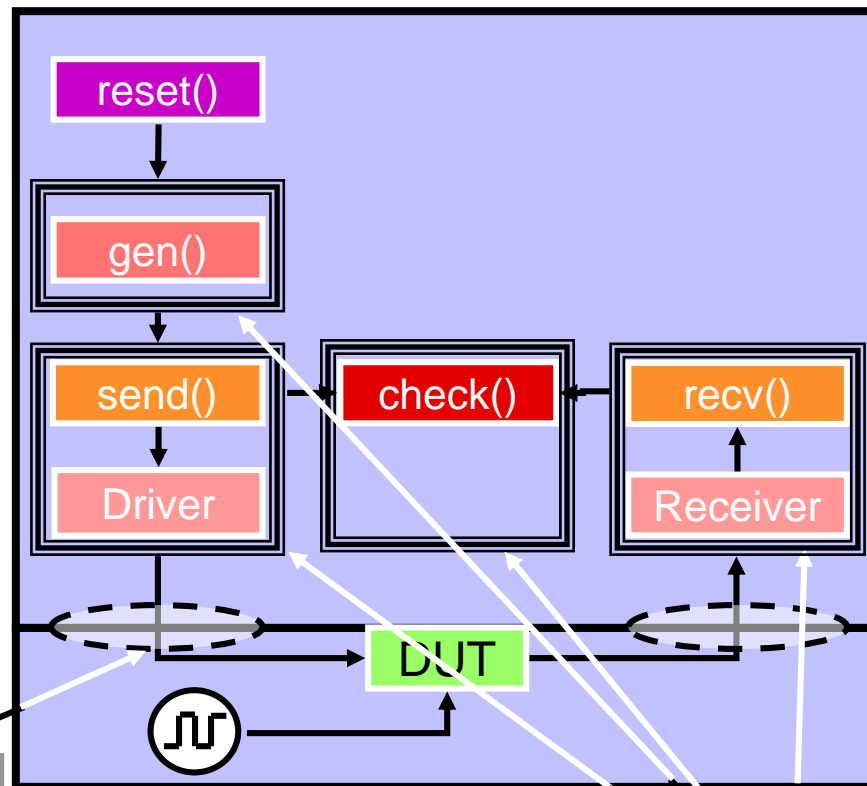


Lab 5 Introduction



120 min

Expand to broad-spectrum verification



Expand to exercise all ports concurrently

Compile & Simulate

semaphore-based resource mgmt

Encapsulate in OOP Class

Unit Objectives Review

Having completed this unit, you should be able to:

- **Establish order of execution using event flag**
- **Avoid resource collision with Semaphores**
- **Pass data via Mailbox**