

Problem Set 7

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Handed In: May 1, 2017

1. Answer to problem 1

a.

$$P(x^{(j)}) = \sum_{z \in \{1,2\}} \prod_{i=0}^n P(x_i^{(j)}|z) \quad (1)$$

$$= \prod_{i=0}^n P(x_i^{(j)}|z=1)P(z=1) + \prod_{i=0}^n P(x_i^{(j)}|z=2)P(z=2) \quad (2)$$

$$P(x^{(j)}) = \alpha \prod_{i=0}^n p_i^{x_i^{(j)}} (1-p_i)^{1-x_i^{(j)}} + (1-\alpha) \prod_{i=0}^n q_i^{x_i^{(j)}} (1-q_i)^{1-x_i^{(j)}} \quad (3)$$

b. By Bayes Theorem,

$$f_z^{(j)} = P(Z=z|x^{(j)}) = \frac{P(Z=z, x^{(j)})}{P(x^{(j)})} = \frac{P(x^{(j)}|Z=z)P(Z=z)}{P(x^{(j)})} \quad (4)$$

$$f_1^{(j)} = \frac{\alpha \prod_{i=0}^n p_i^{x_i^{(j)}} (1-p_i)^{1-x_i^{(j)}}}{\alpha \prod_{i=0}^n p_i^{x_i^{(j)}} (1-p_i)^{1-x_i^{(j)}} + (1-\alpha) \prod_{i=0}^n q_i^{x_i^{(j)}} (1-q_i)^{1-x_i^{(j)}}} \quad (5)$$

$$f_2^{(j)} = \frac{(1-\alpha) \prod_{i=0}^n q_i^{x_i^{(j)}} (1-q_i)^{1-x_i^{(j)}}}{\alpha \prod_{i=0}^n p_i^{x_i^{(j)}} (1-p_i)^{1-x_i^{(j)}} + (1-\alpha) \prod_{i=0}^n q_i^{x_i^{(j)}} (1-q_i)^{1-x_i^{(j)}}} \quad (6)$$

$$(7)$$

c.

$$L = \prod_{j=1}^m P(x^{(j)}|p, q, \alpha) \quad (8)$$

$$LL = \sum_{j=1}^m \log P(x^{(j)}|p, q, \alpha) \quad (9)$$

$$(10)$$

$$E[LL] = E\left[\sum_{j=1}^m \log P(x^{(j)}|p, q, \alpha)\right] = \sum_{j=1}^m E[\log P(x^{(j)}|p, q, \alpha)] \quad (11)$$

$$= \sum_{j=1}^m \sum_{z=1}^2 f_z^{(j)} \log P(Z = z, x^{(j)}|\tilde{p}, \tilde{q}, \tilde{\alpha}) - \sum_{j=1}^m \sum_{z=1}^2 f_z^{(j)} \log f_z^{(j)} \quad (12)$$

$$= \sum_{j=1}^m f_1^{(j)} \log(\alpha \prod_{i=1}^n \tilde{p}_i^{x_i^{(j)}} (1 - \tilde{p})^{1-x_i^{(j)}}) \quad (13)$$

$$+ f_2^{(j)} \log((1 - \alpha) \prod_{i=1}^n \tilde{q}_i^{x_i^{(j)}} (1 - \tilde{q})^{1-x_i^{(j)}}) \quad (14)$$

$$- \sum_{j=1}^m (f_1^{(j)} \log f_1^{(j)} + f_2^{(j)} \log f_2^{(j)}) \quad (15)$$

d. The expected log likelihood $E[LL]$ will be maximized when the derivative is equal to 0.

$$\frac{\partial E}{\partial \tilde{\alpha}} = \sum_{j=1}^m \frac{f_1^{(j)}}{\tilde{\alpha}} - \frac{f_2^{(j)}}{1 - \tilde{\alpha}} = 0 \quad (16)$$

$$\Rightarrow \tilde{\alpha} = \frac{\sum_{j=1}^m f_1^{(j)}}{m} \quad (17)$$

$$\frac{\partial E}{\partial \tilde{p}_i} = \sum_{j=1}^m \frac{f_1^{(j)} x_i}{\tilde{p}_i} - \frac{f_1^{(j)} (1 - x_i)}{1 - \tilde{p}_i} = 0 \quad (18)$$

$$\Rightarrow \tilde{p}_i = \frac{\sum_{j=1}^m f_1^{(j)} x_i^{(j)}}{\sum_{j=1}^m f_1^{(j)}} \quad (19)$$

$$\frac{\partial E}{\partial \tilde{q}_i} = \sum_{j=1}^m \frac{f_2^{(j)} x_i}{\tilde{q}_i} - \frac{f_2^{(j)} (1 - x_i)}{1 - \tilde{q}_i} = 0 \quad (20)$$

$$\Rightarrow \tilde{q}_i = \frac{\sum_{j=1}^m (1 - f_2^{(j)}) x_i^{(j)}}{\sum_{j=1}^m (1 - f_2^{(j)})} \quad (21)$$

- e. $\tilde{\alpha}$ is the estimated probability of generating a sample with $z = 1$.
- \tilde{p} is the estimated probability of getting $x_i = 1$ given that $z = 1$.
- \tilde{q} is the estimated probability of getting $x_i = 1$ given that $z = 2$.

Algorithm 1 Pseudocode:

- 1: **procedure** MYPROCEDURE
 - 2: Initialize parameters p, q, α with random values as estimations $\tilde{p}, \tilde{q}, \tilde{\alpha}$.
 - 3: Calculate the posterior distribution $f_1^{(j)}$ and $f_2^{(j)}$, with the equation from part (b).
 - 4: Use the update rules in part (d) to update estimations $\tilde{p}, \tilde{q}, \tilde{\alpha}$.
 - 5: Repeat step *ii.* and *iii.* until the estimations $\tilde{p}, \tilde{q}, \tilde{\alpha}$ converge.
 - 6: **end procedure**
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- f. The algorithm will predict x_0 as 1 if $\frac{P(X_0=1)}{P(X_0=0)} > 1$, otherwise, the algorithm will predict x_1 as 0. Therefore, $x_0 = \text{sign}(\log \frac{P(X_0=1)}{P(X_0=0)})$.

$$P(X_0 = 0) = P(x_0|x_1, \dots, x_n) \quad (22)$$

$$= P(Z = 1|x_1, \dots, x_n)P(X_0 = 0|Z = 1) \quad (23)$$

$$+ P(Z = 2|x_1, \dots, x_n)P(X_0 = 0|Z = 2) \quad (24)$$

$$= f_1 p_0 + f_2 q_0 \quad (25)$$

$$P(X_0 = 1) = P(x_0|x_1, \dots, x_n) \quad (26)$$

$$= P(Z = 1|x_1, \dots, x_n)P(X_0 = 1|Z = 1) \quad (27)$$

$$+ P(Z = 2|x_1, \dots, x_n)P(X_0 = 1|Z = 2) \quad (28)$$

$$= f_1(1 - p_0) + f_2(1 - q_0) \quad (29)$$

$$f_1 = \frac{\alpha \prod_{i=0}^n p_i^{x_i} (1 - p_i)^{1-x_i}}{\alpha \prod_{i=0}^n p_i^{x_i} (1 - p_i)^{1-x_i} + (1 - \alpha) \prod_{i=0}^n q_i^{x_i} (1 - q_i)^{1-x_i}} \quad (30)$$

$$f_2 = \frac{(1 - \alpha) \prod_{i=0}^n q_i^{x_i} (1 - q_i)^{1-x_i}}{\alpha \prod_{i=0}^n p_i^{x_i} (1 - p_i)^{1-x_i} + (1 - \alpha) \prod_{i=0}^n q_i^{x_i} (1 - q_i)^{1-x_i}} \quad (31)$$

$$x_0 = \text{sign}(\log \frac{P(X_0 = 1)}{P(X_0 = 0)}) \quad (32)$$

$$= \text{sign}(\log(\frac{f_1(1 - p_0) + f_2(1 - q_0)}{f_1 p_0 + f_2 q_0})) \quad (33)$$

$$= \text{sign}(\log(\frac{(1 - p_0)\alpha \prod_{i=0}^n p_i^{x_i} (1 - p_i)^{1-x_i} + (1 - q_0)(1 - \alpha) \prod_{i=0}^n q_i^{x_i} (1 - q_i)^{1-x_i}}{p_0 \alpha \prod_{i=0}^n p_i^{x_i} (1 - p_i)^{1-x_i} + q_0 (1 - \alpha) \prod_{i=0}^n q_i^{x_i} (1 - q_i)^{1-x_i}})) \quad (34)$$

$$(35)$$

- g. According to the result from part (f), the decision surface for this prediction can transform to a linear function. After transforming,

$$x_0 = \text{sign}(\log(\frac{(1-2p_0)\alpha \prod_{i=0}^n p_i^{x_i}(1-p_i)^{1-x_i}}{(1-2q_0)(\alpha-1) \prod_{i=0}^n q_i^{x_i}(1-q_i)^{1-x_i}})) \quad (36)$$

$$= \log((1-2p_0)\alpha \prod_{i=0}^n p_i^{x_i}(1-p_i)^{1-x_i}) - \log((1-2q_0)(\alpha-1) \prod_{i=0}^n q_i^{x_i}(1-q_i)^{1-x_i}) \quad (37)$$

$$= \log(\frac{1-2p_0}{1-2q_0}) + \log(\frac{\alpha}{\alpha-1}) + \sum_{i=0}^n x_i \log(\frac{p_i}{q_i}) + \sum_{i=0}^n (1-x_i) \log(\frac{1-p_i}{1-q_i}) \quad (38)$$

2. Answer to problem 2

- a. The statement means that the probabilities for every event E over variables x_1, \dots, x_n are equal for two directed trees T_0 and T_1 . The joint probability distributions are the same. $P_{T_0}(x_1, \dots, x_n) = P_{T_1}(x_1, \dots, x_n)$.
- b. Assume two directed trees T_i and T_j have different roots x_i and x_j from the undirected tree T . The resulting directed trees are all equivalent if $P_{T_0}(x) = P_{T_1}(x)$. So, $P(x_1|x_2)P(x_2) = P(x_2|x_1)P(x_1) = P(x_1, x_2)$. First, we assume x_i and x_j are nodes in tree T , and they are connected by a path P with length 1.

$$P_{T_i}(x) = P(x_i) \prod_{k \in \{N-i\}}^n P(x_k | \text{Parent}_{x_k}) \quad (39)$$

$$= P(x_i)P(x_j|x_i) \prod_{k \in \{N-P\}} P(x_k | \text{Parent}_{x_k}) \quad (40)$$

$$= P(x_i, x_j) \prod_{k \in \{N-P\}} P(x_k | \text{Parent}_{x_k}) \quad (41)$$

$$= P(x_j)P(x_i|x_j) \prod_{k \in \{N-P\}} P(x_k | \text{Parent}_{x_k}) \quad (42)$$

$$= P(x_j) \prod_{k \in \{N-j\}}^n P(x_k | \text{Parent}_{x_k}) \quad (43)$$

$$= P_{T_j}(x) \quad (44)$$

When the path P has length larger than 1, the hypothesis still holds. $P(x_{\text{root}}) \prod_{k \in \{N-P\}}^n P(x_k | \text{Parent}_{x_k})$ remains the same. We can switch the edges between x_i and x_j according to the chain rule. By switching one edge at a time, the resulting directed trees are all equivalent.