

CHAPTER 1

INTRODUCTION ON COMPUTER NETWORK

ET4230 – 20161

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Chapter 1: introduction

our goal:

- get “feel” and terminology
- more depth, detail *later* in course
- approach:
 - use Internet as example

overview:

- what's the Internet?
- what's a protocol?
- network edge; hosts, access net, physical media
- network core: packet/circuit switching, Internet structure
- performance: loss, delay, throughput
- security
- protocol layers, service models
- history

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Chapter 1: roadmap

1.1 what is the Internet?

1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

1.6 networks under attack: security

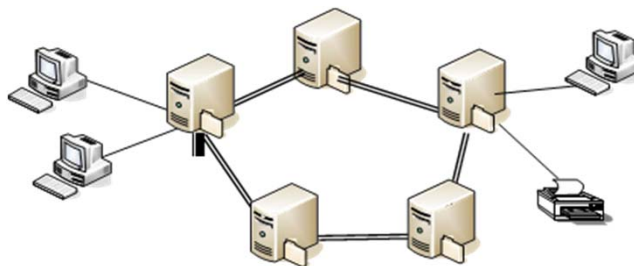
1.7 history

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ĐỊNH NGHĨA

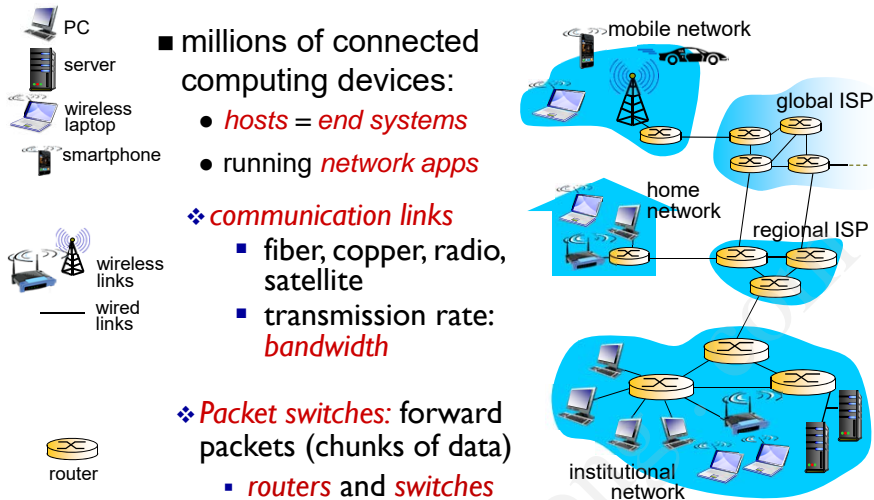
■ Mạng máy tính:

- thiết bị mạng (host, server, network devices)
- phương tiện truyền dẫn vật lý (transmission medium)
- kiến trúc mạng (network architecture)
 - [cấu trúc mạng \(Topology\)](#)
 - [giao thức mạng \(Protocols\)](#)



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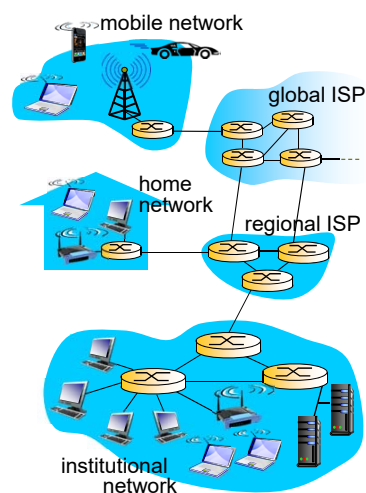
What's the Internet: a component view



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What's the Internet: a protocol view

- *Internet: "network of networks"*
 - Interconnected ISPs
- *protocols* control sending, receiving of msgs
 - e.g., TCP, IP, HTTP, Skype, 802.11
- *Internet standards*
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force



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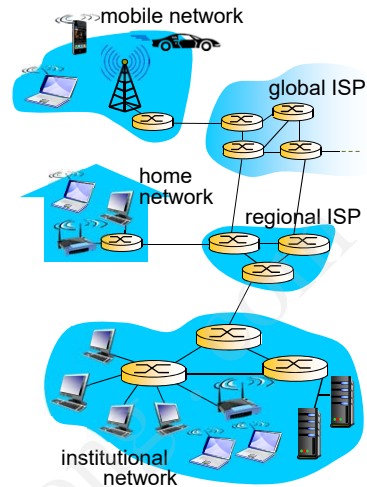
What's the Internet: a service view

- *Infrastructure that provides services to applications:*

- Web, VoIP, email, games, e-commerce, social nets, ...

- *provides programming interface to apps*

- hooks that allow sending and receiving app programs to "connect" to Internet
- provides service options, analogous to postal service



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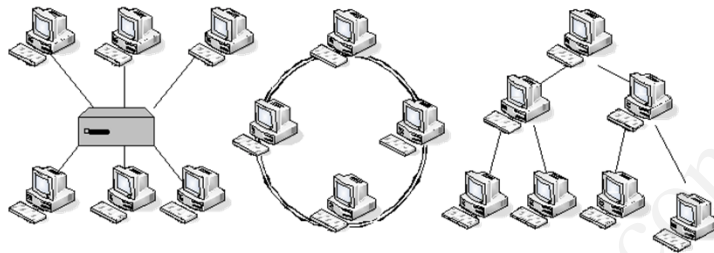
What's the Internet: goal

- Chia sẻ tài nguyên dùng chung
- Nâng cao độ tin cậy
- Môi trường giao tiếp người – máy
- Giảm chi phí đầu tư phần cứng
- Bảo đảm các tiêu chuẩn thống nhất về tính bảo mật, an toàn dữ liệu

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Topology

■ Kiểu điểm - điểm (Point to Point)

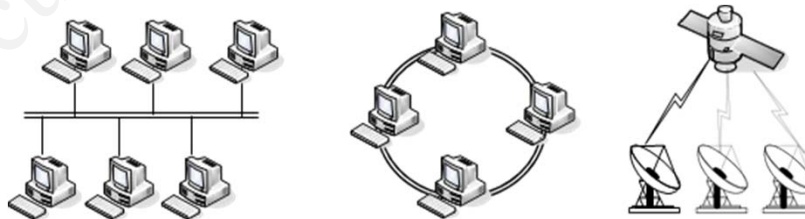


Các mạng có cấu trúc điểm - điểm (Star, Ring, Tree)

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Topology

■ Kiểu đa điểm hay quảng bá (Point to Multipoint, Broadcasting)



Các mạng có cấu trúc quảng bá (Bus, Ring, and Satellite)

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What's a protocol?

human protocols:

- “what’s the time?”
- “I have a question”
- introductions

... specific msgs sent

... specific actions taken
when msgs received,
or other events

network protocols:

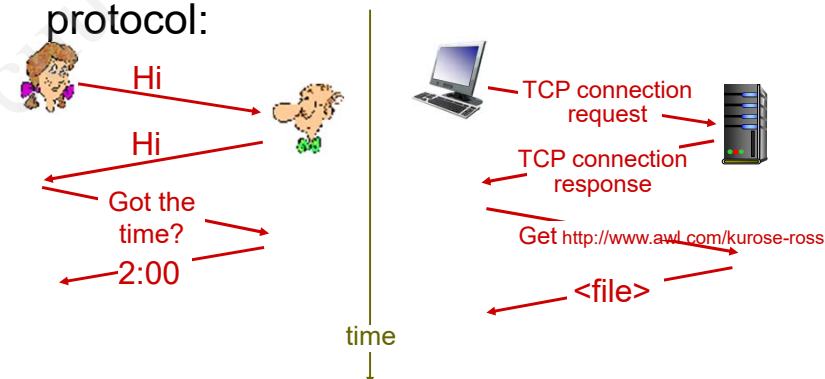
- machines rather than humans
- all communication activity in Internet governed by protocols

protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt

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What's a protocol?

a human protocol and a computer network protocol:



Q: other human protocols?

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What's a protocol?

■ Chức năng giao thức

- Encapsulation
- Fragmentation
- Connection control
- Monitoring
- Flow control
- Error control
- Synchronization
- Addressing

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Chapter 1: roadmap

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1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

1.6 networks under attack: security

1.7 history

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A closer look at network structure

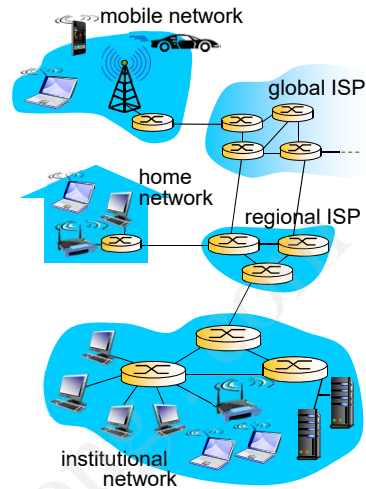
■ *network edge:*

- hosts: clients and servers
- servers often in data centers

❖ *access networks, physical media:* wired, wireless communication links

❖ *network core:*

- interconnected routers
- network of networks



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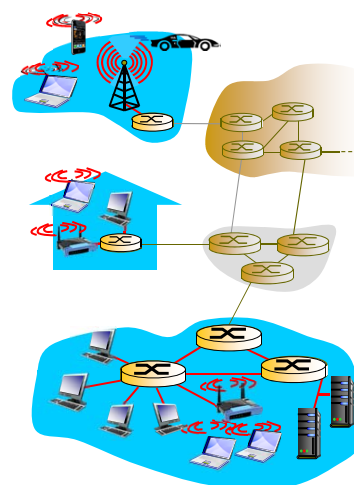
Access networks and physical media

Q: How to connect end systems to edge router?

- residential access nets
- institutional access networks (school, company)
- mobile access networks

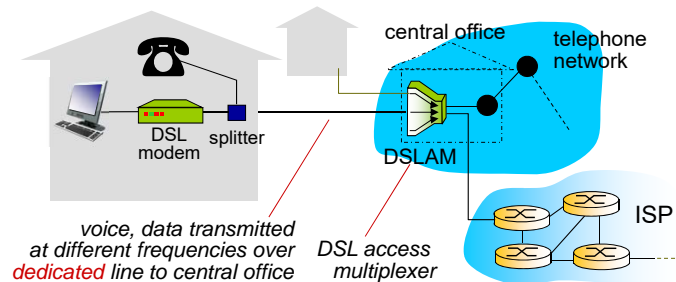
keep in mind:

- bandwidth (bits per second) of access network?
- shared or dedicated?



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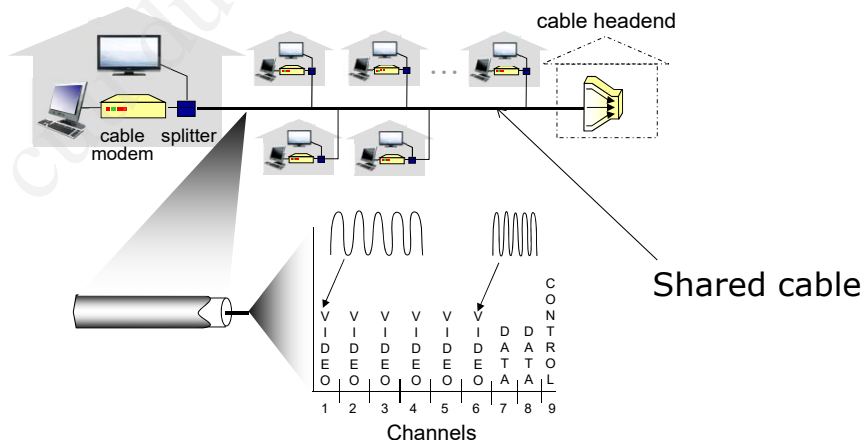
Access net: digital subscriber line (DSL)



- ❖ use **existing** telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- ❖ < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)
- ❖ < 24 Mbps downstream transmission rate (typically < 10 Mbps)

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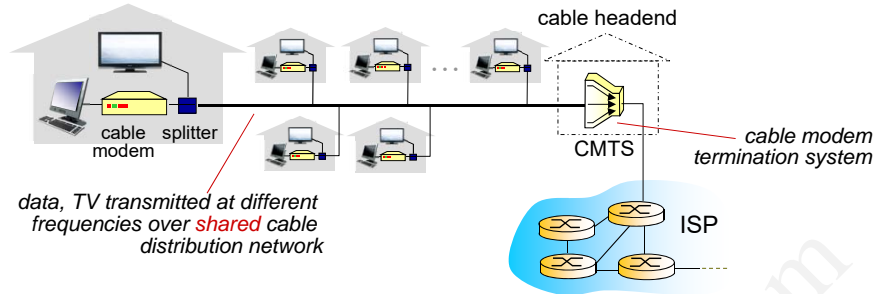
Access net: cable network



frequency division multiplexing: different channels transmitted in different frequency bands

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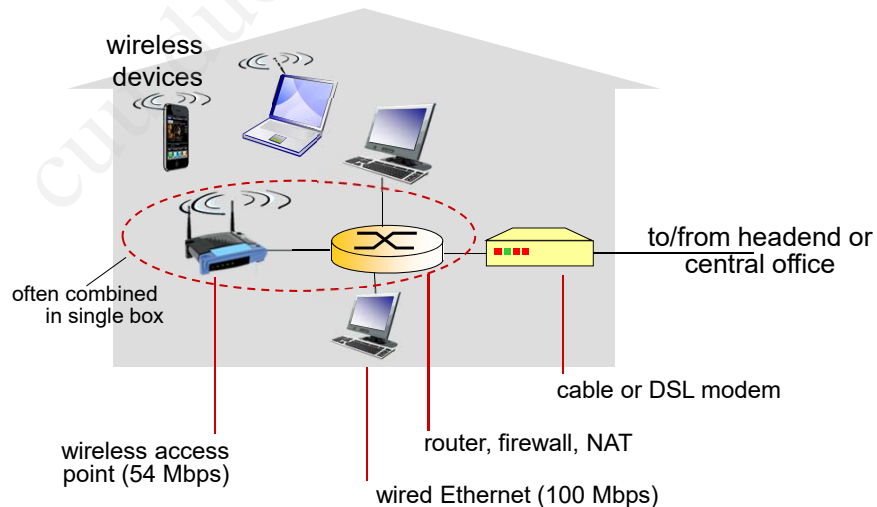
Access net: cable network



- ❖ HFC: hybrid fiber coax
 - asymmetric: up to 30Mbps downstream transmission rate, 2 Mbps upstream transmission rate
- ❖ network of cable, fiber attaches homes to ISP router
 - homes *share access network* to cable headend
 - unlike DSL, which has dedicated access to central office

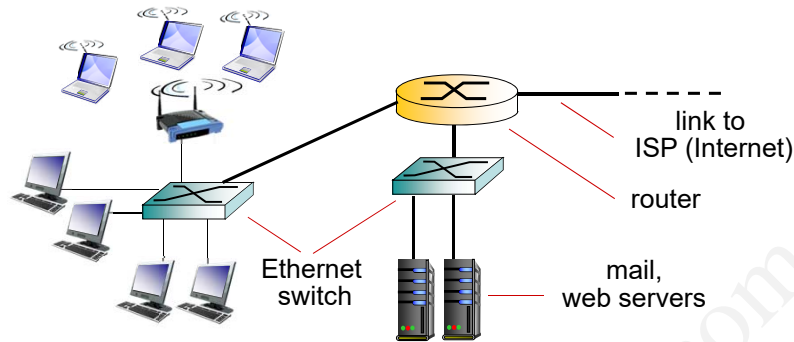
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Access net: home network



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Enterprise access networks (Ethernet)



- typically used in companies, universities, etc
- 10 Mbps, 100Mbps, 1Gbps, 10Gbps transmission rates
- today, end systems typically connect into Ethernet switch

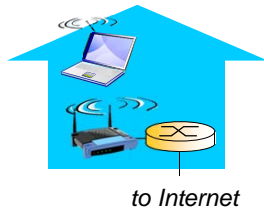
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Wireless access networks

- shared *wireless* access network connects end system to router
 - via base station aka "access point"

wireless LANs:

- within building (100 ft)
- 802.11b/g (WiFi): 11, 54 Mbps transmission rate



to Internet

wide-area wireless access

- provided by telco (cellular) operator; 10's km
- between 1 and 10 Mbps
- 3G, 4G: LTE



to Internet

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Physical media

- **bit**: propagates between transmitter/receiver pairs
- **physical link**: what lies between transmitter & receiver
- **guided media**:
 - signals propagate in solid media: copper, fiber, coax
- **unguided media**:
 - signals propagate freely, e.g., radio

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Physical media: coax, fiber

twisted pair (TP)

- two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps



coaxial cable

- two concentric copper conductors
- bidirectional
- broadband:
 - multiple channels on cable
 - HFC



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Physical media: coax, fiber

fiber optic cable:

- ❖ glass fiber carrying light pulses, each pulse a bit
- ❖ high-speed operation:
 - high-speed point-to-point transmission (e.g., 10' s-100' s Gpbs transmission rate)
- ❖ low error rate:
 - repeaters spaced far apart
 - immune to electromagnetic noise



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Physical media: radio, viba, Infrared

- signal carried in electromagnetic spectrum
- no physical “wire”
- bidirectional
- propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

radio link types:

- ❖ **terrestrial microwave**
 - e.g. up to 45 Mbps channels
- ❖ **LAN** (e.g., WiFi)
 - 11 Mbps, 54 Mbps
- ❖ **wide-area** (e.g., cellular)
 - 3G cellular: ~ few Mbps
- ❖ **satellite**
 - Kbps to 45Mbps channel (or multiple smaller channels)
 - 270 msec end-end delay
 - geosynchronous versus low altitude

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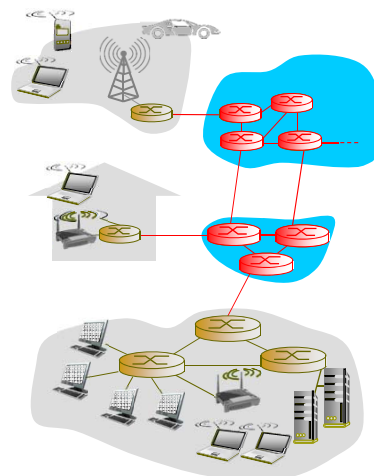
1.6 networks under attack: security

1.7 history

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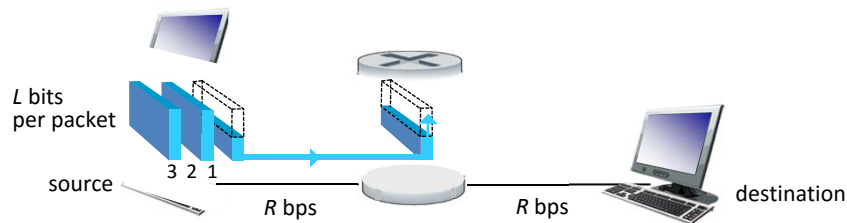
The network core

- mesh of interconnected routers
- packet-switching: hosts break application-layer messages into *packets*
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



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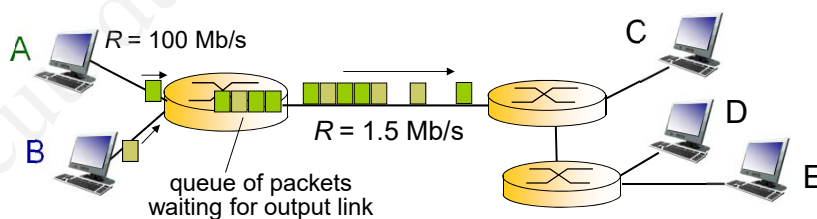
Packet-switching: store-and-forward



- takes L/R seconds to transmit (push out) L -bit packet into link at R bps
 - **store and forward**: entire packet must arrive at router before it can be transmitted on next link
 - ❖ end-end delay = $2L/R$ (assuming zero propagation delay)
- one-hop numerical example:*
- $L = 7.5$ Mbits
 - $R = 1.5$ Mbps
 - one-hop transmission delay = 5 sec
- more on delay shortly ...

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Packet Switching: queueing delay, loss



queueing and loss:

- ❖ If arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
 - packets will queue, wait to be transmitted on link
 - packets can be dropped (lost) if memory (buffer) fills up

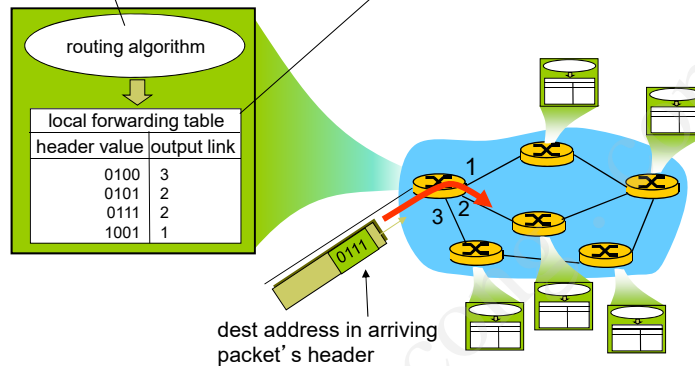
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Two key network-core functions

routing: determines source-destination route taken by packets

- *routing algorithms*

forwarding: move packets from router's input to appropriate router output

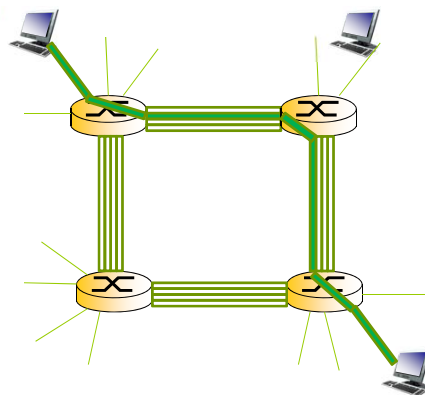


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Alternative core: circuit switching

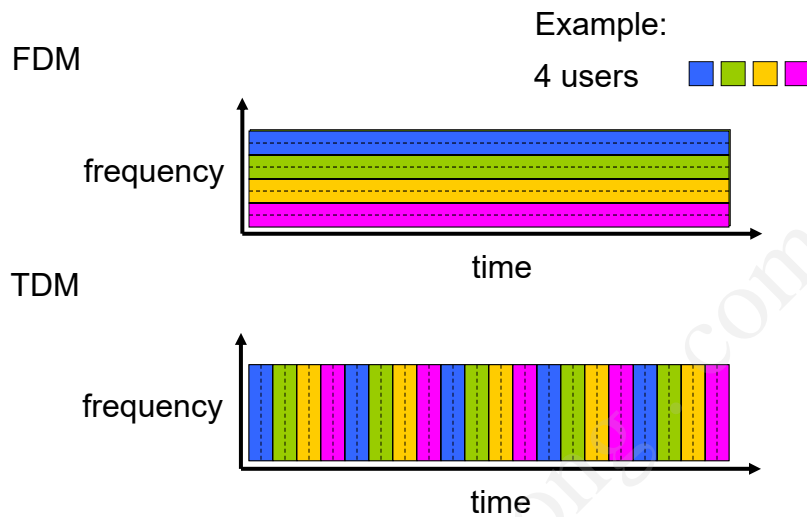
end-end resources allocated to, reserved for "call" between source & dest:

- In diagram, each link has four circuits.
 - call gets 2nd circuit in top link and 1st circuit in right link.
- dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- circuit segment idle if not used by call (*no sharing*)
- Commonly used in traditional telephone networks



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Circuit switching: FDM versus TDM



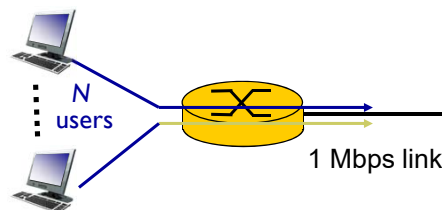
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Packet switching versus circuit switching

packet switching allows more users to use network!

example:

- 1 Mb/s link
- each user:
 - 100 kb/s when “active”
 - active 10% of time



- **circuit-switching:**
 - 10 users
- **packet switching:**
 - with 35 users, probability > 10 active at same time is less than .0004 *

Q: how did we get value 0.0004?

Q: what happens if > 35 users ?

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Packet switching versus circuit switching

is packet switching a “slam dunk winner?”

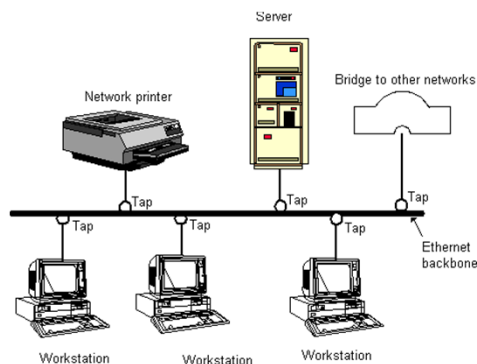
- great for bursty data
 - resource sharing
 - simpler, no call setup
- excessive congestion possible: packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem

Q: human analogies of reserved resources (circuit switching)
versus on-demand allocation (packet-switching)?

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Phân loại mạng máy tính

- Theo khoảng cách
 - Mạng cục bộ LAN (Local Area Networks)

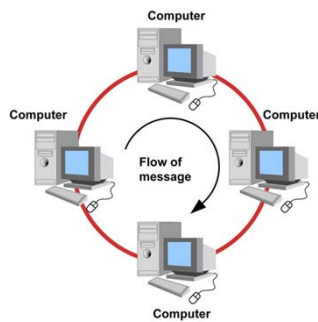


Cấu trúc mạng hình BUS

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Phân loại mạng máy tính

- Theo khoảng cách
 - Mạng cục bộ LAN (Local Area Networks)

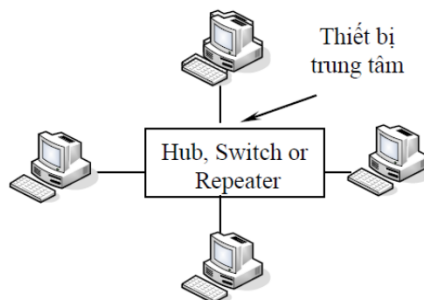


Cấu trúc mạng hình RING

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Phân loại mạng máy tính

- Theo khoảng cách
 - Mạng cục bộ LAN (Local Area Networks)

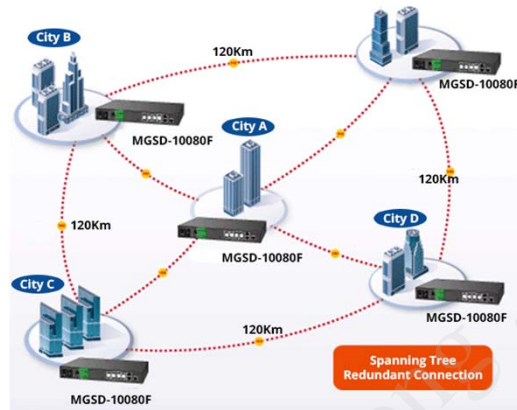


Cấu trúc mạng hình sao

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Phân loại mạng máy tính

- Theo khoảng cách
 - Mạng đô thị MAN (Metropolitan Area Networks)

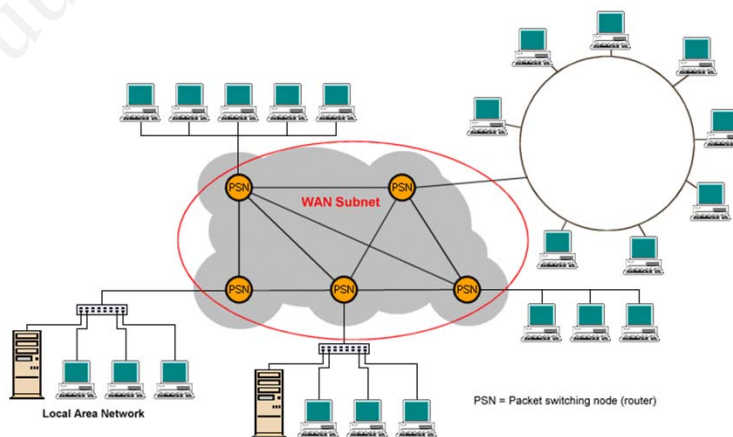


MAN Network Diagram

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Phân loại mạng máy tính

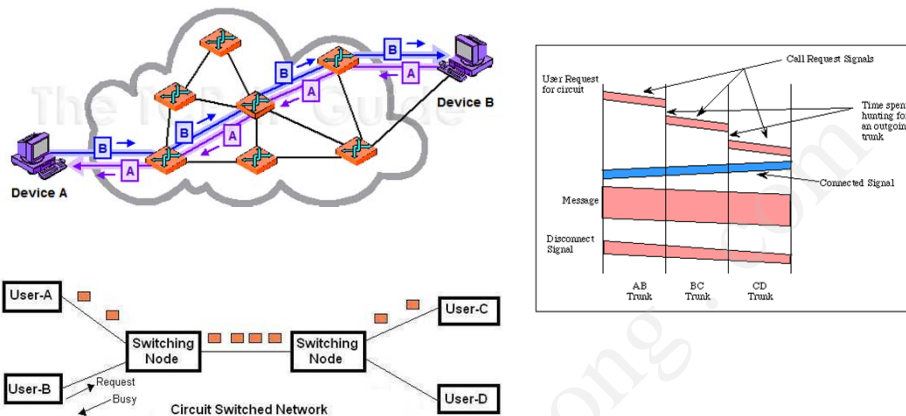
- Theo khoảng cách
 - Mạng diện rộng WAN (Wide Area Networks)



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Phân loại mạng máy tính

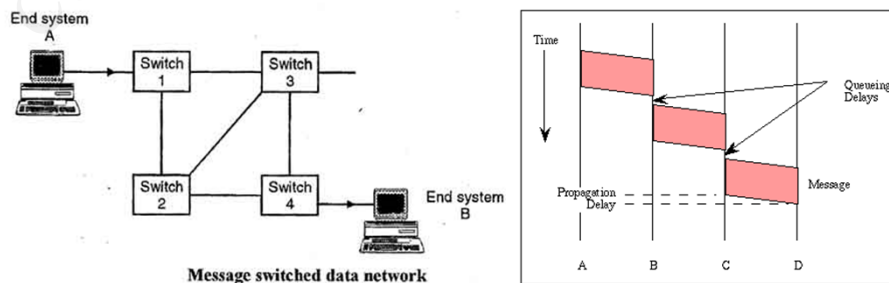
- Phân loại theo cơ chế chuyển mạch
 - Mạng chuyển mạch kênh (Circuit Switched Networks)



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Phân loại mạng máy tính

- Phân loại theo cơ chế chuyển mạch
 - Mạng chuyển mạch thông báo (message-switched network)

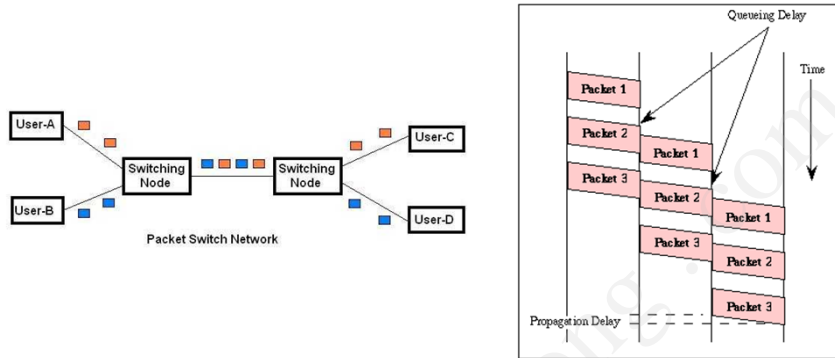


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Phân loại mạng máy tính

■ Phân loại theo cơ chế chuyển mạch

- Mạng chuyển mạch gói (Packet Switched Networks)

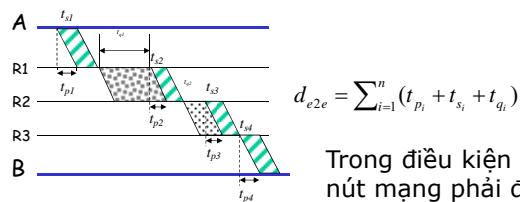
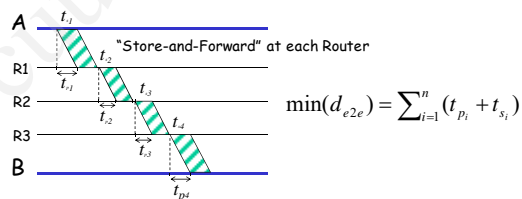


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Phân loại mạng máy tính

■ Phân loại theo cơ chế chuyển mạch

- Mạng chuyển mạch gói (Packet Switched Networks)

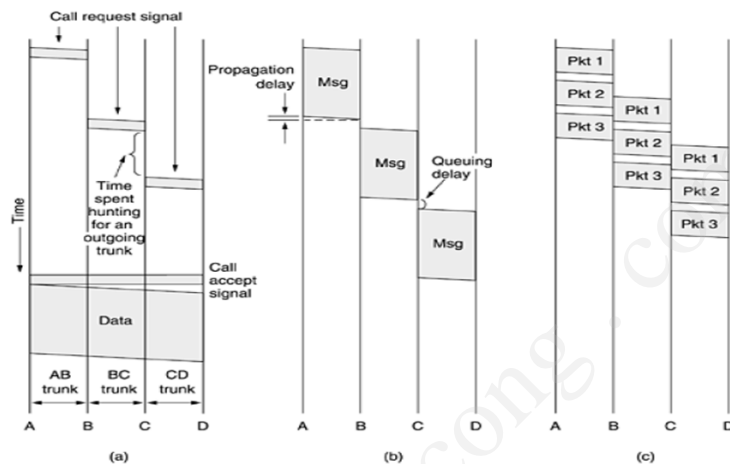


Trong điều kiện tải cao, các gói đi vào nút mạng phải đợi trong hàng đợi trước khi được gửi ra đầu ra

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Phân loại mạng máy tính

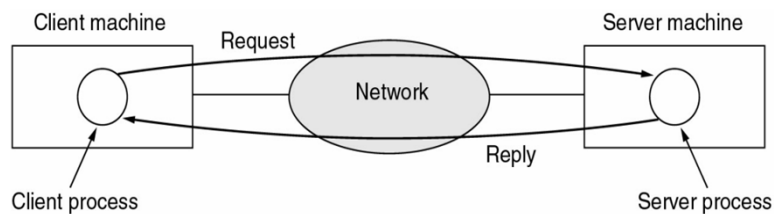
- Phân loại theo cơ chế chuyển mạch
 - So sánh các cơ chế chuyển mạch (bài tập)



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Phân loại mạng máy tính

- Phân loại theo mô hình xử lý dữ liệu
 - Mô hình Client-Server

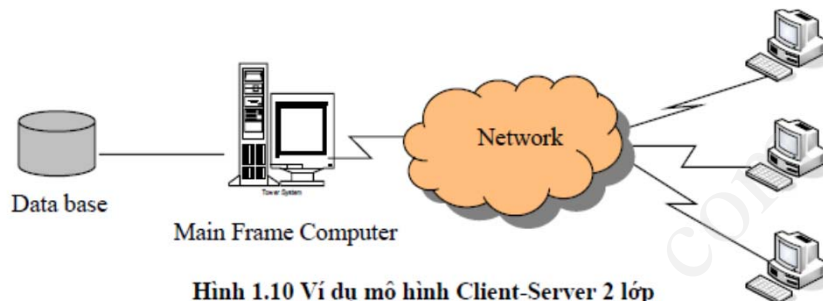


Mô hình chủ /khách (Client / Server)

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Phân loại mạng máy tính

- Phân loại theo mô hình xử lý dữ liệu
 - Mô hình Client-Server

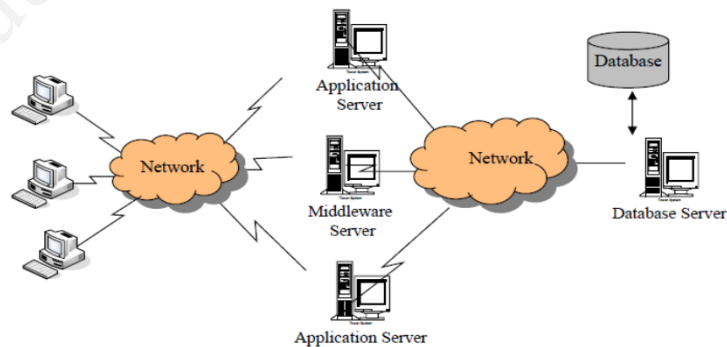


Hình 1.10 Ví dụ mô hình Client-Server 2 lớp
1. Trình duyệt Browser gửi yêu cầu cho Web Server.
2. Web Server trả kết quả về cho trình duyệt

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Phân loại mạng máy tính

- Phân loại theo mô hình xử lý dữ liệu
 - Mô hình Client-Server

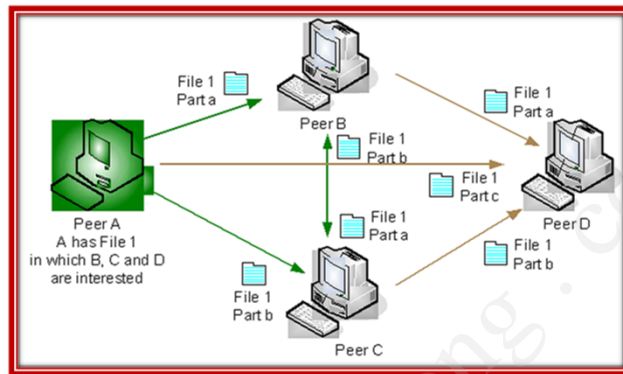


Mô hình Client-Server nhiều lớp

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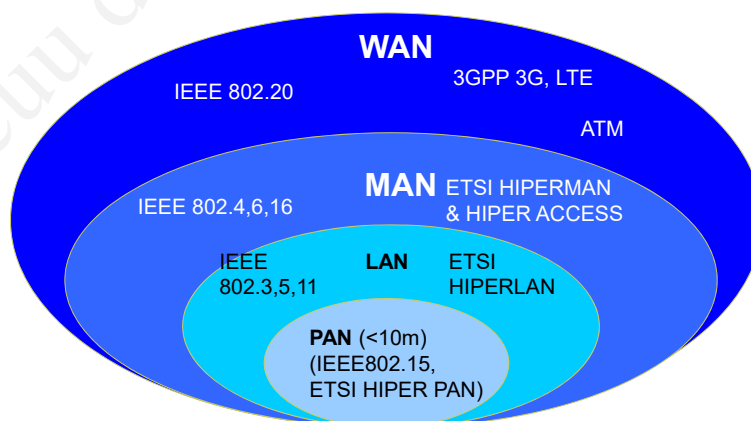
Phân loại mạng máy tính

- Phân loại theo mô hình xử lý dữ liệu
 - Mô hình ngang hàng (Peer-to-Peer)



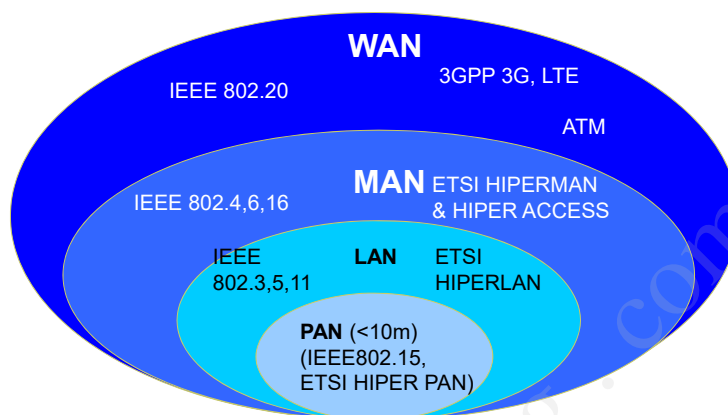
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Kích cỡ mạng



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Một số mạng điển hình

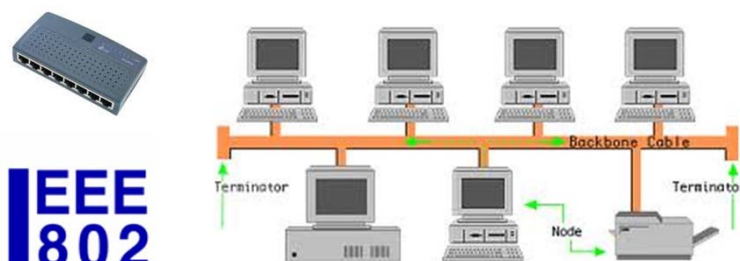


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Một số mạng điển hình

■ IEEE 802

- IEEE 802.3: Chuẩn mạng LAN/MAN – Ethernet
- IEEE 802.4: Chuẩn mạng LAN – Token Bus, chủ yếu được sử dụng trong công nghiệp (GM)

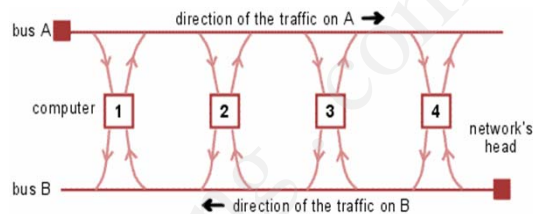
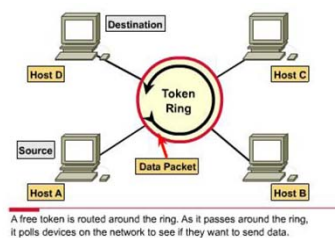


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Một số mạng điện hình

■ IEEE 802

- IEEE 802.5: chuẩn mạng LAN – Token Ring được phát triển bởi IBM
- IEEE 802.6: chuẩn mạng MAN – DQDB (Distributed Queue Dual Bus) với tốc độ 150Mbit/s trên khoảng cách 160km



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Một số mạng điện hình

■ IEEE 802

- IEEE 802.11: chuẩn mạng LAN không dây
- IEEE 802.15: chuẩn mạng cá nhân không dây (Wireless Personal Area Network - WPAN)
 - ▶ IEEE 802.15.1: Bluetooth
 - ▶ IEEE 802.15.3: High rate WPAN (11 – 55Mbit/s): sử dụng cho các ứng dụng multimedia
 - ▶ IEEE 802.15.4: Low rate WPAN/ZigBee: cho các ứng dụng tiêu thụ ít năng lượng, tốc độ thấp (Wireless Sensor Network)



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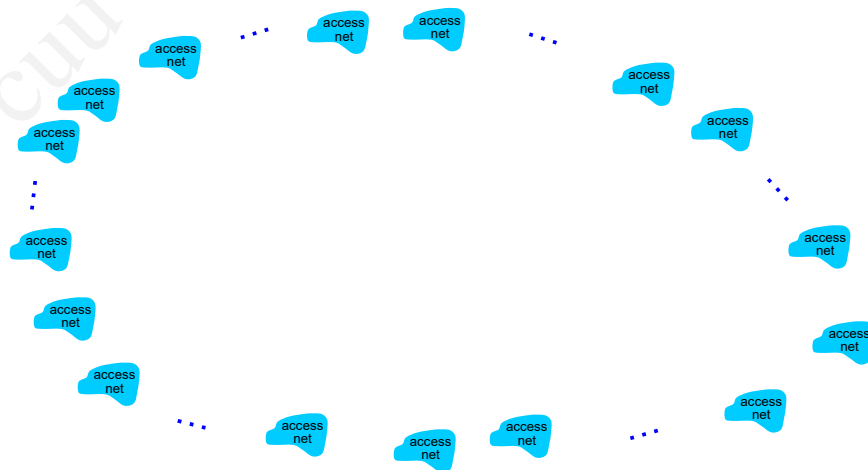
Internet structure: network of networks

- ❖ End systems connect to Internet via **access ISPs** (Internet Service Providers)
 - Residential, company and university ISPs
- ❖ Access ISPs in turn must be interconnected.
 - ❖ So that any two hosts can send packets to each other
- ❖ Resulting network of networks is very complex
 - ❖ Evolution was driven by **economics** and **national policies**
- ❖ Let's take a stepwise approach to describe current Internet structure

57

Internet structure: network of networks

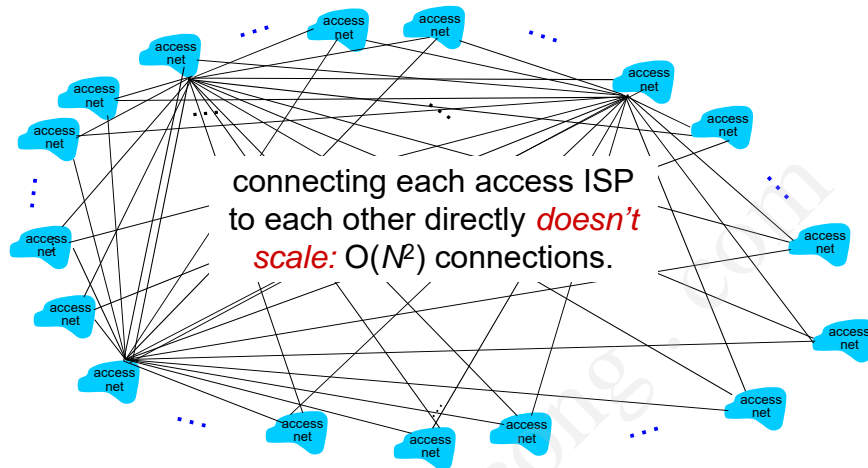
Question: given *millions* of access ISPs, how to connect them together?



58

Internet structure: network of networks

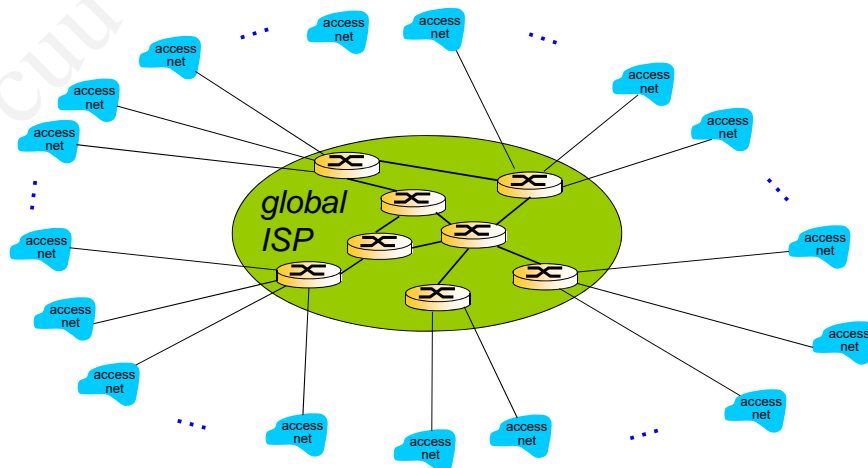
Option: connect each access ISP to every other access ISP?



59

Internet structure: network of networks

Option: connect each access ISP to a global transit ISP? *Customer* and *provider* ISPs have economic agreement.

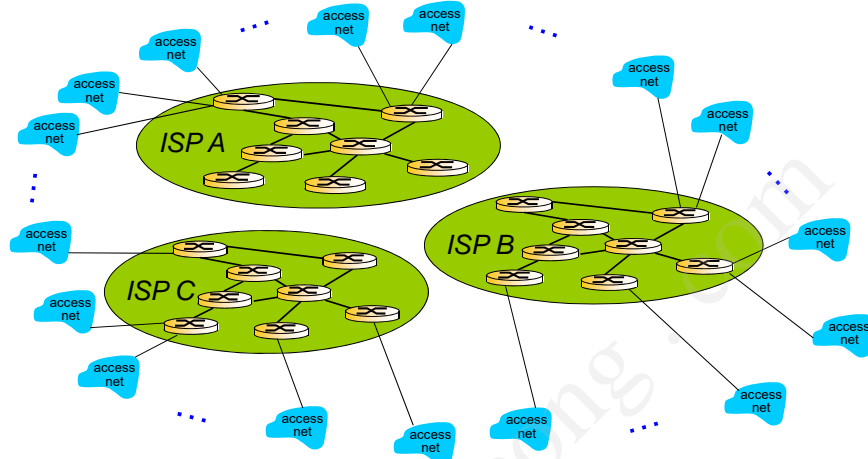


60

Internet structure: network of networks

But if one global ISP is viable business, there will be competitors

....

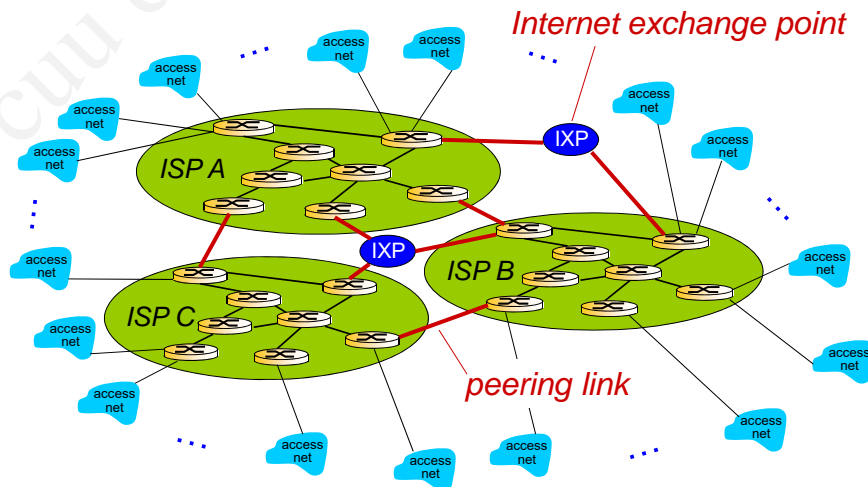


61

Internet structure: network of networks

But if one global ISP is viable business, there will be competitors

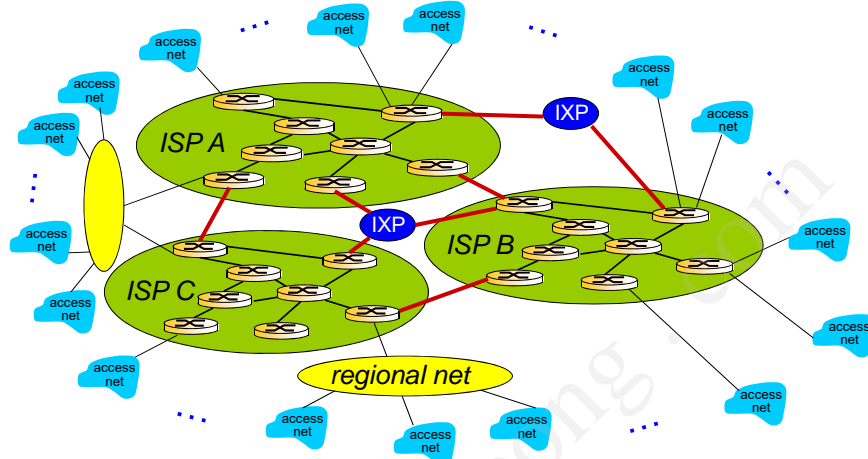
.... which must be interconnected



62

Internet structure: network of networks

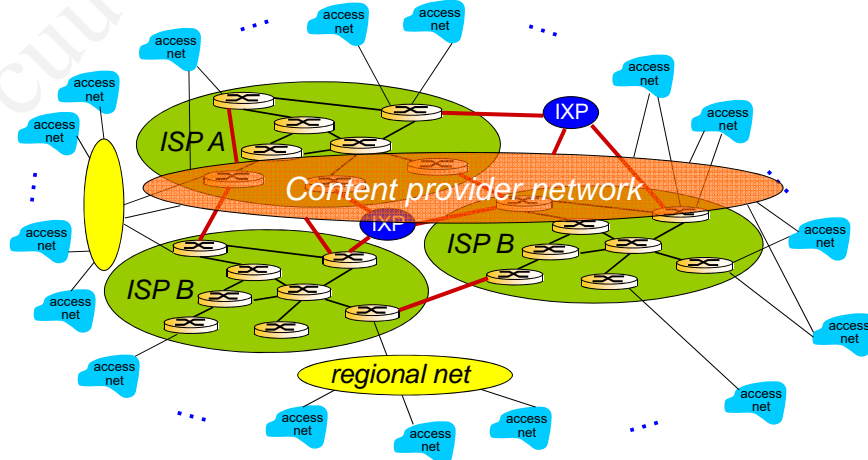
... and regional networks may arise to connect access nets to ISPs



63

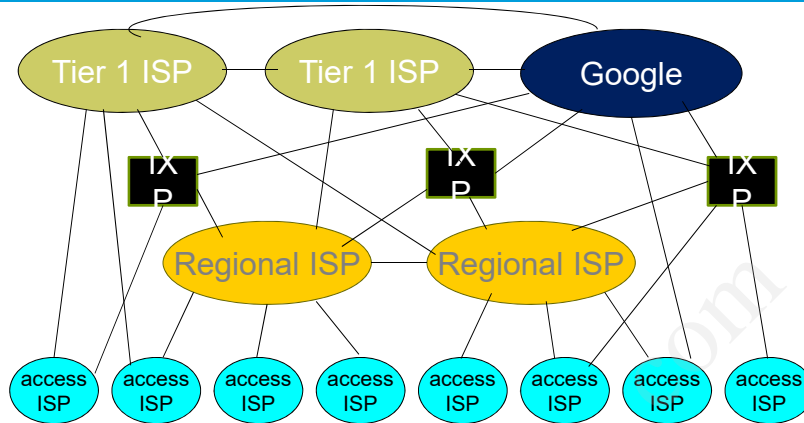
Internet structure: network of networks

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



64

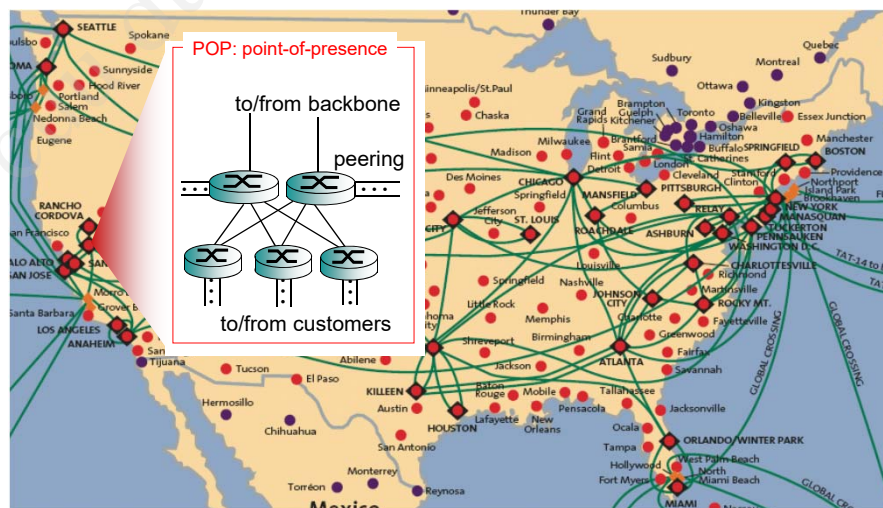
Internet structure: network of networks



- at center: small # of well-connected large networks
 - “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - content provider network (e.g., Google): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

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Tier-1 ISP: e.g., Sprint



66

Một số khái niệm

- **Băng thông:** Số lượng bit có thể truyền đi trên một đơn vị thời gian
 - Thường được biểu diễn qua tốc độ dữ liệu tối đa (maximum data rate)
- **Trễ lan truyền (propagation delay):** thời gian để tín hiệu lan truyền trên kênh vật lý:
 - $t_p = l/v_c$; v_c là vận tốc lan truyền của tín hiệu
 - Phụ thuộc vào môi trường truyền dẫn
 - Electromagnetic signal (light) travels in the medium -- 2×10^8 m/s in fiber.
 - l : Chiều dài của đường truyền vật lý giữa 2 nút mạng (m)

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Một số khái niệm

- **Thời gian phục vụ gói (transmission time):** thời gian gửi hết một gói tin từ bit đầu tiên đến bit cuối cùng lên kênh truyền
 - $t_s = L/C$
 - Dung lượng kênh truyền C (bit/s) (link capacity)
 - A function of bandwidth
 - If bandwidth is B , transmission time is $1/B$.
 - If bandwidth is 10 Mbps, the transmission time is $1/(10 \times 10^6)$ = 1 ms.
- **Trễ hàng đợi t_q (Queuing delay):** thời gian một gói phải lưu lại trong hàng đợi ở nút mạng trung gian
 - How long does it have to wait ?
 - Dependent on the load on the network -- how many packets are traversing that router

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Một số khái niệm

- Trễ từ đầu cuối đến đầu cuối (end-to-end delay d_{e2e} : trễ từ khi gửi một gói tin ở đầu phát cho đến khi nó được nhận ở đầu thu)
- Round Trip Time
 - **RTT** is the length of time it takes for a small packet to be sent plus the length of time it takes for an acknowledgment of that packet to be received (*excluding the transmission time of the packet*)
 - if forward delay = backward delay, $RTT = 2 * \text{Latency}$ (typically assumed -- although not always accurate)

69

Một số khái niệm

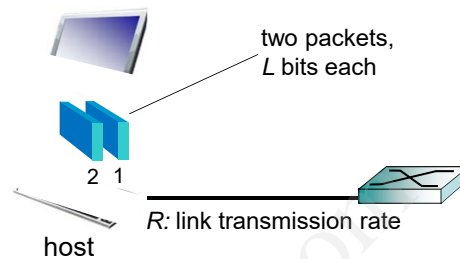
- Thông lượng (Throughput):
 - Định nghĩa hiệu quả sử dụng kênh truyền
 - $\text{Throughput} = \text{Transfer size} / \text{Transfer time}$.
 - What is the transfer time ?
 - $RTT + (\text{Transfer size} / \text{Bandwidth})$
 - (ignoring queuing delays).
 - Ví dụ:
 - 1 MB file over a 1 Gbps network with RTT 100 milliseconds.
 - $\text{Transfer time} = 100 \text{ ms} + (1 \text{ MB} / 1 \text{ Gbps}) = 100 \text{ ms} + 8 \text{ ms} = 108 \text{ ms}$.
 - $\text{Effective throughput} = 1 \text{ MB} / 108 \text{ ms} = 74.1 \text{ Mbps}$.
 - Nhận xét: Impact of data size
 - If data size increases, $(\text{Transfer size} / \text{Bandwidth})$ increases become much larger than RTT, in that case, $\text{Throughput} \sim (\text{Transfer size} / (\text{Transfer size} / \text{Bandwidth})) \sim \text{Bandwidth}$!

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Host: sends *packets* of data

host sending function:

- takes application message
- breaks into smaller chunks, known as *packets*, of length L bits
- transmits packet into access network at *transmission rate* R
 - link transmission rate, aka link *capacity*, aka *link bandwidth*



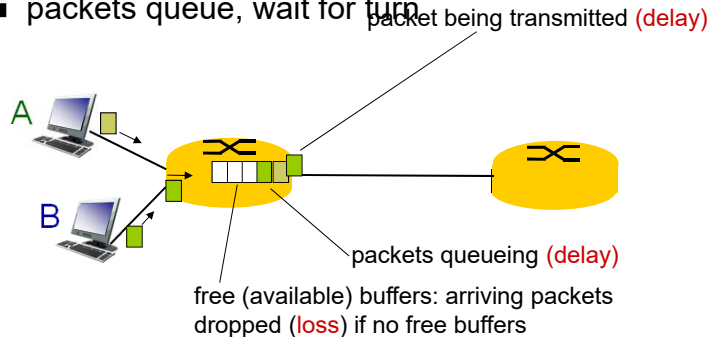
$$\text{packet transmission delay} = \text{time needed to transmit } L\text{-bit packet into link} = \frac{L \text{ (bits)}}{R \text{ (bits/sec)}}$$

71

How do loss and delay occur?

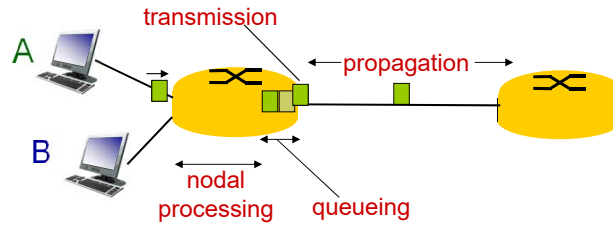
packets *queue* in router buffers

- packet arrival rate to link (temporarily) exceeds output link capacity
- packets queue, wait for turn



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Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{proc} : nodal processing

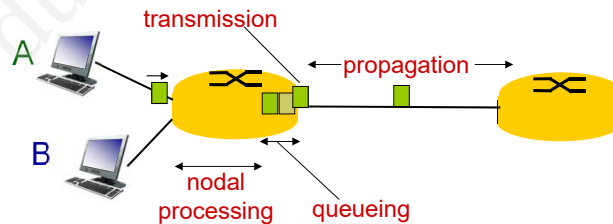
- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

73

Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L : packet length (bits)
- R : link bandwidth (bps)
- $d_{\text{trans}} = L/R$

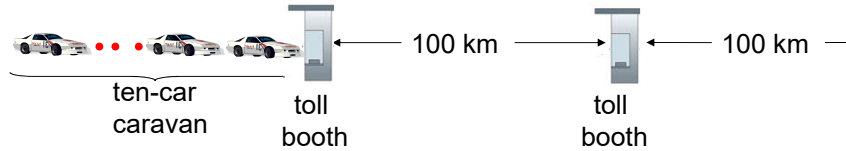
d_{prop} : propagation delay:

- d : length of physical link
- s : propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- $d_{\text{prop}} = d/s$

d_{trans} and d_{prop}
very different

74

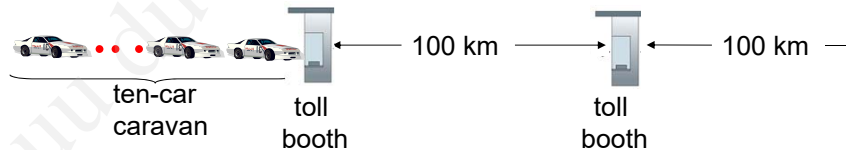
Caravan analogy



- cars “propagate” at 100 km/hr
- toll booth takes 12 sec to service car (bit transmission time)
- car~bit; caravan ~ packet
- **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway = $12 \times 10 = 120$ sec
- time for last car to propagate from 1st to 2nd toll booth: $100\text{km}/(100\text{km/hr}) = 1$ hr
- **A: 62 minutes**

75

Caravan analogy (more)

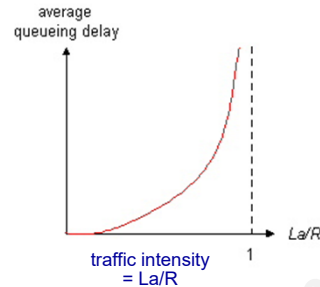


- suppose cars now “propagate” at 1000 km/hr
- and suppose toll booth now takes one min to service a car
- **Q: Will cars arrive to 2nd booth before all cars serviced at first booth?**
- **A: Yes!** after 7 min, 1st car arrives at second booth; three cars still at 1st booth.

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Queueing delay (revisited)

- R : link bandwidth (bps)
- L : packet length (bits)
- a : average packet arrival rate



- ❖ $La/R \sim 0$: avg. queueing delay small
- ❖ $La/R \rightarrow 1$: avg. queueing delay large
- ❖ $La/R > 1$: more “work” arriving than can be serviced, average delay infinite!



$La/R \sim 0$

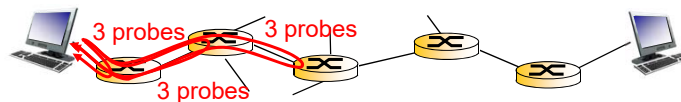


$La/R \rightarrow 1$

77

“Real” Internet delays and routes

- what do “real” Internet delay & loss look like?
- `traceroute` program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



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“Real” Internet delays, routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from
gaia.cs.umass.edu to cs-gw.cs.umass.edu

```

1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 ***
18 ***
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
  
```

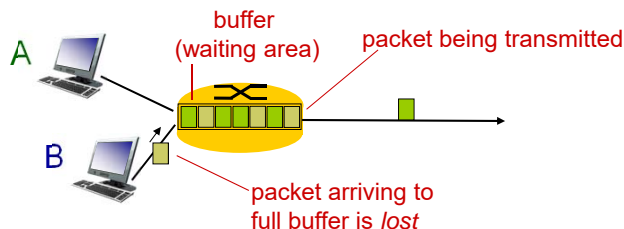
trans-oceanic link

* means no response (probe lost, router not replying)

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Packet loss

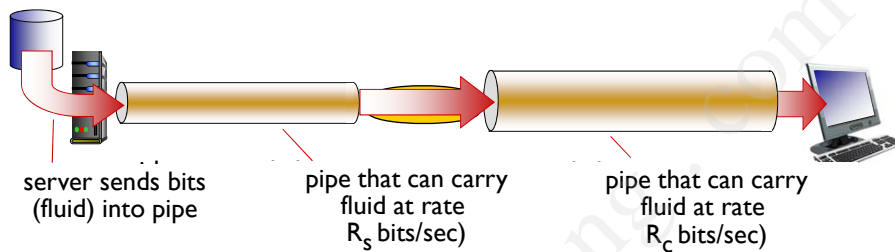
- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all



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Throughput

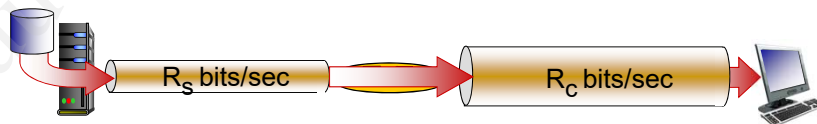
- **throughput**: rate (bits/time unit) at which bits transferred between sender/receiver
 - **instantaneous**: rate at given point in time
 - **average**: rate over longer period of time



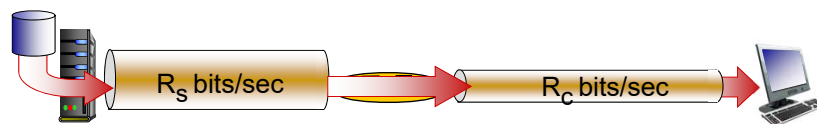
81

Throughput (more)

- $R_s < R_c$ What is average end-end throughput?



- ❖ $R_s > R_c$ What is average end-end throughput?



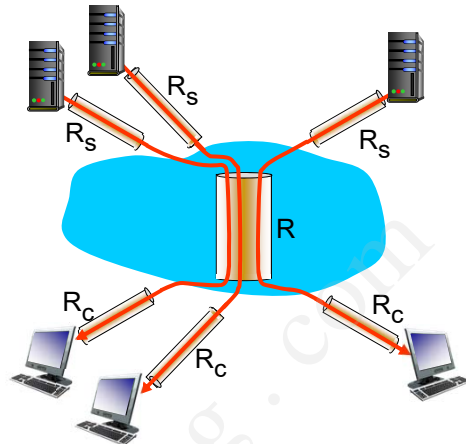
bottleneck link

link on end-end path that constrains end-end throughput

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Throughput: Internet scenario

- per-connection end-end throughput: $\min(R_c, R_s, R/10)$
- in practice: R_c or R_s is often bottleneck



10 connections (fairly) share backbone bottleneck link R bits/sec

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Bài tập

- Compare the delay in sending an x -bit message over a k -hop path in a circuit-switched network and in a (lightly loaded) packet-switched network. The circuit setup time is s sec, the propagation delay is d sec per hop, the packet size is p bits, and the data rate is b bps. Under what conditions does the packet network have a lower delay?
- Suppose that x bits of user data are to be transmitted over a k -hop path in a packet-switched network as a series of packets, each containing p data bits and h header bits, with $x \gg (p + h)$. The bit rate of the lines is b bps and the propagation delay is negligible. What value of p minimizes the total delay?
- A LMDS (Local Multipoint Distribution Service), each sector has its own 36-Mbps channel. According to queuing theory, if the channel is 50% loaded, the queuing time will be equal to the download time. Under these conditions, how long does it take to download a 5-KB Web page? How long does it take to download the page over a 1-Mbps ADSL line? Over a 56-kbps modem?

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Internet history

- Bắt đầu từ một thí nghiệm của dự án của ARPA
- Một liên kết giữa hai nút mạng (IMP tại UCLA và IMP tại SRI).
 - ARPA: Advanced Research Project Agency
 - UCLA: University California Los Angeles
 - SRI: Stanford Research Institute
 - IMP: Interface Message Processor

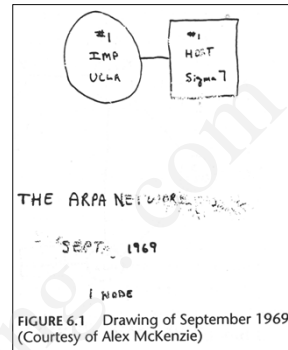


FIGURE 6.1 Drawing of September 1969 (Courtesy of Alex McKenzie)

85

Internet history

- 3 tháng sau: một mạng hoàn chỉnh với 4 nút, 56kbps

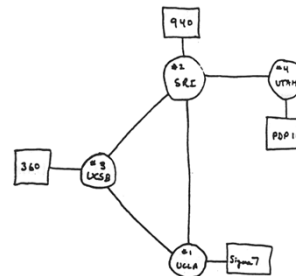
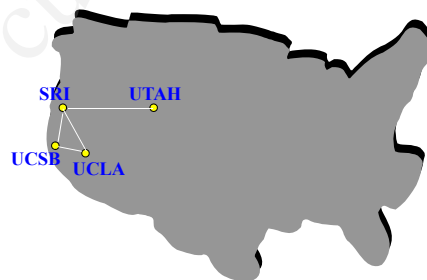


FIGURE 6.2 Drawing of 4 Node Network (Courtesy of Alex McKenzie)

86

Internet history

■ Thập niên 1970:

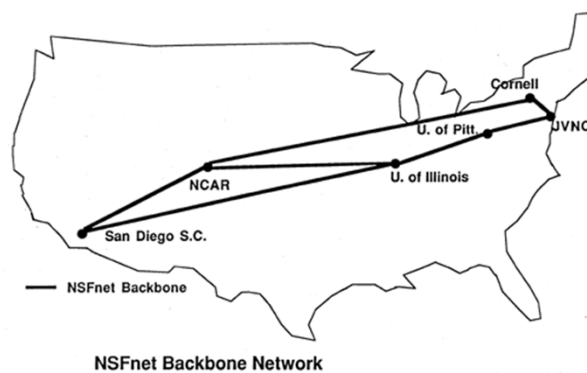
- Từ đầu 1970 xuất hiện các mạng riêng:
 - ALOHAnet tại Hawaii
 - DECnet, IBM SNA, XNA
- 1974: Cerf & Kahn – nguyên lý kết nối các hệ thống mở (Turing Awards)
- 1976: Ethernet, Xerox PARC

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Internet history

■ 1981: Xây dựng mạng NSFNET

- NSF: National Science Foundation: Phục vụ cho nghiên cứu khoa học, do sự quá tải của ARPANET



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Internet history

- 1986: Nối kết USENET& NSFNET

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usenet

Site locations and access exchange paths
Connectivity information from end user to backbone (1986)
Geographic information from CIA World Factbook data
Printed on 100% 12x18 inch high quality paper
Produced with netmap 1.2 by DEC Systems Research Lab

○
Thick lines are backbone paths
Thin solid lines are non-backbone paths
Dashed lines are non-backbone partial links

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Internet history

- Thêm nhiều mạng mới nối vào: MFENET, HEPNET (Dept. Energy), SPAN (NASA), BITnet, CSnet, NSFnet, Minitel ...
- TCP/IP được chuẩn hóa và phổ biến vào 1980
- Berkeley tích hợp TCP/IP vào BSD Unix
- Dịch vụ: FTP, Mail, DNS ...

Cuối 1990's – 2000's:

- Thập niên 90: Web và thương mại hóa Internet
- Đầu 90: ARPAnet chỉ là một phần của Internet
- Đầu 90: **Web**
 - **HTML, HTTP**: Berners-Lee
 - 1994: **Mosaic, Netscape**
- Cuối 90: Thương mại hóa Internet

- Nhiều ứng dụng mới: chat, chia sẻ file **P2P...**
- **E-commerce, Yahoo, Amazon, Google...**
- > 50 triệu máy trạm, > 100 triệu NSD
- **Vấn đề an toàn an ninh thông tin!**
 - Internet dành cho tất cả mọi người
 - Tất cả các dịch vụ phải quan tâm tới vấn đề này

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Internet history

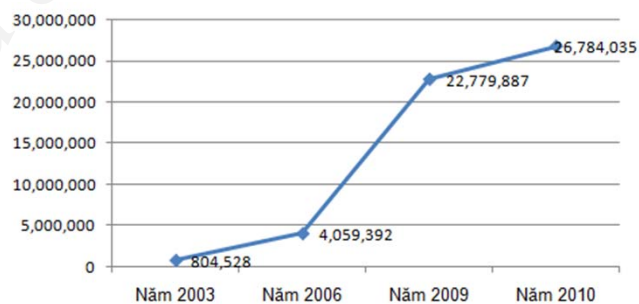
2005-present

- ~750 million hosts
 - Smartphones and tablets
- Aggressive deployment of broadband access
- Increasing ubiquity of high-speed wireless access
- Emergence of online social networks:
 - Facebook: soon one billion users
- Service providers (Google, Microsoft) create their own networks
 - Bypass Internet, providing “instantaneous” access to search, email, etc.
- E-commerce, universities, enterprises running their services in “cloud” (eg, Amazon EC2)

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Internet history

■ Phát triển Internet ở Việt Nam



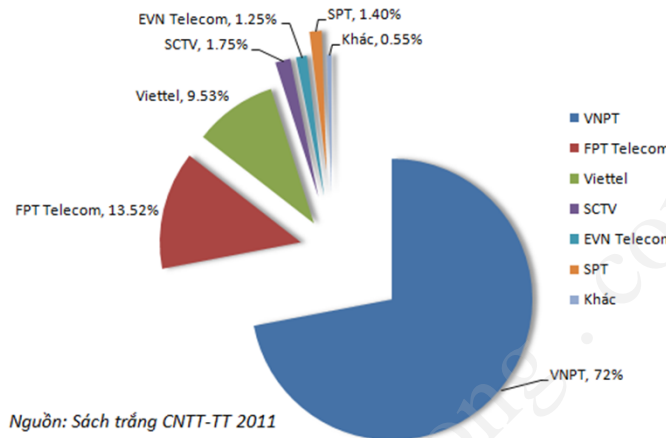
Bản đồ số lượng người dùng Internet ở Việt Nam qua các năm

https://vi.wikipedia.org/wiki/Internet_tại_Việt_Nam

94

Internet history

- Thị phần thuê bao dịch vụ truy cập Internet của các doanh nghiệp (tính đến tháng 12/2010)



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Introduction: summary

covered a “ton” of material!

- Internet overview
- what's a protocol?
- network edge, core, access network
 - packet-switching versus circuit-switching
 - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

you now have:

- context, overview, “feel” of networking
- more depth, detail *to follow!*

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