

$$5.3. (\epsilon = 1, \ell \geq 2) - N_{h=2\ell, r, s=2\ell+1}$$

5.2.2. Find the Upper Bound of  $r_{\max, \text{scalar}}$

there is no unique upper bound

Find  $(\alpha + 1)$  received vectors that span a subspace of dimension  $h$ . This implies that the  $\alpha$  links from the middle layer carry  $\alpha$  vectors which span a subspace of  $\mathbb{F}_{q_s}^h$  whose dimension is at least  $(h - 1)$ , with  $q_s = q^t$ .

Following to Theorem (4.1), we are interested in the following range:  $\ell + \epsilon + 1 \leq h \leq \alpha \ell + \epsilon$ .

(everywhere!) and why is then  $r \leq \begin{bmatrix} \alpha \\ 1 \end{bmatrix}_{q_s}$ ?

For  $3 \leq \alpha < h$ , all  $\alpha$  links must be distinct  $\Rightarrow r \leq \begin{bmatrix} \alpha \\ 1 \end{bmatrix}_{q_s} \Rightarrow r \leq O(q_s^{\alpha-1})$

why? more explanation!

For  $\alpha \geq h \geq 3$ : to achieve  $(h - 1)$ -subspaces of  $\mathbb{F}_{q_s}^h$ , no  $\alpha$  links will contain a vector which is contained in the same  $(h - 2)$ -subspace.

Hence,

for  $\epsilon \geq 1$ , shouldn't it be  $h - \epsilon$  and  $h - \epsilon - 1$ ?

what achieves this and why? This is not a sentence that makes sense.

$$r_{\max, \text{scalar}} \leq (\alpha - 1) \begin{bmatrix} \alpha \\ h - 2 \end{bmatrix}_{q_s} \Rightarrow r_{\max, \text{scalar}} \in O(q_s^{(\alpha - h + 2)(h - 2)t^2})$$

made more details!

seems to be correct

where does it come from?

5.2.3. Calculate Gap

+ a sentence

$$r_{\max, \text{scalar}} = \min \{ r_{\max, \text{vector}} \}$$

state as Theorem?

\* do you mean lin. independent?

state a



Proof. Being similar with the previous subsections, we have:

$$d \leq \alpha \binom{r-1}{\alpha-1} = 2 \frac{(r-1) \dots (r-1)}{1!} \leq 2r$$

□

**Theorem 5.3.** If  $r \leq \Omega(q^{t^2 + \mathcal{O}(t)})$ , then there exists a vector solution for the  $(\epsilon = 1, \ell \geq 2) - \mathcal{N}_{h=2\ell, r, s=2\ell+1}$  network.

*Proof.* As previous, we need  $4 \cdot p \cdot d(r) \leq 1, \forall r \leq r_{\max, \text{vector}}$  so that a vector solution exists. Following to Lemma 5.7, we have  $d \leq 2r \Rightarrow 4 \cdot p \cdot 2r \leq 1 \Rightarrow r \leq \frac{1}{8p}$ . We still need to maximize  $p$  to get lower on  $r_{\max, \text{vector}}$ , and we have  $p \leq \Theta(q^{-t^2 - 2t - 1}), \forall t \geq 2$  in Lemma 5.6. Thus,  $\min \{r_{\max, \text{vector}}\} \in \Omega\left(\frac{1}{8p}\right) = \Omega(q^{t^2 + 2t + 1})$ .

Hence, the Local lemma in 5.1 is satisfied, when  $r \leq \Omega(q^{t^2/2 + \mathcal{O}(t)})$ . None of bad events occurs, so there exists a vector solution for such  $r$ . □

**Lemma 5.8.** A scalar solution for the  $(\epsilon = 1, \ell \geq 2) - \mathcal{N}_{h=2\ell, r, s=2\ell+1}$  network exists, if and only if there exists a Grassmannian code  $\mathcal{G}_q(h = 2\ell, \ell)$  such that any  $\alpha = 2$  subspaces of the set span a subspace of dimension at least  $2\ell - 1$ .

*Proof.* Let's denote any 2 subspaces of  $\mathcal{G}_q(h = 2\ell, \ell)$  as  $\mathcal{U}$  and  $\mathcal{V}$ . If  $\mathcal{U}$  and  $\mathcal{V}$  span a subspace of dimension at least  $2\ell - 1$ , then we have  $\dim(\mathcal{U} + \mathcal{V}) = 2\ell - 1$ , with  $\mathcal{U} + \mathcal{V}$ .

not sure if correct number

See Section VII of Turi/Antonio's paper.

There they discuss the general case  $d=2$

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and derive an upper bound on  $r_{\max, \text{scalar}}$ . (all based on theory of subspace codes)

This is a subspace code!

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## 5. Combinatorial Results

Therefore,  $\dim(\mathcal{U} \cap \mathcal{V}) = \dim(\mathcal{U}) + \dim(\mathcal{V}) - \dim(\mathcal{U} + \mathcal{V}) = 1$ , which leads to the subspace distance  $d_s(\mathcal{U}, \mathcal{V}) = 2\dim(\mathcal{U} + \mathcal{V}) - \dim(\mathcal{U}) - \dim(\mathcal{V}) = 2\ell - 2$ . No  $2\ell$ -dimensional subspaces of  $\mathbb{F}_{q_s}^{2\ell}$  will contain a vector which is contained in the same  $(2\ell - 2)$ -subspace, but  $(2\ell - 1)$  of such subspaces can have such vectors,

$$\Rightarrow r_{\text{scalar}} \leq (2\ell - 1) \left[ \begin{matrix} 2\ell \\ 2\ell - 2 \end{matrix} \right]_{q_s} \Rightarrow r_{\text{scalar}} \leq \mathcal{O} \left( \binom{2\ell}{q_s} \right)$$

There are general upper bounds on the cardinality of subspace codes  $\rightarrow$  see

**Corollary 5.2.** The  $(\epsilon = 1, \ell \geq 2) - \mathcal{N}_{h=2\ell, r, s=2\ell+1}$  network has a vector solution with a gap  $q^{t^2/4 + \mathcal{O}(t)}$ .

Following to Section 3.4.2 and Theorem 5.3, we have the gap size

$$\begin{aligned} r_{\text{max}, \text{scalar}} &= \min \{ r_{\text{max}, \text{vector}} \} \\ \Leftrightarrow q_s^\ell &= q^{t^2/2 + \mathcal{O}(t)} \\ \Leftrightarrow q_s &= q^{t^2/2\ell + \mathcal{O}(t)} \\ \Rightarrow g &= q_s - q_v = q^{t^2/2\ell + \mathcal{O}(t)} \end{aligned} \quad (5.9)$$

This shows us that there exists a better vector solution by comparison with the gap in [EW18, Fig. 4].

$\rightarrow$  see Tavis/Ad. paper & compare to yours.