# PMT Revision (Week 10)

#### Nicholas Sim

#### March 15th, 2018

# 1 Logic

[See handout: C140 2016 Q1c]

# 2 Relations

- 1. [C142 2015 Q1d.v] Let  $R \subseteq A^2$ . Prove that if R is symmetric, then  $R \circ R$  is also symmetric.
- 2. [C142 2005 Q2b.ii] Let R be a binary relation on the natural numbers defined by xRy if and only if y = 2x. Describe the transitive closure of R.
- 3. [C142 2005 Q2b.iii] A binary relation R is antisymmetric if and only if  $(a,b) \in R$  and  $(b,a) \in R$  implies a=b. How many binary relations on  $A=\{1,2,3,4\}$  are both symmetric and antisymmetric? Justify your answer.

[Bonus: What about an arbitrary set?]

## 3 Functions

[C142 2002 Q1c] Let  $f:A\to B$  and  $g:B\to C$  be functions. Prove the following (otherwise give a counterexample):

- 1. f surjective implies  $g \circ f$  surjective;
- 2.  $g \circ f$  injective implies f injective;
- 3.  $g \circ f$  injective implies g injective;
- 4. f, g surjective implies  $g \circ f$  surjective.

Note: these solutions are only outlines.

# 4 Solution to Alternative Elimination

- 1. I would argue that (2-4) are sufficient for  $A \to B$ . From here, we can use EM to get  $A \vee \neg A$ , and  $\vee E$  to get B.
- 2. First use  $\forall I$  to get  $A \lor B \lor C$ . Assume A to get C then  $B \lor C$  (using  $\forall I$ ). Now use  $Alt \lor E$  to get  $B \lor C$ . Use  $Alt \lor E$  again to get C.
- 3. Observe that  $\forall E$  essentially follows from the previous part. (Say: if we have ..., then get  $\vdash^a C$ )
- 4.  $\vdash^a$  is complete if every valid statement expressed in propositional logic can be proven using only the rules of  $\vdash^a$ . It is sound if every statement that can be proven under the system is valid.
- 5. We know from the lectures that  $\vdash$  is sound and complete. Since  $Alt \lor E$  is a derived rule of  $\vdash$ , soundness is preserved (we can't prove anything new, hence nothing invalid). Since  $\lor E$  is a derived rule of  $\vdash$ <sup>a</sup>, completeness is preserved (we can prove at least the same things that  $\vdash$  can).

### 5 Solutions to Relations

- 1. Suppose that R is symmetric, i.e.  $(a,b) \in R \Rightarrow (b,a) \in R$ . Recall that  $R \circ R = \{(a,c) : (\exists b) ((a,b) \in R \land (b,c) \in R)\}$ . Fix  $(a,c) \in R \circ R$ . To show:  $(c,a) \in R \circ R$ . But this is clear as  $\exists b(a,b), (b,c) \in R$ , so  $(b,a), (c,b) \in R$  by symmetry, and  $(c,a) \in R \circ R$ .
- 2.  $x\overline{R}y$  if and only if  $\exists n \in \mathbb{N}_{>0} : y = 2^n x$ .
- 3. R is symmetric if and only if  $(b,a) \in R$  whenever  $(a,b) \in R$ . Let R be both symmetric and antisymmetric. Suppose  $(a,b) \in R$ . Then  $(b,a) \in R$ , so a=b. We easily see that the possible relations R are those which only contain pairs of the form (a,a). So 16.

### 6 Solutions to Functions

- 1. False, let  $C = \{0, 1\}$  and g(x) = 0... (write a surjective f and sets)
- 2. We show the contrapositive. Suppose f not injective. Then  $\exists a_1, a_2 \in A : f(a_1) = f(a_2) \in B$ . Clearly  $g(f(a_1)) = g(f(a_2))$  (as g is a well-defined function), so  $g \circ f$  is not injective.
- 3. False, let  $A = \{0, 1\}, B = \{0, 1, 2\}, C = \{0, 1\}, f = id, g(0) = g(2) = 0, g(1) = 1.$
- 4. Suppose f,g surjective, i.e.  $\forall b \in B \exists a \in A: f(a) = b$  and  $\forall c \in C \exists b \in B: g(b) = c$ . Fix  $c \in C$ . Want:  $\exists a \in A: g(f(a)) = c$ . We have  $b \in B$  s.t. g(b) = c. Similarly  $\exists a \in A: f(a) = b$ , which is what was required.