



# Correction to: The simpliciality of higher-order networks

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The original article can be found online at <https://doi.org/10.1140/epjds/s13688-024-00458-1>

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Following publication of the original article [1], the authors identified errors in the equations in the 2.1 Measures section, Algorithm 2, 3, 4, Figure 1, 2, 3, 4, Table 1, 2 values and a value and wording in the 2.1 Measures, 3.3 Local measures of simpliciality section and the Acknowledgements section.

The incorrect and correct equations and wording in the 2.1 Measures section, Algorithm 2, 3, Table 1, 2, 3.3 Local measures of simpliciality section and Acknowledgements section are indicated hereafter.

The incorrect 2.1. sub-sections [Edit Simpliciality](#) and [Face Edit Simpliciality](#):

**Edit simpliciality** The *edit simpliciality* (ES) is defined as the minimal number (or fraction, in the normalized case) of additional edges needed to make a hypergraph a simplicial complex.

Our formal definition uses the notion of an induced simplicial complex defined in Sect. 1.1. Given a hypergraph  $\mathbf{H} = (V, E)$  for which we want to measure the ES, we find its maximal edges  $\tilde{E}$  and construct the simplicial complex  $\mathbf{S} = (V, C)$  induced on  $\mathbf{H}$ , with  $C = \bigcup_{e \in \tilde{E}} \mathcal{P}(e)$ . The edit simpliciality is then

$$\sigma_{\text{ES}} = \frac{|E|}{|C|}, \quad (2)$$

again satisfying  $\sigma_{\text{ES}} \in [0, 1]$ ; see Fig. 1C. (We note that one can use the induced simplicial complex to define variants of the ES, e.g., a simplicial edit distance  $d_{\text{ES}} = |C| - |E|$  or a normalized distance  $d_{\text{NES}} = (|C| - |E|)/|C| = 1 - |E|/|C| = 1 - \sigma_{\text{ES}}$ .)

The ES answers a slightly different question than the SF does—it counts missing hyperedges that would make the dataset into a simplicial complex, rather than the edges that already satisfy downward closure. It thus offers a complementary, equally interpretable measure of simpliciality. However, the ES has the disadvantage of being sensitive to outliers, as a handful of large hyperedges with few inclusions will rapidly drive  $\sigma_{\text{ES}}$  towards

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0. Indeed, a hyperedge of size  $m$  without any inclusion contributes one edge to  $|E|$  but  $2^m$  edges to  $|C|$  in the denominator of Eq. (2).

*Face edit simpliciality* Finally, building upon the idea of edit simpliciality, we define a more localized notion of simpliciality, using the number of subsurfaces that must be added to the hypergraph to make a particular face a simplex.

Given a hyperedge  $e$ , the number of edges one must add to the hypergraph to make  $e$  a simplex is

$$d_{\text{FES}}(e) = |\mathcal{P}(e)| - |c|,$$

where  $c = \{f \in E \mid f \subseteq e\}$ . We can think of this quantity as an edit distance, or *face edit distance*. We use this quantity to define an average

$$\bar{d}_{\text{FES}} = \frac{1}{|F|} \sum_{e \in F} d_{\text{FES}}(e),$$

where  $F$  is a set of edges—most commonly,  $F = \tilde{E}$  or  $E$ . We exclusively use  $F = \tilde{E}$  in this study. These quantities are on the scale of counts, and to define quantities analogous to previous simpliciality measures, we thus introduce a per-face normalization, either on a distance scale (meaning that the quantity grows as the dataset becomes less simplicial):

$$\bar{d}_{\text{NFES}} = \frac{1}{|F|} \sum_{e \in F} \frac{d_{\text{FES}}(e)}{|\mathcal{P}(e)|},$$

or, similarly to previous definitions, on a simpliciality scale:

$$\sigma_{\text{FES}} = \frac{1}{|F|} \sum_{e \in F} \left( 1 - \frac{d_{\text{FES}}(e)}{|\mathcal{P}(e)|} \right). \quad (3)$$

We call this last measure the *face edit simpliciality* (FES).

The FES normalizes the face edit distance as a fraction of its maximal simpliciality. This normalization removes the dominance of large edges in the calculation of  $\sigma_{\text{ES}}$  and, in fact, exponentially down-weights the contribution of these edges. In addition, because this metric is computed on faces, this is an averaged local metric.

The correct 2.1. sub-sections Edit Simpliciality and Face Edit Simpliciality:

*Edit simpliciality* The *edit simpliciality* (ES) is defined as the minimal number (or fraction, in the normalized case) of additional edges needed to make a hypergraph a simplicial complex.

Our formal definition uses the notion of an induced simplicial complex defined in Sect. 1.1. Given a hypergraph  $\mathbf{H} = (V, E)$  for which we want to measure the ES, we find its maximal edges  $\tilde{E}$  and construct the simplicial complex  $\mathbf{S} = (V, C)$  induced on  $\mathbf{H}$ , with  $C = \bigcup_{e \in \tilde{E}} \mathcal{P}(e)$ . The edit simpliciality is then

$$\sigma_{\text{ES}} = \frac{|E| - |\tilde{E}|}{|C| - |\tilde{E}|}, \quad (2)$$

again satisfying  $\sigma_{\text{ES}} \in [0, 1]$ ; see Fig. 1C. We subtract the number of maximal edges from the total number of edges so that the edit simpliciality is zero in the case of a hypergraph with no subsfaces. (We note that one can use the induced simplicial complex to define variants of the ES, e.g., a simplicial edit distance  $d_{\text{ES}} = |C| - |E|$  or a normalized distance  $d_{\text{NES}} = (|C| - |E|)/(|C| - |\tilde{E}|) = 1 - (|E| - |\tilde{E}|)/(|C| - |\tilde{E}|) = 1 - \sigma_{\text{ES}}$ .)

The ES answers a slightly different question than the SF does—it counts missing hyperedges that would make the dataset into a simplicial complex, rather than the edges that already satisfy downward closure. It thus offers a complementary, equally interpretable measure of simpliciality. However, the ES has the disadvantage of being sensitive to outliers, as a handful of large hyperedges with few inclusions will rapidly drive  $\sigma_{\text{ES}}$  towards 0. Indeed, a hyperedge of size  $m$  without any inclusion contributes one edge to  $|E|$  but  $2^m$  edges to  $|C|$  in the denominator of Eq. (2).

*Face edit simpliciality* Finally, building upon the idea of edit simpliciality, we define a more localized notion of simpliciality, using the number of subsfaces that must be added to the hypergraph to make a particular face a simplex.

Given a hyperedge  $e$ , the number of edges one must add to the hypergraph to make  $e$  a simplex is

$$d_{\text{FES}}(e) = |\mathcal{P}(e)| - |c|,$$

where  $c = \{f \in E \mid f \subseteq e\}$ . We can think of this quantity as an edit distance, or *face edit distance*. We use this quantity to define an average

$$\bar{d}_{\text{FES}} = \frac{1}{|F|} \sum_{e \in F} d_{\text{FES}}(e),$$

where  $F$  is a set of edges—most commonly,  $F = \tilde{E}$  or  $E$ . We exclusively use  $F = \tilde{E}$  in this study. These quantities are on the scale of counts, and to define quantities analogous to previous simpliciality measures, we thus introduce a per-face normalization, either on a distance scale (meaning that the quantity grows as the dataset becomes less simplicial):

$$\bar{d}_{\text{NFES}} = \frac{1}{|F|} \sum_{e \in F} \frac{d_{\text{FES}}(e)}{|\mathcal{P}(e)| - 1},$$

or, similarly to previous definitions, on a simpliciality scale:

$$\sigma_{\text{FES}} = \frac{1}{|F|} \sum_{e \in F} \left( 1 - \frac{d_{\text{FES}}(e)}{|\mathcal{P}(e)| - 1} \right). \quad (3)$$

We call this last measure the *face edit simpliciality* (FES). We subtract one in the denominators of both expressions so that when an edge has no subsfaces, its normalized face edit distance is one.

The FES normalizes the face edit distance as a fraction of its maximal simpliciality. This normalization removes the dominance of large edges in the calculation of  $\sigma_{\text{ES}}$  and, in fact, exponentially down-weights the contribution of these edges. In addition, because this metric is computed on faces, this is an averaged local metric.

### The incorrect Table 1:

**Table 1** Properties of empirical datasets and their simpliciality.  $|V|$ ,  $|E|$ ,  $\langle k \rangle$ ,  $\langle s \rangle$ ,  $\sigma_{SF}$ ,  $\sigma_{ES}$ , and  $\sigma_{FES}$  denote the number of nodes, the number of hyperedges, the mean degree, the mean edge size, the simplicial fraction (SF), edit simpliciality (ES), and the face edit simpliciality (FES), respectively

Dataset	$ V $	$ E $	$\langle k \rangle$	$\langle s \rangle$	$\sigma_{SF}$	$\sigma_{ES}$	$\sigma_{FES}$
<i>Proximity datasets</i>							
contact-primary-school	242	12,704	52.50	2.42	0.85	0.92	0.94
contact-high-school	327	7,818	23.91	2.33	0.81	0.93	0.92
hospital-lyon	75	1,824	24.32	2.43	0.91	0.95	0.97
<i>Email datasets</i>							
email-enron	143	1,442	10.08	2.97	0.31	0.05	0.50
email-eu	967	23,729	24.54	3.12	0.32	0.05	0.52
<i>Biological datasets</i>							
diseasome	516	314	0.61	3.00	0.00	0.05	0.04
disgenenet	1,982	760	0.38	5.14	0.00	0.00	0.01
ndc-substances	2,740	4,754	1.74	5.16	0.02	0.01	0.07
<i>Other</i>							
congress-bills	1,715	58,788	34.28	4.95	0.03	0.01	0.10
tags-ask-ubuntu	3,021	145,053	48.01	3.43	0.15	0.25	0.46

### The correct Table 1:

**Table 1** Properties of empirical datasets and their simpliciality.  $|V|$ ,  $|E|$ ,  $\langle k \rangle$ ,  $\langle s \rangle$ ,  $\sigma_{SF}$ ,  $\sigma_{ES}$ , and  $\sigma_{FES}$  denote the number of nodes, the number of hyperedges, the mean degree, the mean edge size, the simplicial fraction (SF), edit simpliciality (ES), and the face edit simpliciality (FES), respectively

Dataset	$ V $	$ E $	$\langle k \rangle$	$\langle s \rangle$	$\sigma_{SF}$	$\sigma_{ES}$	$\sigma_{FES}$
<i>Proximity datasets</i>							
contact-primary-school	242	12,704	52.50	2.42	0.85	0.88	0.94
contact-high-school	327	7,818	23.91	2.33	0.81	0.91	0.92
hospital-lyon	75	1,824	24.32	2.43	0.91	0.94	0.97
<i>Email datasets</i>							
email-enron	143	1,442	10.08	2.97	0.31	0.04	0.50
email-eu	967	23,729	24.54	3.12	0.32	0.04	0.52
<i>Biological datasets</i>							
diseasome	516	314	0.61	3.00	0.00	0.02	0.04
disgenenet	1,982	760	0.38	5.14	0.00	0.00	0.01
ndc-substances	2,740	4,754	1.74	5.16	0.02	0.00	0.07
<i>Other</i>							
congress-bills	1,715	58,788	34.28	4.95	0.03	0.00	0.10
tags-ask-ubuntu	3,021	145,053	48.01	3.43	0.15	0.11	0.46

### The incorrect 3.3 Local measures of simpliciality section:

The sentence currently reads:

(The correlation drops to  $\rho = 0.69$  when comparing the SF and ES).

The sentence should read:

(The correlation drops to  $\rho = 0.6$  when comparing the SF and ES).

The sentence currently reads:

For tags-ask-ubuntu, FES is weakly assortative, whereas the other two measures are weakly disassortative.

The sentence should read:

For tags-ask-ubuntu, FES is weakly assortative, SF is weakly disassortative, and ES is strongly disassortative.

The incorrect Table 2:

**Table 2** The simplicial assortativity of each dataset filtered to only include interactions of sizes two and three for computational tractability

Dataset	$\rho_{SF}$	$\rho_{ES}$	$\rho_{FES}$
<i>Proximity datasets</i>			
contact-primary-school	0.15	0.15	0.14
contact-high-school	0.22	0.34	0.24
hospital-lyon	-0.02	-0.02	-0.01
<i>Email datasets</i>			
email-enron	0.29	0.29	0.24
email-eu	0.19	0.16	0.16
<i>Biological datasets</i>			
ndc-substances	0.57	0.65	0.72
diseasome	N/A	0.46	0.75
disgenenet	N/A	0.55	0.89
<i>Other</i>			
congress-bills	0.78	0.48	0.75
tags-ask-ubuntu	-0.03	-0.08	0.04

The correct Table 2:

**Table 2** The simplicial assortativity of each dataset filtered to only include interactions of sizes two and three for computational tractability

Dataset	$\rho_{SF}$	$\rho_{ES}$	$\rho_{FES}$
<i>Proximity datasets</i>			
contact-primary-school	0.15	0.17	0.14
contact-high-school	0.22	0.37	0.24
hospital-lyon	-0.02	-0.01	-0.01
<i>Email datasets</i>			
email-enron	0.29	0.29	0.24
email-eu	0.19	0.16	0.16
<i>Biological datasets</i>			
ndc-substances	0.56	0.54	0.69
diseasome	N/A	0.28	0.68
disgenenet	N/A	0.28	0.78
<i>Other</i>			
congress-bills	0.78	0.33	0.75
tags-ask-ubuntu	-0.03	-0.24	0.04

The incorrect Algorithm 2:

---

**Algorithm 2:** Exhaustive edit simpliciality

---

**Input:**  $K$ , a set of edge sizes

$m$ , the minimum acceptable simplex size

$\mathbf{H} = (V, E)$ , a hypergraph

$T$ , a trie constructed from the edges in  $\mathbf{H}$

**Output:**  $\sigma_{\text{ES}}$

$\sigma_{\text{ES}} = 0$

// Construct the set of maximal faces.

$F = \{e \in E \mid e \notin f, \forall f \in E, |e| \geq m\}$

//  $D$  stores the unique missing subfaces.

$D = \emptyset$ , is a set of sets

// Iterate over all maximal faces.

**for**  $f \in F$  **do**

// For each maximal face of the hypergraph, add all of  
its missing subfaces not already present in the  
global set of missing faces.

**for**  $e \in \mathcal{P}_K(f)$  **do**

**if**  $e \notin T$  **then**  
     $| D \leftarrow D \cup e$   
    **end**

**end**

**end**

// The number of edges in the minimal simplicial complex  
is the sum of the number of edges in the original  
hypergraph and the number of missing subfaces.

$\sigma_{\text{ES}} = |E|/(|E| + |D|)$

**return**  $\sigma_{\text{ES}}$

---

The correct Algorithm 2:

---

**Algorithm 2:** Exhaustive edit simpliciality
 

---

**Input:**  $K$ , a set of edge sizes  
 $m$ , the minimum acceptable simplex size  
 $\mathbf{H} = (V, E)$ , a hypergraph  
 $T$ , a trie constructed from the edges in  $\mathbf{H}$

**Output:**  $\sigma_{\text{ES}}$

```

 $\sigma_{\text{ES}} = 0$ 
 $E = \{e \in E \mid |e| \in K, |e| > m\}$ 
// Construct the set of maximal faces.
 $F = \{e \in E \mid e \notin f, \forall f \in E\}$ 
// D stores the unique missing subfaces.
 $D = \emptyset$ , is a set of sets
// Iterate over all maximal faces.
for  $f \in F$  do
    // For each maximal face of the hypergraph, add all of
    // its missing subfaces not already present in the
    // global set of missing faces.
    for  $e \in \mathcal{P}_K(f)$  do
        if  $e \notin T$  then
             $|D \leftarrow D \cup e$ 
        end
    end
end
// The number of edges in the minimal simplicial complex
// is the sum of the number of edges in the original
// hypergraph and the number of missing subfaces.
 $\sigma_{\text{ES}} = (|E| - |\tilde{E}|) / (|E| - |\tilde{E}| + |D|)$ 
return  $\sigma_{\text{ES}}$ 

```

---

The incorrect Algorithm 3:

---

**Algorithm 3:** Memory-efficient edit simpliciality
 

---

**Input:**  $K$ , a set of edge sizes  
 $m$ , the minimum acceptable simplex size  
 $\mathbf{H} = (V, E)$ , a hypergraph  
 $T$ , a trie constructed from the edges in  $\mathbf{H}$

**Output:**  $\sigma_{\text{ES}}$

```

 $\sigma_{\text{ES}} = 0$ 
// Construct the set of maximal faces.
 $F = \{e \in E \mid e \notin f, \forall f \in E, |e| \geq m\}$ 
 $d = 0$ 
// Iterate over all enumerated maximal faces.
for  $i = 1 \dots |F|$  do
   $f = F_i$ 
  // First, calculate the number of missing faces for a given
  // maximal face.
   $\tilde{d} = |\mathcal{P}_K(f)|$ 
  for  $e \in \mathcal{P}_K(f)$  do
    if  $e \in T$  then
      |  $\tilde{d} \leftarrow \tilde{d} - 1$ 
    end
  end
  // Update the total number of missing subfaces
   $d \leftarrow d + \tilde{d}$ 
  // Calculate the number of redundant missing subfaces
  // counted for the maximal faces already seen. To prevent
  // looping over all previous maximal edges, we iterate only
  // over the previous maximal faces, which are also
  // neighbors of the current maximal face.
   $D = \emptyset$ 
  for  $j = \{1 \dots i - 1\} \cap \{k \mid e_k \cap f \neq \emptyset\}$  do
    // For each prior maximal face, we add the missing edges
    // formed by the powerset of the intersection of that
    // edge and the current maximal edge to the complete set
    // of redundant missing edges.
     $e = F_j$ 
     $g = e \cap f$ 
    for  $h \in \mathcal{P}_{K \cup |g|}(g)$  do
      if  $g \notin T$  then
        |  $D \leftarrow D \cup g$ 
      end
    end
  end
  // Subtract the redundant missing subfaces
   $d \leftarrow d - |D|$ 
end
 $\sigma_{\text{ES}} = |E| / (|E| + d)$ 
return  $\sigma_{\text{ES}}$ 

```

---

The correct Algorithm 3:

---

**Algorithm 3:** Memory-efficient edit simpliciality
 

---

**Input:**  $K$ , a set of edge sizes  
 $m$ , the minimum acceptable simplex size  
 $\mathbf{H} = (V, E)$ , a hypergraph  
 $T$ , a trie constructed from the edges in  $\mathbf{H}$

**Output:**  $\sigma_{ES}$

```

 $\sigma_{ES} = 0$ 
 $E = \{e \in E \mid |e| \in K, |e| > m\}$ 
// Construct the set of maximal faces.
 $F = \{e \in E \mid e \notin f, \forall f \in E\}$ 
 $d = 0$ 
// Iterate over all enumerated maximal faces.
for  $i = 1 \dots |F|$  do
   $f = F_i$ 
  // First, calculate the number of missing faces for a given
  // maximal face.
   $\tilde{d} = |\mathcal{P}_K(f)|$ 
  for  $e \in \mathcal{P}_K(f)$  do
    if  $e \in T$  then
       $\tilde{d} \leftarrow \tilde{d} - 1$ 
    end
  end
  // Update the total number of missing subfaces
   $d \leftarrow d + \tilde{d}$ 
  // Calculate the number of redundant missing subfaces
  // counted for the maximal faces already seen. To prevent
  // looping over all previous maximal edges, we iterate only
  // over the previous maximal faces, which are also
  // neighbors of the current maximal face.
   $D = \emptyset$ 
  for  $j = 1 \dots i - 1 \cap \{k \mid e_k \cap f \neq \emptyset\}$  do
    // For each prior maximal face, we add the missing edges
    // formed by the powerset of the intersection of that
    // edge and the current maximal edge to the complete set
    // of redundant missing edges.
     $e = F_j$ 
     $g = e \cap f$ 
    for  $h \in \mathcal{P}_{K \cup g}(g)$  do
      if  $g \notin T$  then
         $D \leftarrow D \cup g$ 
      end
    end
  end
  // Subtract the redundant missing subfaces
   $d \leftarrow d - |D|$ 
end
// The number of edges in the minimal simplicial complex is
// the sum of the number of edges in the original hypergraph
// and the number of missing subfaces.
 $\sigma_{ES} = (|E| - |\tilde{E}|) / (|E| - |\tilde{E}| + d)$ 
return  $\sigma_{ES}$ 

```

---

The incorrect Algorithm 4:

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**Algorithm 4:** Face edit simpliciality

---

**Input:**  $K$ , a set of edge sizes  
     $m$ , the minimum acceptable simplex size  
     $\mathbf{H} = (V, E)$ , a hypergraph  
     $T$ , a trie constructed from the edges in  $\mathbf{H}$

**Output:**  $\sigma_{\text{FES}}$

```
 $\sigma_{\text{FES}} = 0$ 
// Construct the set of maximal faces.
 $F = \{e \in E \mid e \notin f, \forall f \in E, |e| \geq m\}$ 
 $\sigma_{\text{FES}} = 0$ 
// Iterate over all maximal faces.
for  $f \in F$  do
    // For each maximal face, calculate the fraction of
    // missing faces.
     $s = 0$ 
    for  $e \in \mathcal{P}_K(f)$  do
        if  $e \in T$  then
             $| s \leftarrow s + 1/|\mathcal{P}_K(f)|$ 
        end
    end
    // Update the running average.
     $\sigma_{\text{FES}} \leftarrow \sigma_{\text{FES}} + s/|F|$ 
end
return  $\sigma_{\text{FES}}$ 
```

---

The correct Algorithm 4:

---

**Algorithm 4:** Face edit simplicity

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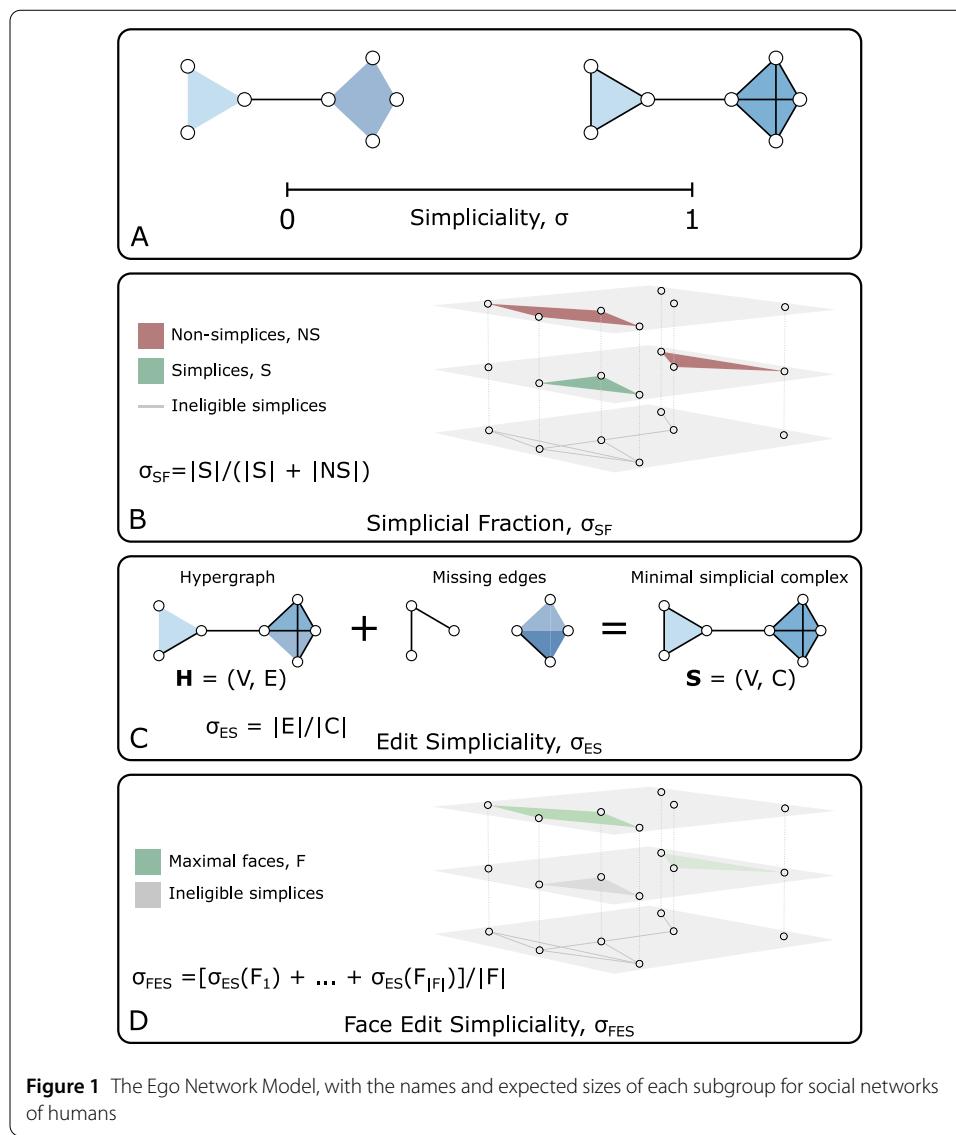
**Input:**  $K$ , a set of edge sizes  
     $m$ , the minimum acceptable simplex size  
     $\mathbf{H} = (V, E)$ , a hypergraph  
     $T$ , a trie constructed from the edges in  $\mathbf{H}$

**Output:**  $\sigma_{\text{FES}}$

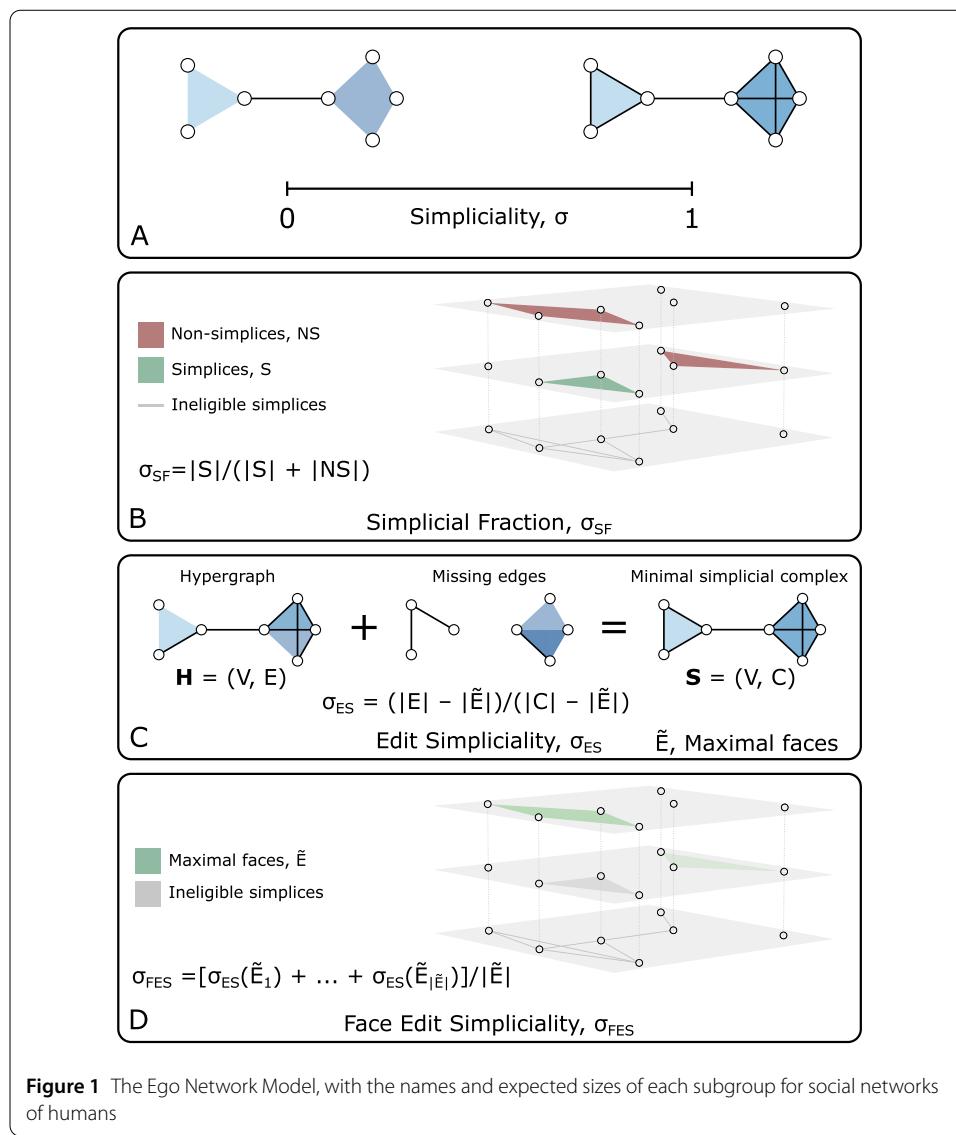
```
 $\sigma_{\text{FES}} = 0$ 
 $E = \{e \in E \mid |e| \in K, |e| > m\}$ 
// Construct the set of maximal faces.
 $F = \{e \in E \mid e \notin f, \forall f \in E\}$ 
 $\sigma_{\text{FES}} = 0$ 
// Iterate over all maximal faces.
for  $f \in F$  do
    // For each maximal face, calculate the fraction of
    // missing faces.
     $s = 0$ 
    for  $e \in \mathcal{P}_K(f)$  do
        if  $e \in T$  then
             $| s \leftarrow s + 1/(|\mathcal{P}_K(f)| - 1)$ 
        end
    end
    // Update the running average.
     $\sigma_{\text{FES}} \leftarrow \sigma_{\text{FES}} + s/|F|$ 
end
return  $\sigma_{\text{FES}}$ 
```

---

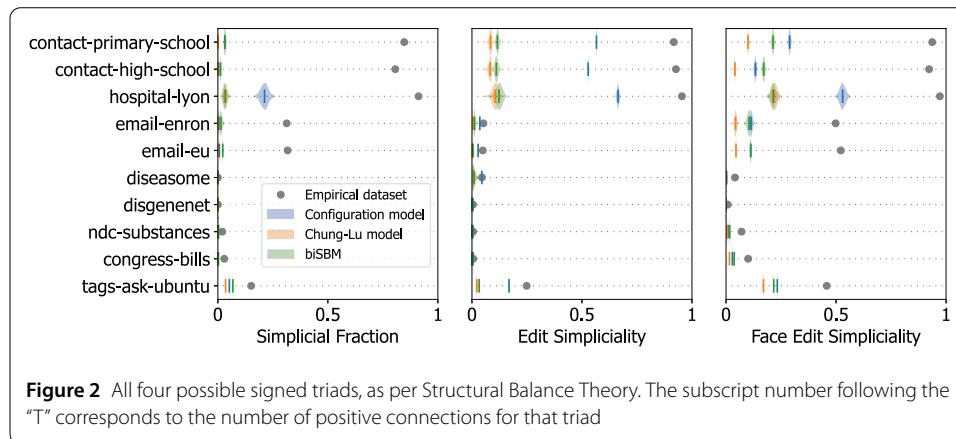
The incorrect Figure 1:



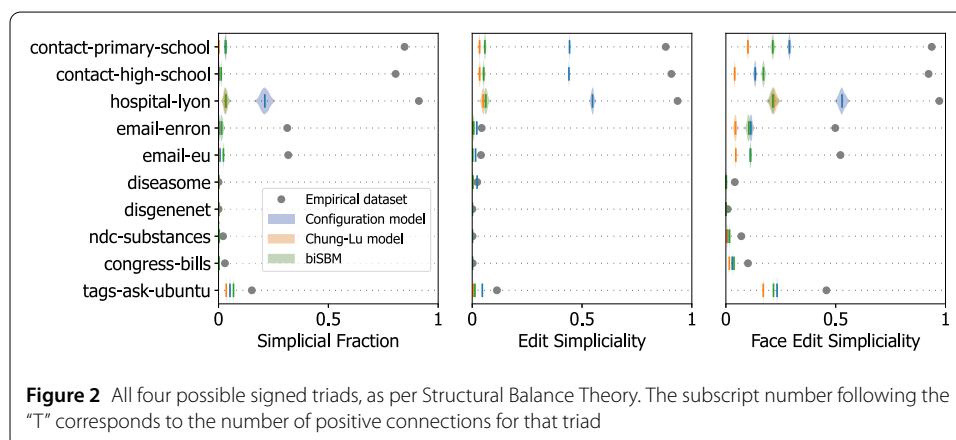
The correct Figure 1:



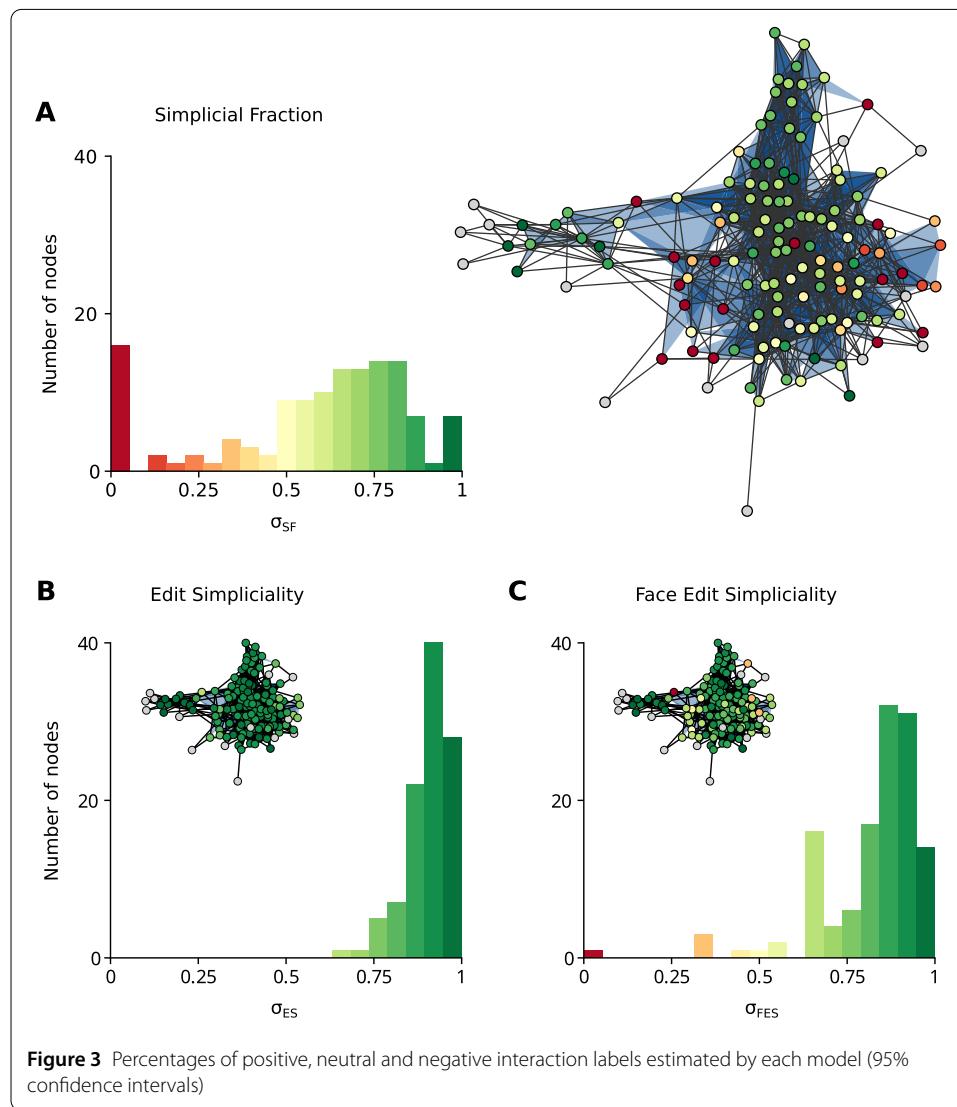
The incorrect Figure 2:



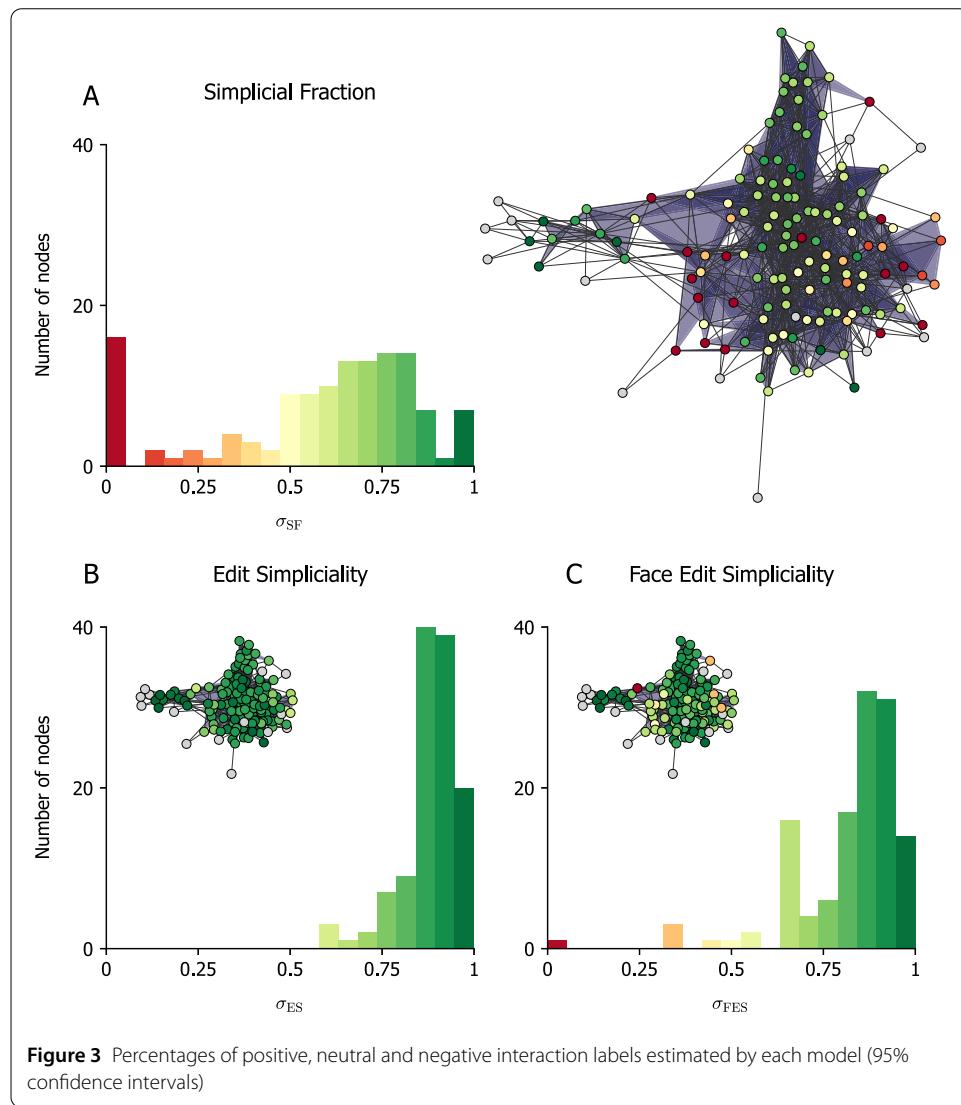
The correct Figure 2:



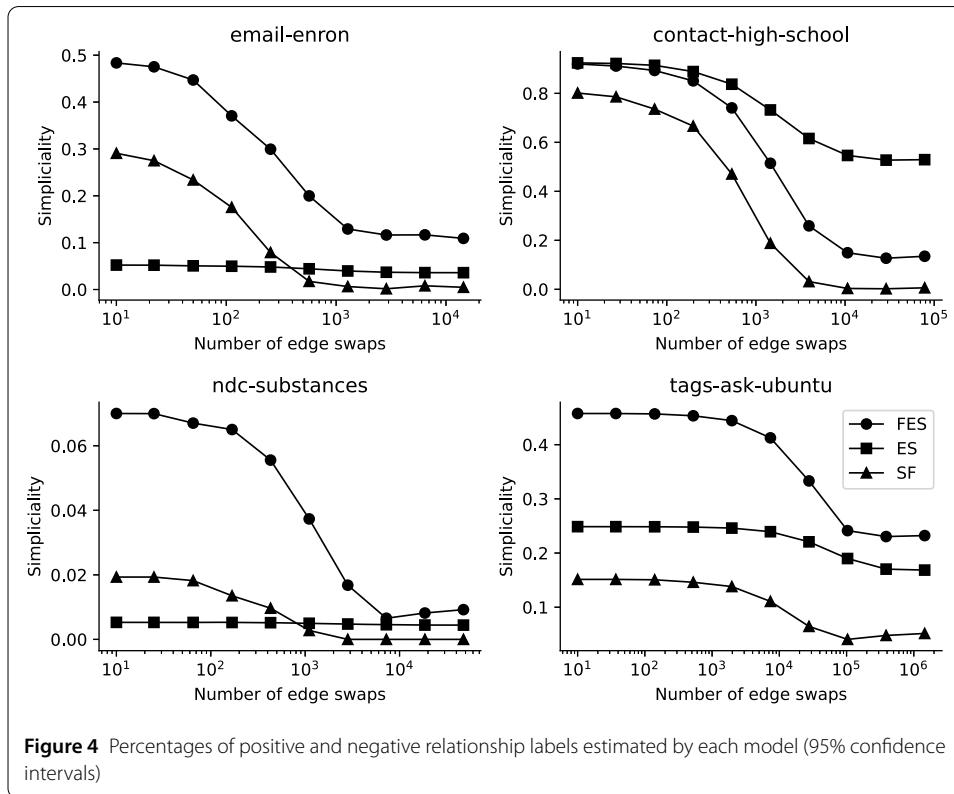
The incorrect Figure 3:



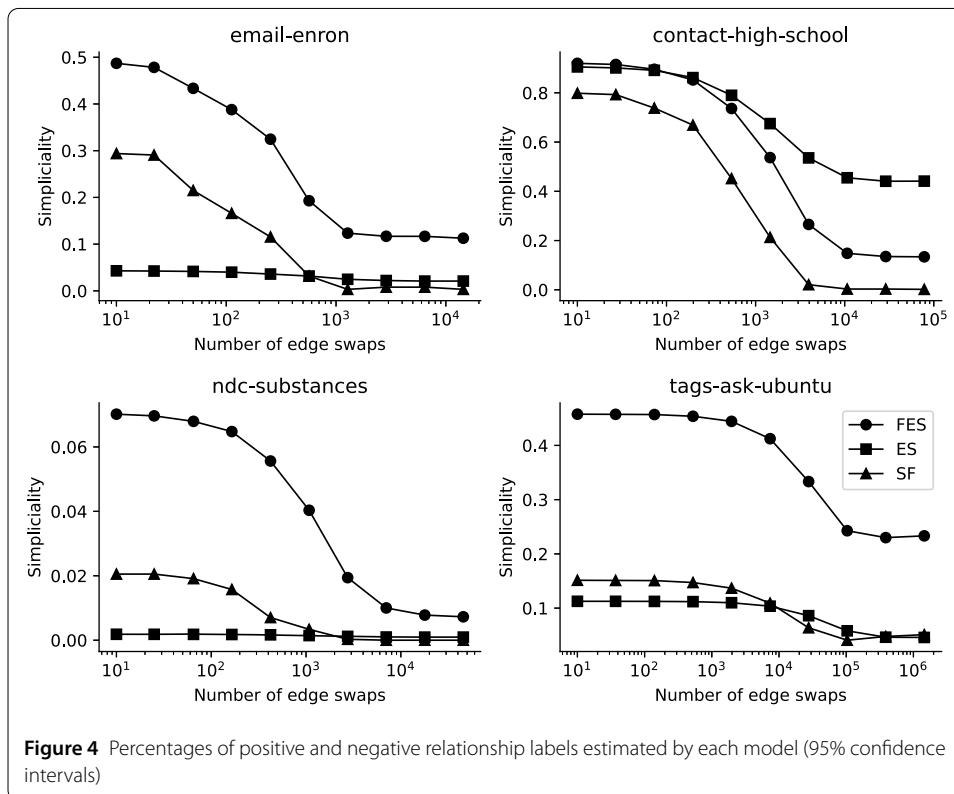
The correct Figure 3:



The incorrect Figure 4:



The correct Figure 4:



Ref. 63 is updated and now reads: Landry, N.: nwlandry/the-simplicity-of-higher-order-networks: v0.3. <https://doi.org/10.5281/zenodo.15707834>

The incorrect Acknowledgements section reads:

N.W.L. would like to acknowledge the participants of the “Workshop on Modelling and Mining Complex Networks as Hypergraphs” at Toronto Metropolitan University and Tim LaRock for helpful conversations. N.W.L. would also like to thank Tzu-Chi Yen for lending his expertise on the biSBM inference.

The correct Acknowledgements section should read:

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Equations in the 2.1 Measures section, Algorithm 2, 3, 4, Figure 1, 2, 3, 4, Table 1, 2, 3.3 Local measures of simplicity section and Acknowledgements section have been updated above and the original article [1] has been corrected.

## Declarations

### Competing interests

The authors declare no competing interests.

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