

# ECON501 Problem Set 6

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## Problem 2

**8.E.1** We write out the normal form of this game. The pure strategies are  $AA$ ,  $AN$ ,  $NA$ ,  $NN$ , where  $A$  means attack and  $N$  means not attack, and the first letter denotes the action if the general sees strong type.

	AA	AN	NA	NN
AA	$\frac{M}{4} - \frac{s+w}{2}, \frac{M}{4} - \frac{s+w}{2}$	$\frac{M}{2} - \frac{s+w}{4}, \frac{M}{4} - \frac{s}{2}$	$\frac{3M}{4} - \frac{s+w}{4}, -\frac{w}{2}$	$M, 0$
AN	$\frac{M}{4} - \frac{s}{2}, \frac{M}{2} - \frac{s+w}{4}$	$\frac{M}{4} - \frac{s}{4}, \frac{M}{4} - \frac{s}{4}$	$\frac{M}{2} - \frac{s}{4}, \frac{M}{4} - \frac{w}{4}$	$\frac{M}{2}, 0$
NA	$-\frac{w}{2}, \frac{3M}{4} - \frac{s+w}{4}$	$\frac{M}{4} - \frac{w}{4}, \frac{M}{2} - \frac{s}{4}$	$\frac{M}{4} - \frac{w}{4}, \frac{M}{4} - \frac{w}{4}$	$\frac{M}{2}, 0$
NN	$0, M$	$0, \frac{M}{2}$	$0, \frac{M}{2}$	$0, 0$

We note  $NA$  is strictly dominated by  $AN$  since  $w > s$ , so we can restrict our attention to

	AA	AN	NN
AA	$\frac{M}{4} - \frac{s+w}{2}, \frac{M}{4} - \frac{s+w}{2}$	$\frac{M}{2} - \frac{s+w}{4}, \frac{M}{4} - \frac{s}{2}$	$M, 0$
AN	$\frac{M}{4} - \frac{s}{2}, \frac{M}{2} - \frac{s+w}{4}$	$\frac{M}{4} - \frac{s}{4}, \frac{M}{4} - \frac{s}{4}$	$\frac{M}{2}, 0$
NN	$0, M$	$0, \frac{M}{2}$	$0, 0$

If  $M < s$ , then  $NN, NN$  is the only Nash equilibrium.

Note that the best reply to  $NN$  is always  $AA$ . If  $M/2 < s$  then the best reply to  $AA$  is  $NN$ . So if  $M/2 < s$ , then  $(AA, NN)$  and  $(NN, AA)$  are Nash equilibria.

If  $w > M > s$ , then the best reply to  $AN$  is  $AN$ . Hence, if this holds (note this is not mutually exclusive with  $M/2 < s$ ), then  $(AN, AN)$  is a Nash equilibrium.

Note that the best reply to  $AA$  is never  $AA$ , so the only potential other Nash equilibria are  $(AA, AN)$  and  $(AN, AA)$ . In order for  $AA$  to be a best response to  $AN$ , we require  $M > w$ . In order for  $AN$  to be the best response to  $AA$ , we need  $M/2 > s$ . Hence, for  $M > w$  and  $M/2 > s$ ,  $(AA, AN)$  and  $(AN, AA)$  are Nash equilibria.

**8.E.3**

**9.C.2**

**9.C.3**

#### 9.C.4

### Problem 3

The strategies for the sender are  $LL, LR, RL, RR$ , denoting the actions after observing  $t_1$  and  $t_2$ . The receiver strategies are  $uu, ud, du, dd$ , the actions after observing  $L$  and  $R$ . In the subsequent parts, we denote  $\mu_L$  the receiver belief in  $t_1$  after observing  $L$  and  $\mu_R$  the receiver belief in  $t_1$  after observing  $R$ .

(a) We start with pooling equilibria. Note that in no case after observing  $t_2$  will the sender want to play  $L$ , so the pooling equilibria must have the sender send  $R$  always. In this case, the belief of the receiver on  $t_1$  versus  $t_2$  is just  $(0.5, 0.5)$ , and hence  $u$  is optimal. Hence,  $(RR, uu)$  is a PBE, supported by belief  $\mu_L \geq 0.5$  and  $\mu_R = 0.5$ . Additionally,  $(RR, du)$  is also a PBE, supported by  $\mu_L \leq 0.5$  and  $\mu_R = 0.5$ . (since the beliefs off-path can be arbitrary).

For separating equilibria, we note that the sender always plays  $R$  after realization  $t_2$ . Then after realizing  $t_1$ , the separating pure equilibria dictates the sender plays  $L$ . Hence, the only separating PBE strategy for the sender is  $LR$ . The only consistent beliefs are then  $\mu_R = 0$ ,  $\mu_L = 1$ , and the resulting strategy for the receiver is  $ud$ . So the only separating PBE is  $(LR, ud)$  with  $\mu_R = 0$ ,  $\mu_L = 1$ .

(b) We start with pooling equilibria. If player 1 plays  $LL$ , then  $\mu_L = 1/2$ , so player 2 plays  $u$  after observing  $L$ . Off-path, player 2 cannot want to play  $d$  after observing  $R$ , else  $t_1$  type would want to deviate to  $R$ . Hence, a PBE here is  $(LL, uu)$ , where  $\mu_L = 1/2$  and  $\mu_R \leq 2/3$ .

If player 1 plays  $RR$ , then  $\mu_R = 1/2$ , so player 2 plays  $u$  after observing  $R$ . But this means the  $t_1$  type would rather deviate to  $L$ , so player 1 cannot play  $RR$  in a PBE. Hence the only pooling PBE is the  $(LL, uu)$  one we described earlier.

Now, for separating PBE. If player 1 plays  $LR$ , player 2's beliefs are  $\mu_L = 1$ ,  $\mu_R = 0$ . So player 2 plays  $du$ . Note that player 1 does not want to deviate for either type, so this is a PBE:  $(LR, du)$  with  $\mu_L = 1$ ,  $\mu_R = 0$ . Now, if player 1 plays  $RL$ , player 2's beliefs are  $\mu_L = 0$ ,  $\mu_R = 1$ , so player 2's best response is  $ud$ . Once again, we observe that neither type of player 1 would like to deviate. Hence this is also a PBE:  $(RL, ud)$  with  $\mu_L = 0$ ,  $\mu_R = 1$ .

### Problem 4

First, we find the Nash equilibria. These are  $(L, R)$ ,  $(L, pL + (1 - p)R)$  and  $(M, L)$ , where  $0 \leq p \leq 0.7$ .

We find a PBE (no restrictions on beliefs off path) first for the first two equilibria. Since player 1 playing  $M, R$ , is off path, we can allow player 2's belief over the information set to be what we want. If it is  $\alpha M + (1 - \alpha)R$  for  $\alpha < 1/2$ , then  $(L, R)$  is supported. If it is  $(1/2)M + (1/2)R$ , then player 2 can take play  $pL + (1 - p)R$ ,  $p \in [0, 0.7]$  at his information set, since that is optimal given his belief and player 1 playing  $L$  is still optimal.

Finally, for  $(M, L)$ , we note the belief at player 2's information set is  $M$  with certainty. Then,  $L$  is supported for player 2 under this belief. Hence, this is also a PBE.

## Problem 5

We first look for Nash equilibria where player 1 plays a pure strategy. Suppose player 1 plays  $R$  for sure. Then player 3 will always play  $L$  for sure. Then player 2 can mix between  $L$  and  $R$ . Let player 2 play  $pL + (1-p)R$ . In order for player 1 to not want to deviate,  $(1-p) \leq 1/4$  or  $p \geq 3/4$ . Hence,  $(R, pL + (1-p)R, L)$  is a Nash equilibrium, for  $p \geq 3/4$ . Now suppose player 1 plays  $L$  for sure. Player 2 will always play  $R$ , since it always gives a strictly higher payout. Player 3 will then also play  $R$  since it gives the highest payout. Hence,  $(L, R, R)$  is also a Nash equilibrium.

Finally, we look for Nash equilibria where player 1 mixes. Suppose player 1 plays  $(1-a)L + aR$ . In order to want to mix, we know that player 1 always gets payoff 1 when playing  $R$ . So player 1 must have expected payoff at least 1 playing  $L$ . Let player 2 play  $R$  with probability  $b$ , and player 3 play  $R$  with probability  $c$ . The mixture indifference condition for player 1 implies

$$4b(1-c) + bc = 1$$

$$4b - 3bc = 1$$

$$b(4 - 3c) = 1$$

If  $b = 1$ , then we must have  $c = 1$ . If  $c = 1$ , then  $b = 1$ . So if players 2, 3 play  $(R, R)$ , then player 1 can mix. In order for it to be optimal for 2 to play  $R$ , the expected value from playing  $R$  must exceed the expected value from playing  $L$ , or

$$1 \geq (1-a) + 4a = 1 + 3a$$

This requires  $a = 0$ , but we already counted  $(L, R, R)$ . Hence, if player 1 mixes, we must have players 2 and 3 also both mix.

Now, we find the Nash equilibrium where all 3 players mix. We already have the indifference condition for player 1 is :

$$b(4 - 3c) = 1$$

The indifference condition for mixing for player 2 is:

$$1 = 4ac + a(1-c) = 3ac + a = a(3c + 1)$$

The indifference condition for player 3 is that players 1 and 2 play  $(L, R)$  and  $(R, L)$  with equal probability, or

$$a(1-b) = b(1-a)$$

$$a - ab = b - ab$$

$$a = b$$

So we get

$$a(4 - 3c) = 1$$

$$a(3c + 1) = 1$$

This implies

$$4 - 3c = 3c + 1$$

$$c = 1/2$$

$$a = 2/5$$

$$b = 2/5$$

Hence, the other Nash equilibrium is

$$\left( \frac{3}{5}L + \frac{2}{5}R, \frac{3}{5}L + \frac{2}{5}R, \frac{1}{2}L + \frac{1}{2}R \right)$$

Now, for PBE, we need to construct beliefs for players 2 and 3 that support the equilibrium. Consider  $(R, pL + (1 - p)R, L)$ , for  $p \geq 3/4$ . To support this, we have player 2's belief be that player 1 played  $R$  for sure, and player 3's belief is that players 1 and 2 played  $(R, L)$ . Alternatively, for  $p = 0$ , the equilibrium  $(R, R, L)$  is also supported for player 2 believing player 1 played  $R$  and player 3 believing players 1 and 2 played some mixture of  $L, L$  and  $R, L$  (this is fine since player 3's belief is off-path). Similarly, for  $(L, R, R)$ , we have player 2's belief is that player 1 played  $L$  and player 3's belief is that players 1 and 2 played  $(L, R)$ .

Now, for the mixed-strategy Nash equilibrium, we can make this a PBE by the following beliefs: player 2 must believe with probability  $3/5$  that player 1 played  $L$ , and player 3 has the following belief:

$$\frac{3}{7}(L, L) + \frac{2}{7}(L, R) + \frac{2}{7}(R, L)$$

## Problem 6

For pure-strategy Nash equilibria, we note that we have  $(R, C, L)$  and  $(R, S, R)$ . Note that player 1 guarantees maximum payoff by playing 1: the only way player 1 could mix is if players 2 and 3 always play  $C, L$ . Clearly, if player 3 plays  $L$ , player 2 is best off playing  $C$ . Further, in order for player 3 to want to play  $L$ , we need player 1 playing  $R$  with higher probability than  $L$ . Hence,  $(\alpha L + (1 - \alpha)R, C, L)$  is Nash for  $\alpha \leq 1/2$ . If player 1 doesn't mix, the only way for player 2 to want to mix is if player 3 also mixes, but player 3 only wants to mix when player 2 plays  $S$ . Hence, for players 1 and 2 playing  $R, S$ , player 3 can play  $\beta L + (1 - \beta)R$  with  $\beta \leq 1/2$ , and this will be a Nash eq. Hence  $(R, S, \beta L + (1 - \beta)R)$  with  $\beta \leq 1/2$  are the other Nash equilibria.

Now, for PBE, we need beliefs of players 2 and 3 to support the strategies. Let  $\mu_2$  be the belief of player 2 in the left side, and  $\mu_3$  be the belief of player 3 in the left side. For  $(R, C, L)$ , this is supported by  $\mu_2 = 0, \mu_3 = 0$ . We can also support  $(R, S, R)$  with  $\mu_2 = 0, \mu_3 \geq 1/2$  (in order for player 3 to play  $R$ ). Further,  $(\alpha L + (1 - \alpha)R, C, L)$  for  $\alpha \leq 1/2$  is supported for  $\mu_2 = \alpha, \mu_3 = \alpha$ . Lastly,  $(R, S, \beta L + (1 - \beta)R)$  with  $\beta \leq 1/2$  is supported by beliefs  $\mu_2 = 0, \mu_3 = 1/2$ .