Neighbor Management Policy for 6LoWPAN

Signaling and Policy guidelines

https://tools.ietf.org/html/draft-jadhav-lwig-nbr-mgmt-policy-00

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Why Neighbor Management?

Challenges

- Unknown network size, unknown Node density
- · Constrained networks with limited neighbor cache
 - Density is higher than neighbor cache size

Expectation of neighbor management

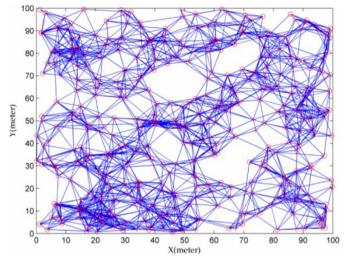
- Improved network stability, reduced churn in routing adjacencies
- Once the neighbor is accepted, the associated resources are guaranteed

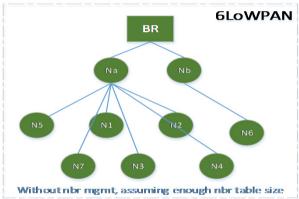
Trivial Neighbor Management policies

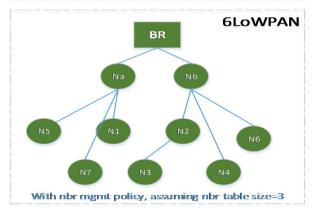
- Evict LRU entry for new insertion when table is full
- First come first serve

Protocol agnostic policy

 Even though the draft references RPL and PANA extensively, the proposed policy is routing protocol and key management protocol agnostic

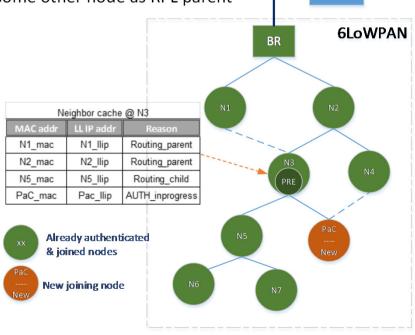






Holistic approach towards neighbor management

- An example security-enabled 6LoWPAN/RPL network
 - Key management protocols before RPL network formation
 - PANA as example, used by Wi-SUN
 - Draft explains neighbor management differences with respect to RPL storing as well as non-storing mode of operations
- Cases where neighbor table update happens
 - Relay based signaling during authentication
 - PRE selection by PaC, usually involves discovery messaging. No std procedure for PRE discovery.
 - PRE needs to add PaC as nbr since it will act as relay till the auth process completes
 - Note that post-auth-success, PaC may choose some other node as RPL parent
 - RPL's parent selection using DIO messaging
 - RPL's routing child node
 - Implicit vs explicit signaling
 - Implicit signaling in Storing MOP for NCE
 - Explicit signaling required in Non-Storing MOP



PRE = PANA Relay Element PaC = PANA Client MOP = Mode of Operation

Neighbor Management Operations

Insertion

- Problem with simple logic (If table space is available: insert)
 - RPL's DIO storm in dense network may overwhelm neighbor cache
 - Same parent chosen by all the nodes resulting in nbr cache containing only routing child entries
 - Similarly PRE discovery may result in the same PRE been made use of by several PaCs.

Eviction

- Issues with eviction
 - · An routing child eviction may have ripple effect on all grand-childs
 - Similarly if a PaC NCE is added on PRE, then early eviction may result in neighbor churn.
- Evicting non-preferred parent NCE is usually possible without much implications
 - For e.g. on receiving DAO, one can evict a "low-priority" parent entry from neighbor cache

Reinforcement

- NCEs needs to be reinforced
 - Reinforcement can be done by passive/active hearing or by explicit probing. *Draft does not define how to do this.*
- Reinforcement allows the link quality estimation to be updated, eventually helping in eviction decision

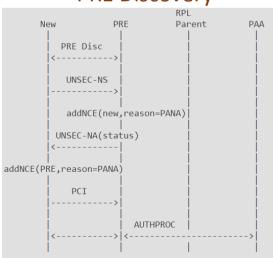
Clearing unused Neighbor table entries

- Important that unused NCEs be reclaimed soon
- For storing MOP, route invalidation is important since routing entries are mapped to NCEs
- For Non-storing MOP,
 - since there is no route invalidation procedure, the child node needs to deregister using NS(lifetime=0)
- PRE neighbors
 - After authentication is successful, the PRE auth entries can be removed
 - However there is no way of explicit identification of auth finish
 - Usually reachability timeout will remove such entries. For neighbors added for authentication, the reachability timer can be reduced to a lower value.

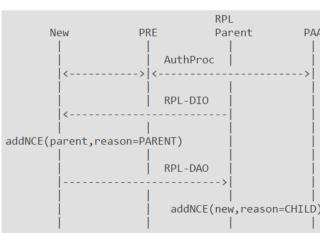
Signaling recommendation for Neighbor management

- As far as possible use implicit mechanism for neighbor entry addition
 - Use DIO/DAO messaging to populate NCEs in case of storing MOP
 - In case of Non-storing MOP, DAO flows end to end, thus explicit NDP signaling in the form of NS/NA is required.
- Implicit mechanism works only if there is a way to send negative status if NCE addition fails
 - For e.g. in case of PANA, there is no way (currently) for PRE to respond back with negative status
 - Thus explicit NDP signaling is involved to populate neighbor cache entries which can also signal failure if needed.

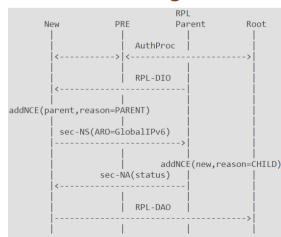
PRE Discovery



Storing MOP

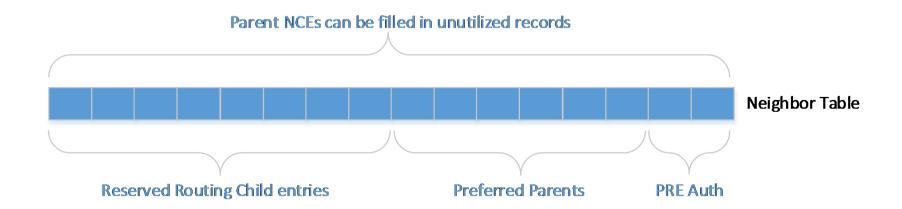


Non-Storing MOP



Proposed guidance for reservation based policy

- Basic principles
 - Reservation of routing direct child entries
 - Reservation of relay element entries
 - Parent node's entries can be inserted at will and can occupy reserved entries
 - Because parent entries could be evicted if necessary, unlike routing direct childs and relay element entries
 - Insertion reason (RPL_parent, RPL_child, Other) is attached with every NCE
- Graceful rejection of DAO/PANA messages
 - NACK for rejecting DAO
 - Negative status in NDP NA response



Issues with implicit/reactive policy

Limitations of reactive policy

- Scenario: A parent whose nbr cache is full sends a DIO ...
- A child node may still select this parent node since DIO does not signal NCE metric
- Thus there would be an additional signaling to reject this parent node
- Worse, in the future, the child node may again select the same parent based on new DIO from the parent node.
- The same problem applies while PRE discovery...
- Guidance:
 - A proactive approach to signal NCE metric
 - For example, metric containers can be shared between RPL and PRE discovery messaging
 - Can RPL metrics containers (RFC 6551) be reused by another protocol?

Discussions

- WG Adoption
 - As a general protocol agnostic guidance for nbr mgmt ...
- Contiki implementation ongoing...

Thank You