

Numerical Analysis: Midterm (50 marks)

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Fall 11111100101

You are not required to complete every question. Although the total marks on the exam sum to 55, your midterm grade will be calculated out of 50.

✿✿ Question 1 (Floating point arithmetic, 10 marks). True or false?

1. Let $(\bullet)_2$ denote binary representation. Then $(0.1101)_2 + (0.0011)_2 = (1.0)_2$.
2. Let ε_{64} denote the machine epsilon for the `Float64` format, i.e. `eps(Float64)`. Then the number 2 is representable exactly in this format, and the next representable number is $2 + \varepsilon_{64}$.
3. It holds that $(10000)_2 \times (0.1010)_2 = (1010)_2$.
4. In Julia, `Float64(.6) == Float32(.6)` evaluates to `true`.
5. The spacing (in absolute value) between successive single-precision (`Float32`) floating point numbers is constant.
6. Infinitely many distinct real numbers can be represented exactly in the `Float64` format, but only finitely many can be represented exactly in the `Float32` format.
7. It holds that $(0.\overline{101})_2 = \frac{5}{7}$.
8. Machine addition $\hat{+}$ is an operation that is *associative* but not *commutative*.
9. The machine epsilon is the smallest number of the form 2^{-n} with $n \in \mathbb{N}$ that can be represented exactly in a floating point format.
10. In Julia, the expression `1 + eps()/3 == 1 + eps()` evaluates to `true`.
11. **Bonus.** In Julia, the expression `exp(eps()/2) == 1 + eps()` evaluates to `true`.

Explain briefly:

12. **Bonus.** In Julia, the expression `cos(eps()) == 1` evaluates to `true`.

Explain briefly:

Question 2 (Interpolation and approximation, **10 marks**). Are the following statements true or false? Prove or disprove. Recall that \mathcal{P}_d denotes the set of polynomials of degree at most d .

1. Assume that $x_0 < x_1 < x_2 < x_3$ and y_0, y_1, y_2, y_3 are given real numbers. Then there exists a polynomial $p \in \mathcal{P}_3$ such that $p(x_i) = y_i$ for all $i \in \{0, 1, 2, 3\}$.

Justification:

2. Assume that $x_0 < x_1 < x_2$ and y_0, y_1, y_2 are given real numbers. Then, there can exist *at most one* polynomial $p \in \mathcal{P}_2$ such that $p(x_i) = y_i$ for all $i \in \{0, 1, 2\}$.

Justification:

3. Let $p \in \mathcal{P}_d$ be a polynomial of degree $d > 0$, and let $q: \mathbf{R} \rightarrow \mathbf{R}$ be given by $q(x) = p(x+1) - p(x)$. Then it holds that $q \in \mathcal{P}_{d-1}$.

Justification:

4. For $n \in \mathbf{N}$, let $x_i^n = i/n$ for $i = 0, 1, \dots, n$. Assume that $u: \mathbf{R} \rightarrow \mathbf{R}$ is the smooth function given by $u(x) = \sin(3x) + x^3$, and let $p_n \in \mathcal{P}_n$ denote the polynomial interpolation of u at the points $x_0^n, x_1^n, \dots, x_n^n$. Then it holds that

$$\lim_{n \rightarrow \infty} \left(\max_{x \in [0,1]} |u(x) - p_n(x)| \right) = 0.$$

Justification:

✿✿✿ **Question 3** (Interpolation, open-ended question, **5 marks**). In polynomial interpolation, the error depends both on the function being interpolated and the choice of interpolation nodes. Consider two families of nodes on the interval $[-1, 1]$:

- Equally spaced nodes,
- Chebyshev nodes.

Discuss qualitatively (and illustrate with examples or plots if you wish) how the interpolation error behaves as the degree n increases in each case. Why does one choice of nodes perform better for large n , and what is the mathematical motivation for using Chebyshev nodes? *You may refer to Runge's phenomenon, but go beyond merely stating it.*

❑ **Implementation exercise 1** (Interpolation, 5 marks). Write Julia code that computes and plots the interpolating polynomial $p \in \mathcal{P}_3$ through the following points: $(0, 0)$, $(1, 4)$, $(2, 15)$, $(3, 40)$. The plot should display both the interpolation points and the graph of the interpolating polynomial over an appropriate range. Do not use any other library than the ones already imported.

```
using LinearAlgebra
using Plots
# Write your code here
```

Question 4 (Numerical integration, 5 marks). Are the following statements true or false? Justify briefly.

1. The degree of precision of the following quadrature rule is 2:

$$\int_{-1}^1 u(x) dx \approx 2u(0).$$

Justification:

2. The degree of precision of the following rule is equal to 3:

$$\int_{-1}^1 u(x) dx \approx u\left(-\frac{1}{3}\right) + u\left(\frac{1}{3}\right).$$

Justification:

3. For any natural number $N > 0$, there exists a quadrature rule with a degree of precision equal to $2N - 1$ of the form

$$\int_{-1}^1 u(x) dx \approx \sum_{n=1}^N w_n u(x_n).$$

Justification:

4. Let $x_i^N = i/N$ and consider the following approximation of $\int_0^1 u(x) dx$:

$$\widehat{I}_N = \frac{1}{2N} \left(u(x_0^N) + 2u(x_1^N) + 2u(x_2^N) + \dots + 2u(x_{N-2}^N) + 2u(x_{N-1}^N) + u(x_N^N) \right). \quad (1)$$

Suppose first that u is the Runge function, given by $u(x) = (1 + 25x^2)^{-1}$. Then \widehat{I}_N diverges in the limit $N \rightarrow \infty$.

Justification:

5. Let $u(x) = \cos(3x)$ and let \widehat{I}_N be as in (1). Then it holds that

$$\lim_{N \rightarrow +\infty} \left(\left| \widehat{I}_N - \int_0^1 u(x) dx \right| \right) = 0.$$

Justification:

6. (**Bonus.**) Fix $u(x) = 2x - 1$ and let \widehat{I}_N be as in (1). Then $\widehat{I}_N = 0$ for all $N \geq 2$.

Justification:

Question 5 (Gaussian–Hermite numerical integration, **10 marks**). The Gauss–Hermite quadrature formula with n nodes is an approximation of the form

$$I(u) := \int_{-\infty}^{\infty} u(x) e^{-x^2} dx \approx \sum_{i=1}^n w_i u(x_i) =: \hat{I}_n(u),$$

which is exact when u is a polynomial of degree $\leq 2n - 1$. Note that the nodes are numbered $1, \dots, n$. For this question, we take for granted that, for integers $i \geq 0$, it holds that

$$\int_{-\infty}^{\infty} x^i e^{-x^2} dx = \begin{cases} 0, & \text{if } i \text{ is odd,} \\ (i-1)!! \sqrt{\frac{\pi}{2^i}}, & \text{if } i \text{ is even,} \end{cases}$$

where $(i-1)!! := 1 \times 3 \times 5 \times \dots \times (i-1)$. In particular, with all the integrals being over $(-\infty, \infty)$, the following special cases may be useful in your computations:

$$\int e^{-x^2} dx = \sqrt{\pi}, \quad \int x^2 e^{-x^2} dx = \frac{\sqrt{\pi}}{2}, \quad \int x^4 e^{-x^2} dx = \frac{3}{4} \sqrt{\pi}, \quad \int x^6 e^{-x^2} dx = \frac{15}{8} \sqrt{\pi}.$$

1. **(5 marks)** Find the nodes and weights of the Gauss–Hermite rule with $n = 3$ nodes. By symmetry, we expect nodes of the form $(-z, 0, z)$ and weights (w_1, w_2, w_1) , which reduces the number of unknowns to three.

Your answer:

2. (5 marks) Let $\{H_0, H_1, \dots\}$ denote orthogonal polynomials for the inner product

$$\langle f, g \rangle := \int_{-\infty}^{\infty} f(x)g(x) e^{-x^2} dx$$

which, in addition, satisfy the following two conditions:

- For all $i \in \mathbf{N}$, the polynomial H_i is of degree i .
- The leading coefficient of H_i , which multiplies x^i , is equal to 1.

Calculate H_0 , H_1 , H_2 and H_3 . What is the relationship between H_3 and the quadrature rule found in the first item?

Your answer:

3. (Bonus, **2 marks**) Calculate H_4 and, using this result, deduce the nodes and weights of the Gauss–Hermite quadrature with 4 points.

Your answer:

❑ **Implementation exercise 2** (Numerical integration, 10 marks). The midpoint quadrature rule reads

$$\int_{-1}^1 u(x) dx \approx 2u(0).$$

- (3 marks) Write a function `midpoint(u, a, b)` that returns, using this quadrature rule, an approximation of the integral

$$\int_a^b u(x) dx. \quad (2)$$

```
function midpoint(u, a, b)
    # Write your code here
end
```

- (4 marks) Write a function `composite_midpoint(u, a, b, N)` that returns an approximation of the integral (2), this time using a composite version of the midpoint rule. More precisely, the approximation should be obtained by partitioning the integration interval $[a, b]$ into N cells, and applying the midpoint rule within each cell.

```
function composite_midpoint(u, a, b)
    # Write your code here
```

```
end
```

- (3 marks) Take $u(x) = \cos(x)$, $a = -1$ and $b = 1$. In this setting, plot the evolution of the error for N varying from 1 to 1000.

```
using Plots
# Write your code here
```