

1. The computational cost of the forward substitution algorithm, for a matrix of size $n \times n$, scales as $\mathcal{O}(n)$.
2. The computational cost to calculate the LU decomposition of an invertible matrix scales as $\mathcal{O}(n^2)$.
3. Suppose that A is a symmetric, positive definite matrix. Calculating the Cholesky decomposition of A requires fewer floating point operations than calculating the LU decomposition of A .
4. The condition number $\kappa_2(A)$ of a square invertible matrix A is always strictly larger than 1.
5. The spectral radius of a matrix square A is zero if and only if A is zero.
6. If the condition number of a matrix A is very large ($\gg 1$), then it is possible that rounding errors arising from floating point arithmetic will have a large impact on the accuracy of the numerical solution to the linear system $Ax = b$ (calculated by LU decomposition followed by forward and backward substitution, for example).
7. Suppose that $A \in \mathbf{R}^{n \times n}$ is symmetric and positive definite, and let $b \in \mathbf{R}^n$. Consider the following iterative method for solving the linear system $Ax = b$:

$$x^{(k+1)} = x^{(k)} + \omega (b - Ax^{(k)}). \quad (1)$$

This iteration converges to the exact solution of the linear system for all $\omega \in \mathbf{R}$.

8. The convergence speed of the iteration (1), for the optimal value of ω , is independent of the condition number $\kappa_2(A)$.
9. If A is symmetric positive definite, there always exists $\omega \in \mathbf{R}$ such that the iteration (1) converges.
10. If A is a nonzero matrix, then its norm $\|A\|_2$ is strictly positive.
11. Suppose that we want to solve the linear systems $Ax = b$ and $Ax = c$ by a direct method. To this end, the LU decomposition of A can be calculated only once.
12. Assuming that the matrix A and the vector b are already defined, write (on paper) Julia code implementing 100 iterations of (1).