

Recitation 12

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Subtyping

- A type defines a set of objects. A subtype ($S <: T$) defines a set of objects that have at least the methods and fields of T .
- A subtype is a subset of the set defined by its parent type.
 - e.g. $\text{Nat} <: \text{Int}$
- If the values of S are a subset of T , then an expression expecting T values can also receive S values.
 - e.g. $x+2$, $f(x)$, $\text{if true then } x \text{ else } y$

Function subtyping

- Covariant subtyping: If $B <: A$, then $C \rightarrow B <: C \rightarrow A$.
 - Example: let B be Nat and A be Int , then $\text{Nat} <: \text{Int}$
 - Let $f1: \text{String} \rightarrow \text{Nat}$, $f2: \text{String} \rightarrow \text{Int}$ each interpret the string as a number
 - In a context where I want to interpret a string into a Int , I can use either $f1$ or $f2$
 - In a context where I want to interpret a string into an Nat , I can only use $f1$
 - $f1$ can be used in all contexts where $f2$ is expected
 - Therefore $f1 <: f2$

Function subtyping

- Contravariant subtyping: If $B <: A$, then $B \rightarrow C >: A \rightarrow C$.
 - Example: let B be Nat and A be Int, then $\text{Nat} <: \text{Int}$
 - Let $f1: \text{Nat} \rightarrow \text{String}$, $f2: \text{Int} \rightarrow \text{String}$ each convert the argument as a string
 - In a context where I want to convert a Nat, I can use either f1 or f2
 - In a context where I want to convert an Int, I can only use f2
 - f2 can be used in all contexts where f1 is expected
 - Therefore $f2 <: f1$

Function subtyping

- “ \rightarrow type constructor is contravariant in argument type”
 - Given some type S , consider a type constructor C that takes type $T1$ and returns the type $T1 \rightarrow S$
 - If $T1 <: T2$, then $T1 \rightarrow S >: T2 \rightarrow S$
 - $C(T1) >: C(T2)$, analogous to `Queue[Duck] >: Queue[Bird]`
 - Type parameter for C is contravariant

Function subtyping

- “ \rightarrow type constructor covariant in return type”
 - Given some type S , consider a type constructor C that takes type $T1$ and returns the type $S \rightarrow T1$
 - If $T1 <: T2$, then $S \rightarrow T1 <: S \rightarrow T2$
 - $C(T1) <: C(T2)$, analogous to `Queue[Duck] <: Queue[Bird]`
 - Type parameter for C is covariant

Function subtyping

- Arguments \sim = contravariant positions, return values \sim = covariant positions
- Function type operator is contravariant in argument type and covariant in return type

Example

```
class CoVar[+T](x: T):  
  def method1: T = x  
  def method2(y: T): List[T] = List(x,y)  
  
class ContraVar[-T](x: T):  
  def method1: T = x  
  def method2(y: T): List[T] = List(x,y)
```

- At least one method in each class violates the variance annotation of the class' type parameter. Which method?

Example

```
var c := CoVar[Car]
var sc := CoVar[SportsCar]
// SportsCar <: Car

// Because T is covariant
// CoVar[SportsCar] <: CoVar[Car]
// sc must support all methods c supports

c.method2(porsche)
sc.method2(porsche)
c.method2(limo)
sc.method2(limo) // Not allowed!
```

```
class CoVar[+T](x: T):
  def method1: T = x
  def method2(y: T): List[T] = List(x,y)

class ContraVar[-T](x: T):
  def method1: T = x
  def method2(y: T): List[T] = List(x,y)
```

Example

```
var c := ContraVar[Car](limo)
var sc := ContraVar[SportsCar](porsche)
// SportsCar <: Car

// Because T is contravariant
// ContraVar[SportsCar] >: ContraVar[Car]
// c must support all methods sc supports

def race_sportscar(x: sportscar) = {...}
race_sportscar(sc.method1)
race_sportscar(c.method1) // Not allowed!
```

```
class CoVar[+T](x: T):
  def method1: T = x
  def method2(y: T): List[T] = List(x,y)

class ContraVar[-T](x: T):
  def method1: T = x
  def method2(y: T): List[T] = List(x,y)
```