## Sample Solution for Homework 8

## Problem 1 Type Checking and Type Inference (16 Points)

(a) Use the type inference rules of the type system to determine for each of the following expressions whether the expression is well-typed. If an expression is not well-typed, explain why. If it is well-typed, give the inferred type. You do not need to show the individual inference steps.

The program is well-typed. The inferred type is **Num**.

(ii)

```
const x = true;
const y = x + 1;
x * y
```

The program is not well-typed. The + operator on line 2 expects two expressions of type **Num** as arguments. However, the inferred type of x is **Bool**.

(iii)

```
const f = function f(x: Num): Num => Num (
          (y: Num) => (x === y ? 1 : y * f(x)(y + 1))
          );
f(3)
```

The program is well-typed. The inferred type is  $Num \Rightarrow Num$ .

(iv)

```
const f = function f(x: Num): Num => Num (
      (y: Num) => (x === y ? 1 : y * f(x)(y + 1))
      );
f(3) === f(3)
```

The program is not well-typed. The problem is that line 3 checks equality between the expression f(3) and itself. However, the inferred type for f(3) is  $Num \Rightarrow Num$  and the type system does not allow expressions of function types to be compared using the equality operator.

```
(v)

1 const f = function f(x: Num) (
2    (y: Num) => (x === y ? 1 : y * f(x)(y + 1))
3    );
4    f(3)(1)
```

The program is not well-typed. The problem is that line 1 declares a recursive function f. However, the function expression is missing the annotation of the return type. The type system does not support recursive function expressions without return type annotations.

(b) For each of the following programs, find concrete types for the missing parameter type annotations  $\tau_1$ ,  $\tau_2$ , and  $\tau_3$  such that the given program is well-typed according to the typing rules. If no such types exist, explain why. If you can find type annotations that make the program well-typed, what is the inferred type of f for your annotations? Are your chosen types the only annotations that work? If not, give at least one other choice of annotations that also makes the program well-typed.

```
(i)

1 const f = (x: \tau_1) => (y: \tau_2) => (z: \tau_3) => x(y(z));

2 const g = (x: Num) => x + 1;

3 const h = (x: Num) => x * 2;

4 f(h)(g)(3)
```

The program is well-typed for the following type annotations:

$$\begin{array}{ll} \tau_1 &= \operatorname{Num} \Rightarrow \operatorname{Num} \\ \tau_2 &= \operatorname{Num} \Rightarrow \operatorname{Num} \\ \tau_3 &= \operatorname{Num} \end{array}$$

These are the only possible type annotations for which the program is well-typed. The inferred type of f is

$$\tau_1 \Rightarrow \tau_2 \Rightarrow \tau_3 \Rightarrow \text{Num}$$

(ii)

1 const f = (x:  $\tau_1$ ) => (y:  $\tau_2$ ) => (z:  $\tau_3$ ) => x(y(z));

2 const g = (x:  $\tau_3$ ) => x;

3 f(g)(g)

The program is well-typed for the following type annotations:

$$\begin{array}{ll} \tau_1 &= \operatorname{Num} \Rightarrow \operatorname{Num} \\ \tau_2 &= \operatorname{Num} \Rightarrow \operatorname{Num} \\ \tau_3 &= \operatorname{Num} \end{array}$$

Another possible annotation is:

$$egin{array}{ll} au_1 &= \mathtt{Bool} \Rightarrow \mathtt{Bool} \ au_2 &= \mathtt{Bool} \Rightarrow \mathtt{Bool} \ au_3 &= \mathtt{Bool} \end{array}$$

In general, any annotation that satisfies the following equalities would work:

$$\tau_1 = \tau_2 = (\tau_3 \Rightarrow \tau_3)$$

The inferred type of f is as in (i) for the respective values of  $\tau_1$ ,  $\tau_2$ , and  $\tau_3$ .

(iii)

```
const f = (x: \tau_1) => (y: \tau_2) => x(y);
const g = (x: Bool) => x ? 1 : 0;
const h = (x: Num) => x + x;
f(g)(true) + f(h)(1)
```

There exists no type annotation for which this program is well-typed. The problem is that calling f on g and h forces  $\tau_1 = \text{Num} \Rightarrow \text{Num}$  and  $\tau_1 = \text{Bool} \Rightarrow \text{Num}$ . We thus must satisfy Bool = Num which is impossible. To make the program well-typed we need a more expressive type system that supports type parameterization.