

#### IFDS Framework

CS-UH 3260 Static Program Analysis

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#### Previously

Context sensitivity

Call strings

Functional approach

## Interprocedural Finite Distributive Subset Problems

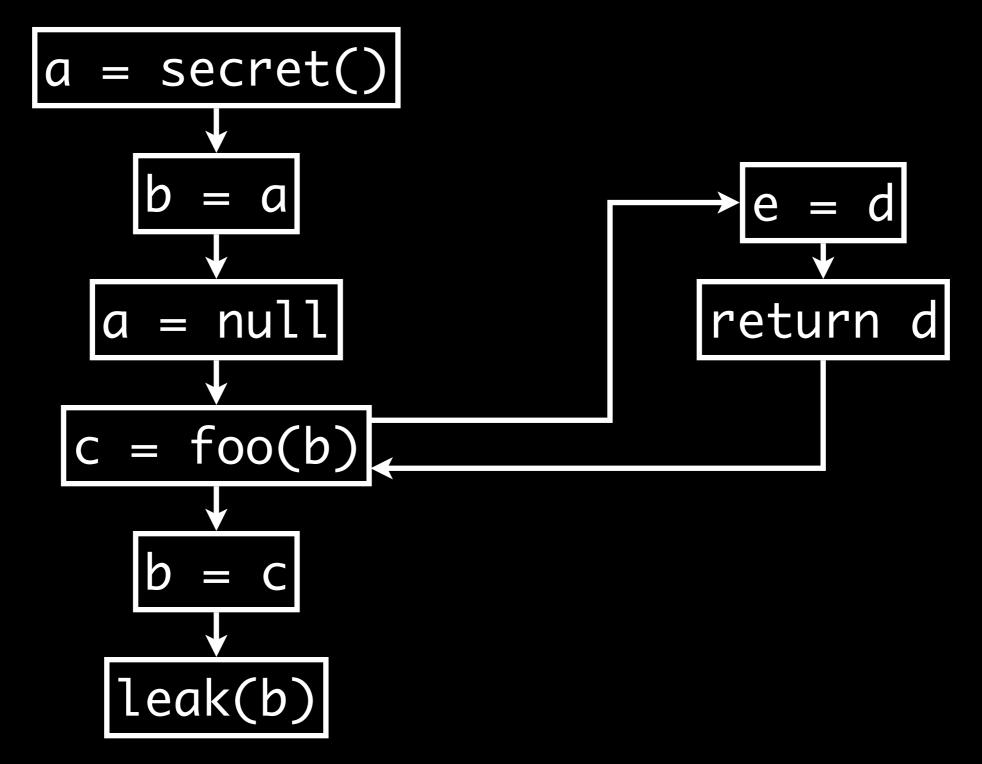
Thomas Reps, Susan Horwitz, and Mooly Sagiv. Precise Interprocedural Dataflow Analysis via Graph Reachability. ACM SIGPLAN-SIGACT Symposium on Principles of Programming Languages (POPL '95), pages 49–61.

# Interprocedural Finite Distributive Subset Problems

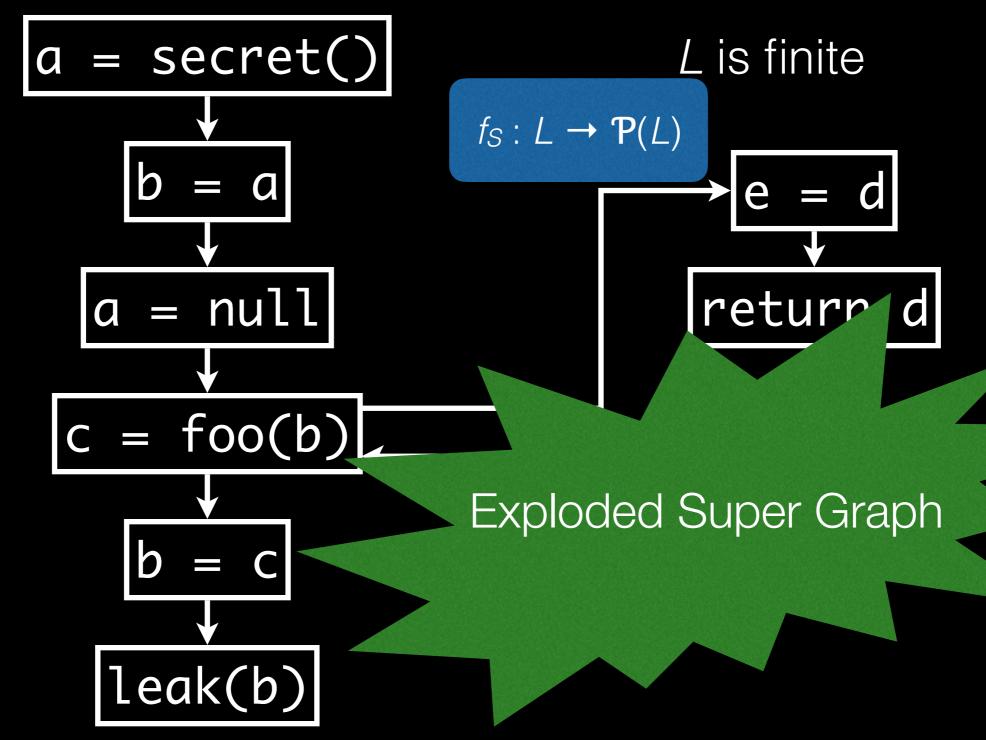
#### Interprocedural

```
main() {
    a = secret();
    b = a;
    a = null;
    c = foo(b);
    b = c;
    leak(b);
}
```

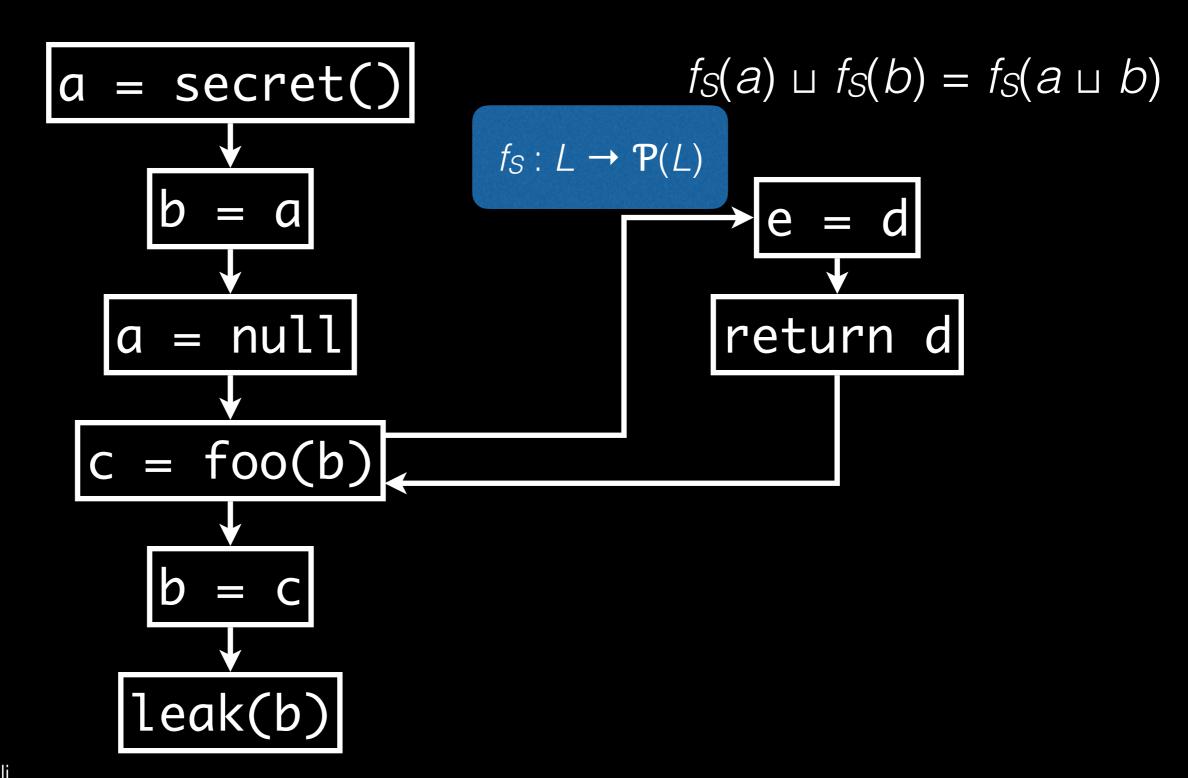
#### Interprocedural



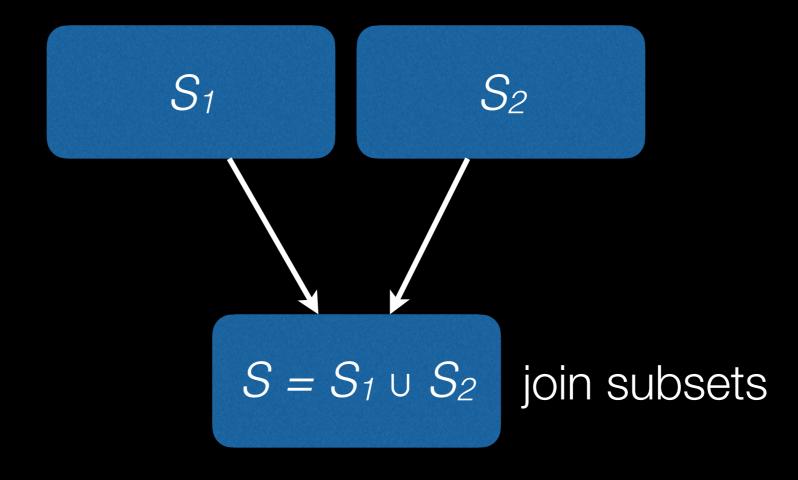
#### **F**inite



#### **D**istributive



#### Subset



alias taint typestate

IFDS

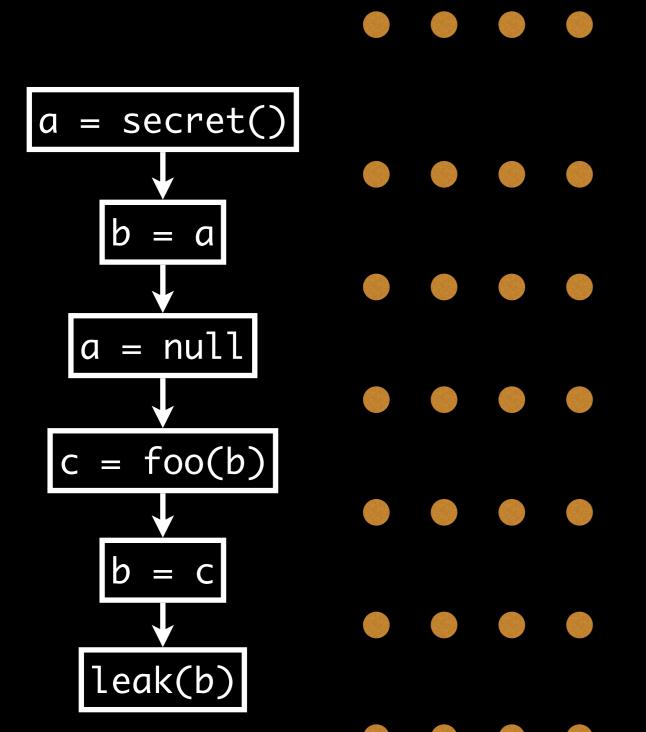
constant propagation

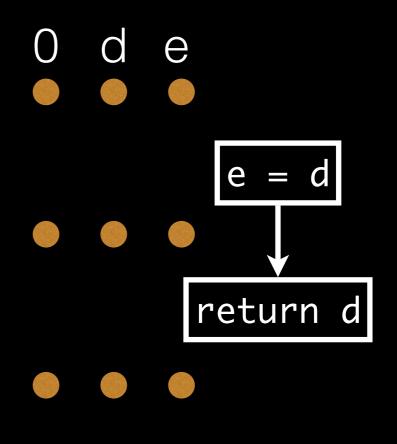
reaching definitions

uninitialized variables

### Graph Reachability

#### Graph Reachability

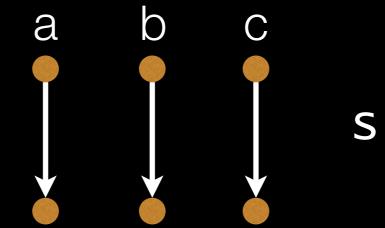




- each node is a fact
- fact holds at a statement => node is reachable in ESG

## Examples of Flow Functions

#### Identity Flow Functions

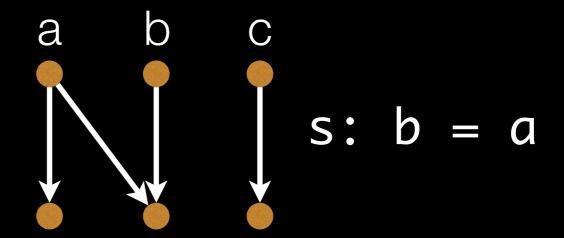


out(s) = facts after s

every fact is reachable iff it was previously reachable

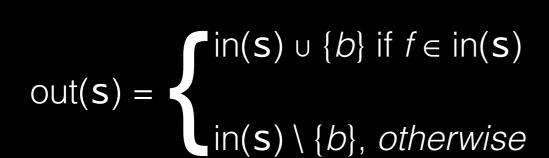
#### Flow Functions

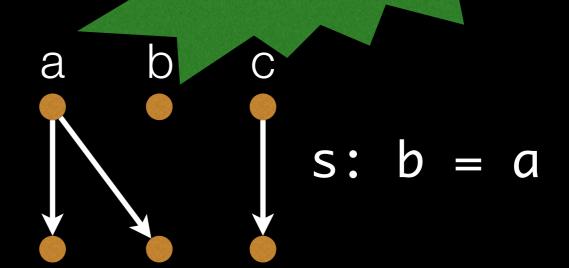
out(s) = 
$$\begin{cases} in(s) \cup \{b\} & \text{if } f \in in(s) \\ in(s), & \text{otherwise} \end{cases}$$



every fact is reachable iff it was previously reachable, and b is also reachable if a was reachable

#### Flow Functions

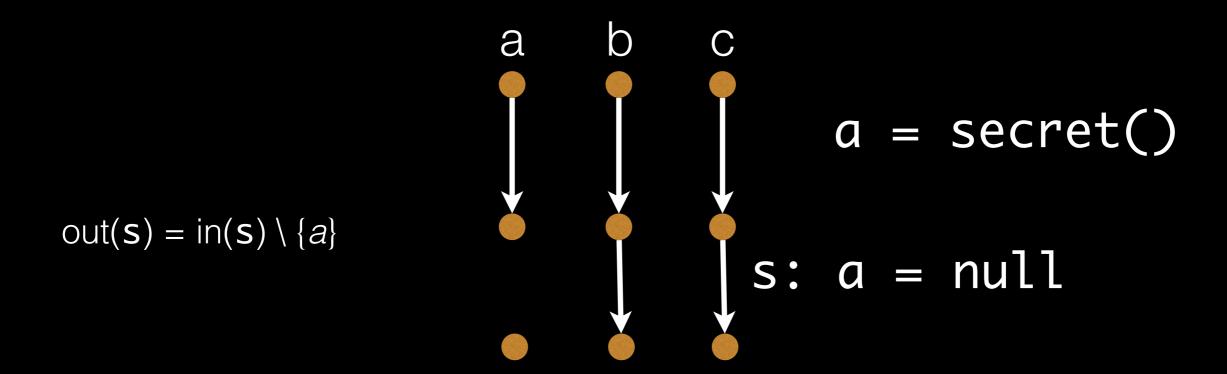




Taint Analysis

every fact except *b* is reachable iff it was previously reachable; *b* is reachable if *a* was reachable

#### "Killing" Facts

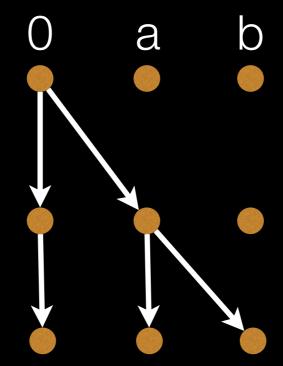


a is not reachable, even if it was before

#### Generating Facts

$$\mathsf{out}(\mathbf{s}) = \mathsf{in}(\mathbf{s}) \cup \{a\}$$

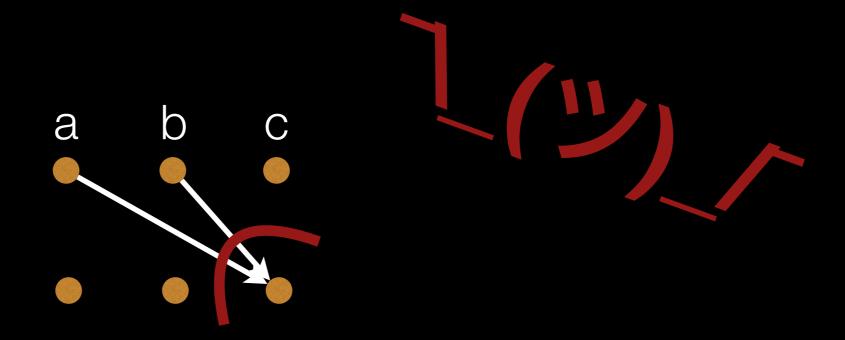
$$out(s) = in(s) \cup \{b\}$$



$$b = a$$

0 is the tautological fact => always reachable

### Non-Distributive Flow Functions



e.g., full constant propagation c = a + b

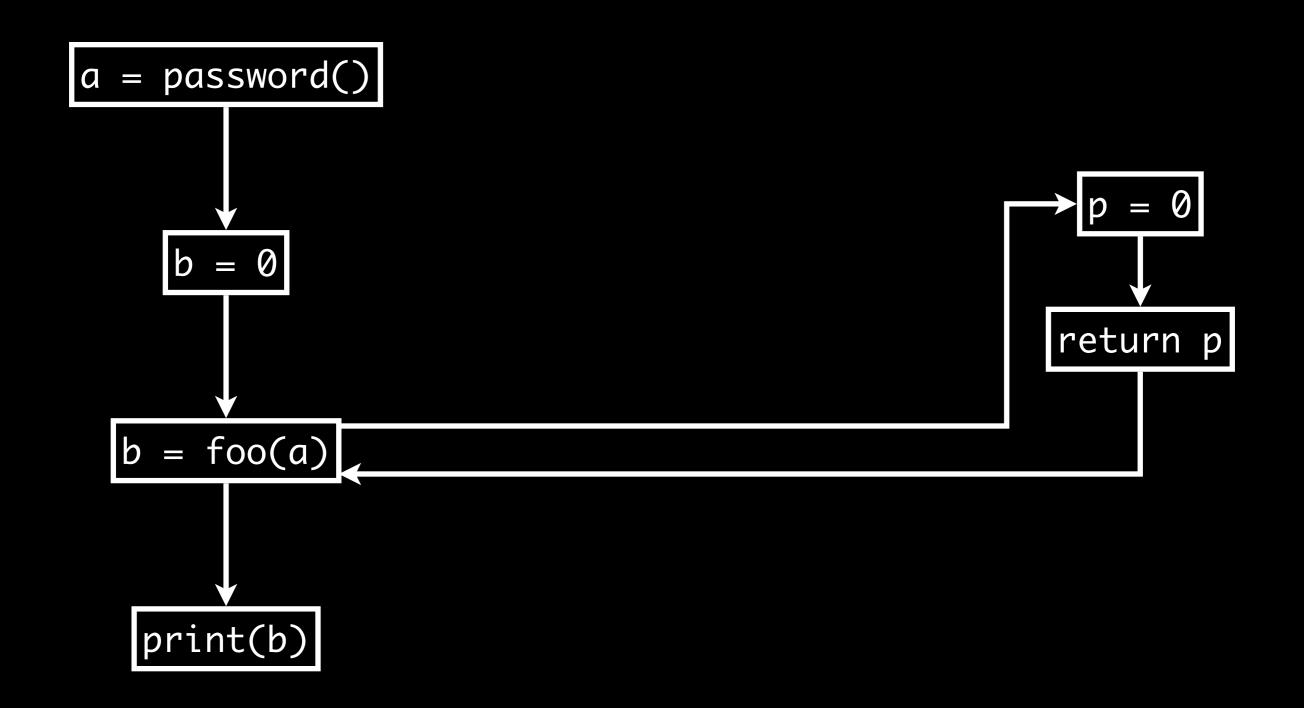
```
void main() {
   int a = password();
   int b = 0;

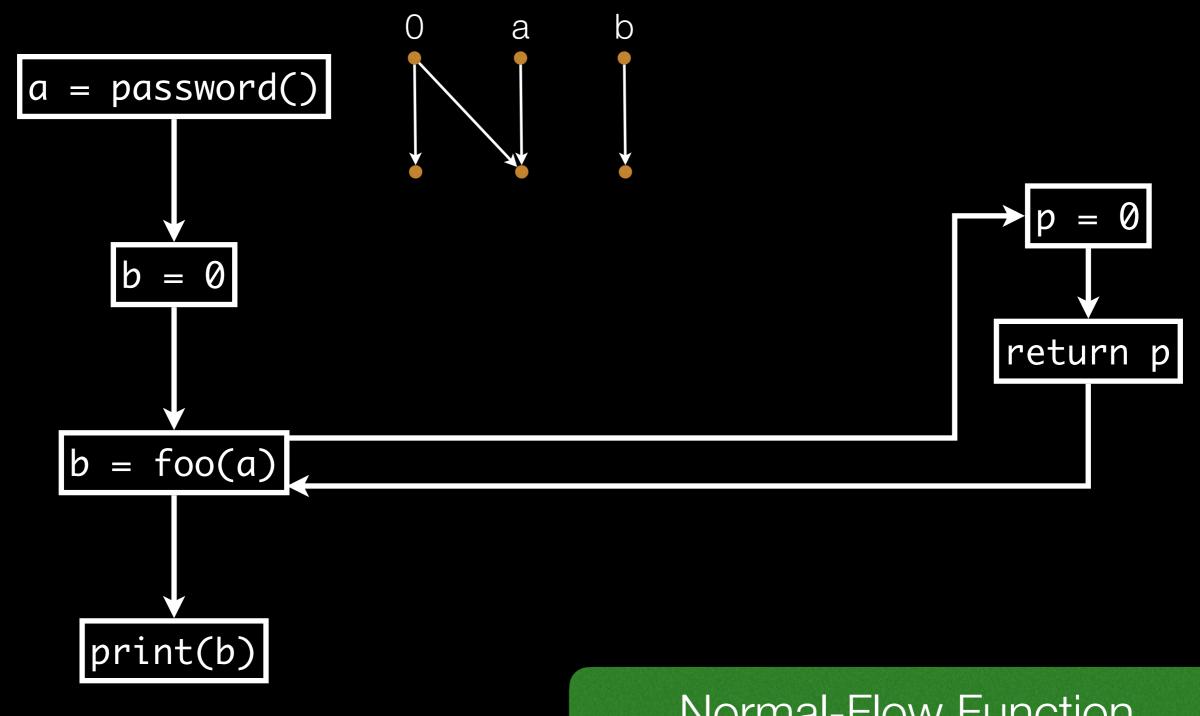
b = foo(a);

print(b);
}
```

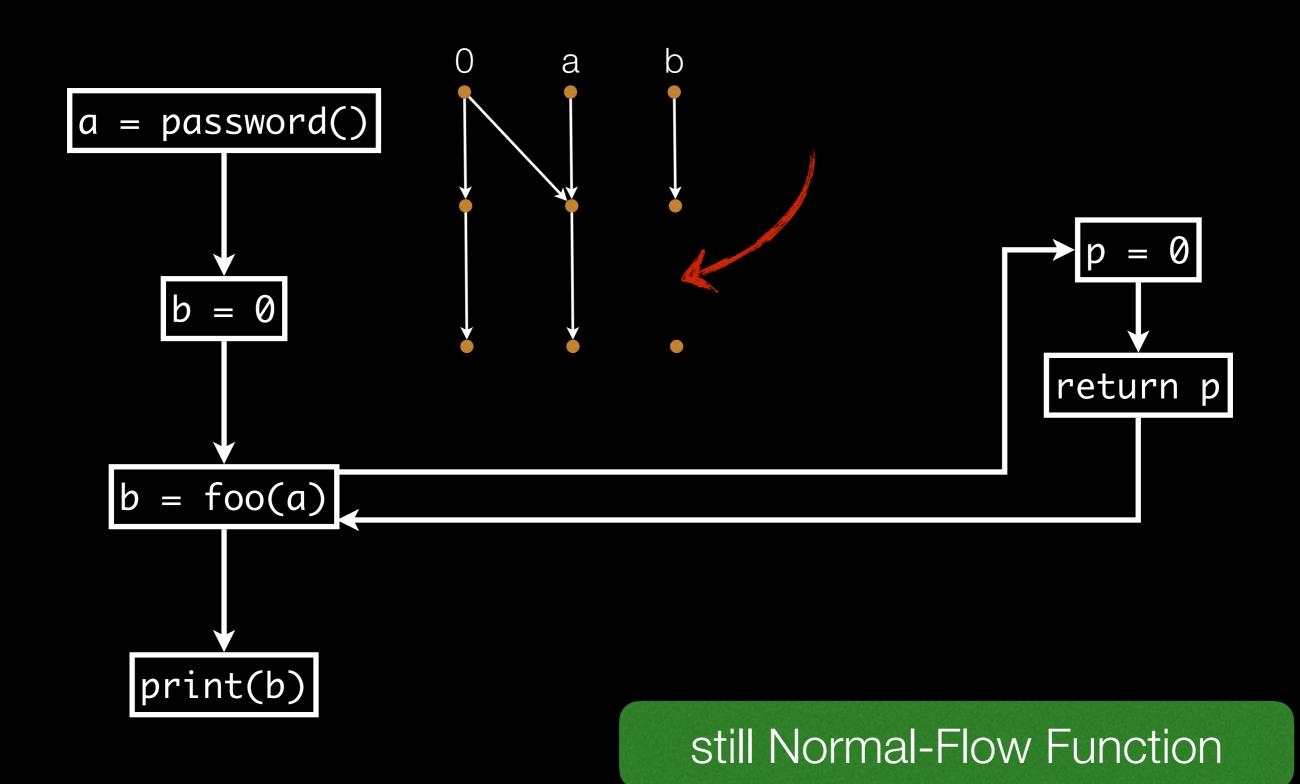
```
int foo(int p) {
   p = 0;

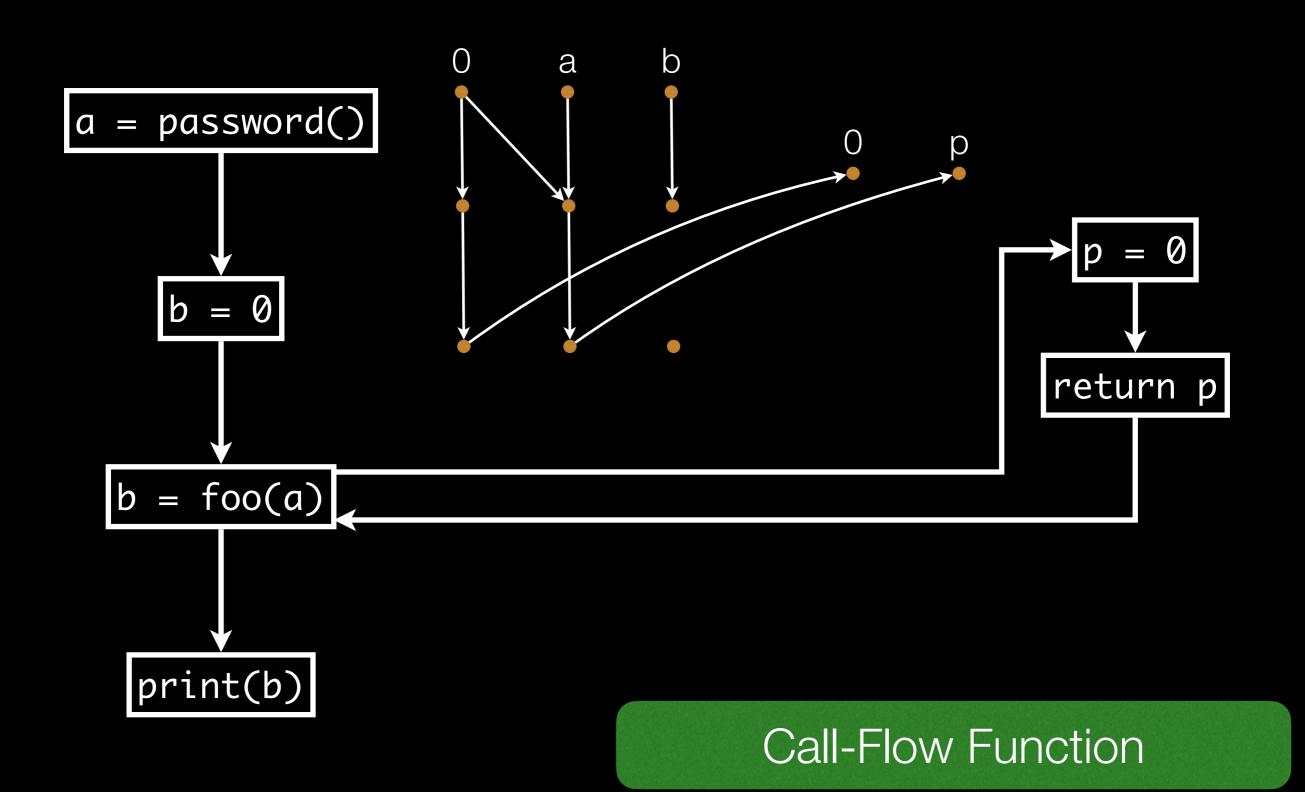
   return p;
}
```

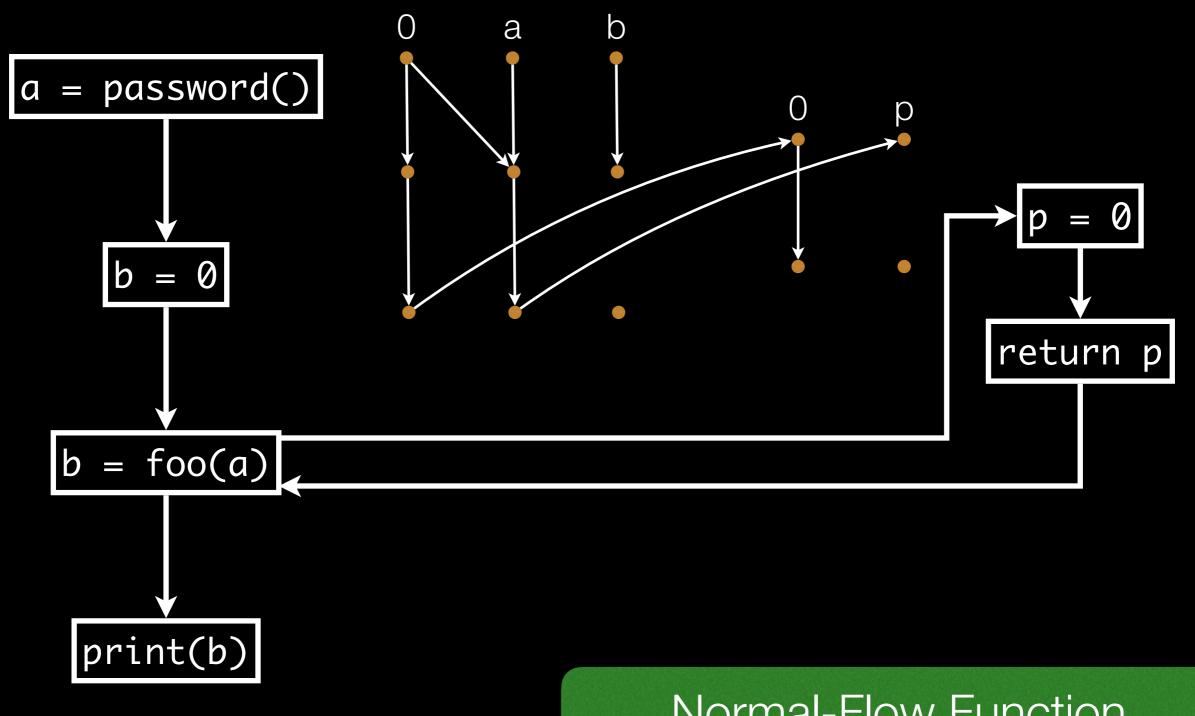




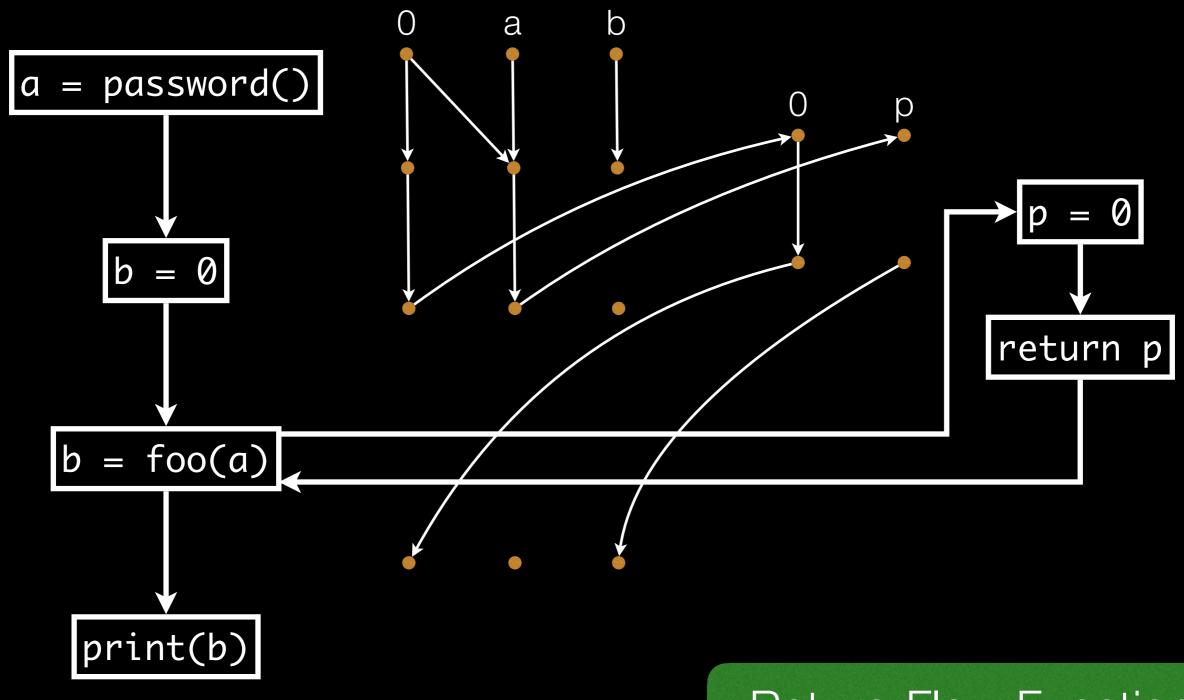
Normal-Flow Function



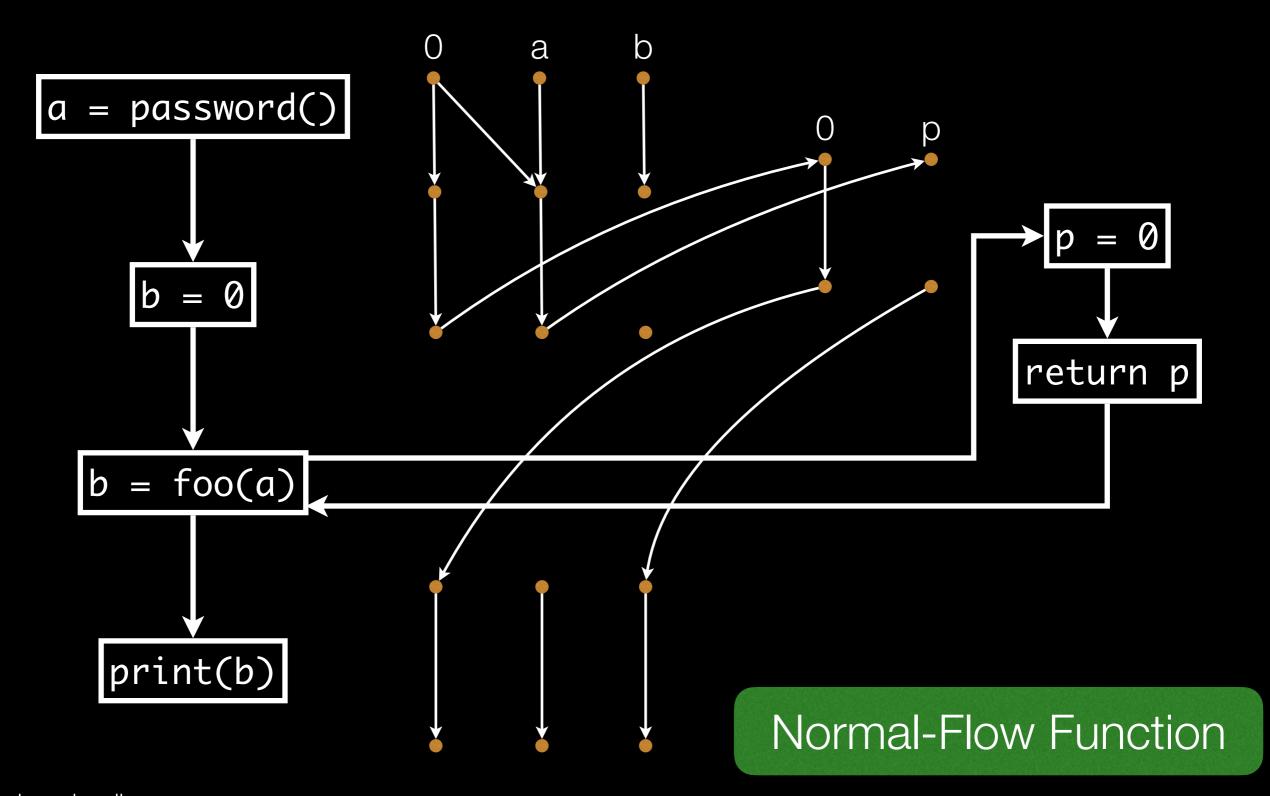


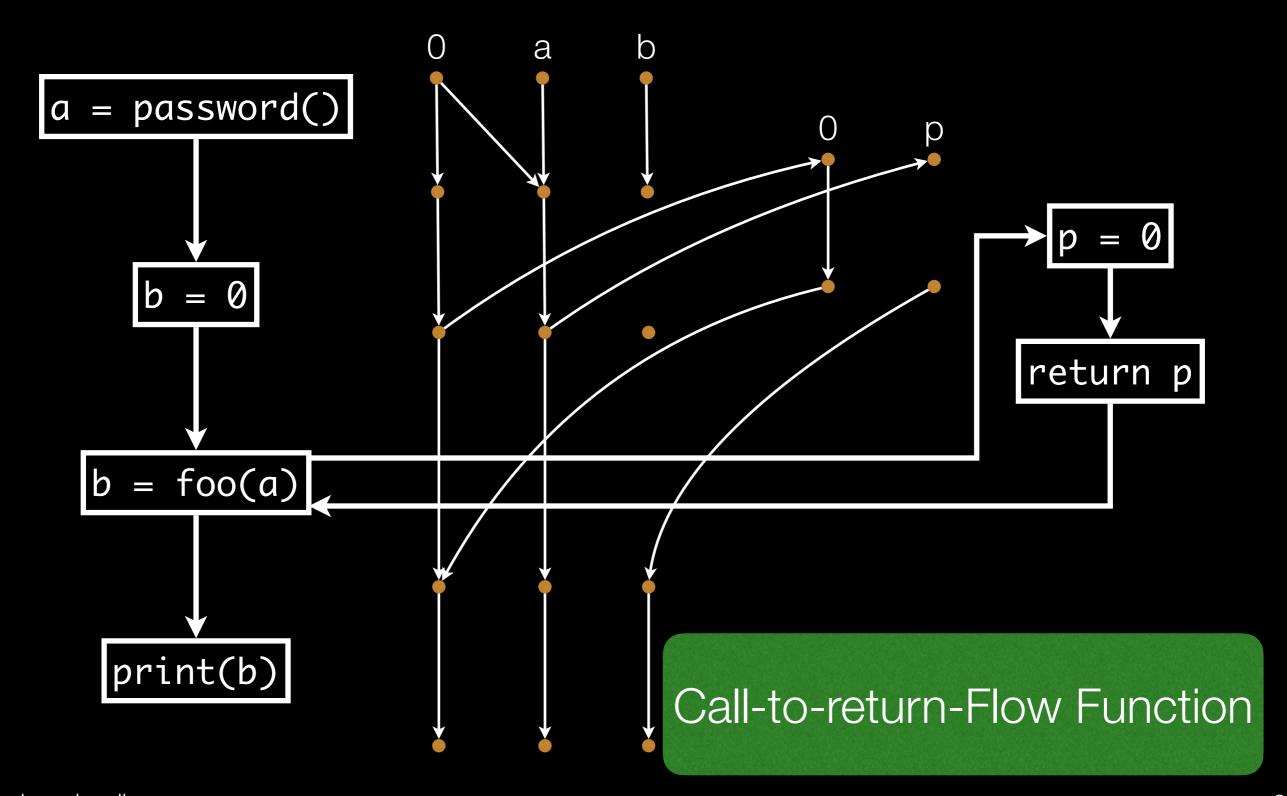


Normal-Flow Function



Return-Flow Function





#### Flow Functions

Call-Flow Function

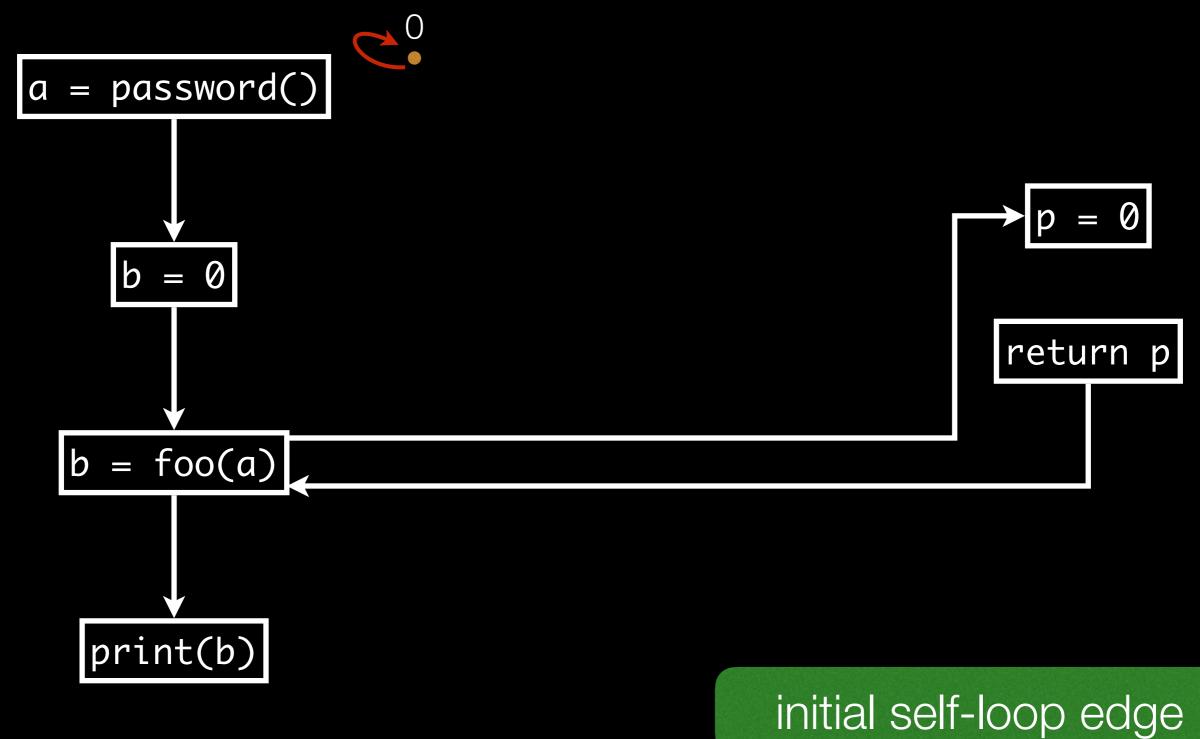
Normal-Flow Function

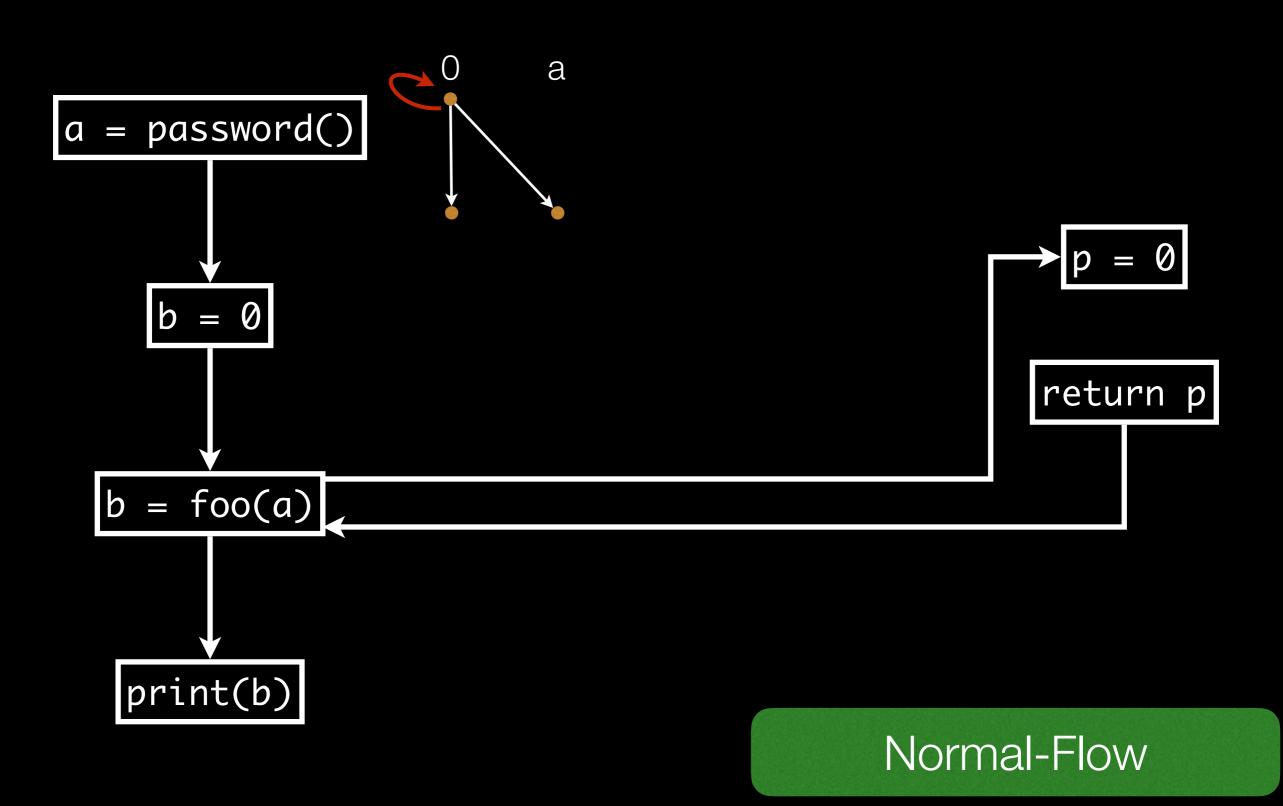
Return-Flow Function

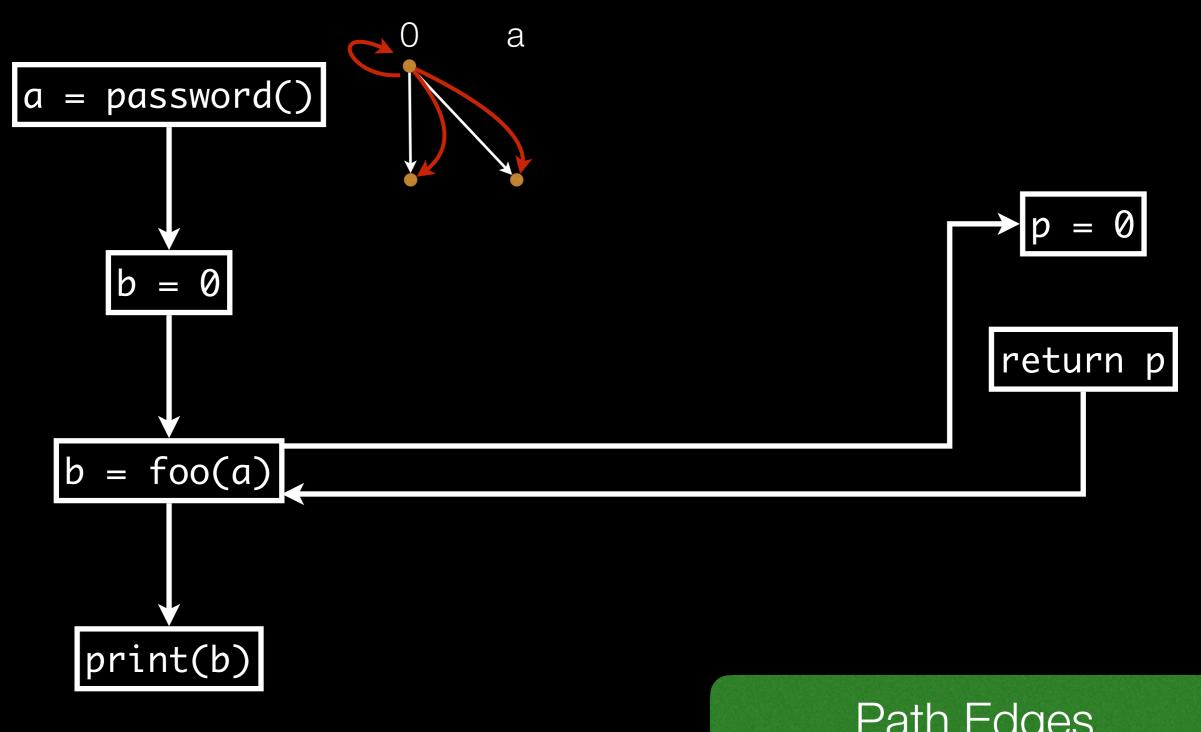
id Ca

Call-to-return-Flow Function

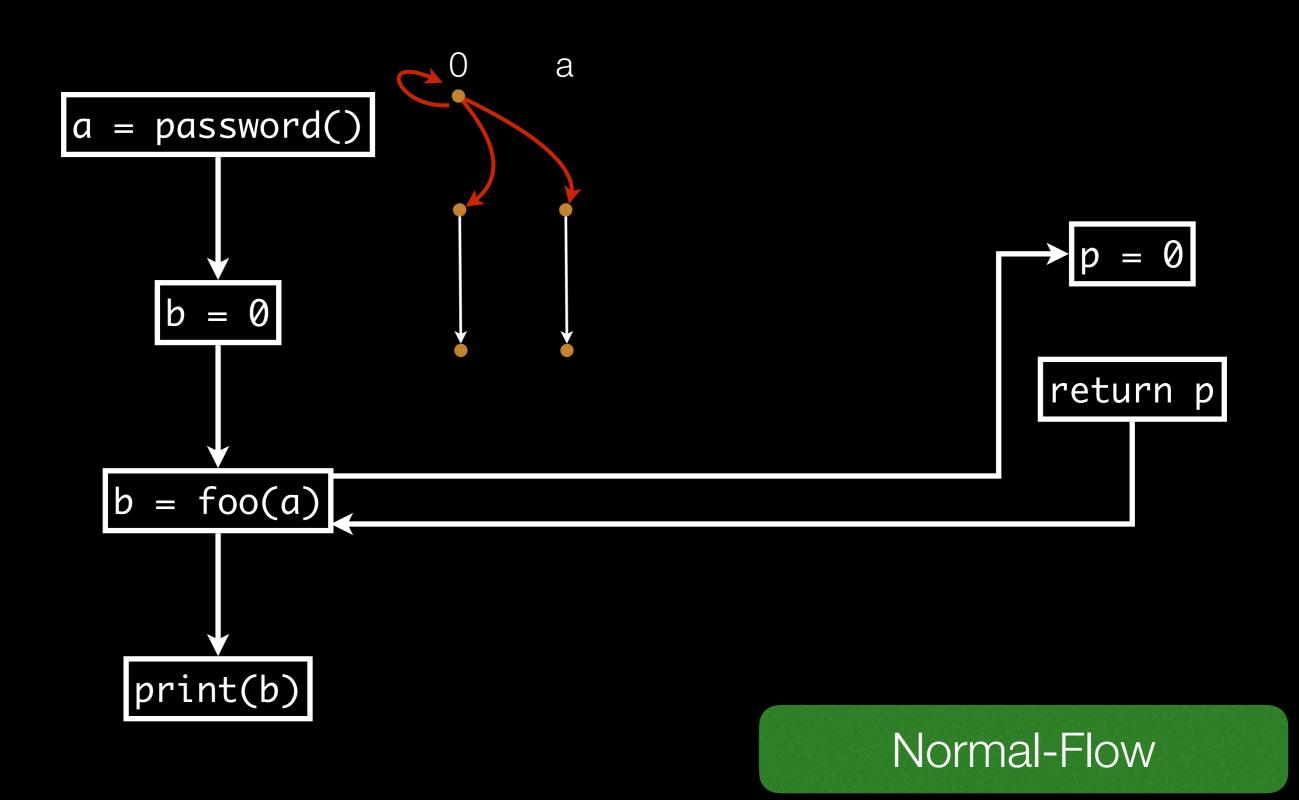


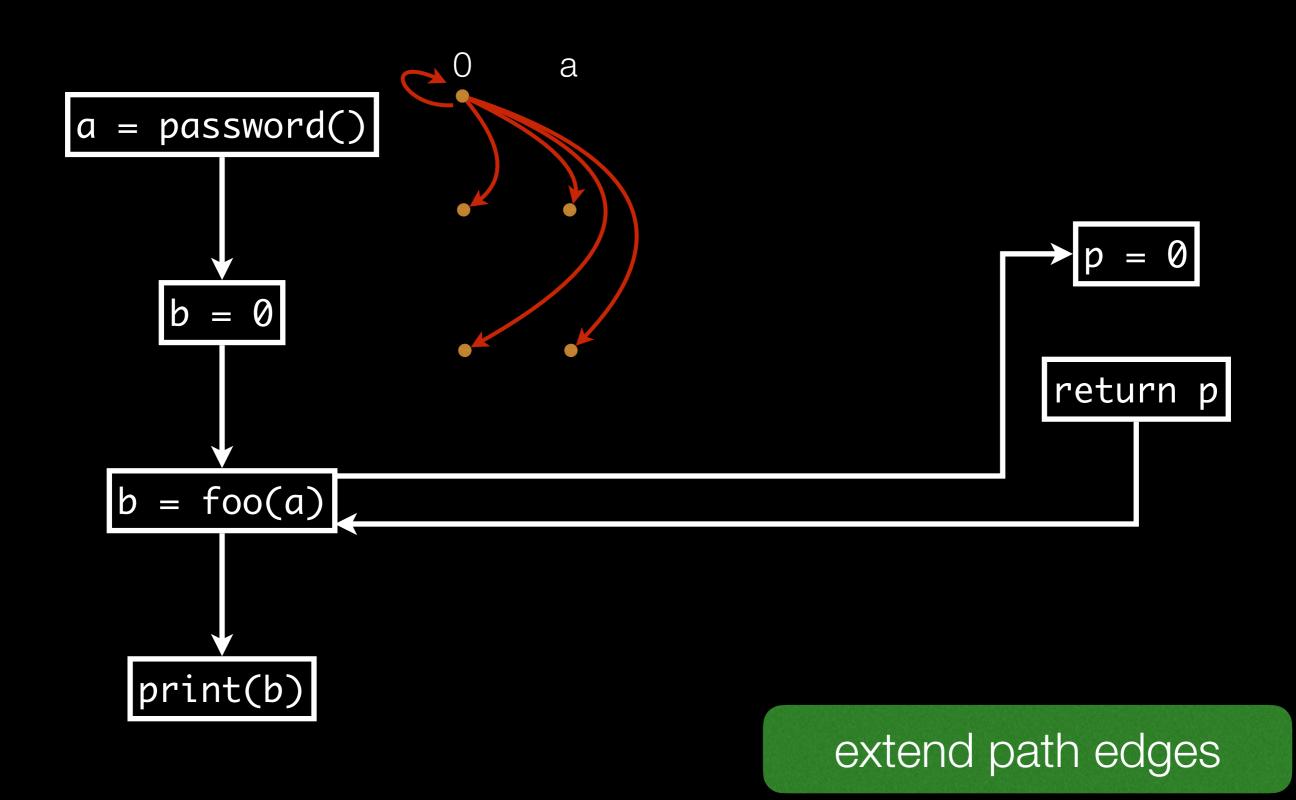


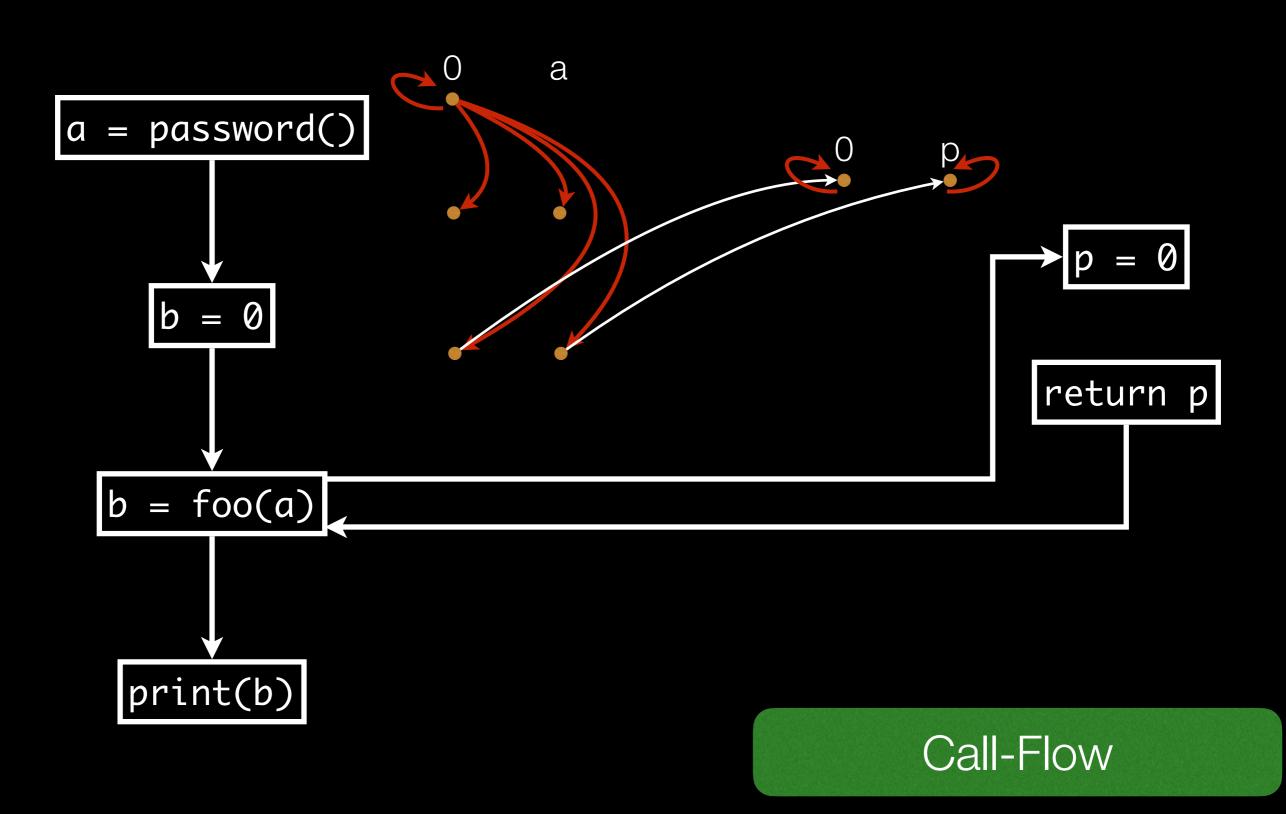


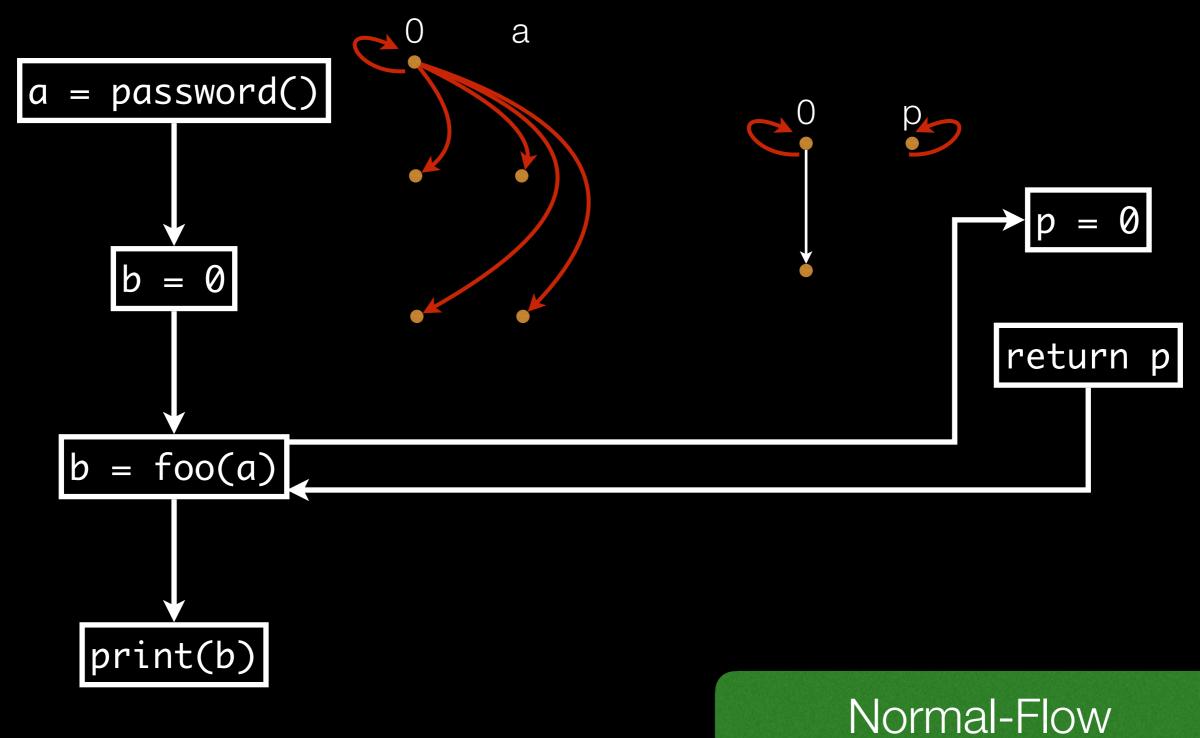


Path Edges

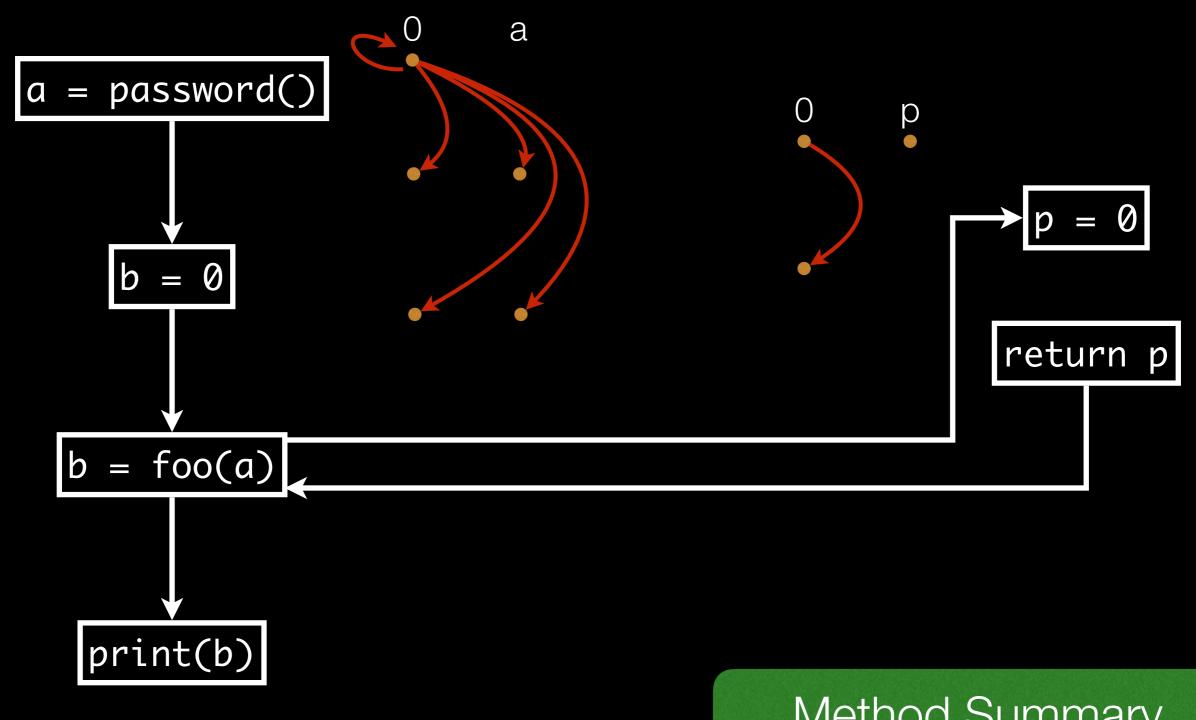




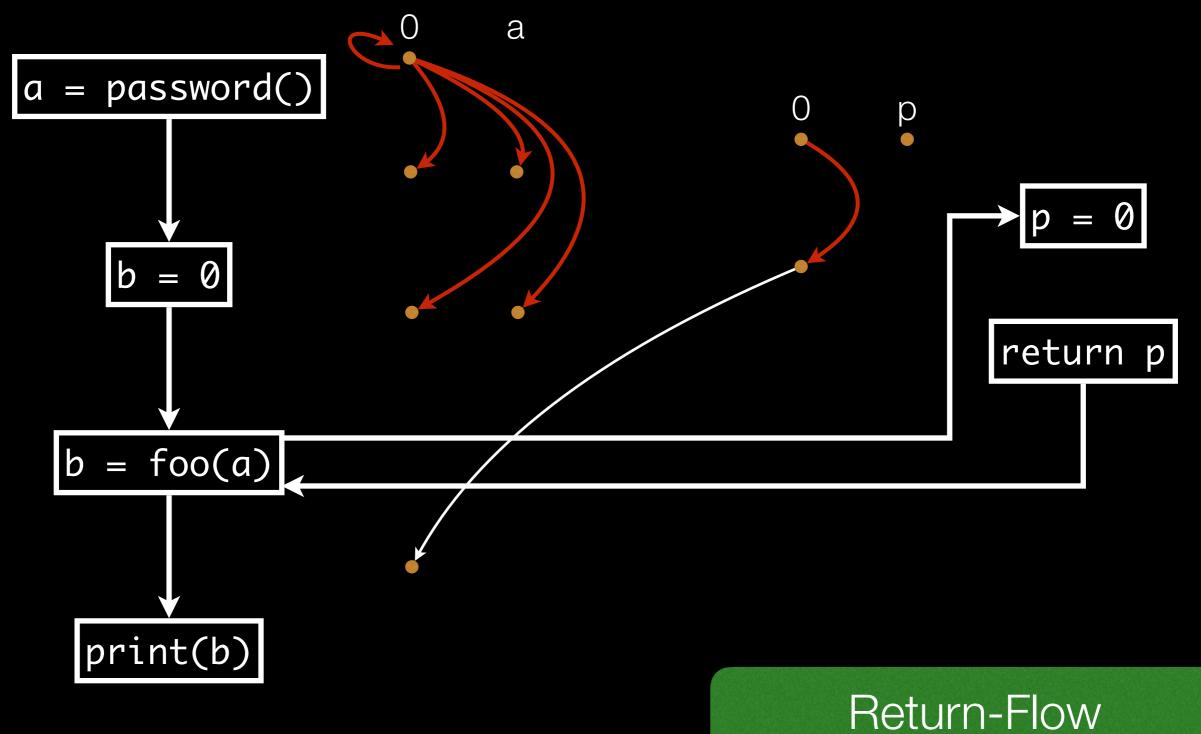


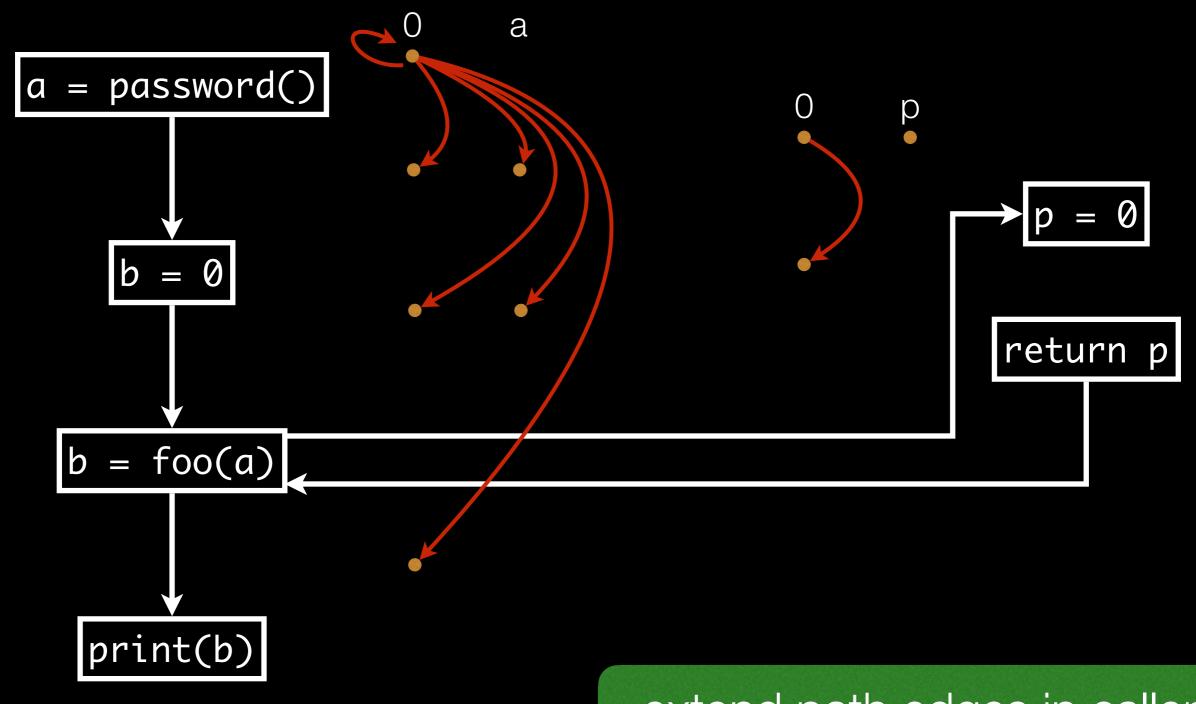


INOTITIAL LIOW

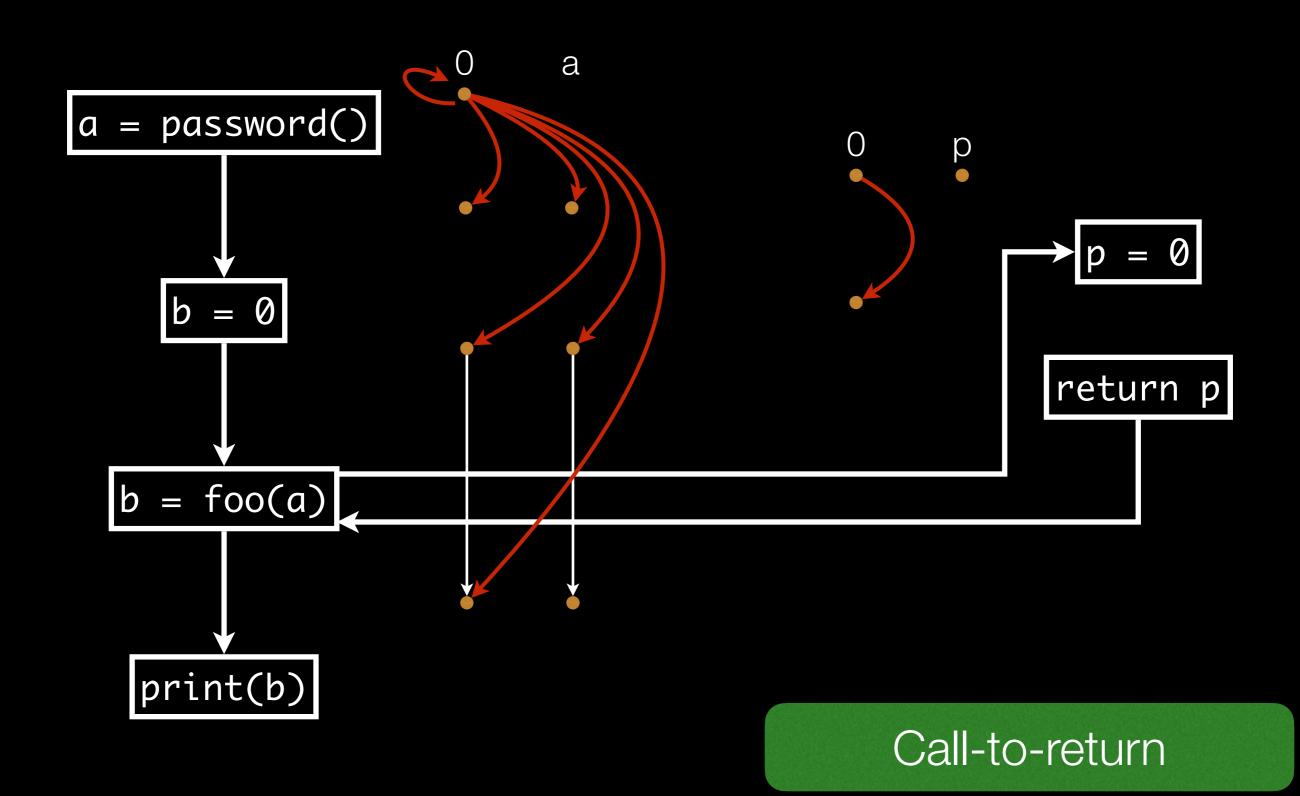


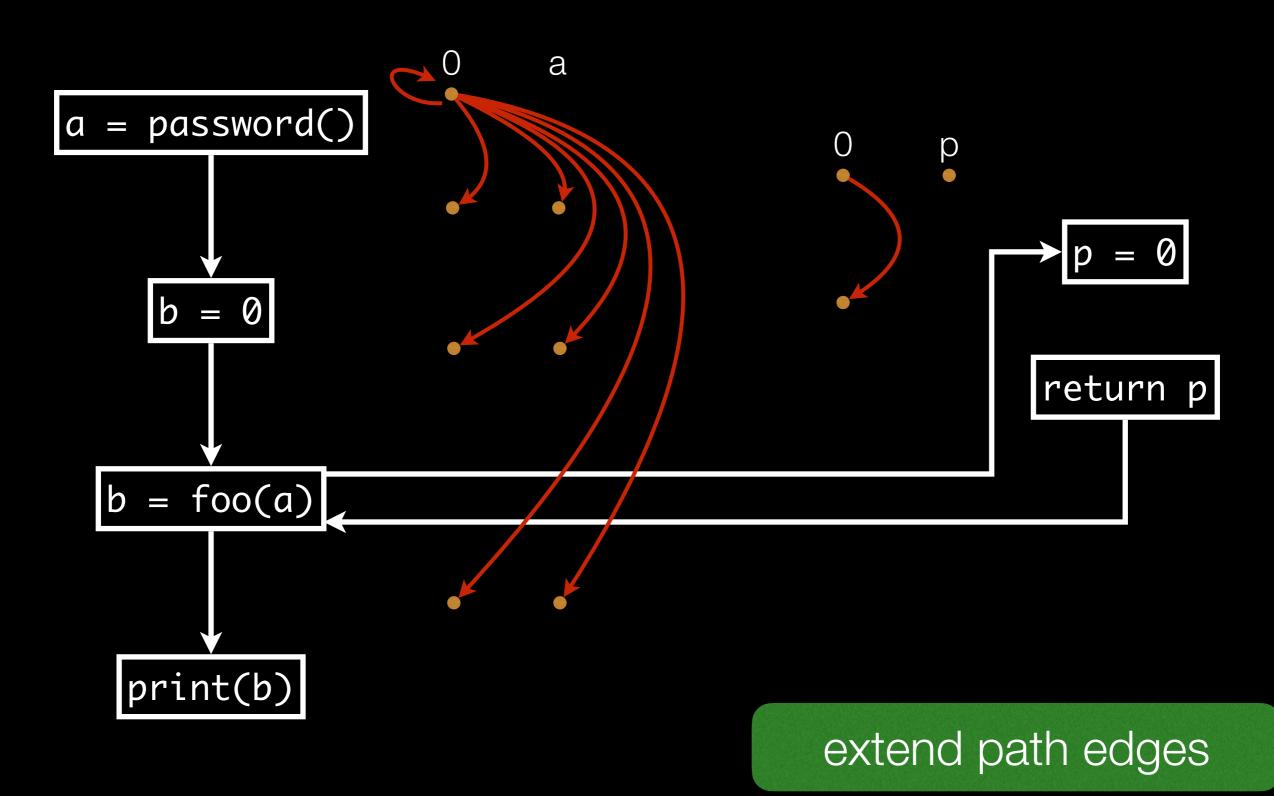
Method Summary

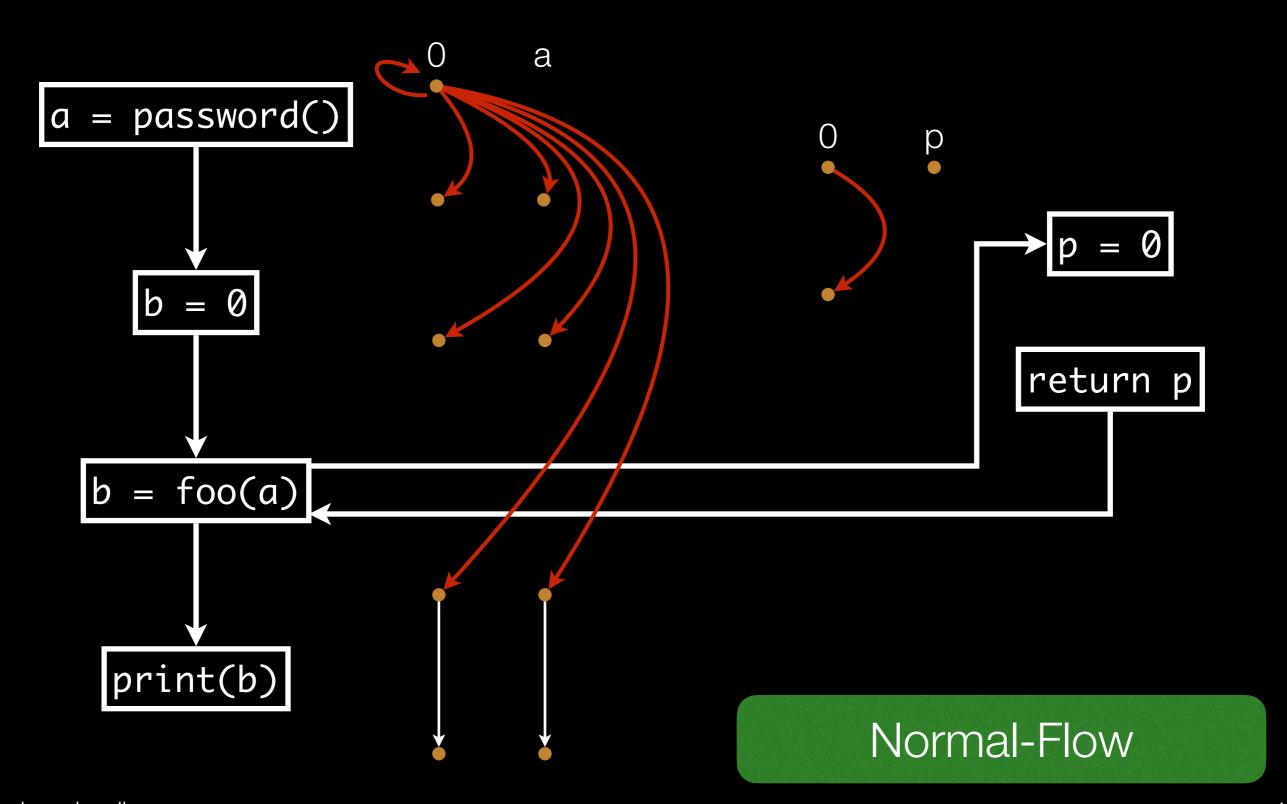


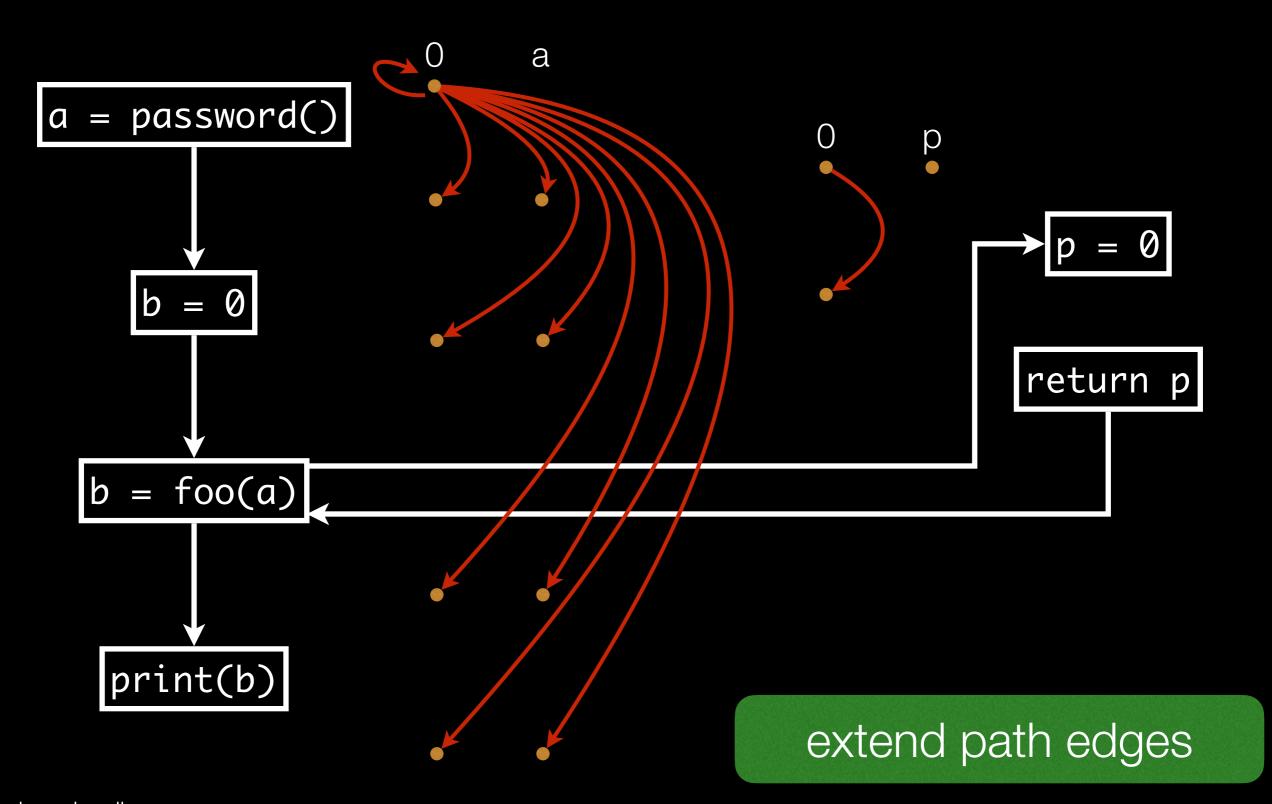


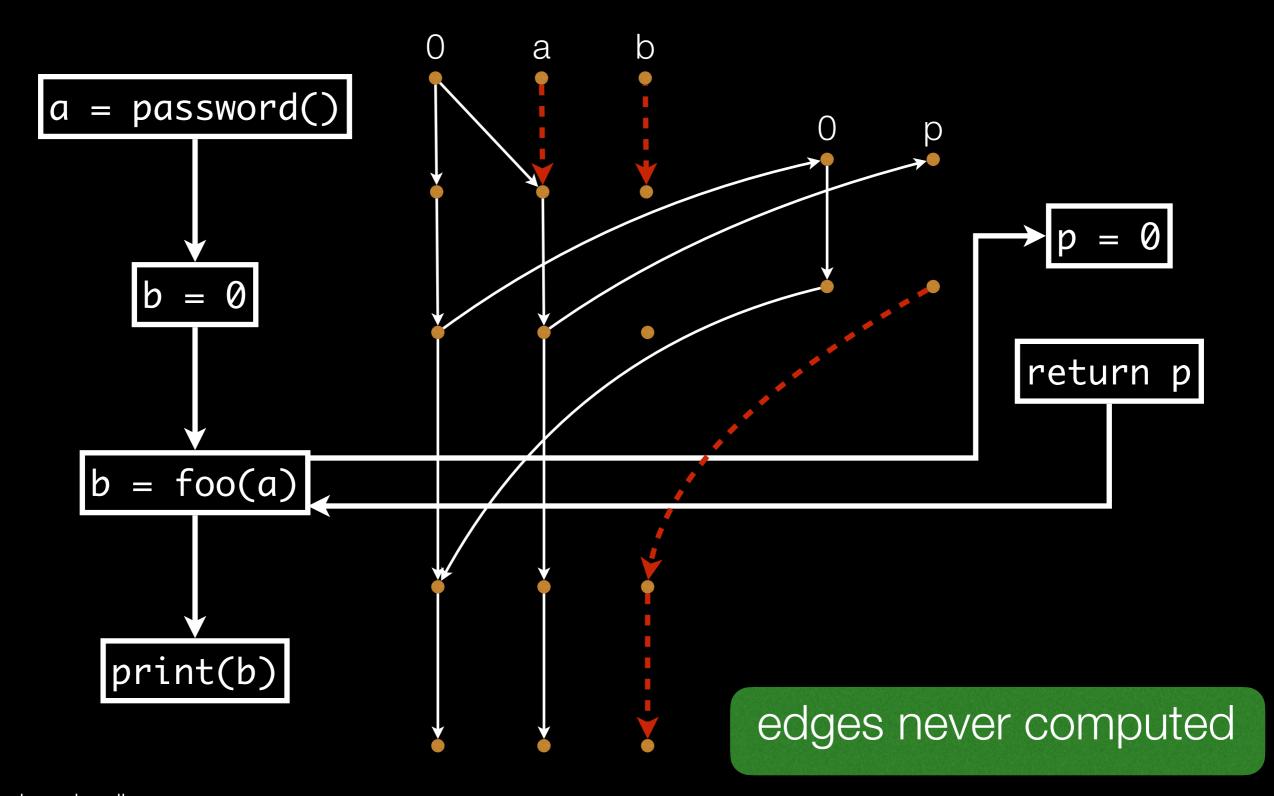
extend path edges in caller











# Tidbits!

## Field-Sensitivity

a.f

a.\*

A.f

Field-Sensitive

Field-Insensitive

Field-Based

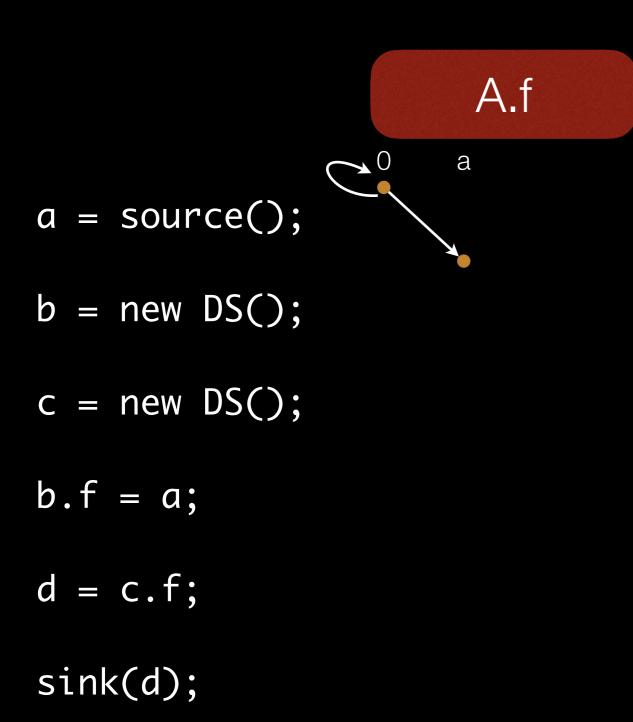
A.f

$$b = new DS();$$

$$c = new DS();$$

$$b.f = a;$$

$$d = c.f;$$





a



$$b = new DS();$$

$$c = new DS();$$

$$b.f = a;$$

$$d = c.f;$$



а



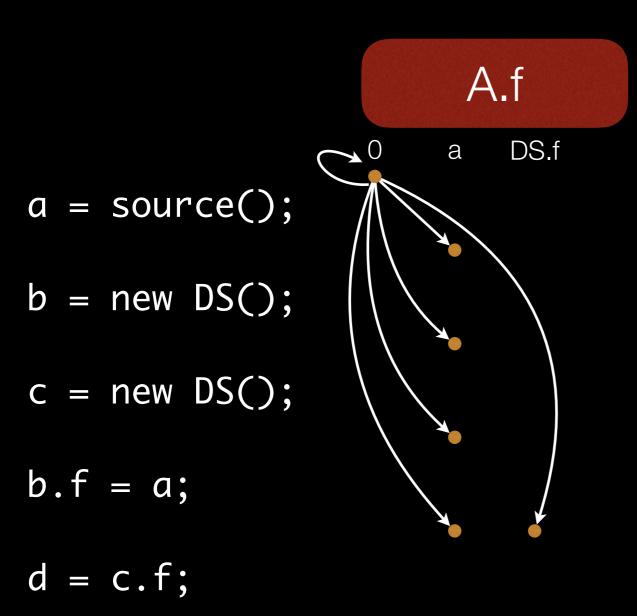
b = new DS();

c = new DS();

$$b.f = a;$$

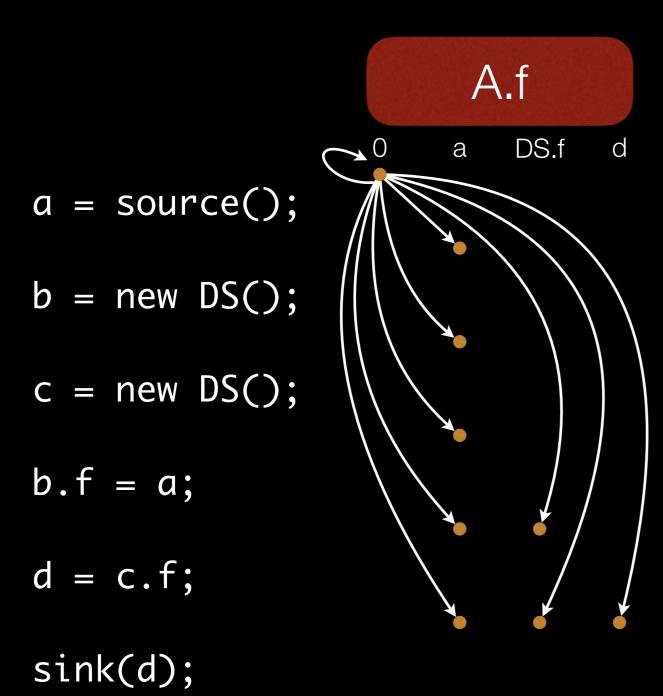
d = c.f;

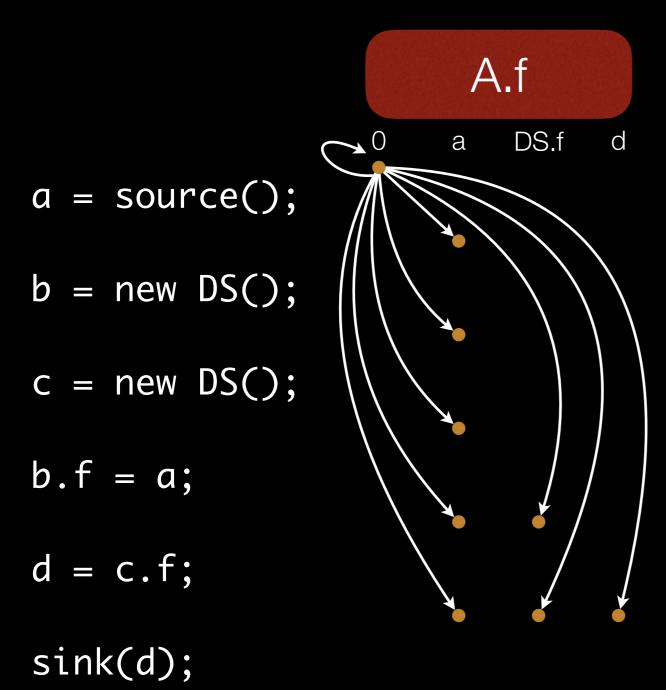
sink(d);

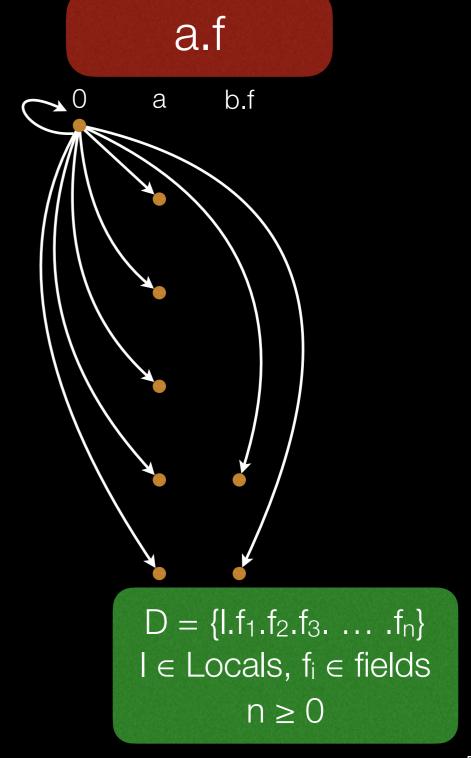


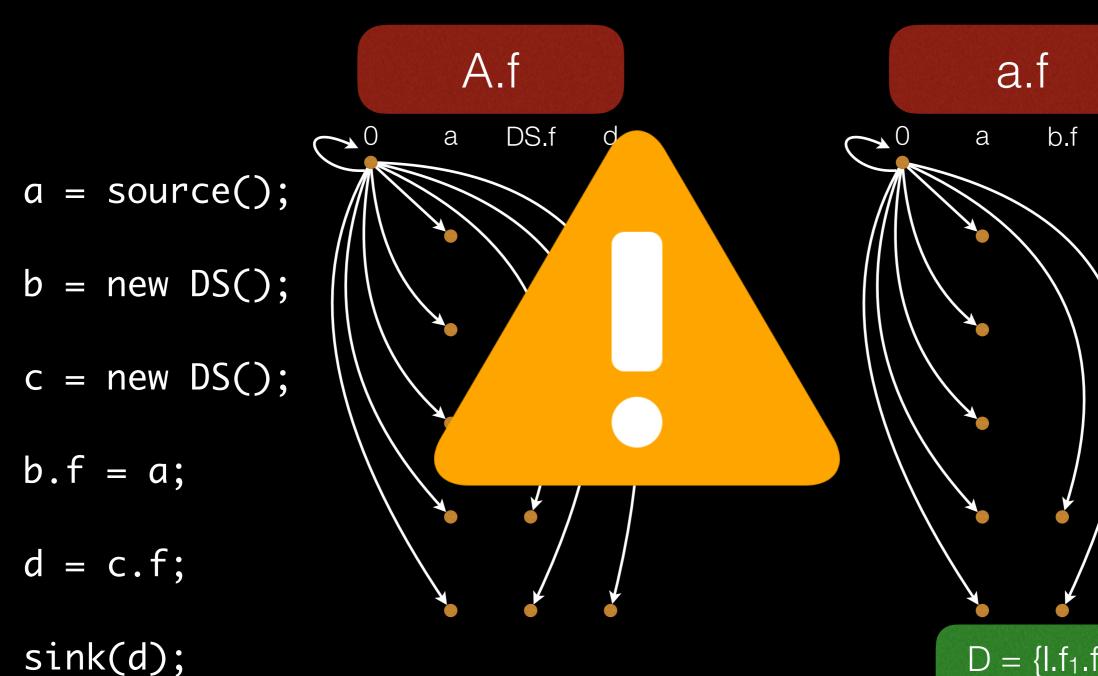
D = Locals u Fields

sink(d);







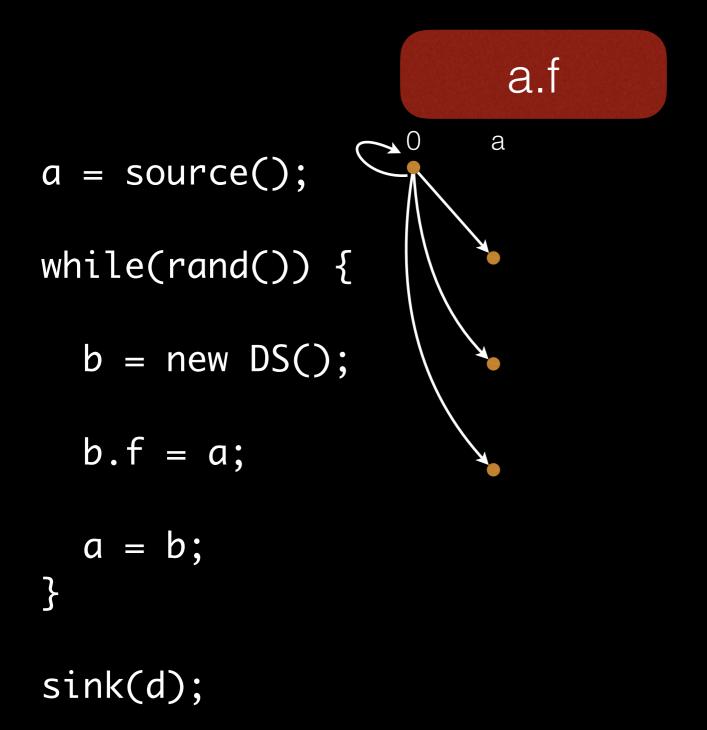


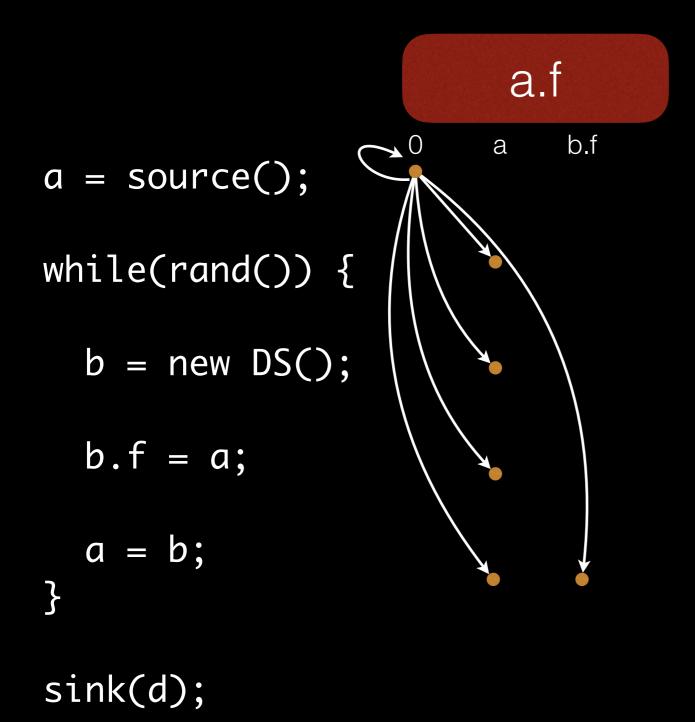
Infinite domain!!

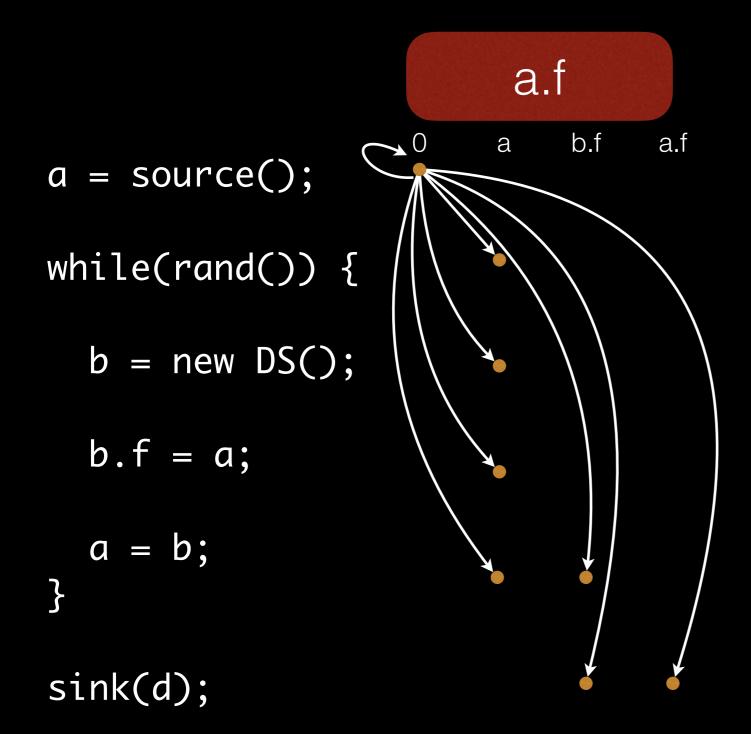
$$\begin{split} D &= \{l.f_1.f_2.f_3. \dots .f_n\} \\ I &\in Locals, \ f_i \in fields \\ n &\geq 0 \end{split}$$

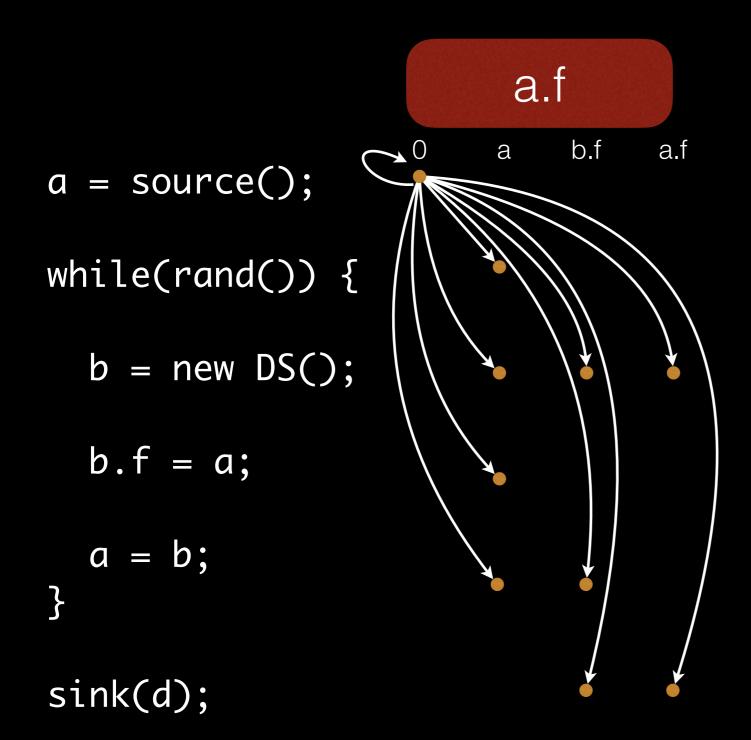
```
a.f
a = source();
while(rand()) {
  b = new DS();
  b.f = a;
  a = b;
sink(d);
```

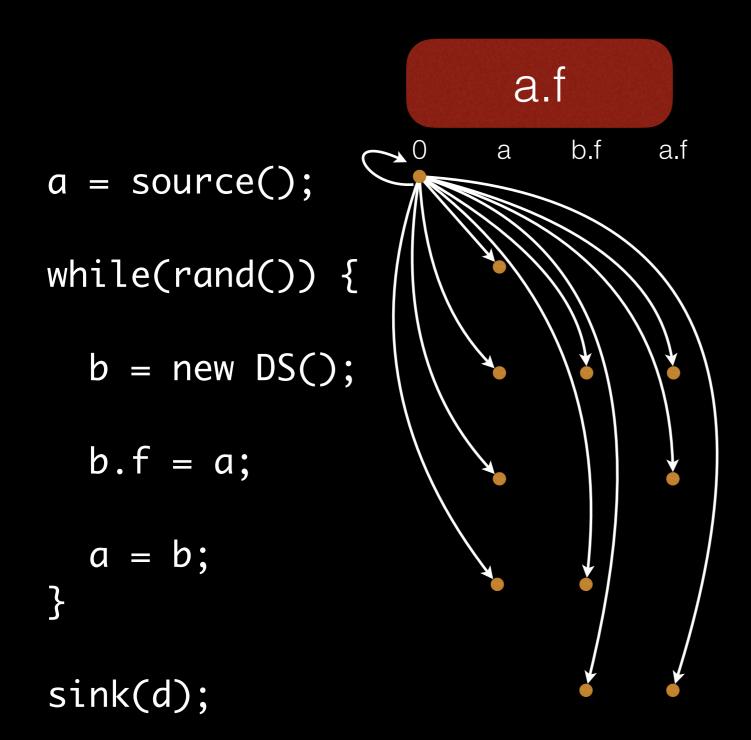
a.f a = source();while(rand()) { b = new DS(); b.f = a;a = b;sink(d);

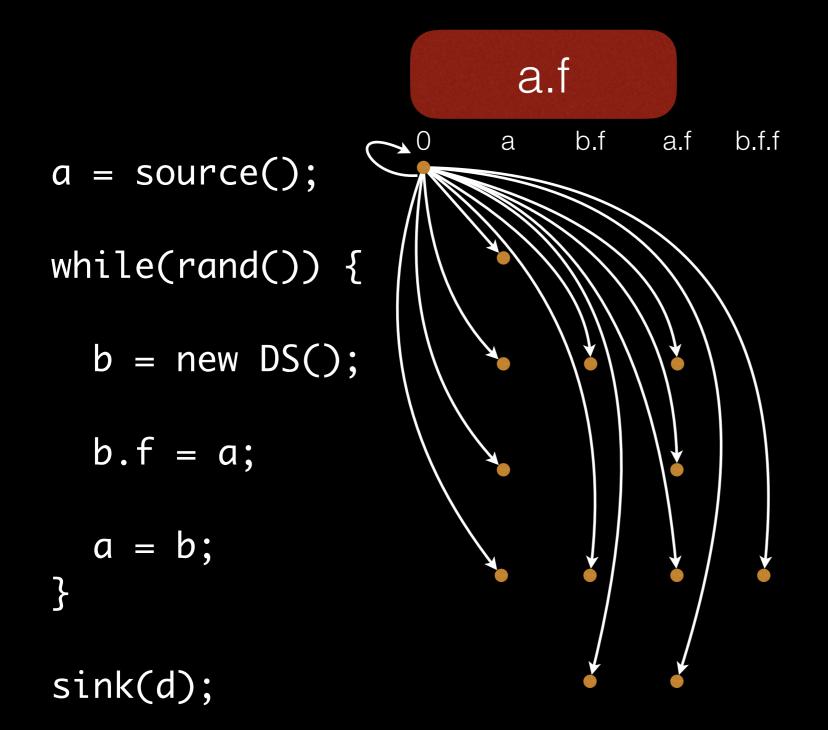


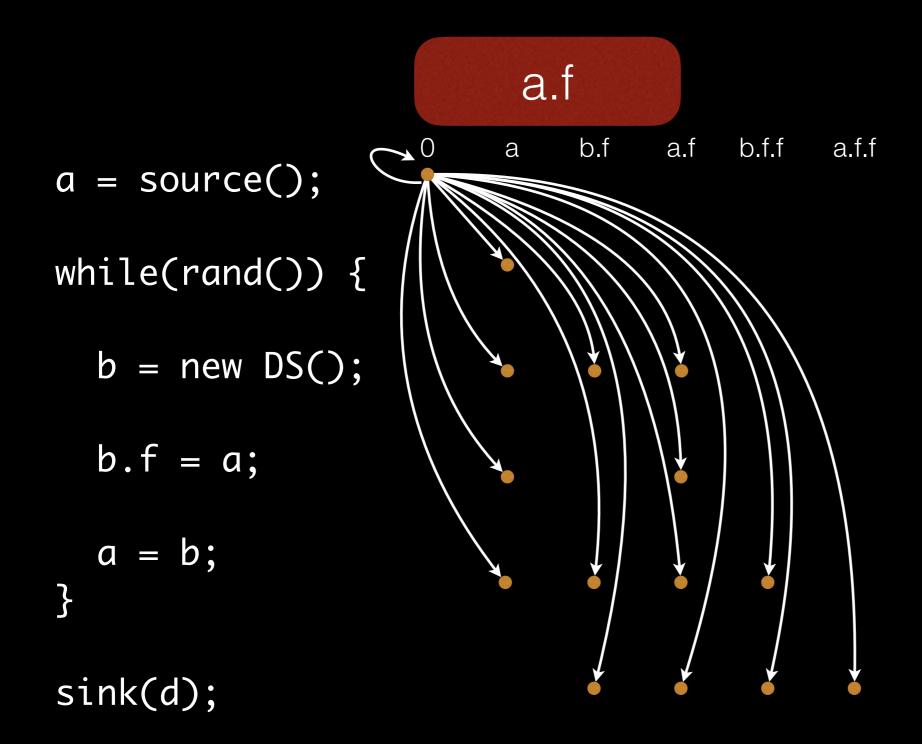


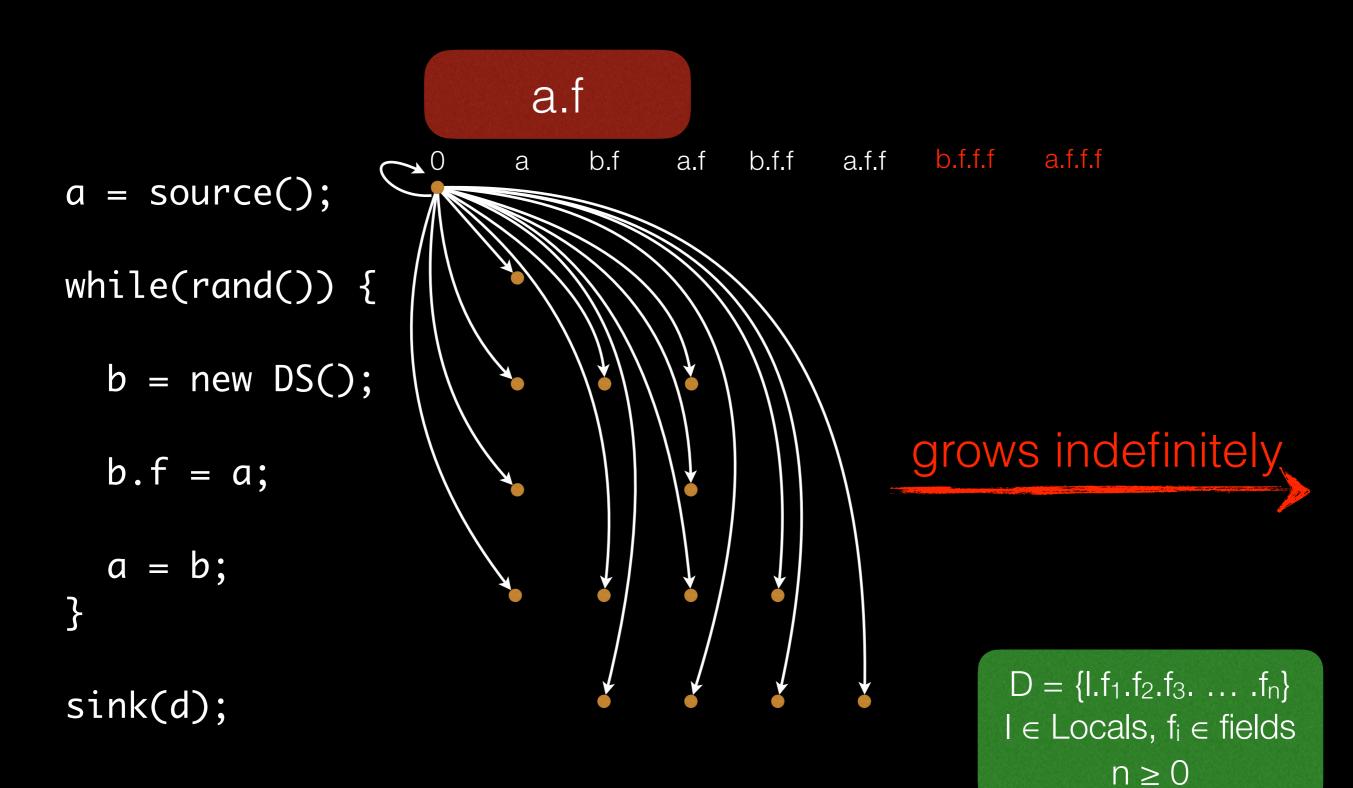


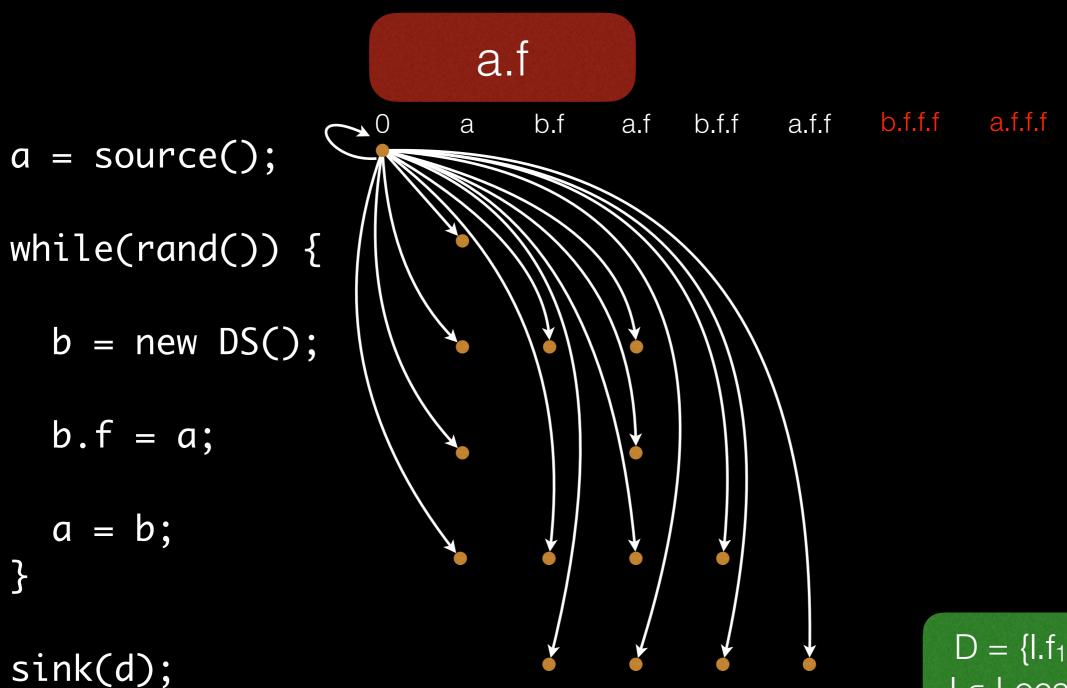


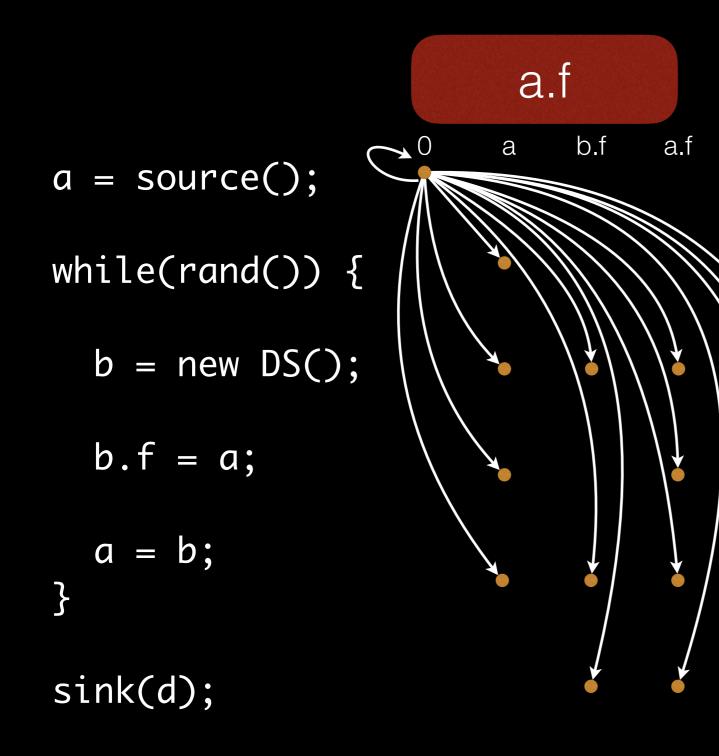






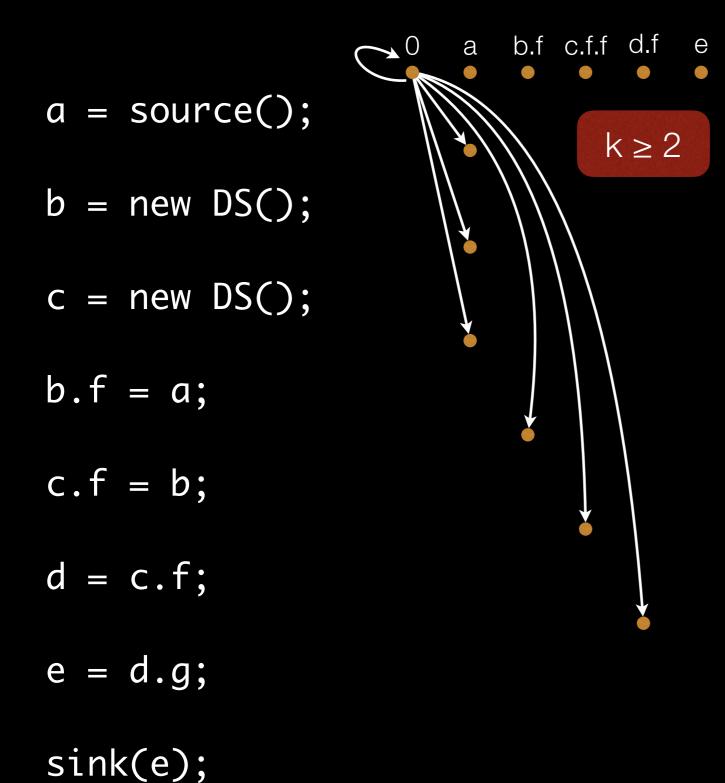


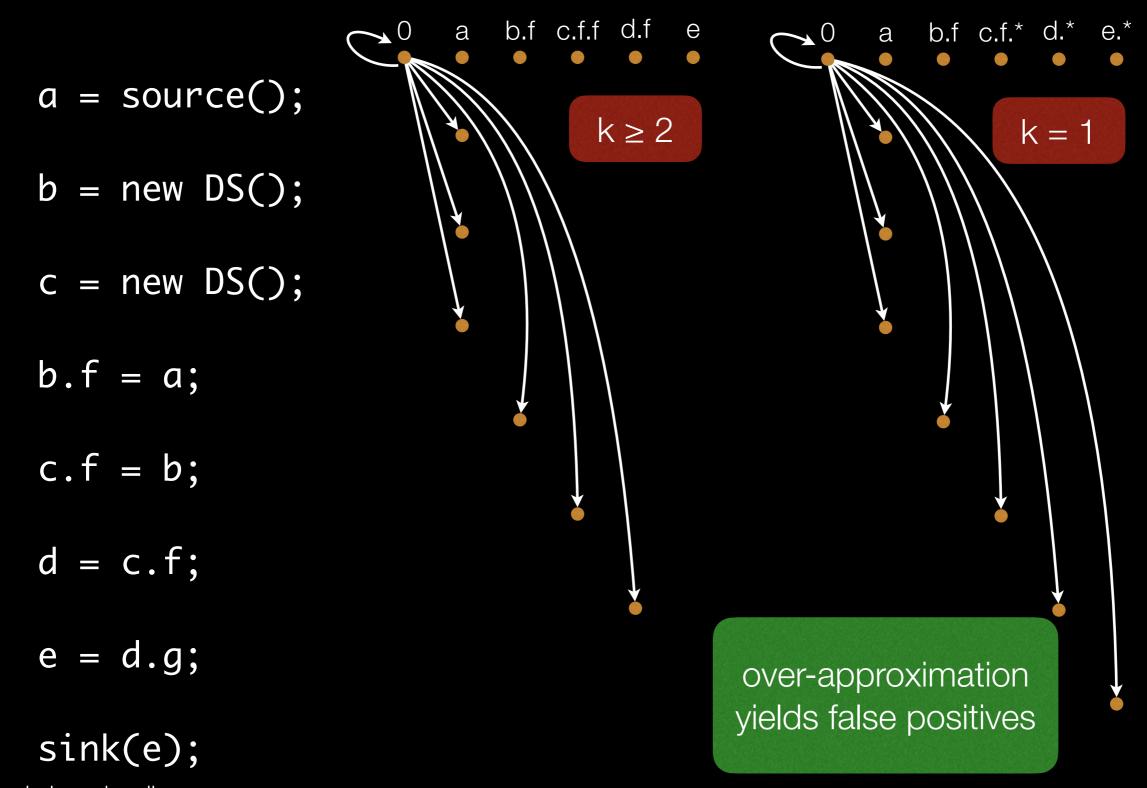




unsound domain beyond k => imprecise







#### **Tidbits**

Flow Sensitivity?

Context Sensitivity?

#### **Tidbits**

Flow Sensitivity?

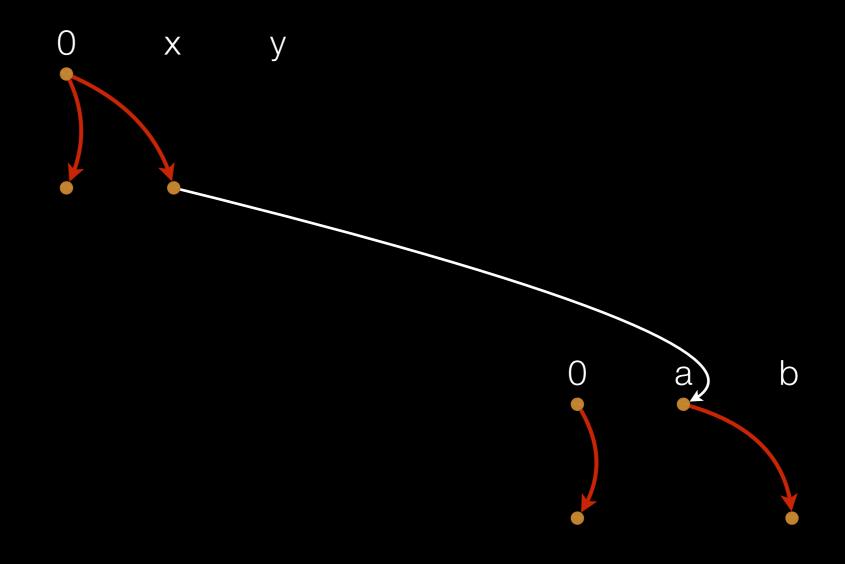
Method Summaries

Context Sensitivity?

• • •

call1();

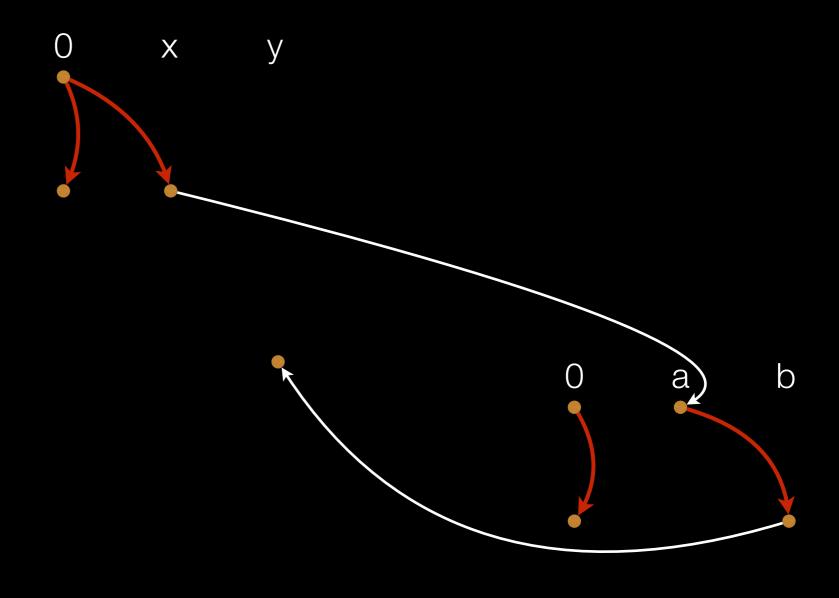
• •



• •

call1();

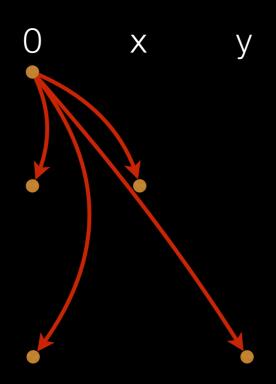
• •

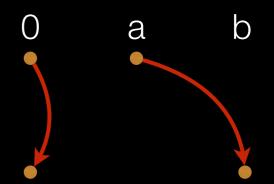


• • •

call1();

• •

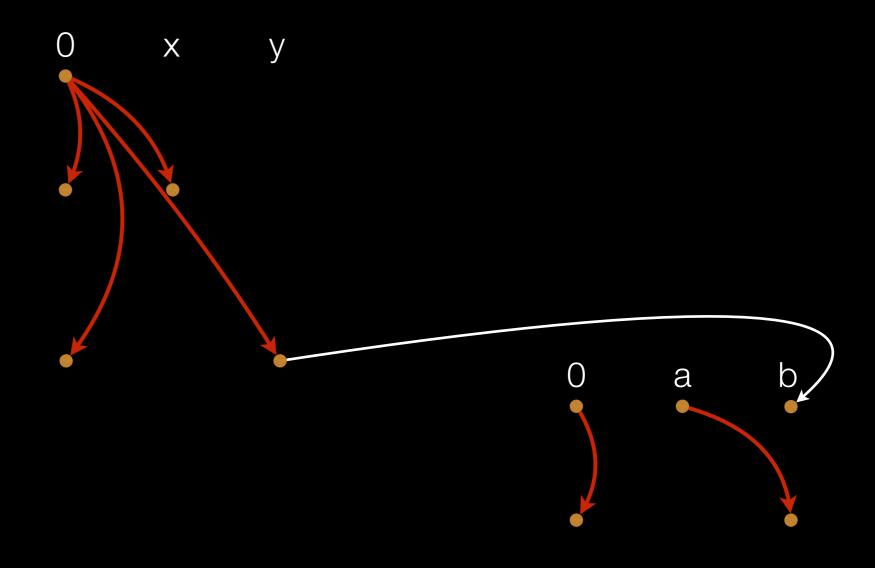




• •

call1();

• •

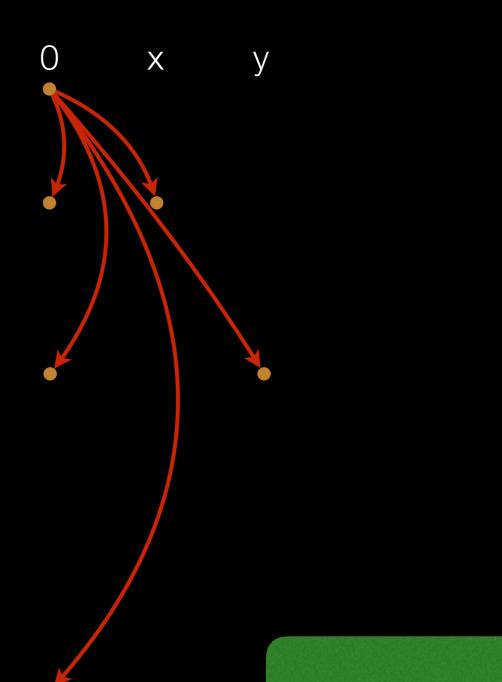


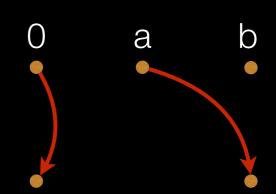
• •

call1();

• •

call2();





same summary different contexts => different results

#### **Tidbits**

- Flow Sensitivity
- Context Sensitivity
- Large Domains => performance problems

$$r = v + 1$$
 $0 \quad (v,0) \quad (r,1)$ 

In any context, if v is 0 before the call, then it is true that r is 1 after the call

```
main() {
x = inc(1);
y = inc(x);
z = inc(y);
 print(z);
      (v,0) (r,1)
```

```
inc(v) {
 r = v + 1;
 return r;
```

```
main() {
                        inc(v) {
x = inc(1);
                          r = v + 1;
y = inc(x);
                          return r;
z = inc(y);
 print(z);
       (v,0) (r,1) (v,1) (r,2)
```

```
main() {
                              inc(v) {
 x = inc(1);
                                 = \vee + 1.
 V - inc(V)
           Solution: IDE instead of IFDS
          (v,0) (r,1) (v,1) (r,2) (v,2)
```

### Recap

- Interprocedural Finite Distributive Subset
- Flow Functions
- Taint Analysis Example
- On-the-fly ESG
- Tidbits: field-sensitivity, contextsensitivity, large domains

#### Next

• Distributive Environments (IDE)