### Outline

This topic looks at another problem solved by breadth-first traversals

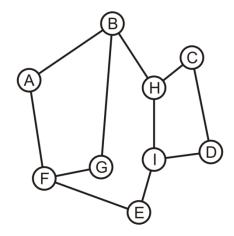
- Determining if a graph is bipartite
- Definition of a bipartite graph
- The algorithm
- An example

### **Definition**

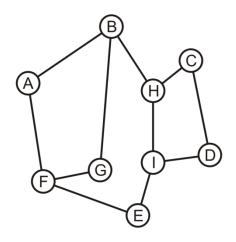
#### **Definition**

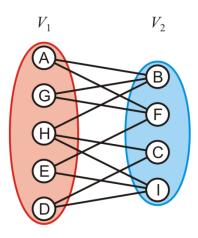
- A bipartite graph is a graph where the vertices V can be divided into two disjoint sets  $V_1$  and  $V_2$  such that **every** edge has one vertex in  $V_1$  and the other in  $V_2$ 

Consider this graph: is it bipartite?

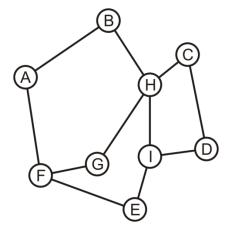


Yes: With a little work, it is possible to determine that we can decompose the vertices into two disjoint sets



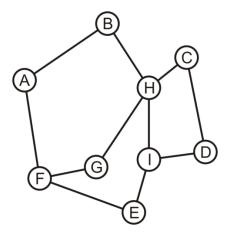


Is this graph bipartite?



In this case, it is not a bipartite graph

– Can we find a traversal that will determine if a graph is bipartite?

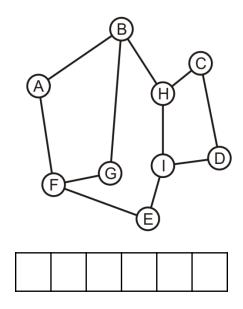


Consider using a breadth-first traversal for a connected graph:

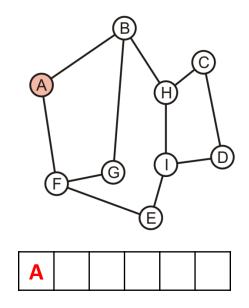
- Choose a vertex, mark it belonging to  $V_1$  and push it onto a queue
- While the queue is not empty, pop the front vertex v and
  - Any adjacent vertices that are already marked must belong to the set not containing v, otherwise, the graph is not bipartite (we are done); while
  - Any unmarked adjacent vertices are marked as belonging to the other set and they are pushed onto the queue
- If the queue is empty, the graph is bipartite

With the first graph, we can start with any vertex

We will use colours to distinguish the two sets

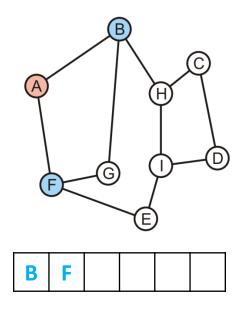


Push A onto the queue and colour it red



Pop A and its two neighbours are not marked:

Mark them as blue and push them onto the queue

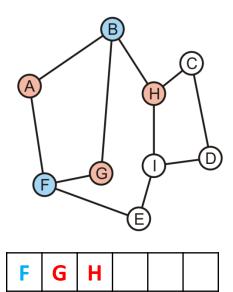


#### Pop B—it is blue:

- Its one marked neighbour, A, is red

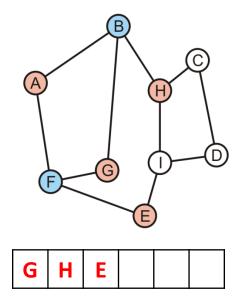
Its other neighbours G and H are not marked: mark them red and push

them onto the queue



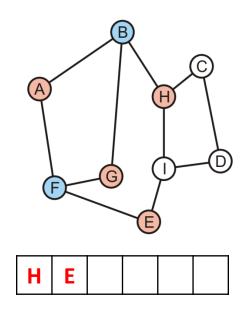
#### Pop F—it is blue:

- Its two marked neighbours, A and G, are red
- Its neighbour E is not marked: mark it red and pus it onto the queue



### Pop G—it is red:

Its two marked neighbours, B and F, are blue

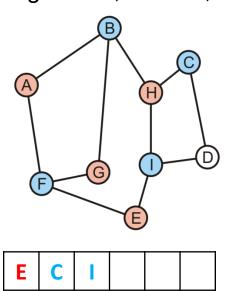


#### Pop H—it is red:

- Its marked neighbours, B, is blue

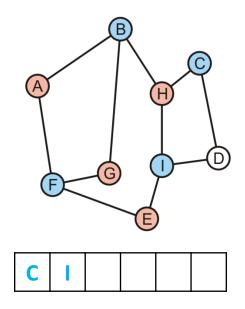
It has two unmarked neighbours, C and I; mark them blue and push

them onto the queue



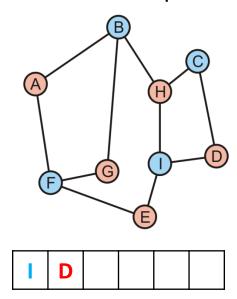
### Pop E—it is red:

Its marked neighbours, F and I, are blue



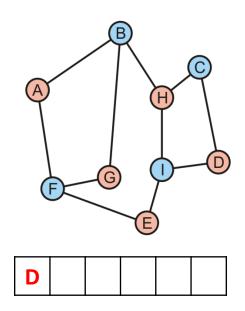
#### Pop C—it is blue:

- Its marked neighbour, H, is red
- Mark D as red and push it onto the queue



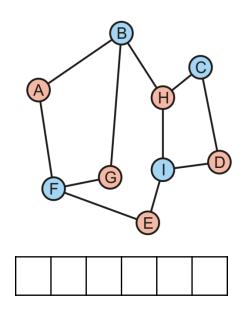
#### Pop I—it is blue:

Its marked neighbours, H, D and E, are all red

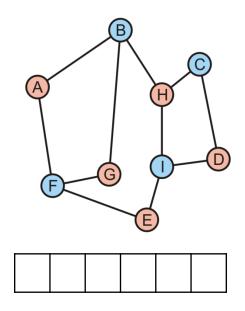


#### Pop D—it is red:

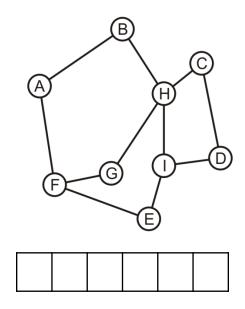
Its marked neighbours, C and I, are both blue



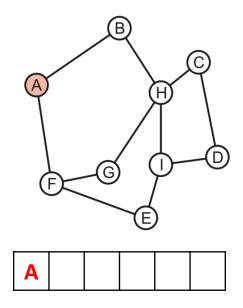
The queue is empty, the graph is bipartite



Consider the other graph which was claimed to be not bipartite

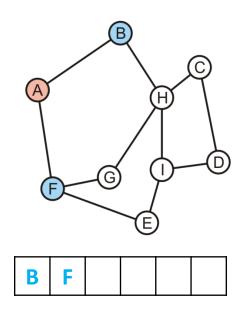


Push A onto the queue and colour it red



#### Pop A off the queue:

 Its neighbours are unmarked: colour them blue and push them onto the queue

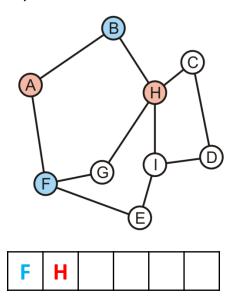


#### Pop B off the queue:

- Its one neighbour, A, is red

The other neighbour, H, is unmarked: colour it red and push it onto the

queue

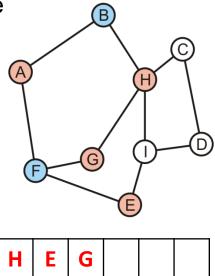


### Pop F off the queue:

- Its one neighbour, A, is red

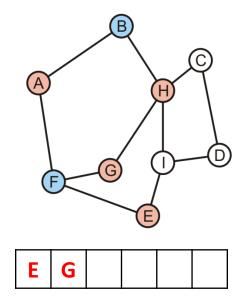
The other neighbours, E and G, are unmarked: colour them red and

push it onto the queue



Pop H off the queue—it is red:

- Its one neighbour, G, is already red
- The graph is not bipartite



#### **Definition**

Cycles that contains either an even number or an odd number of vertices are said to be *even cycles* and *odd cycles*, respectively

#### Theorem

A graph is bipartite if and only if it does not contain any odd cycles

Reference: Kleinberg and Tardos

### Sumary

This topic looked at identifying bipartite graphs

- Perform a breadth-first traversal
- Each vertex is given one of two identifiers (we used color)
- The first vertex is identified as one color and pushed onto the queue
- When a vertex is popped:
  - Each unvisited neighbor is pushed onto the tree with the opposite color
  - Each visited neighbor must be the opposite color
    - If one is not, the graph is not bipartite

### References

Wikipedia, http://en.wikipedia.org/wiki/Breadth-first\_search#Testing\_bipartiteness http://en.wikipedia.org/wiki/Breadth-first\_search http://en.wikipedia.org/wiki/Bipartite\_graph

[1] Jon Kleinberg and Éva Tardos, *Algorithm Design*, Addison Wesley, 2006, §§3.2-5, pp.78-99.

These slides are provided for the ECE 250 *Algorithms and Data Structures* course. The material in it reflects Douglas W. Harder's best judgment in light of the information available to him at the time of preparation. Any reliance on these course slides by any party for any other purpose are the responsibility of such parties. Douglas W. Harder accepts no responsibility for damages, if any, suffered by any party as a result of decisions made or actions based on these course slides for any other purpose than that for which it was intended.