

Botnet

# Fodcha Is Coming Back, Raising A Wave of Ransom DDoS



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## Background

On April 13, 2022, 360Netlab first disclosed the Fodcha botnet. After our article was published, Fodcha suffered a crackdown from the relevant authorities, and its authors quickly responded by leaving "Netlab pls leave me alone I surrender" in an updated sample. No surprise, Fodcha's authors didn't really stop updating after the fraudulent surrender, and soon a new version was released.

In the new version, the authors of Fodcha redesigned the communication protocol and started to use xxtea and chacha20 algorithms to encrypt sensitive resources and network communications to avoid detection at the file & traffic level; at the same time, a dual-C2 scheme with OpenNIC domain as the primary C2 and ICANN domain as the backup C2 was adopted.

Relying on the strong N-day vulnerability integration capabilities, the comeback of Focha is just as strong as the previous ones. In our data view, in terms of scale, Fodcha has once again developed into a massive botnet with more than 60K daily active bots and 40+ C2 IPs, we also observed it can easily launch more than 1Tbps DDos traffic; in terms of attacks, Fodcha has an average of 100+ daily attack targets and more than 20,000 cumulative attacks, on October 11, Fodcha hit its record and attacked 1,396 unique targets in that single day.

While Fodcha was busy attacking various targets, it has not forgot to mess with us, we saw it using `N3t1@bG@Y` in one of it scan payload.

## Timeline

Backed by our `BotMon` systems, we have kept good track of Fodcha's sample evolution and DDoS attack instructions, and below are the sample evolution and some important DDoS attack events we have seen. (Note: The Fodcha sample itself does not have a specific flag to indicate its version, this is the version number we use internally for tracking purposes)

- On January 12, 2022, the first Fodcha botnet sample was captured.
- April 13, 2022, Disclosure of the [Fodcha](#) botnet, containing versions V1, V2.
- April 19, 2022, captured version V2.x, using `OpenNIC's TLDs` style C2
- April 24, 2022, version V3, using xxtea algorithm to encrypt configuration information, adding `ICANN's TLDs` style C2, adding `anti-sandbox` & `anti-debugging` mechanism.
- June 5, 2022, version V4, using structured configuration information, `anti-sandboxing` & `anti-debugging` mechanism were removed.
- July 7, 2022, version V4.x with an additional set of `ICANN C2`.
- On September 21, 2022, a well-known cloud service provider was attacked with traffic exceeding `1Tbps`.

## Botnet Size

In April, we confirmed that the number of Fodcha's global daily live bots was about `60,000` ([refer to our other article](#)). We don't have accurate number of the current

size, but suspect that the number of current active bots has not dropped, maybe more than **60,000** now.

There is a positive relationship between the size of a botnet and the number of C2 IPs, and the most parsimonious view is that "the larger the botnet, the more C2 infrastructure it requires. In April, there were **10 c2s** to control the **60,000** bots; After that, we observed that the IPs corresponding to its C2 domains continued to increase. Using a simple dig command to query the latest C2 domain name **yellowchinks.dyn**, we can see it resolves to **44 IPs**.

```

]$/ dig yellowchinks.dyn @opennic2.eth-services.de
;; Truncated, retrying in TCP mode.

; <>> DiG 9.7.3-P3-RedHat-9.7.3-8.P3.el6 <>> yellowchinks.dyn @opennic2.eth-services.de
;; global options: +cmd
;; Got answer:
;; →HEADER← opcode: QUERY, status: NOERROR, id: 3711
;; flags: qr aa rd ra; QUERY: 1, ANSWER: 44, AUTHORITY: 0, ADDITIONAL: 0

;; QUESTION SECTION:
;yellowchinks.dyn.           IN      A

;; ANSWER SECTION:
yellowchinks.dyn.    300    IN      A      185.45.192.97
yellowchinks.dyn.    300    IN      A      46.17.47.212
yellowchinks.dyn.    300    IN      A      185.117.73.115
yellowchinks.dyn.    300    IN      A      194.36.189.157
yellowchinks.dyn.    300    IN      A      194.147.87.242
yellowchinks.dyn.    300    IN      A      91.206.93.243
yellowchinks.dyn.    300    IN      A      101.41.102.132
yellowchinks.dyn.    300    IN      A      185.117.73.116
yellowchinks.dyn.    300    IN      A      185.183.98.205
yellowchinks.dyn.    300    IN      A      185.198.57.95
yellowchinks.dyn.    300    IN      A      46.17.47.54
yellowchinks.dyn.    300    IN      A      193.233.253.93
yellowchinks.dyn.    300    IN      A      193.124.24.42
yellowchinks.dyn.    300    IN      A      185.143.220.100
yellowchinks.dyn.    300    IN      A      46.17.42.190
yellowchinks.dyn.    300    IN      A      185.183.98.228
yellowchinks.dyn.    300    IN      A      46.17.43.237
yellowchinks.dyn.    300    IN      A      185.117.75
yellowchinks.dyn.    300    IN      A      46.29.160.14
yellowchinks.dyn.    300    IN      A      185.117.75.45
yellowchinks.dyn.    300    IN      A      195.133.52.70
yellowchinks.dyn.    300    IN      A      101.49.222.133
yellowchinks.dyn.    300    IN      A      195.45.192.96
yellowchinks.dyn.    300    IN      A      194.87.197.3
yellowchinks.dyn.    300    IN      A      185.117.75.117
yellowchinks.dyn.    300    IN      A      185.183.96.7
yellowchinks.dyn.    300    IN      A      91.149.232.128
yellowchinks.dyn.    300    IN      A      185.117.75.34
yellowchinks.dyn.    300    IN      A      46.17.41.79
yellowchinks.dyn.    300    IN      A      185.45.192.212
yellowchinks.dyn.    300    IN      A      91.149.232.129
yellowchinks.dyn.    300    IN      A      185.183.96.60
yellowchinks.dyn.    300    IN      A      185.117.73.109
yellowchinks.dyn.    300    IN      A      185.141.27.157
yellowchinks.dyn.    300    IN      A      185.183.96.8
yellowchinks.dyn.    300    IN      A      185.141.27.235
yellowchinks.dyn.    300    IN      A      193.233.253.10
yellowchinks.dyn.    300    IN      A      193.233.253.220
yellowchinks.dyn.    300    IN      A      193.38.50.197
yellowchinks.dyn.    300    IN      A      194.147.84.28
yellowchinks.dyn.    300    IN      A      194.147.86.193
yellowchinks.dyn.    300    IN      A      195.133.52.29
yellowchinks.dyn.    300    IN      A      194.156.121.87
yellowchinks.dyn.    300    IN      A      194.156.120.36

;; Query time: 287 msec
;; SERVER: 195.10.195.195#53(195.10.195.195)
;; WHEN: Fri Oct 21 15:24:54 2022
;; MSG SIZE  rcvd: 738

```

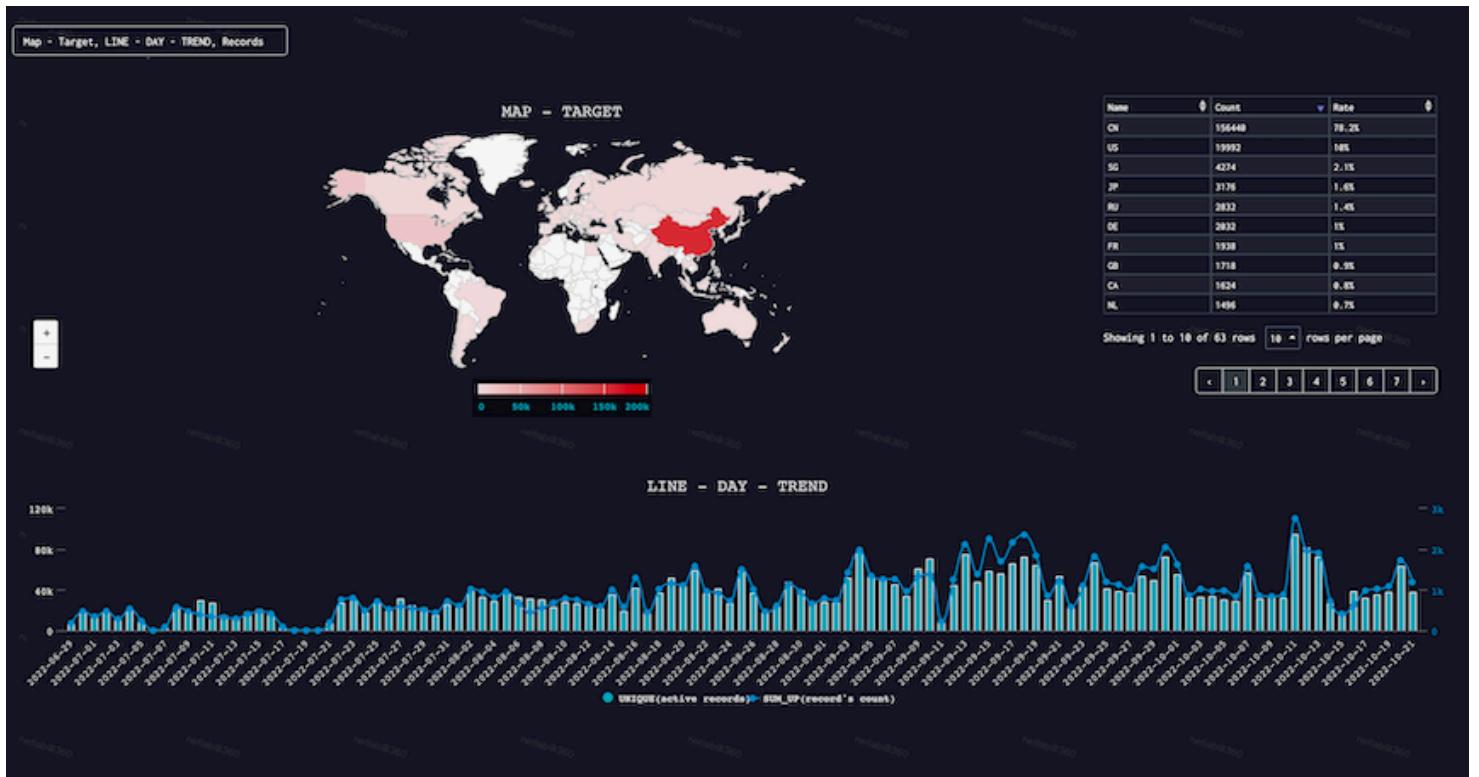
**44 C2 IPs**

**Fodcha C2 Infrastructure**

One likely reason for this is that their botnet is so large that they need to invest more IP resources in order to have a reasonable ratio between Bots and C2s to achieve load balancing.

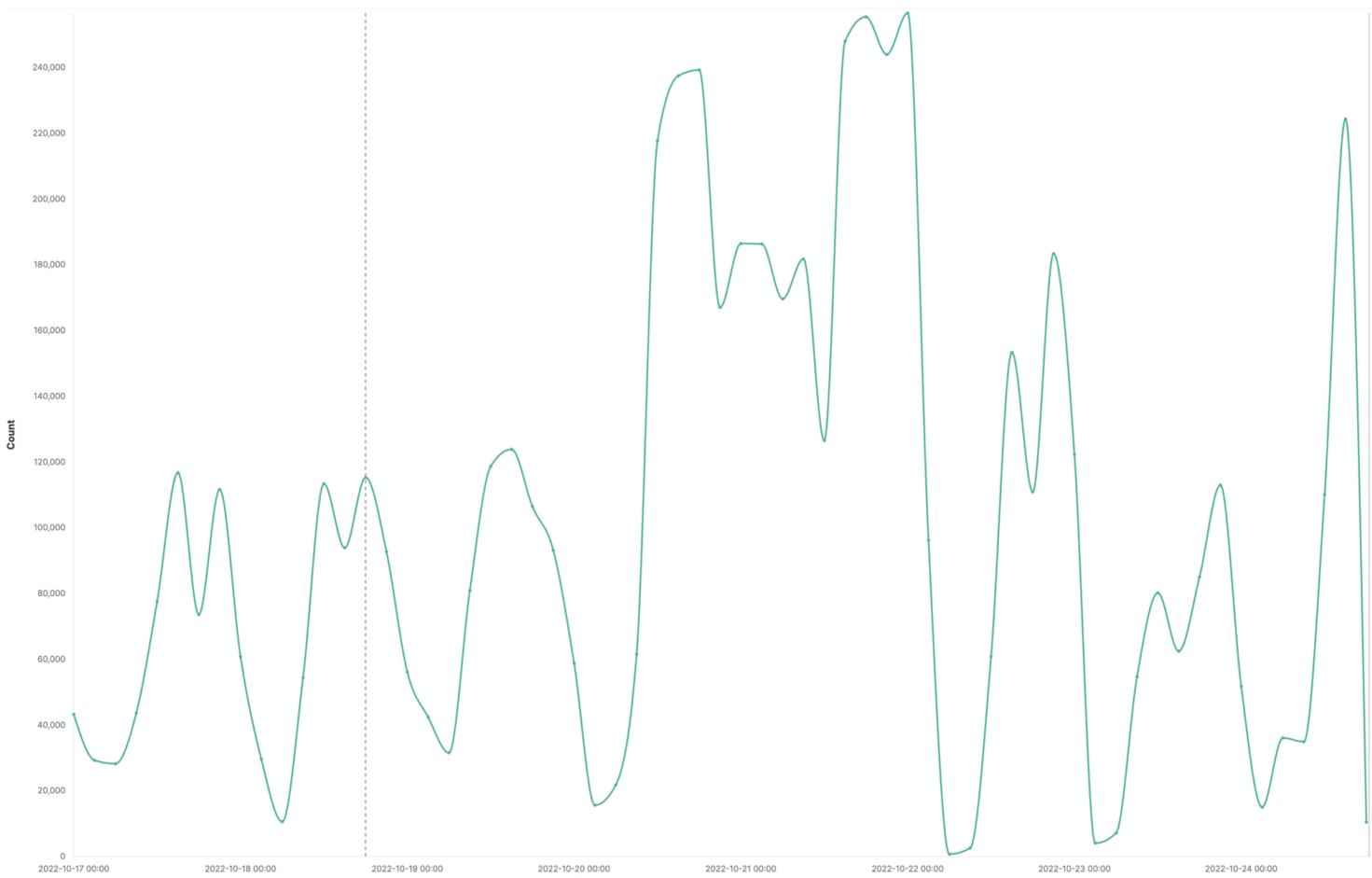
# DDoS Statistics

More C2 IPs cost more money, and it seems that the business is good as it has been very active launching ddos attacks. We have excerpted the data from June 29, 2022 to the present, and the attack trends and target area distribution are as follows.



We can see that the ddos attacks has been non-stop, and China and US have the most targets.

The time distribution of the attack instructions within 7 days is shown below, which shows that Fodcha launched DDoS attacks throughout  $7 * 24$  hours, without any obvious working time zone.



# Sample Analysis

We have divided the captured samples into four major versions, of which V1 and V2 have been analyzed in the previous blog, here we select the latest V4 series samples as the main object of analysis, their basic information is shown below.

```
MD5: ea7945724837f019507fd613ba3e1da9
ELF 32-bit LSB executable, ARM, version 1, dynamically linked (uses shared libs), sti
LIB: uclibc
PACKER: None
version: V4

MD5: 899047ddf6f62f07150837aef0c1ebfb
ELF 32-bit LSB executable, ARM, version 1 (SYSV), statically linked, stripped
Lib: uclibc
Packer: None
Version: V4.X
```

When Fodcha's Bot executes, it will first check the operating parameters, network connectivity, whether the "LD\_PRELOAD" environment variable is set, and whether it is debugged. These checks can be seen as a simple countermeasure to the typical emulator & sandbox.

When the requirements are met, it first decrypts the configuration information and print “snow slide” on the Console, then there are some common host behaviors, such as single instance, process name masquerading, manipulating watchdog, terminating specific port processes, reporting specific process information, etc. The following will focus on the decryption of configuration information, network communication, DDoS attacks and other aspects of Fodcha.

## Decrypting configuration information (Config)

Fodcha uses a side-by-side Config organization in V2.X and V3, and a structured Config organization in V4 and V4.X. The following figure clearly shows the difference.

```
int prepare_config() V2.X and V3
{
    int result; // r0
    __int16 v1[10]; // [sp+4h] [bp-14h] BYREF

    dword_25B34 = dec_str(&unk_1C734, 16, off_258E4, v1);
    word_25B38 = v1[0];
    dword_25B3C = dec_str(&unk_1C748, 12, off_258E4, v1);
    word_25B40 = v1[0];
    dword_25B44 = dec_str(&unk_1C758, 12, off_258E4, v1);
    word_25B48 = v1[0];
    dword_25B4C = dec_str(&unk_1C768, 20, off_258E4, v1);
    word_25B50 = v1[0];
    dword_25B54 = dec_str(&unk_1C780, 12, off_258E4, v1);
    word_25B58 = v1[0];
    dword_25B5C = dec_str(&unk_1C790, 12, off_258E4, v1);
    word_25B60 = v1[0];
    dword_25B64 = dec_str(&unk_1C7A0, 8, off_258E4, v1);
    word_25B68 = v1[0];
    dword_25B6C = dec_str(&unk_1C7AC, 8, off_258E4, v1);
    word_25B70 = v1[0];
    dword_25B74 = dec_str(&unk_1C7B8, 12, off_258E4, v1);
    word_25B78 = v1[0];
    dword_25B7C = dec_str(&unk_1C7C8, 8, off_258E4, v1);
    word_25B80 = v1[0];
    dword_25B84 = dec_str(&unk_1C7D4, 72, off_258E4, v1);
    word_25B88 = v1[0];
    dword_25B8C = dec_str(&unk_1C820, 44, off_258E4, v1);
    word_25B90 = v1[0];
    result = dec_str(&unk_1C850, 20, off_258E4, v1);
    ...
}
```

Multiple Configs

```
int prepare_config() V4 and V4.X
{
    int v0; // r6
    int v1; // r4
    char *v2; // r4
    void *v3; // r5
    int v4; // r0
    unsigned __int16 v5; // r1
    _BYTE *v6; // r0
    int result; // r0
    __int16 v8; // [sp+2h] [bp+0h] BYREF

    v0 = 0;
    v1 = 0;
    do
    {
        v2 = (char *)&unk_1F1A0 + 16 * v1;
        v3 = calloc((unsigned __int8)v2[6] | ((unsigned __int8)v2[7] << 8)) + 1, 1u);
        v4 = (unsigned __int8)v2[2] | ((unsigned __int8)v2[3] << 8);
        v5 = ((unsigned __int8)v2[4] | ((unsigned __int8)v2[5] << 8)) - v4;
        *((__DWORD *)v2 + 3) = v3;
        v6 = dec_str((int)&unk_1F2B0 + v4, v5, (int)aPjbinbbeasdddf, &v8);
        ++v0;
        result = sub_149F0(v3, v6, (unsigned __int8)v8 | (HIBYTE(v8) << 8));
        v1 = v0;
    }
    while ( v0 != 17 );
    return result;
}
```

Structural Configs

Although the organization of Config is different, their encryption methods are the same. As we can see by the constants referenced in the code snippet below, they use the xxtea algorithm with the key PJbiNbbeasddDfsc .

```

if ( (unsigned __int16)v4 != 1 )
{
    v21 = 0x9E3779B9 * (0x34u / (unsigned __int16)v4 + 6);
    if ( v21 )
    {
        v22 = &v13[v4];
        do
        {
            v23 = v39;
            v24 = (v21 >> 2) & 3;
            v25 = v22 - 2;
            do
            {
                v20 = v25[1]
                - (((4 * v20) ^ (*v25 >> 5)) + ((16 * *v25) ^ (v20 >> 3))) ^ ((v20 ^ v21)
                    + (*v25 ^ v17[v24 ^ v23-- & 3])));
                v25[1] = v20;
                --v25;
            }
            while ( v23 );
            v26 = (((4 * v20) ^ (v13[v39] >> 5)) + ((16 * v13[v39]) ^ (v20 >> 3))) ^ ((v20 ^ v21) + (v13[v39] ^ v17[v24]));
            v21 += 0x61C88647;
            v20 = *v13 - v26;
            *v13 = v20;
        }
        while ( v21 );
    }
}

```

After inversion, we wrote the following [IDAPYTHON](#) script to decrypt the configuration information.

```

# md5: ea7945724837f019507fd613ba3e1da9
# requirement: pip install xxtea-py
# test: ida7.6_python3

import ida_bytes
import xxtea

BufBase=0x1F2B0
ConfBase=0x0001F1A0
key=b"PJbiNbbeasddDfsc"
for i in range(17):
    offset=ida_bytes.get_word(i*16+ConfBase+2)
    leng=ida_bytes.get_word(i*16+ConfBase+4)-offset
    buf=ida_bytes.get_bytes(BufBase+offset,leng)
    print("index:%d, %s" %(i,xxtea.decrypt(buf,key)))

```

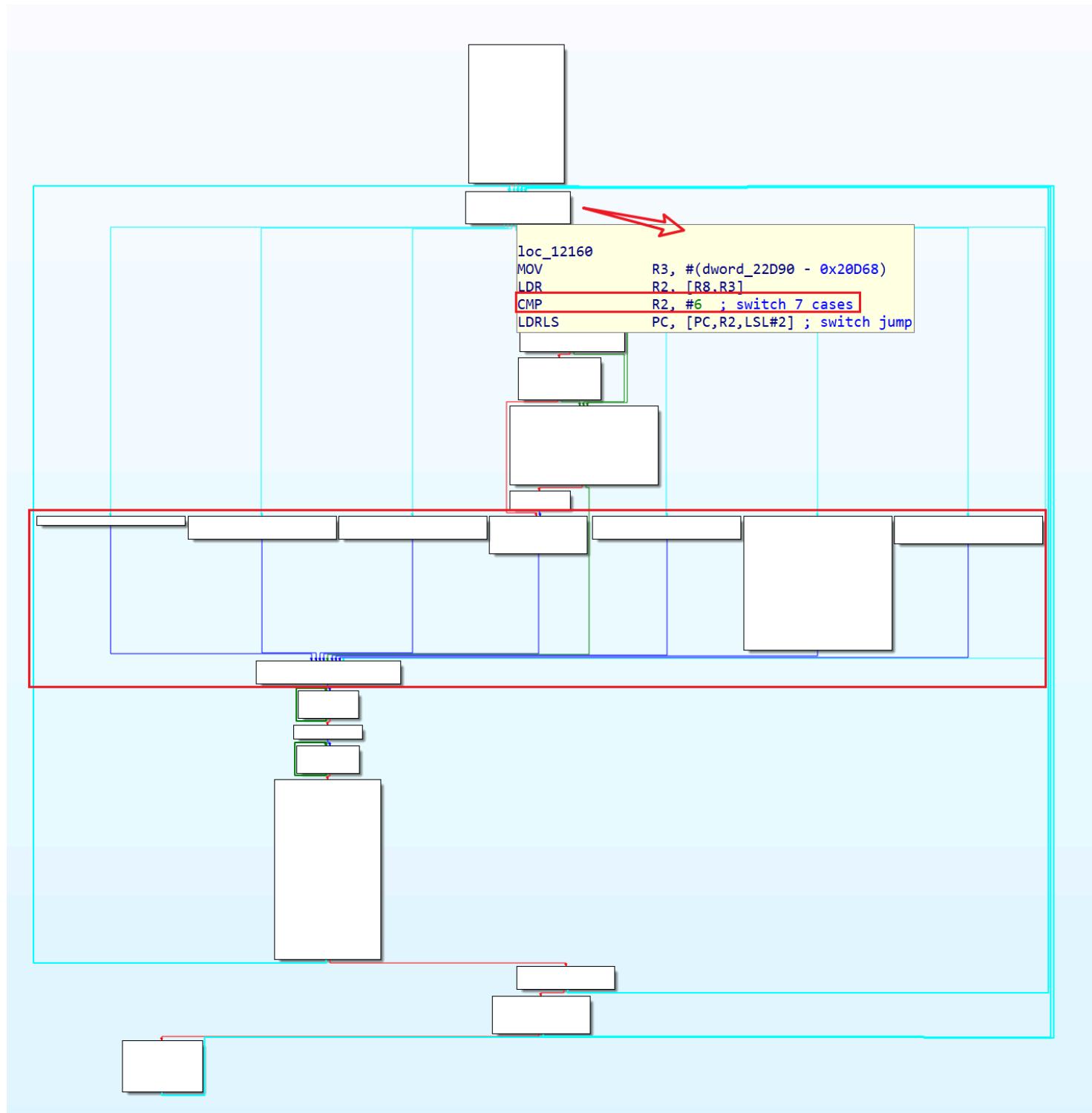
The decrypted Config information is shown in the following table. You can see that index 11 still retains the aforementioned "surrender" egg, index 12 is worth mentioning, as it is the reporter server address to which Fodcha reports some process-specific information.

INDEX	VALUE
0	snow slide

INDEX	VALUE
1	/proc/
2	/stat
3	/proc/self/exe
4	/cmdline
5	/maps
6	/exe
7	/lib
8	/usr/lib
9	.ri
10	GET /geoip/?res=10&r HTTP/1.1\r\nHost: 1.1.1.1\r\nConnection: Close\r\n\r\n
11	Netlab pls leave me alone I surrender
12	kvsolutions.ru
13	api.opennicproject.org
14	watchdog
15	/dev/
16	TSource Engine Query

## Network communication

Fodcha's network communication has a very fixed feature at the code level: a perpetual While loop, with switch-case processing at each stage, so that the CFG graphs generated by each version of Fodcha's network protocol processing functions are highly similar in IDA.



In summary, Fodcha's network communication goes through the following four steps.

1. Decryption C2
  2. DNS query
  3. Establishing communication
  4. Execute the command

## 0x1: Decrypting C2

Different versions of Fodcha support different types of C2. V2.X has only 1 group of OpenNIC C2; V3 & V4 have 1 group of OpenNIC C2 and 1 group of ICANN C2; and V4.X has the most, 1 group of OpenNIC C2 and 2 groups of ICANN C2, the following diagram shows the difference very clearly.

```

V2.X
do
{
    v3 = C2_GET(&unk_1F400, 1452, 0);
    v4 = (void *)v3;
    if ( v3 )
    {
        v5 = (unsigned __int8 *)DNS_QUERY(v3, 1);
        free(v4);
        if ( v5 )
        {
            LABEL_14:
                v17 = *((__DWORCD *)v5 + 1);
                v18 = sub_12EFC(v6, v7, v8, v9) % (unsigned int)*v5;
                *v1 = *(__DWORD *)v17 + 4 * v18;
                sub_13BAE(v5, v19);
                return 1;
            }
            v10 = sub_12EFC(v6, v7, v8, v9) % 5;
            sleep(v10 + 2);
        }
        v11 = v2++ == 10;
    }
    while ( !v11 );
    v12 = 1;
    do
    {
        v13 = C2_GET(&unk_1FDAC, 180, 0);
        v14 = (void *)v13;
        if ( v13 )
        {
            v5 = (unsigned __int8 *)DNS_QUERY(v13, 0);
            free(v14);
            if ( v5 )
                goto LABEL_14;
            v15 = sub_12EFC(v6, v7, v8, v9) % 5;
            sleep(v15 + 2);
        }
        v11 = v12++ == 10;
    }
    while ( !v11 );
    return 0;
}

```

```

V3&V4
do
{
    v3 = C2_GET(&unk_21F14, 1452, 0);
    v4 = v3;
    if ( v3 )
    {
        v5 = (unsigned __int8 *)DNS_QUERY(v3, 1);
        v6 = sub_15ASC(v4);
        if ( v5 )
            goto LABEL_11;
        v10 = sub_10918(v6, v7, v8, v9);
        sub_164C0(v10 % 5 + 2);
    }
}
while ( v2++ != 10 );
while ( 1 )
{
    do
    {
        v12 = C2_GET(&unk_224C0, 180, 0);
        v13 = v12;
    }
    while ( !v12 );
    v5 = (unsigned __int8 *)DNS_QUERY(v12, 0);
    v6 = sub_15ASC(v13);
    if ( v5 )
        break;
    while ( 1 )
    {
        v14 = sub_10918(v6, v7, v8, v9);
        sub_164C0(v14 % 5 + 2);
        v15 = C2_GET(&unk_224C0, 180, 0);
        v16 = v15;
        if ( v15 )
            break;
        v5 = (unsigned __int8 *)DNS_QUERY(v15, 0);
        v6 = sub_15ASC(v16);
        if ( v5 )
            goto LABEL_11;
    }
}
LABEL_11:
    v17 = *((__DWORD *)v5 + 1);
    v18 = sub_10918(v6, v7, v8, v9);
    sub_12B10(v18, *v5);
    *v1 = *(__DWORD *)v17 + 4 * v19;
    return sub_10CA8(v5);
}

```

```

V4.X

```

Although the types & numbers of C2 are different, their processing logic is almost the same as shown in the figure below.

Firstly, a C2 domain name is obtained through the C2\_GET function, and then the corresponding IP of C2 is obtained through the DNS\_QUERY function, where the first parameter of C2\_GET is the C2 ciphertext data, and the second parameter is the length, while the second parameter of DNS\_QUERY implies the type of C2.

```

do
{
    v3 = C2_GET(&unk_1CA6C, 1568, 0);
    v4 = v3;
    if ( v3 )
    {
        v5 = (unsigned __int8 *)DNS_QUERY(v3, 1);
        v6 = sub_17590(v4);
        if ( v5 )
            goto LABEL_11;
        v10 = sub_10924(v6, v7, v8, v9);
        sub_18D78(v10 % 5 + 2);
    }
}
while ( v2++ != 10 );
while ( 1 )
{
    do
    {
        v12 = C2_GET(&unk_1D090, 192, 0);
        v13 = v12;
    }
    while ( !v12 );
    v5 = (unsigned __int8 *)DNS_QUERY(v12, 0);
    v6 = sub_17590(v13);
    if ( v5 )
        break;
    while ( 1 )
    {
        v14 = sub_10924(v6, v7, v8, v9);
        sub_18D78(v14 % 5 + 2);
        v15 = C2_GET(&unk_1D090, 192, 0);
        v16 = v15;
        if ( !v15 )
            break;
        v5 = (unsigned __int8 *)DNS_QUERY(v15, 0);
        v6 = sub_17590(v16);
        if ( v5 )
            goto LABEL_11;
    }
}

```

Ciphertext of OpenNic C2

Flag 1 indict OpenNic

Ciphertext of ICCAN C2

FLAG 0 indict ICCAN

A valid C2 domain name can be obtained through C2\_GET, and its internal implementation can be divided into 2 steps.

- First, the C2 ciphertext data must be decrypted.
- Then they are constructed into a legitimate domain name.

## Decrypting C2 ciphertext data

The C2 ciphertext data uses the same encryption method as the configuration information, i.e. xxtea, and the key is also

**PJbiNbbeasddDfsc**. The OpenNic C2 data can be decrypted by the following simple IDAPYTHON script.

```
#md5: 899047DDF6F62F07150837AEF0C1EBFB
import xxtea
import ida_bytes
import hexdump
key=b"PJbiNbbeasddDfsc"
buf=ida_bytes.get_bytes(0x0001CA6C,1568) # Ciphertext of OpenNic C2
plaintext=xxtea.decrypt(buf,key)
print(plaintext)
```

The decrypted C2 data is shown below, you can see that the C2 data consists of 2 parts, the front is the domain names, the back is the TLDs, they are separated by the "/" symbol in the red box.

```
"14\nwehateyellow,s3x4fun,techsupporthelpars,s3x4fun,cookiedough,s3x4fun,parasjha,botnetskid,parasjha,cookiedough,parasjha,cooki edough,parasjha,parasjha,parasjha,cookiedough,cookiedough,yellowchinks,cookiedough,s3x4fun,wehateyellow,s3x4fun,s3x4fun,botnetsk id,s3x4fun,parasjha,wehateyellow,s3x4fun,botnetskid,wehateyellow,yellowchinks,parasjha,s3x4fun,botnetskid,botnetskid,s3x4fun,coo kiedough,s3x4fun,botnetskid,botnetskid,cookiedough,wearelegal,s3x4fun,botnetskid,parasjha,s3x4fun,wehateyellow,parasjha,cookiedough,s3x4fun,funnyyellowpeople,wehateyellow,botnetskid,s3x4fun,parasjha,wehateyellow,botnetskid,chinksdogeaters,wehateyellow,para sjha,s3x4fun,parasjha,blackpeeps,botnetskid,cookiedough,pepperfan,botnetskid,parasjha,parasjha,cookiedough,wehateyellow,botnetsk id,botnetskid,parasjha,wehateyellow,parasjha,wehateyellow,cookiedough,parasjha,wehateyellow,chinkchink,cookiedough,cookiedough,b otnetskid,cookiedough,wehateyellow,cookiedough,s3x4fun,parasjha,parasjha,wehateyellow,s3x4fun,botnetskid,peepeepoo,botnetskid,bo tnetskid,botnetskid,wehateyellow,parasjha,wehateyellow,cookiedough,s3x4fun,s3x4fun,botnetskid,respectkkk,s3x4fun,wehateyellow,we hateyellow,parasjha,botnetskid,s3x4fun,wehateyellow,wehateyellow,bladderfull,wehateyellow,botnetskid,parasjha,cookiedough,botnet skid,cookiedough,tsengtsing,cookiedough,botnetskid,wehateyellow,parasjha,obamalover,/pirate,at,geek,libre,ozz,geek,indy,libre,at,geek,libre,oss,libre,geek,libre,libre,at,geek,oss,libre,dyn,oss,geek,dyn,at,ozz,libre,at,geek,geek,geek,oss,dyn,libre,oss,geek,oss,oss,oss,at,oss,"
```

## Constructing a domain name

Fodcha has a specific domain name construction method, and the equivalent Python implementation is shown below.

```
# md5: 899047ddf6f62f07150837aef0c1ebfb
# requirement: pip install xxtea-py
# test: ida7.6_python3

import xxtea
import ida_bytes

def getcnt(length):
    cnt=1
    while True:
        cnt +=1
        calc=2
```

```

        for i in range(1,cnt):
            calc+=2+12*i%cnt

        if calc +cnt==length-1:
            return cnt

key=b"PJbiNbbeasddDfsc"
buf=ida_bytes.get_bytes(0x0001CA6C,1568) # Ciphertext of OpenNic C2
plaintext=xxtea.decrypt(buf,key)

domains,tlds=plaintext.split(b'/')
domainList=domains.split(b',')
tldList=tlds.split(b',')

cnt=getcnt(len(domainList))

print("-----There're %d C2-----" %cnt)
coff=2
for i in range(0,cnt):
    if i ==0:
        c2Prefix=domainList[i+coff]
    else:
        coff+=12*i %cnt+2
        c2Prefix=domainList[i+coff]
c2Tld=tldList[(cnt-i-1)*3]
print(c2Prefix + b'.' + c2Tld)

```

Taking the C2 data obtained above as input, the following 14 OpenNIC C2s are finally constructed.

```

techsuporthelpars.oss
yellowchinks.geek
yellowchinks.dyn
wearelegal.geek
funnyyellowpeople.libre
chinksdogeaters.dyn
blackpeeps.dyn
pepperfan.geek
chinkchink.libre
peepeepoo.libre
respectkkk.geek
bladderfull.indy
tsengtsing.libre
obamalover.pirate

```

Readers familiar with the ICANN domain name system may think at first glance that our decryption is wrong, because the ICANN domain name system does not support these TLDs, they would be "unresolvable", but in fact they are the domain names under the OpenNIC system, OpenNIC is independent of the OpenNIC, which supports the TLDs shown in the figure below. The domain names of OpenNIC cannot be resolved by common DNS (such as `8.8.8.8`, `101.198.198.198`) and must use the specified NameServer. Readers can check out [OpenNIC's official website](#) for more details.

## New Top-Level Domains!

OpenNIC's TLDs grant you access to a whole new space on the web. These domains can only be accessed **using our democratic nameservers**. Once you're in, click a button below to register your free domain!



Using the same method, we can get the following 4 ICANN C2s.

```
cookiemonsterboob[.]com  
forwardchinks[.]com  
doodleleching[.]com  
milfsfors3x[.]com
```

## 0X2: DNS lookup

When the C2 domain name is successfully obtained, Bot performs the domain name resolution through the function `DNS_QUERY`, its 2nd parameter is a FLAG, which implies the different processing of OpenNIC/ICANN C2, and the corresponding code snippet is shown below.

```

if ( flag_a2 )
{
    v13 = OpenNicDNS_via_API(&v60);
    if ( !v13 )
    {
        v13 = OpenNicDNS_via_HardCode();
        v14 = &v57[1010];
        v60 = v13;
    }
}
else
{
    v13 = ICCANDNS_via_HardCode();
    v15 = &v57[1010];
    v60 = v13;
}

```

**A2==1**

**A2==0**

It can be seen that there are 2 options for the resolution of OpenNIC C2.

- Option 1: Request from `api.opennicproject.org` through API interface to get nameserver dynamically

```
GET /geoip/?res=10&r HTTP/1.1
```

```
Host: 1.1.1.1
```

Fake

```
Connection: Close
```

```
HTTP/1.1 200 OK
```

```
Server: nginx
```

```
Date: Thu, 21 Jul 2022 18:01:52 GMT
```

```
Content-Type: text/plain; charset=UTF-8
```

```
Content-Length: 143
```

```
Connection: close
```

```
Vary: Accept-Encoding
```

```
Strict-Transport-Security: max-age=15768000
```

```
Allow: GET, HEAD
```

```
X-Upstream-Cache-Status: HIT
```

```
X-Cache-Key: geoip 180.163.225.13 res=10&r
```

```
144.24.181.253
```

```
91.217.137.37
```

```
94.247.43.254
```

```
194.36.144.87
```

```
103.1.206.179
```

```
130.61.117.123
```

```
51.77.149.139
```

```
195.10.195.195
```

```
94.16.114.254
```

```
94.16.114.254
```

- Option 2: Use the hard-coded nameserver shown in the figure below

```
int __fastcall OpenNicDNS_via_HardCode()
{
    int result; // r0

    switch ( sub_134C8() % 7u )
    {
        case 0u:
            result = 0xC3C30AC3;
            break;
        case 1u:
            result = 0x579024C2;
            break;
        case 2u:
            result = 0x4FDE5097;
            break;
        case 3u:
            result = 0x57C7854F;
            break;
        case 4u:
            result = 0x7B753D82;
            break;
        case 5u:
            result = 0x25E2B35F;
            break;
        default:
            result = 0x8B954D33;
            break;
    }
    return result;
}
```



195.10.195.195
194.36.144.87
151.80.222.79
79.133.199.87
130.61.117.123
95.179.226.37
51.77.149.139

For ICANN C2, there is only one option, i.e., use the hard-coded nameserver shown in the figure below.

```

int __fastcall ICCANDNS_via_HardCode()
{
    int result; // r0

    switch ( sub_134C8() % 7 )
    {
        case 1u:
            result = 0x4040808;
            break;
        case 2u:
            result = 0x1010101;
            break;
        case 3u:
            result = 0x2010101;
            break;
        case 4u:
            result = 0xDEDE43D0;
            break;
        case 5u:
            result = 0xDCDC43D0;
            break;
        case 6u:
            result = 0x9090909;
            break;
        default:
            result = 0x8080808;
            break;
    }
    return result;
}

```

8.8.4.4  
1.1.1.1  
1.1.1.2  
208.67.222.222  
208.67.220.220  
9.9.9.9  
8.8.8.8

The actual resolution of C2 `techsupporthelpars.oss`, for example, is reflected in the network traffic as follows.

Source	Destination	Protocol	Destination Port	Info
172.19.108.150	8.8.4.4	DNS	53	Standard query 0xe64d A api.opendnsproject.org
8.8.4.4	172.19.108.150	DNS	43953	Standard query response 0xe64d A api.opendnsproject.org CNAME api.opendns.org A 116.203.98.109
172.19.108.150	116.203.98.109	TCP		50892 → 80 [SYN] Seq=0 Win=64240 Len=0 MSS=1460 SACK_PERM TSval=1855618149 TSecr=0 WS=128
1.1.1.1	172.19.108.150	TCP		53 → 36574 [FIN, ACK] Seq=1 Ack=2 Win=65536 Len=0
172.19.108.150	1.1.1.1	TCP		36574 → 53 [ACK] Seq=2 Ack=2 Win=64256 Len=0
116.203.98.109	172.19.108.150	TCP		80 → 50892 [SYN, ACK] Seq=0 Ack=1 Win=28960 Len=0 MSS=1460 SACK_PERM TSval=1058099480 TSecr=1855618149 WS=128
172.19.108.150	116.203.98.109	TCP		50892 → 80 [ACK] Seq=1 Ack=1 Win=64256 Len=0 TSval=1855618393 TSecr=1058099480
172.19.108.150	116.203.98.109	HTTP		GET /geoip/?res=10&r HTTP/1.1
116.203.98.109	172.19.108.150	TCP		80 → 50892 [ACK] Seq=68 Win=29056 Len=0 TSval=1058099541 TSecr=1855618393
116.203.98.109	172.19.108.150	TCP		80 → 50892 [PSH, ACK] Seq=1 Ack=68 Win=29056 Len=266 TSval=1058099541 TSecr=1855618393 [TCP segment of a reassembled packet]
172.19.108.150	116.203.98.109	TCP		50892 → 80 [ACK] Seq=68 Ack=267 Win=64128 Len=0 TSval=1855618639 TSecr=1058099541
116.203.98.109	172.19.108.150	HTTP		HTTP/1.1 200 OK (text/plain)
172.19.108.150	138.197.140.189	DNS	53	Standard query 0xb03 A techsupporthelpars.oss
172.19.108.150	116.203.98.109	TCP		50892 → 80 [ACK] Seq=68 Ack=413 Win=64128 Len=0 TSval=1855618681 TSecr=1058099541
<				
> Frame 16: 211 bytes on wire (1688 bits), 211 bytes captured (1688 bits)				
> Ethernet II, Src: 02:42:a8:18:97:eb (02:42:a8:18:97:eb), Dst: 02:42:a8:18:97:eb (02:42:a8:18:97:eb)				
> Internet Protocol Version 4, Src: 116.203.98.109, Dst: 172.19.10.10				
> Transmission Control Protocol, Src Port: 80, Dst Port: 50892, Seq: 142(226), Dst Seq: 145(225)				
> [2 Reassembled TCP Segments (411 bytes): #14(266), #16(145)]				
> Hypertext Transfer Protocol				
> Line-based text data: text/plain (10 lines)				
>				

Why use OpenNIC / ICANN dual C2?

Fodcha's author has built a redundant OpenNIC / ICANN dual-C2 architecture, why did he do so?

From a C2 infrastructure perspective, after Fodcha was exposed, its C2 was added to various security lists. Quad9DNS (9.9.9.9), for example, had sent a tweet about Fodcha domain traffic spike



Quad9  
@Quad9DNS

...

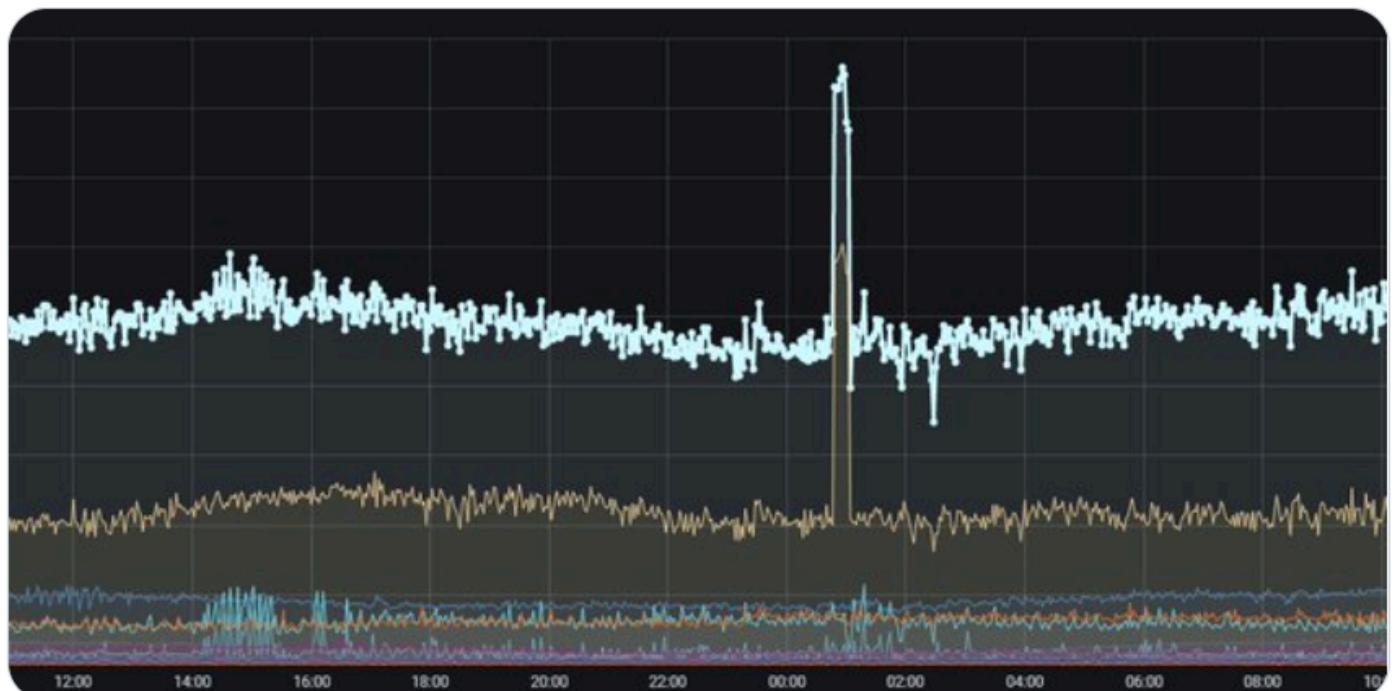
Wow so big spike in blocks this past couple of days and mostly due to **fridgexperts[.]cc**

**Fodcha v2 C2**

Quad9 has got you in the interim, but please Internet - please patch your routers, DVRs, and servers!

#Fodcha #cybersecurity #dns

[cyware.com/news/enemybot-...](http://cyware.com/news/enemybot-...)



8:49 PM · Apr 20, 2022 · TweetDeck

After Fodcha was cracked down, its author, when reselecting the C2 infrastructure, looked at the **DNS Neutrality** feature touted by OpenNIC to eliminate the

possibility of C2 being regulated & taken over.

At the same time, OpenNIC based C2 has its own problems, such as the NameServer of OpenNIC may not be accessible in some regions, or there are efficiency or stability problems in domain name resolution. For the sake of robustness, Fodcha authors re-added ICANN C2 as the backup C2 after V3 to form a redundant structure with the main C2.

## 0x3: Establishing communication

Fodcha Bot establishes a connection to C2 through the following code snippet, which has a total of 22 ports.

```
int __fastcall establist_connect(int a1, int a2, int a3, int a4)
{
    unsigned int v4; // r0
    int result; // r0
    sockaddr_in v6; // [sp+0h] [bp-20h] BYREF

    memset(&v6.sin_port, 0, 14);
    v6.sin_family = 2;
    v4 = sub_10924();
    v6.sin_port = HIBYTE(portlist[v4 % 0x15]) | ((unsigned __int8)portlist[v4 % 0x15] << 8);
    sub_FC54(&v6.sin_addr.s_addr);
    dword_26CB8[0] = socket(2, 1, 0);
    sub_15068((void *)dword_26CB8[0], (void *)4, (void *)0x800);
    result = connect(dword_26CB8[0], &v6, 16);
    dword_28C90 = 1;
    return result;
}

; _WORD portlist[22]
portlist    DCW 2348, 12381, 8932, 8241, 38441, 23845, 8745, 6463
            ; DATA XREF: sub_FD94+50↑o
            ; .text:off_FE58↑o
            DCW 7122, 1114, 6969, 1337, 4200, 3257, 7214, 2474, 4444
            DCW 2222, 3333, 5555, 24811, 0
```

Once the connection with C2 is successfully established, the Bot and C2 must go through 3 phases of interaction before communication is actually established.

- Stage 1: Bot requests the key&nonce of the chacha20 encryption algorithm from C2.
- Stage 2: Bot and C2 use the key&nonce from stage 1 for identity confirmation.

- Stage 3: Bot sends the encrypted upline & group information to C2.

To aid in the analysis, we ran the Bot sample within a restricted environment and used `fsdsad` as the packet string to generate the network traffic shown in the figure below, and the details of how this traffic was generated are described below.

00000000 06 00 00 f0 70 00 16 36 93 93 b7 27 5c 9a 2a 16 ....p..6 ...'\\.*.	Stage 1
00000010 09 d8 13 32 01 d2 69 1d 25 f3 42 00 32 ....2..i. %B.2	
00000000 80 6d 88 06 cd 54 60 d8 99 63 39 fb f7 ba c3 4b .m...T`..c9....K	
00000010 a1 e2 0f 79 28 72 ba 0e 05 d0 96 ad 92 a5 53 5e ...y(r... .....S^	
00000020 60 e5 5b 8d `.[.	
00000024 22 c8 a3 bb 31 9c 5b 25 12 e7 6a 47 24 18 f9 ee " 1 [% jG\$	
0000001D dc 23 c5 69 43 43 10 18 b6 12 62 48 1c e5 a2 19 .#.iCC...bH....	Stage 2
0000002D da 94 80 93 0f 08 71 4e 01 7e dc 56 bf 90 3d 32 .....qN ..~.V..=2	
0000003D ac 5d ae b8 31 4f 1b f7 e6 .]..10... .	
00000034 dc 23 c5 0d 5a 43 16 c0 72 88 16 cc 38 29 48 ae .#.ZC.. r...8)H.	
00000044 74 ad 43 f8 4c 1f 54 5b fe 77 5b 22 bc fa 31 6a t.C.L.T[ .w["..1j	
00000054 f4 75 ac d5 7e f4 86 6f 1c e1 5e .u...~.o ..^	
00000046 dd 23 c3 94 a5 43 2e e3 a1 a9 7d 32 e0 86 3b 36 .#...C... }2..;6	Stage 3
00000056 3f 6d 0d 2e f6 a8 f7 13 23 1d 9d 71 39 f5 fa 4b ?m..... #..q9..K	
00000066 7f aa 2b ..+	
00000069 be 50 a1 e0 ae 07 .P.....	
0000005F 01 00 16 6c 33 00 1d df e8 19 4b 45 b0 b1 50 38 ...13... KE..P8	CMD
0000006F 8d 28 3e 78 7c 4d cc 3e 2a 96 48 f1 88 78 95 2b .(>x M.> *.H..x.+	
0000007F 96 43 ef 07 .C..	
00000083 d8 23 c5 af c8 42 eb 68 6e 02 d2 fc 53 f4 eb fe .#...B.h n....S...	
00000093 f5 5f f0 8b c5 66 . ...f	

### Stage 1: Bot ---> C2 ,formatted as head(7 bytes) + body( random 20-40 bytes)

Bot actively sends an initialization message with `netstage=6` to C2, this message has the format of head+body, and the meaning of each field is shown below.

00000000 06 00 00 f0 70 00 16 36 93 93 b7 27 5c 9a 2a 16 ....p..6 ...'\\.*.	
00000010 09 d8 13 32 01 d2 69 1d 25 f3 42 00 32 ....2..i. %B.2	

### head

The length of head is 7 bytes, and the format is shown below.

06	----->netstage,1byte,06 means init
f0 70	----->tcpip checksum, 2byte,
00 16	----->length of body, 2 bytes

### checksum

## checksum

The checksum in head uses the tcp/ip checksum, which is calculated for the whole payload, and the original value of the checksum offset is `\x00\x00`, and the python implementation of the checksum is as follows.

```
def checksum(data):
    s = 0
    n = len(data) % 2
    for i in range(0, len(data)-n, 2):
        s+= ord(data[i]) + (ord(data[i+1]) << 8)
    if n:
        s+= ord(data[-1])
    while (s >> 16):
        s = (s & 0xFFFF) + (s >> 16)
        s = ~s & 0xffff
    return s

buf="\x06\x00\x00\x00\x00\x00\x00\x00\x16\x36\x93\x93\xb7\x27\x5c\x9a\x2a\x16\x09\xd8\x13\x32\x01\x02\x69
print(hex(checksum(buf)))

#hex(checksum(buf))
#0x70f0
```

## body

body is a randomly generated content, meaningless.

```
00000000  36 93 93 b7 27 5c 9a 2a 16 09 d8 13 32 01 d2 69
00000010  1d 25 f3 42 00 32
```

## Stage 1: C2-->Bot, 2 rounds

When C2 receives the message `netstage=6` from Bot, it will send 2 rounds of data to BOT.

00000000	80 6d 88 06 cd 54 60 d8 99 63 39 tb t/ ba c3 4b	.m....T`.. .c9....K
00000010	a1 e2 0f 79 28 72 ba 0e 05 d0 96 ad 92 a5 53 5e	...y(r.. .....S^
00000020	60 e5 5b 8d	`.[.
00000024	22 c8 03 bb 31 0c 5b 25 12 e7 6a 47 24 18 f9 ee	"...1.[% ..jG\$...

- The first round, 36 bytes , the original message is encrypted by xxtea and decrypted as the key of chacha20, with the length of 32bytes

```

import hexdump
import xxtea
key=b"PJbiNbbeasddDfsc"
keyBuf=bytes.fromhex("806d8806cd5460d8996339fbf7bac34ba1e20f792872ba0e05d096")
chaKey=xxtea.decrypt(keyBuf, key)
hexdump.hexdump(chaKey)

#chaKey
00000000: E6 7B 1A E3 A4 4B 13 7F 14 15 5E 99 31 F2 5E 3A
00000010: D7 7B AB 0A 4D 5F 00 EF 0C 01 9F 86 94 A4 9D 4B

```

- Second round, 16 bytes , the original message is encrypted by xxtea, decrypted as the nonce of chacha20, length 12bytes

```

import hexdump
import xxtea
key=b"PJbiNbbeasddDfsc"
nonBuf=bytes.fromhex("22c803bb310c5b2512e76a472418f9ee")
chaNonce=xxtea.decrypt(nonBuf, key)
hexdump.hexdump(chaNonce)

#chaNonce
00000000: 98 79 59 57 A8 BA 7E 13 59 9F 59 6F

```

## Stage 2: Bot--->C2, chacha20 encryption

Once Bot receives the key and nonce of chacha20, it sends the message

`netstage=4` to C2, this time the message is encrypted using chacha20, the key&nonce is obtained from the previous stage, the number of rounds encrypted is 1.

```

0000001D dc 23 c5 69 43 43 10 18 b6 12 62 48 1c e5 a2 19 .#.iCC.. ..bH....
0000002D da 94 80 93 0f 08 71 4e 01 7e dc 56 bf 90 3d 32 .....qN .~.V..=2
0000003D ac 5d ae b8 31 4f 1b f7 e6 .]..10.. .

```

We can decrypt the above traffic using the following python code that

```

from Crypto.Cipher import ChaCha20
cha=ChaCha20.new(key=chaKey,nonce=chaNonce)
cha.seek(64)
tmp=bytes.fromhex('dc23c56943431018b61262481ce5a219da9480930f08714e017edc56bf903d32ac')
rnd3=cha.decrypt(tmp)

```

The decrypted traffic is shown below, it still has the format of head (7 bytes) + body as described before, where the value of the netstage field of head is 04, which represents the authentication.

```
00000000: 04 00 00 FA 8C 00 22 64 4B 11 90 B6 4F 11 4B E6 ....."dK...0.K.  
00000010: 94 CA 70 1A CA 6F 81 38 FB 3F 9B 16 7B 92 5A 06 ..p..o.8.?..{.Z.  
00000020: B9 DB 7E DD 93 1B 81 FD 5E .....^
```

## Stage 2: C2 ---> Bot, chacha20 encryption

After receiving the authentication message from Bot, C2 also sends a message with `netstage=4` to Bot's data, also using chacha20 encryption, and the key,nonce,round number is the same as that used by Bot.

```
00000034 dc 23 c5 0d 5a 43 16 c0 72 88 16 cc 38 29 48 ae .#.ZC.. r...8)H.  
00000044 74 ad 43 f8 4c 1f 54 5b fe 77 5b 22 bc fa 31 6a t.C.L.T[."..1j  
00000054 f4 75 ac d5 7e f4 86 6f 1c e1 5e .u..~..o ..^
```

Using the same code as Bot to decrypt the traffic, we can see that its format is also head+body, and the value of netstage is also 04.

```
00000000: 04 00 00 9E 95 00 24 BC 8F 8B E4 32 6B DD A1 51 .....$....2k..Q  
00000010: 3A F3 B3 71 89 78 A4 2D 04 36 1C 62 78 F8 56 5E :..q.x.-.6.bx.V^  
00000020: E1 F3 7C B0 DC A0 1C 65 A4 1B B0 ...|....e...
```

After Bot and C2 send each other the message `netstage=4`, the chacha20 key&nonce representing stage 1 is recognized by both parties, and the authentication of each other is completed, and Bot enters the next stage to prepare to go online.

## Stage 3: Bot--->C2, 2 rounds, chacha encryption

Bot sends `netstage=5` message to C2 to indicate that it is ready to go online, and then reports its own grouping information to C2, these 2 rounds of messages also use chacha20 encryption.

- First round

00000046	dd 23 c3 94 a5 43 2e e3 a1 a9 7d 32 e0 86 3b 36	.#...C... ..}2..;6
00000056	3f 6d 0d 2e f6 a8 f7 13 23 1d 9d 71 39 f5 fa 4b	?m..... #..q9..K
00000066	7f aa 2b	..+

- Second round

00000069 be 50 a1 e0 ae 07

.P....

After the above two rounds of data decryption, we can see that the content of the group is exactly the preset `fsdsad`, which means our analysis is correct, so the Bot is successfully online and starts to wait for the execution of the command sent by C2.

00000000: 05 00 06 07 6A 00 1C 9F 5C AA 8F CC B3 72 D2 C9 ....j...\\....r...
00000010: 71 33 FD A7 33 CF 07 65 D9 5C DA 31 FD F7 9D 7F q3..3..e.\.1....
00000020: 6A 2C FB i..
00000000: 66 73 64 73 61 44 fsdsad

## 0x4: Execute command

Bot successfully online, support the netstage number, as shown in the figure below, the most important is the `netstage = 1` on behalf of the DDoS task, Fodcha reuse a large number of Mirai attack code, a total of `17` kinds of attack methods are supported.

```

if ( (unsigned __int8)netstage == 2 )
    exit(1);
if ( (unsigned __int8)netstage == 3 )
    return heartbeat(dword_20D68);
if ( (unsigned __int8)netstage != 1 )           if netstage == 1
{
    v5 = dword_20D68[0];
    dword_22D90 = 0;
LABEL_14:
    close(v5);
    v6 = rand() % 10;
    return sleep(v6 + 3);
}
return ddos_task(dword_20D68, &unk_20D6C, v4);
}

```

Take the following `DDos_Task` traffic (`netstage=01`) as an example.

0000005F	01 00 16 6c 33 00 1d df e8 19 4b 45 b0 b1 50 38	...13... .KE..P8
0000006F	8d 28 3e 78 7c 4d cc 3e 2a 96 48 f1 88 78 95 2b	.(>x M.> * .H..x.+
0000007F	96 43 ef 07	.C..
00000083	d8 23 c5 af c8 42 eb 68 6e 02 d2 fc 53 f4 eb fe	.#...B.h n....S..
00000093	f5 5f f0 8b c5 66	._...f

ddos instruction

The attack instructions are still encrypted using `chacha20`, and the decrypted instructions are shown below, which might ring a bell for readers who are familiar with `Mirai`.

```
00000000: 00 00 00 3C 07 01 xx 14 93 01 20 02 00 00 02 01
00000010: BB 01 00 02 00 01
```

The format and parsing of the above attack instructions are shown in the following table.

OFFSET	LEN (BYTES)	VALUE	MEANING
0x00	4	00 00 00 3c	Duration
0x04	1	07	Attack Vector, 07
0x05	1	1	Attack Target Cnt
0x06	4	xx 14 93 01	Attack Target, xx.20.147.1
0x0a	1	20	Netmask
0x0b	1	02	Option Cnt
0x0c	5	00 00 02 01 bb	OptionId 0, len 2, value 0x01bb ---> (port 443)
0x11	5	01 00 02 00 01	OptionId 1, len 2, value 0x0001---> (payload len 1 byte)

When Bot receives the above instruction, it will use the tcp message with a payload of 1 byte to conduct a DDoS attack on the target `xx.20.147.1:443`, which corresponds to the actual packet capture traffic.

Destination	Protocol	Destination Port	Info
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588180 Ack=3719926464 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588181 Ack=3719926465 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588182 Ack=3719926466 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588183 Ack=3719926467 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588184 Ack=3719926468 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588185 Ack=3719926469 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588186 Ack=3719926470 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588187 Ack=3719926471 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588188 Ack=3719926472 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588189 Ack=3719926473 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588190 Ack=3719926474 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588191 Ack=3719926475 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588192 Ack=3719926476 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588193 Ack=3719926477 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588194 Ack=3719926478 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588195 Ack=3719926479 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588196 Ack=3719926480 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588197 Ack=3719926481 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588198 Ack=3719926482 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588199 Ack=3719926483 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588200 Ack=3719926484 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588201 Ack=3719926485 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588202 Ack=3719926486 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588203 Ack=3719926487 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588204 Ack=3719926488 Win=5969408 Len=1 [
.20.147.1	TCP	443	40458 → 443 [PSH, ACK, CWR, AE, Reserved] Seq=588205 Ack=3719926489 Win=5969408 Len=1 [

# Misc

## 0x01: Racism

From some of the OpenNIC C2 constructs, it seems that the author of Fodcha is more hostile to people from some regions.

```
yellowchinks.geek
wearelegal.geek
funnyyellowpeople.libre
chinksdogeaters.dyn
blackpeeps.dyn
bladderfull.indy

wehateyellow
```

## 0x02: Ransom DDoS

Fodcha had the following string attached to the UDP attack command it issued.

```
send 10 xmr to 49UnJhpvRRxDXJHYcz0UEiK3EKCQZorZWaV6HD7axKGQd5xpUQeNp7Xg9RATFpL4u8dzP-
```

The corresponding attack traffic is shown below, the wallet address appears to be illegal, perhaps the operators behind Fodcha are experimenting with the attack-as-ransom business model, we will see how it evolve.

Destination	Protocol	Destination Port	Info
.241.237.245	UDP	13258	56855 → 13258 Len=1400
.241.237.61	UDP	13258	24614 → 13258 Len=1400
.197.100.215	UDP	13258	35840 → 13258 Len=1400
.197.96.162	UDP	13258	57099 → 13258 Len=1400
.233.151.207	UDP	13258	51056 → 13258 Len=1400
.233.24.50	UDP	13258	58783 → 13258 Len=1400
.241.237.245	UDP	13258	56855 → 13258 Len=1400
.241.237.61	UDP	13258	24614 → 13258 Len=1400
.197.100.215	UDP	13258	35840 → 13258 Len=1400
.197.96.162	UDP	13258	57099 → 13258 Len=1400
.233.151.207	UDP	13258	51056 → 13258 Len=1400
.233.24.50	UDP	13258	58783 → 13258 Len=1400
.241.237.245	UDP	13258	56855 → 13258 Len=1400
.241.237.61	UDP	13258	24614 → 13258 Len=1400
.197.100.215	UDP	13258	35840 → 13258 Len=1400
.197.96.162	UDP	13258	57099 → 13258 Len=1400
197.96.162	UDP	13258	57099 → 13258 Len=1400

Ransom Message From Fodcha

```

> Frame 310: 1442 bytes on wire (11536 bits), 1442 bytes radiated (11536 bits)
> Ethernet II, Src: 02:42:ac:13:64:70 (02:42:ac:13:64:70), Dst: 02:42:a8:10:97:eb
> Internet Protocol Version 4, Src: 172.19.100.11, Dst: 120.197.96.162
> User Datagram Protocol, Src Port: 57099, Dst Port: 13258
> Data (1400 bytes)

0020 60 a2 df 0b 33 ca 05 80 ef 7c 73 65 6e 64 20 31
0030 30 20 78 6d 72 20 74 6f 20 34 39 55 6e 4a 68 70
0040 76 52 52 78 44 58 4a 48 59 63 7a 6f 55 45 69 4b
0050 33 45 4b 43 51 5a 6f 72 5a 57 61 56 36 48 44 37
0060 61 78 4b 47 51 64 35 78 70 55 51 65 4e 70 37 58
0070 67 39 52 41 54 46 70 4c 34 75 38 64 7a 50 66 41
0080 6e 75 4d 59 71 73 32 4b 63 68 31 73 6f 61 66 35
0090 42 35 6d 64 66 4a 31 62 20 6f 72 20 77 65 20 77
00a0 69 6c 6c 20 73 68 75 74 64 6f 77 6e 20 79 6f 75
00b0 72 20 62 75 73 69 6e 65 73 73 00 00 00 00 00 00
00c0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

...3... |send 1
0 xmr to 49UhJhp
vRRxDXJH YczolEik
3EKCQZor ZWaV6HD7
axKGd5x pUqeNp7X
g9RATFpL 4u8dzPfA
nuMyqs2K ch1soaf5
B5mdfJ1b or we w
ill shut down you
r busine ss.....
..... .

```

# Contact us

Readers are always welcomed to reach us on [twitter](#) or email us to netlab[at]360.cn.

loC

C2

```
v1,v2:
folded[.]in
fridgexperts[.]cc
```

ICANN C2:

```
forwardchinks[.]com
doodleleching[.]com
cookiemonsterboob[.]com
milfsfors3x[.]com
```

OpenNIC C2:

yellowchinks.geek  
yellowchinks.dyn  
wearelegal.geek  
tsengtsing.libre  
techsupporthelpars.oss  
respectkkk.geek  
pepperfan.geek  
peepeepoo.libre  
obamalover.pirate  
funnyyellowpeople.libre  
chinksdogeaters.dyn  
chinkchink.libre  
bladderfull.indy  
blackpeeps.dyn  
91.206.93.243  
91.149.232.129  
91.149.232.128  
91.149.222.133  
91.149.222.132  
67.207.84.82  
54.37.243.73  
51.89.239.122  
51.89.238.199  
51.89.176.228  
51.89.171.33  
51.161.98.214  
46.17.47.212  
46.17.41.79  
45.88.221.143  
45.61.139.116  
45.41.240.145  
45.147.200.168  
45.140.169.122  
45.135.135.33  
3.70.127.241  
3.65.206.229  
3.122.255.225  
3.121.234.237  
3.0.58.143  
23.183.83.171  
207.154.206.0  
207.154.199.110  
195.211.96.142  
195.133.53.157  
195.133.53.148  
194.87.197.3  
194.53.108.94  
194.53.108.159  
194.195.117.167  
194.156.224.102

194.147.87.242  
194.147.86.22  
193.233.253.93  
193.233.253.220  
193.203.12.157  
193.203.12.156  
193.203.12.155  
193.203.12.154  
193.203.12.151  
193.203.12.123  
193.124.24.42  
192.46.225.170  
185.45.192.96  
185.45.192.227  
185.45.192.212  
185.45.192.124  
185.45.192.103  
185.198.57.95  
185.198.57.105  
185.183.98.205  
185.183.96.7  
185.143.221.129  
185.143.220.75  
185.141.27.238  
185.141.27.234  
185.117.75.45  
185.117.75.34  
185.117.75.119  
185.117.73.52  
185.117.73.147  
185.117.73.115  
185.117.73.109  
18.185.188.32  
18.136.209.2  
178.62.204.81  
176.97.210.176  
172.105.59.204  
172.105.55.131  
172.104.108.53  
170.187.187.99  
167.114.124.77  
165.227.19.36  
159.65.158.148  
159.223.39.133  
157.230.15.82  
15.204.18.232  
15.204.18.203  
15.204.128.25  
149.56.42.246  
139.99.166.217  
139.99.153.49  
139.99.142.215  
139.162.69.4

138.68.10.149  
137.74.65.164  
13.229.98.186  
107.181.160.173  
107.181.160.172

## Reporter

kvsolutions[.]ru  
icarlyfanss[.]com

## Samples

ea7945724837f019507fd613ba3e1da9  
899047ddf6f62f07150837aef0c1ebfb  
0f781868d4b9203569357b2dbc46ef10

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## Botnet



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## P2P 僵尸网络：回顾·现状·持续监测

缘起 P2P结构的网络比传统的 C/S结构具有更好的可扩展性和健壮性，这些优点很早就为 botnet的作者所认识到并被用到他们的僵尸网络中。从时间上看，2007年出现的Storm可以算是这方面的鼻祖，那时 botnet这种网络威胁刚为大众所知。Storm之后，陆续又有 Karen、ZeroAccess、GameOver、Hijime、mozi等20来种P2P botnet先后出现...



Nov 2, 2022

16 min



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Botnet

## 卷土重来的DDoS 狂魔：Fodcha僵尸 网络再次露出獠牙

背景 2022年4月13日，360Netlab首次向社区披露了Fodcha僵尸网络，在我们的文章发表之后，Fodcha遭受到相关部门的打击，其作者迅速做出回应，在样本中留下Netlab pls leave me alone I surrender字样向我们投降。本以为Fodcha会就此淡出江湖，没想到这次投降只是一个不讲武德的假动作，Fodcha的作者在诈降之后并没有停下更新的脚步...



Oct 27, 2022

23 min



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