



Parallel Hashing

Program for Undergraduate Research

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Contents

1	Introduction	i
2	Intel SIMD Instructions	i
3	Hashing Functions	ii
	3.1 Model 1	ii
	3.2 Model 2	ii
	3.3 Multiply-Shift Hash	ii
	3.4 MurMurHash3	iii
	3.5 CityHash	iv
	3.6 Tabular Hash	iv
4	Experiment and Results	iv
5	Future Work	iv
	5.1 HyperLogLog	iv
	5.2 Bloom Filter	iv
	5.3 Count-Min Sketch	iv
6	Conclusions	iv

1 Introduction

Hash functions are

2 Intel SIMD Instructions

SIMD is an instruction set available mostly on all current processors. In this project we used AVX(Advanced Vector Extension) and AVX2 instructions which are available for Intel processors since Sandy Bridge architecture. With AVX instructions, it is possible to process 128 bits of data in registers on parallel, with AVX2 this increased to 256 bits. The SIMD instructions we used on this project and their descriptions could be found in the following paragraph.

```
__m256i _mm256_add_epi32 ( __m256i a, __m256i b)
```

Add 8 packed to 256-bit side by side 32-bit integers in a and b, and store the results in dst.

```
__m256i _mm256_add_epi64 ( __m256i a, __m256i b)
```

Add 4 packed to 256-bit side by side 64-bit integers in a and b, and store the results in dst.

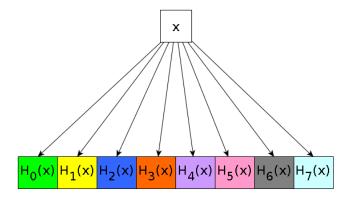
```
<u>__int32</u> _mm256_extract_epi32 ( __m256i a, const int b)
```

Add 4 packed to 256-bit side by side 64-bit integers in a and b, and store the results in dst.

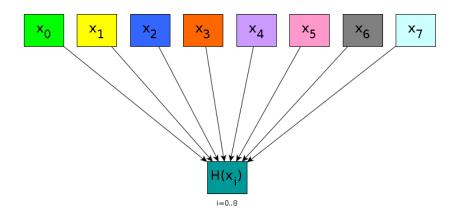
3 Hashing Functions

We have developed our work under 2 models; Model 1, one data multiple hash, where we generate multiple hash values from a single data sample, and Model 2, one data one hash, where we generate only one hash value using the same random seed for a single data sample. As our hash functions, we used Multiply-Shift Hash, MurMurHash3 and Tabular Hash.

3.1 Model 1



3.2 Model 2



3.3 Multiply-Shift Hash

- 1: **function** MULTIPLY-HASH(x)
- 2: 2 randomly sampled 64 bit integers m_a, m_b ; return $(m_a * (uint_64) * x + m_b) >> 32$

Conclusions _____ iii

3.4 MurMurHash3

```
1: function MURMUR3(key, len, seed)
      c1 = 0xcc9e2d51
 2:
      c2 = 0x1b873593
 3:
 4:
      r1 = 15
      r2 = 13
 5:
      m = 5
 6:
      n = 0xe6546b64
 7:
      hash = seed
 8:
 9:
      for each fourByteChunk of key do
          k \leftarrow fourByteChunk
10:
          k = k * c1
11:
          k = (k \text{ ROL } r1)
12:
          k = k * c2
13:
          hash = hash XOR k
14:
          hash = (hash ROL r2)
15:
          hash = hash * m + n
16:
                                     ▶ Endian swapping is only necessary on big-endian machines.
      for each remaining byte in key do
17:
          remainingBytes = SwapEndianOrderOf(remainingBytesInKey)
18:
          remainingBytes = remainingBytes * c1
19:
          remainingBytes = (remainingBytes ROL r1)
20:
          remainingBytes = remainingBytes * c2
21:
22:
          hash = hash XOR remainingBytes
      hash = hash XOR len
23:
      hash = hash XOR (hash SHR 16)
24:
      hash = hash * 0x85ebca6b
25:
       hash = hash XOR (hash SRH 13)
26:
      hash = hash * 0xc2b2ae35
27:
      hash = hash XOR (hash SHR 16)
```

3.5 Tabular Hash

4 Experiment and Results

5 Future Work

- 5.1 HyperLogLog
- 5.2 Bloom Filter
- 5.3 Count-Min Sketch

6 Conclusions