

Functional and Logic Programming

Lecture 9 — Algebraic data types

Pedro Vasconcelos
DCC/FCUP

2025

Type synonyms

We can define a new name for an existing type using a **type synonym** declaration.

Example from the Prelude:

```
type String = [Char]
```

Type synonyms (cont.)

Type synonym declarations are used to improve program readability.

Example: in a 2D game we might define

```
type Pos = (Int,Int)      -- x, y coordinates  
type Ship = Pos  
type Obstacles = [Pos]    -- list of positions
```

Then we can write

```
colides :: Ship → Obstacles → Bool
```

instead of

```
colides :: (Int,Int) → [(Int,Int)] → Bool
```

Type synonyms (cont.)

Type synonym declarations can also have parameters.

Example: association lists of key-value pairs.

```
type Assoc k v = [(k,v)]
```

```
ages :: Assoc String Int
```

```
ages = [("Anne", 21), ("Bob", 18), ("Carol", 19)]
```

```
emails :: Assoc String String
```

```
emails = [("Anne", "anne@domain.org"),  
          ("Bob", "bob@example.biz"),  
          ("Carol", "carol@bigtech.com")]
```

Type synonyms (cont.)

Type synonyms may be used in other synonym declarations.

```
type Pos = (Int,Int)
type Ship = Pos           -- OK
type Obstacles = [Pos]    -- OK
```

But they cannot be used recursively:

```
type Stream = (Int,Stream)  -- ERROR
```

For the later, we need to declare a **new data type**.

Declaring new data types

We can define new data types using a `data` declaration.

Example:

```
data Season = Winter | Spring | Summer | Autumn
```

defines a new type named `Season` with four possible values (Winter, Spring, Summer, Autumn).

Declaring new data types (cont.)

- ▶ The `data` declares enumerates all possible cases for the new type
- ▶ `Winter`, `Spring`, etc. are called the **constructors** of `Season`
- ▶ Constructors must be unique (i.e. you can't use the same name constructor in two distinct types or two alternatives of the same type)
- ▶ Names of types and constructors must start with an upper-case letter

Declaring new data types (cont.)

We can use this new type just like pre-defined language types.

```
next :: Season → Season
next Winter = Spring
next Spring  = Summer
next Summer  = Autumn
next Autumn  = Winter
```

```
isWarm :: Season → Bool
isWarm Summer = True
isWarm Autumn  = True
isWarm _       = False
```

```
allSeasons :: [Season]
allSeasons = iterate next Winter
-- [Winter, Spring, Summer, Autumn, Winter ...]
```


Constructors with parameters

The constructors can may also take any number of parameters.

Example:

```
data Shape = Circ Float          -- radius
           | Rect Float Float    -- width, height
```

```
square :: Float → Shape
square h = Rect h h
```

```
area :: Shape → Float
area (Circ radius)      = pi*radius^2
area (Rect width height) = width*height
```

Constructors with parameters (cont.)

- ▶ Constructors may take any number of parameters
- ▶ The parameters may be of distinct types
- ▶ We use constructors both to **build** values and as patterns to **decompose** values

```
Circ 1.0 :: Shape
```

```
Rect 2.0 3.0 :: Shape
```

```
area :: Shape → Float
```

```
area (Circ radius)      = pi*radius^2
```

```
area (Rect width height) = width*height
```

```
-- alternative, using case
```

```
area shape = case shape of
```

```
  Circ radius      → pi*radius^2
```

```
  Rect width height → width*radius
```

New data types with parameters

Data declarations can also have type parameters.

Example (from the Prelude):

```
data Maybe a = Nothing | Just a
```

Maybe is useful to represent an optional value or result:

```
safeHead :: [a] → Maybe a  
safeHead []      = Nothing  
safeHead (x:xs) = Just x
```

Maybe is often used instead of a nullable value (which does not exist in Haskell).

Algebraic data types

- ▶ Data declarations allow us to build new types from simpler ones using **disjunctions** and **conjunctions**
- ▶ These behave like sums and products, hence the designation of **algebraic**
- ▶ Pairs are the special case of a **binary product**:

```
data Pair a b = MkPair a b
  -- isomorphic to (a,b)
```

- ▶ **Either** is the special case of a **binary sum**:

```
data Either a b = Left a | Right b
  -- defined in the Prelude
```

Typeclasses

- ▶ By default, new types do not implement the type classes `Eq`, `Ord` or `Show`
- ▶ We can ask for automatic definitions using a `deriving` clause in the declaration:

```
data Shape = Circ Float
           | Rect Float Float
           deriving (Eq, Show)
```

- ▶ The derived operations are implemented structurally:

```
> show (Circ 2)
"Circ 2.0"
> Rect 1 (1+1) == Rect 1 2
True
```

Recursive types

Unlike type synonyms, data declarations may be recursive.

For example, we can define the type of lists as follows:

```
data List a
  = Empty                -- empty list []
  | Cons a (List a)      -- add an element (:)

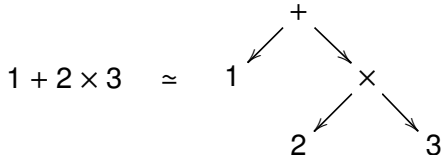
```

This type is **isomorphic** to `[a]`, i.e. you can convert any value of this type into a value of `[a]` and vice-versa (exercise).

Syntax Trees

We can use a recursive type to represent expressions as a **syntax tree**: the operators are the nodes in the tree and the constants are the leaves.

Example



Syntax Trees (cont.)

The type of expressions is:

```
data Expr
  = Val Int           -- constant
  | Add Expr Expr     -- add 2 expressions
  | Mult Expr Expr    -- multiply 2 expressions
```

The previous example expression is:

```
Add (Val 1) (Mult (Val 2) (Val 3))
```


Recursion over trees

Because syntax trees are recursive types, it is natural to define functions over syntax trees using pattern matching and recursion.

Evaluation

The **evaluation** of an expression should produce its value.

Example:

```
eval (Add (Val 1) (Mult (Val 2) (Val 3))) == 7
```

Evaluation (cont.)

We can define `eval` using recursion.

```
eval :: Expr → Int
eval (Val n) = n
eval (Add e1 e2) = eval e1 + eval e2
eval (Mult e1 e2) = eval e1 * eval e2
```

Try this using Haskellite:

```
eval (Add (Val 1) (Mult (Val 2) (Val 3)))
```

Pretty-printing

Pretty-printing transforms a syntax tree into a text string suitable for users.

Example:

```
pretty (Add (Val 1) (Mult (Val 2) (Val 3)))  
== " (1+ (2*3)) "
```

(Note that this is *not* what `show` does.)

Pretty-printing (cont.)

Again we can define `pretty` recursively.

```
pretty :: Expr → String
pretty (Val n)
  | n >= 0      = show n
  | otherwise   = "(" ++ show n ++ ")"
pretty (Add e1 e2)
  = "(" ++ pretty e1 ++ "+" ++ pretty e2 ++ ")"
pretty (Mult e1 e2)
  = "(" ++ pretty e1 ++ "*" ++ pretty e2 ++ ")"
```

The above definition may introduce redundant parenthesis; removing some of them is left as an exercise.

Parsing

Parsing transforms unstructured text into a syntax tree.

Example:

```
parse "1+2*3"  
== (Add (Val 1) (Mult (Val 2) (Val 3)))
```

Parsing (cont.)

- ▶ Unlike pretty-printing, parsing is inherently a **partial function**: not all text strings are valid syntax trees
- ▶ Example: " $1++2*$ (" does not correspond a valid expression
- ▶ This means we have to allow for the possibility of failure
- ▶ We will dedicate a future lecture to **parsing using combinators**

Newtype declarations

- ▶ Newtype declarations allow us to define a new type that is isomorphic to another existing type
- ▶ A newtype must have a single constructor and single field, e.g.:

```
newtype UserID = UID Int
```

- ▶ Unlike a type synonym, the newtype is distinct from the old type
- ▶ Unlike a data declaration, the newtype has the same runtime representation as the old type
- ▶ Unlike a data declaration, the newtype constructor is strict (i.e. not lazy); `UID undefined = undefined`

Record syntax

- ▶ We can use a data type with a single constructor to define a custom product type
- ▶ We can also use record syntax to define names for each field
- ▶ Example: type for user contact info

```
data Contact = Contact
    { name :: String,
      , email :: String
      , address :: String
    }
```

(NB: `Contact` is both a type name and a constructor name.)

Record syntax (cont.)

We can use the field names:

- ▶ to create a contact;
- ▶ as functions to get the field values;
- ▶ to create a new record.

```
>>> let p = Contact { name="Pedro"  
                      , address="My home"  
                      , email="my@email.com" }
```

```
>>> name p  
"Pedro"
```

```
>>> address p  
"My home"
```

```
>>> p { email= "new@email.com" }  
Contact { name="Pedro", address="My home",  
email="new@email.com" }
```