

Calculating the Cell Polynomial of Catacondensed Polycyclic Hydrocarbons

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In the paper we establish a simple algorithm of low complexity for calculating the so-called cell polynomials for a catacondensed polycyclic unsaturated hydrocarbon H , which enables the Kekulé structure count (KSC) and the algebraic structure count (ASC) of H to be calculated.

1. INTRODUCTION

During the past 2 decades applications of graph theory to polycyclic aromatic hydrocarbons have raised great interest in resonance theory. In this respect, Herndon's resonance theory^{1,2} and the conjugated circuit model introduced by Randić^{3,4} are of considerable importance. The sextet polynomial found by Hosoya and Yamaguchi⁵ allows a systematic combinatorial enumeration of Kekulé structures of (catacondensed) aromatic hydrocarbons. This polynomial was shown to possess a number of interesting properties and to reflect Clar's resonant sextet theory.⁶ Various further developments of the sextet polynomial concept can be found, e.g., in refs 7–9.

In the present paper we establish a simple algorithm for calculating the so-called cell polynomials for a catacondensed polycyclic unsaturated hydrocarbon H , which enable the Kekulé structure count (KSC) and the algebraic structure count (ASC) of H to be calculated. In connection with the ASC of H see, e.g., the papers by Wilcox,¹⁰ Herndon,¹¹ Klein et al.,¹² Dias,¹³ and refs 14 and 15.

2. DEFINITIONS AND NOTATION

Let ϵ denote the Euclidean plane and let \underline{G} be the set of all two-connected finite planar graphs (without loops and multiple edges). Let $G = (V, E) \in \underline{G}$ with vertex set $V = V(G)$ and edge set $E = E(G)$, and let $n = n(G) = |V|$, $m = m(G) = |E|$ denote the numbers of vertices and edges of G , respectively. Graph G_ϵ is an embedding of G into ϵ . G_ϵ subdivides ϵ in $m - n + 1 =: c = c(G)$ finite open domains D_i , $i = 1, 2, \dots, c$, the infinite open domain D_0 and the union of the boundaries of all these domains (Figure 1).

The boundaries of the D_i are denoted by $B_i = B(D_i)$, $i = 1, 2, \dots, c$; B_0 is also called the *contour* (or periphery) of G_ϵ . G_ϵ is called a *map* denoted by $M = M(G)$. Cell C_i is the union $D_i \cup B_i$, $i = 1, 2, \dots, c$. Note that $B(C_i) =: B_i$. The *inner dual* $D = D(M)$ of M is the dual of M without the vertex that corresponds to the infinite domain. Map M is a *contour map* if $D(M)$ is a tree (Figure 1).

A *perfect matching* (PM) or linear factor of a graph G is a set of pairwise disjoint edges that cover all vertices of G (Figure 5).

Case $c(M) = 0$. Then $M = \emptyset$ is the zero map (without edges and vertices). By convention, map \emptyset has precisely one PM.

Case $c(M) = 1$. Then M has exactly one cell which is called an *isolated cell*.

Case $c(M) > 1$. Edge $e \in E(M)$ is called a *contour edge* or an *internal edge* of M if the whole edge, or only its endpoints,

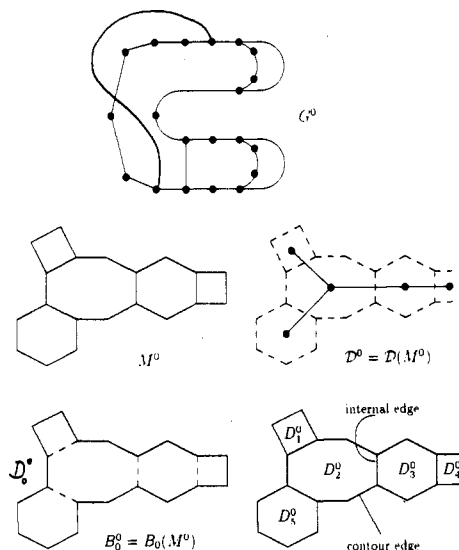


Figure 1.

lie on B_0 (Figure 1). Cell C is an end cell of M if C corresponds to an end vertex of $D(M)$ (in particular, an isolated cell is an end cell).

Let \underline{M} denote the set of all contour maps.

Map $M \in \underline{M}$ is *bipartite* if and only if $V = V(M)$ can be divided into two disjoint subsets $\tilde{V} = \tilde{V}(M)$ and $\bar{V} = \bar{V}(M)$ ($V = \tilde{V} \cup \bar{V}$, $\tilde{V} \cap \bar{V} = \emptyset$) such that vertices from the same subset are never adjacent.

Let \underline{M}^b denote the set of all bipartite contourmaps.

Note that for every $M \in \underline{M}^b$ $\tilde{n} = \bar{n}$ and M has a PM.

Observation 1. Every $M \in \underline{M}^b$ with $c(M) > 1$ has at least two end cells.

Observation 2. Every $M \in \underline{M}^b$ is the last of a (finite) sequence of bipartite contour maps $\{M_i\}$, $i = 1, 2, \dots, c$, where M_1 is an isolated cell and for $i = 2, 3, \dots, c$ we obtain M_i by adding cell C_i (C_i is an end cell of M_{i-1}) to M_{i-1} (see Figure 2).

Any plane image (i.e., an embedding in ϵ) of a (bipartite) planar graph G is called a (bipartite) *pattern* (of G). The cells of a pattern are defined in an analogous way as above.

Let $\underline{\mathcal{M}}^b$ denote the set of all (bipartite) subpatterns of all maps from \underline{M}^b .

3. DEFINITION OF THE CELL POLYNOMIAL OF A BIPARTITE PATTERN

For $\mathcal{M} \in \underline{\mathcal{M}}^b$ let $\underline{C} = \underline{C}(\mathcal{M}) = \{C_{i_1}, C_{i_2}, \dots, C_{i_N}\}$ and $N = N(\mathcal{M}) = \{i_1, i_2, \dots, i_N\}$ denote the set of all cells and of all cell indices (labels), respectively (note that if $\underline{C} = \emptyset$, then $N = \emptyset$).

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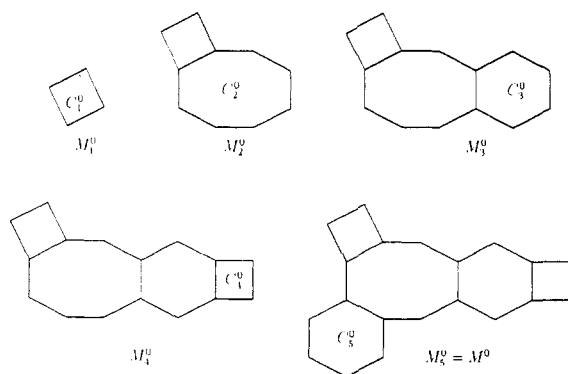


Figure 2.

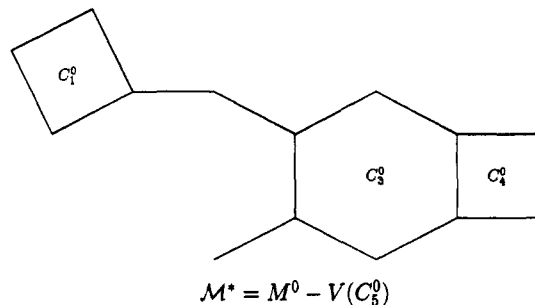


Figure 3.

To every subset $I \subseteq N$ there corresponds a cell set $\underline{C}_I = \{C_i | C_i \in \underline{C} \text{ and } i \in I\}$ of \mathcal{M} ; I is called an *index set* of \mathcal{M} .

The pattern \mathcal{M}_I is obtained from \mathcal{M} by deleting in \mathcal{M} all vertices of all cells from \underline{C}_I and all edges which are incident with these vertices. Index set $I \subseteq N$ is *feasible* (in \mathcal{M}) if (i) any two cells $C', C'' \in \underline{C}_I$ are disjoint ($B(C') \cap B(C'') = \emptyset$) and (ii) \mathcal{M}_I has a PM. Note that $\mathcal{M}_\emptyset = \mathcal{M}$; this implies that $I = \emptyset$ is feasible if and only if \mathcal{M} has a PM.

To every cell $C_i \in \underline{C}(\mathcal{M})$ is assigned a weight $w_i = w(C_i)$ which is an element of an (algebraic) ring (here it suffices to assume that w_i is a real number). Cell set \underline{C}_I of \mathcal{M} has *weight*

$$w(\underline{C}_I) = \begin{cases} \prod_{i \in I} w_i, & \text{if } I \text{ is feasible} \\ 0 & \text{otherwise} \end{cases}$$

(this includes $w(\underline{C}_\emptyset) = 1$ if \emptyset is feasible). The *cell polynomial* $f_{\mathcal{M}} = f_{\mathcal{M}}(w_1, w_2, \dots, w_c)$ of \mathcal{M} is defined as

$$f_{\mathcal{M}} = \sum_{I \subseteq N(\mathcal{M})} w(\underline{C}_I)$$

For example, consider the pattern $\mathcal{M}^* =: \mathcal{M}^0 - V(C_5^0)$, given in Figure 3.

Here $N(\mathcal{M}^*) = \{1, 3, 4\}$; the only feasible sets are \emptyset and $\{1\}$; thus $f_{\mathcal{M}^*} = 1 + w_1$.

Note that for any bipartite pattern \mathcal{M} the cell polynomial $f_{\mathcal{M}}$ can be defined as described above.

4. CALCULATION OF THE CELL POLYNOMIAL OF A BIPARTITE PATTERN WHICH IS PART OF A BIPARTITE CONTOUR MAP

Let $\mathcal{M} \in \underline{\mathcal{M}}^b$ with cell set \underline{C} , where cell $C_i \in \underline{C}$ has weight $w_i = w(C_i)$. The cell polynomial $f_{\mathcal{M}}$ can be calculated applying the following recursive procedure.

Algorithm f:

(f.0) If \mathcal{M} consists of a single (isolated) vertex,

put $f_{\mathcal{M}} = 0$;

(f.1) $f_\emptyset = 1$;

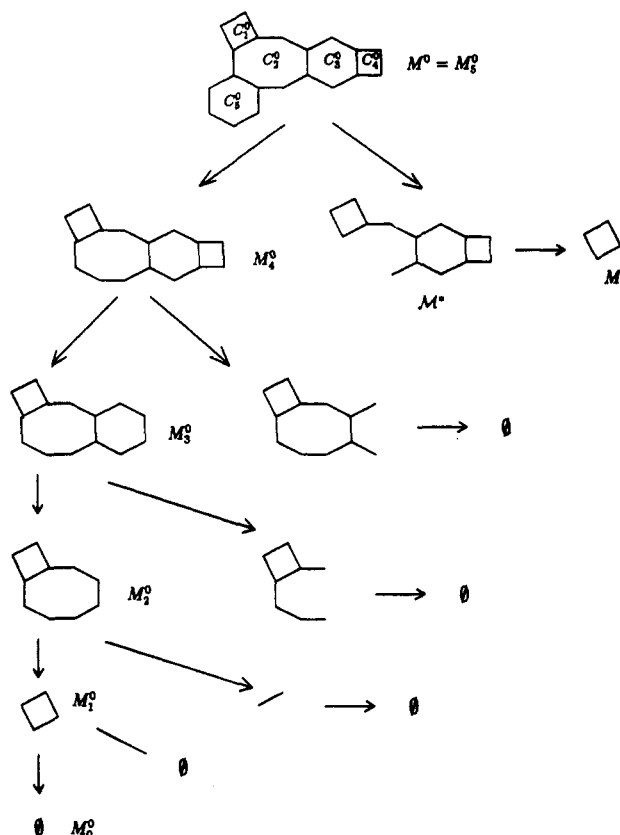


Figure 4.

(f.2) If \mathcal{M} has components $\mathcal{M}_1, \mathcal{M}_2, \dots, \mathcal{M}_p$,

put $f_{\mathcal{M}} = f_{\mathcal{M}_1} f_{\mathcal{M}_2} \dots f_{\mathcal{M}_p}$;

(f.3) If $\mathcal{M} \in \underline{\mathcal{M}}^b$, \mathcal{M} is connected, and $c(\mathcal{M}) > 1$, consider the following cases.

(i) \mathcal{M} has no cut vertex ($\mathcal{M} \in \underline{\mathcal{M}}^b$).

Then (because of observation 2) \mathcal{M} can be obtained from $\mathcal{M}_c \in \underline{\mathcal{M}}^b$ by adding cell C to \mathcal{M}_c (C is an end cell of \mathcal{M}).

Delete in \mathcal{M} all vertices of $V(C)$ (and all edges which are incident with these vertices). Denote the resulting pattern by \mathcal{M}^c . Put

$$f_{\mathcal{M}} = f_{\mathcal{M}_c} + w(C)f_{\mathcal{M}_c}$$

(ii) \mathcal{M} has a cut vertex ($\mathcal{M} \in \underline{\mathcal{M}}^b$).

(ii.1) \mathcal{M} has a hanging edge $(u, v) \in E(\mathcal{M})$.

Delete vertices u, v and all edges which are incident with them. Denote the resulting pattern by \mathcal{M}' . Put

$$f_{\mathcal{M}} = f_{\mathcal{M}'}$$

(ii.2) \mathcal{M} has no hanging edge.

Then there exists an induced subpattern \mathcal{M}'' of \mathcal{M} such that $\mathcal{M}'' \in \underline{\mathcal{M}}^b$, \mathcal{M}'' contains exactly one cut vertex of \mathcal{M} , and \mathcal{M}'' is contained in no other submap of \mathcal{M} with these properties. Delete in \mathcal{M} all vertices of \mathcal{M}'' (and all edges incident with them); this results in a pattern which we call \mathcal{M}''' . Put

$$f_{\mathcal{M}} = f_{\mathcal{M}''} f_{\mathcal{M}'''}$$

By successively applying f.0–f.3, eventually the polynomial $f_{\mathcal{M}}$ of $\mathcal{M} \in \underline{\mathcal{M}}^b$ is found. For example consider the map $\mathcal{M}^0 \in \underline{\mathcal{M}}^b$ of Figure 2. \mathcal{M}^0 has five cells C_i^0 with weights $w_i^0 = w(C_i^0)$, $i = 1, 2, 3, 4, 5$. Here we find (see Figure 4)

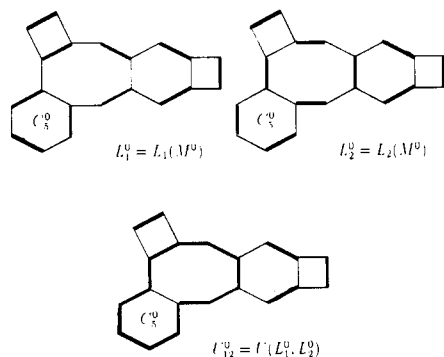


Figure 5.

$$f_{M^0} = f_{M_3^0} = f_{M_4^0} + w_3^0 f_{M^*} = \dots = f_{M_4^0} + w_3^0 f_{M_1^0}$$

and for $i = 3, 2, 1, 0$ by use of (f.3(ii))

$$f_{M_{i+1}^0} = f_{M_i^0} + w_{i+1}^0 f_{\emptyset}$$

Therefore,

$$f_{M^0} = 1 + w_1^0 + w_2^0 + w_3^0 + w_4^0 + w_5^0 + w_1^0 w_5^0$$

It can easily be checked that this result is in accordance with the definition of the cell polynomial (section 3).

5. COARSENEO CELL POLYNOMIALS

The cell polynomial $f_{\mathcal{M}} = f_{\mathcal{M}}(w_1, w_2, \dots, w_c)$ of $\mathcal{M} \in \underline{\mathcal{M}}^b$ can be "coarsened". Let $n_i = n(C_i)$ denote the number of vertices of $B_i = B(C_i)$. Inserting

$$w_i = \begin{cases} x, & \text{if } n_i \equiv 2, \text{ mod } 4 \\ y, & \text{if } n_i \equiv 0, \text{ mod } 4 \end{cases}$$

into $f_{\mathcal{M}}$, we find the *first coarsened cell polynomial* $f^*_{\mathcal{M}} = f^*_{\mathcal{M}}(x, y)$. With $x = y =: z$ we obtain the *second coarsened cell polynomial* $f^{**}_{\mathcal{M}} = f^{**}_{\mathcal{M}}(z) = f^*_{\mathcal{M}}(z, z)$. As an example we use the contour map M^0 of Figure 2:

$$f^*_{M^0} = 1 + 2x + 3y + xy$$

$$f^{**}_{M^0} = 1 + 5z + z^2$$

If there is no danger of confusion, we shall directly transfer concepts and symbols, originally defined for planar graphs, to their plane images (embeddings, patterns).

Let $\underline{L} = \underline{L}(G)$ and $l = l(G)$ denote the set and the number of all PMs of G , respectively.

Observation 3: For every $\mathcal{M} \in \underline{\mathcal{M}}^b$

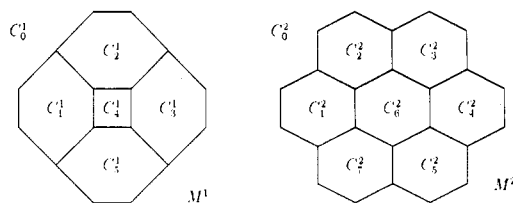
$$f^{**}_{\mathcal{M}} = l(\mathcal{M})$$

This formula follows immediately from the way the algorithm is constructed; it is well-known for \mathcal{M} being a catacondensed hexagonal system.⁵

A *perfect basic figure* (PBF) of graph G is a subgraph U of G with the following properties: (i) Every component of U is a circuit or a dumbbell $\bullet-\bullet$; (ii) U covers all vertices of G (Figure 5).

Clearly, the union of any two PMs $L^*, L^* \in \underline{L}(G)$ of G is the edge set of a PBF; conversely, every PBF U can be obtained this way. We shall denote the PBF U determined by the pair $[L^*, L^*]$ briefly by (L^*, L^*) . Let $q(U) = q(L^*, L^*)$ denote the number of circuits contained in $U = (L^*, L^*)$ whose length is a multiple of 4.

Lemma: There is a unique partition $\underline{L}, \underline{L}^{\parallel}$ of \underline{L} (i.e., $\underline{L} = \underline{L} \cup \underline{L}^{\parallel}$, $\underline{L} \cap \underline{L}^{\parallel} = \emptyset$ where \underline{L} or $\underline{L}^{\parallel}$ may be empty) such that



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