

A simple SSD-efficiency test

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Abstract A linear programming SSD-efficiency test capable of identifying a dominating portfolio is proposed. It has $T + n$ variables and at most $2T + 1$ constraints, whereas the existing SSD-efficiency tests are either unable to identify a dominating portfolio, or require solving a linear program with at least $O(T^2 + n)$ variables and/or constraints.

Keywords Stochastic dominance · Portfolio analysis · Linear programming

1 Introduction

The concept of second-order stochastic dominance (SSD), introduced in economics by Hadar and Russell [7] and Rothschild and Stiglitz [18], has become one of the central concepts in risk modelling. We say that portfolio rate of return X , modelled as a random variable (r.v.) on some probability space, dominates r.v. Y by SSD, and write $X \succ_2 Y$, if X is preferred to Y for any risk-averse expected utility maximizer (that is, for any agent with increasing and concave utility function, see [13]). Thus, the notion of SSD allows one to compare some of the investment opportunities without knowing exact utility function of a particular agent. This is particularly important, because identifying a utility function is a difficult task, which resorts to an extensive questionnaire procedure.

Given a (convex) set \mathcal{V} of admissible portfolio rates of return, r.v. $Y \in \mathcal{V}$ is called SSD-dominated within \mathcal{V} , if $X \succ_2 Y$ for some $X \in \mathcal{V}$ (in this case, X will be called a dominating portfolio), and SSD-efficient otherwise. The SSD definition implies that an optimal investment for a risk-averse expected-utility maximizer belongs to the

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set of SSD-efficient portfolios. However, this fact goes beyond the expected utility theory: only SSD-efficient portfolios may be optimal for an agent who maximizes convex Yaari dual utility function (see [19, Theorem 2]), minimizes a law-invariant convex risk measure (see [4, Corollary 4.59]), or uses the mean-deviation model [5]. SSD-efficiency is also central to solving the inverse portfolio problem for identifying investor's risk preferences, see [6]. This motivates the following question: determine whether the given portfolio Y is SSD-efficient within a given set \mathcal{V} , and if not, find a dominating portfolio.

Assuming that the underlying probability space is a finite T -element set $\Omega = \{\omega_1, \dots, \omega_T\}$, and \mathcal{V} consists of linear combinations of rates of return of n assets, Post [15] developed a linear program with $O(T + n)$ variables and constraints¹, which tests whether a given $Y \in \mathcal{V}$ is SSD-efficient subject to the additional assumption

$$Y(\omega_i) \neq Y(\omega_j), \quad i \neq j, \quad (1)$$

i.e. that ties do not occur in the distribution of Y . This assumption holds with probability 1 if the distribution of Y is an approximation of a continuous one using T Monte-Carlo simulations. However, assumption (1) is rarely encountered in practical applications. As explained in [15, Section II-C], tied returns may occur, for example, when analysing bootstrap pseudo-samples or evaluating a riskless alternative. Even if (1) holds for base assets, it might fail for some mixtures of them, or for derivative securities. Kopa and Post [9] show how the analysis can be generalized using a weakly increasing ranking to account for ties.²

Based on his test, Post [15, Section IV] made a surprising conclusion that some of the popular financial indices are SSD-dominated and hence cannot be optimal investments for a risk-averse agent. Post's test, however, fails to identify a dominating portfolio X if Y is SSD-dominated. In other words, for an agent holding a portfolio with rate of return Y , this test may show the *existence* of better investment opportunities within set \mathcal{V} , but does not identify them.

Assuming (1), Post [16] shows that a given portfolio is dominated by its mixture with the dual solution portfolio of Post [15], provided that the mixture lies in the local neighborhood with the same strictly increasing ranking as the evaluated portfolio. For a general case, several SSD-efficiency linear programming tests capable to identify a dominating portfolio have been developed, see e.g. [8–10]. Moreover, methods developed by Dentcheva and Ruszczyński [1, 2] and Kopa and Chovanec [8] can be used to find an *optimal* dominating portfolio under various definitions of optimality. However, all those tests use $O(T^2 + n)$ variables and constraints, which can be computationally intense, because the typical values of T are above 10^2 or even 10^3 . Fabian et al. [3] introduced a technique for solving optimization problems with SSD constraints which uses $O(n)$ variables, but with cuts from an exponential number of inequali-

¹ Throughout the paper, we do not count the bounds on decision variables (such as $x_j \geq 0$) as constraints, because the linear programming algorithms can efficiently deal with such bounds. This does not make any significant difference for our linear program, because it has at most $T + n$ such bounds.

² A method for treatment of ties has been outlined already in [15, Section II-C]. However, no SSD-test for the case with ties has been explicitly formulated in that paper.

ties added algorithmically. In contrast, Luedtke [12] suggested a test with $O(T + n)$ constraints but $O(T^2 + n)$ variables. The existence of a linear programming SSD-efficiency test with $O(T + n)$ variables and constraints, and to be capable of finding a dominating portfolio, was an open question. Such a test is the main result of this work.

We present a linear programming test, with $T + n$ variables and $2T + 1$ constraints, which, given any portfolio return $Y \in \mathcal{V}$, possibly with ties, tests whether Y is SSD-efficient within \mathcal{V} , and if not, finds a dominating portfolio X . A possible limitation is that our solution portfolio X may itself be inefficient, that is, dominated by a third portfolio Z . If this limitation is of concern, an additional test may be needed, e.g., the full Kopa and Post [9] test.

2 The SSD-efficiency test

Let $\Omega = \{\omega_1, \dots, \omega_T\}$ be a finite probability space, with probability measure \mathbb{P} such that $\mathbb{P}[\omega_i] = p_i$, $i = 1, \dots, T$. A random variable (r.v.) is any function $X : \Omega \rightarrow \mathbb{R}$. $F_X(x) = \mathbb{P}[X \leq x]$ and $q_X(\alpha) = \inf\{x | F_X(x) > \alpha\}$ will denote the cumulative distribution function (CDF) and quantile function of an r.v. X , respectively. We say, that r.v. X dominates r.v. Y by SSD, and write $X \succ_2 Y$, if $Eu(X) \geq Eu(Y)$ for every increasing concave function $u : \mathbb{R} \rightarrow \mathbb{R}$, with inequality being strict for some u . Equivalently, $X \succ_2 Y$ if and only if

$$\int_0^\alpha q_X(\beta) d\beta \geq \int_0^\alpha q_Y(\beta) d\beta, \quad \alpha \in (0, 1], \quad (2)$$

with inequality being strict for some $\alpha \in (0, 1]$, see Theorem 2.58 in [4].

Let \mathcal{V} be a convex set of r.v.s, and $Y \in \mathcal{V}$ be fixed. This paper presents a test which determines if Y is SSD-efficient within \mathcal{V} , and if not, finds a dominating portfolio.

Because Y is fixed and given, we can assume without loss in generality that

$$Y(\omega_i) \leq Y(\omega_j), \quad i < j. \quad (3)$$

Given any r.v. U , we denote $(s(1), \dots, s(T))$ a permutation of set $(1, \dots, T)$ such that $Y(\omega_{s(i)}) \leq Y(\omega_{s(j)})$, $i < j$, and $U(\omega_{s(i)}) \leq U(\omega_{s(j)})$ whenever $Y(\omega_{s(i)}) = Y(\omega_{s(j)})$, $i < j$.

Proposition 1 *Y is SSD-dominated if and only if there exists a non-zero r.v. $U \in (\mathcal{V} - Y)$ such that $\sum_{i=1}^t u_{s(i)} p_{s(i)} \geq 0$, $t = 1, \dots, T$. Moreover, in this case $X = Y + \varepsilon U$ is a dominating portfolio for any $\varepsilon \in (0, \min\{1, \min_{i \in J} \frac{y_{s(i+1)} - y_{s(i)}}{u_{s(i)} - u_{s(i+1)}}\})$ where $J \subset \{1, \dots, T\}$ is the set of indices³ for which $y_{s(i+1)} - y_{s(i)} > 0$ and $u_{s(i)} - u_{s(i+1)} > 0$.*

³ The set of such indices may be an empty set. Throughout the paper, we will use the convention that the minimum over an empty set is equal to $+\infty$.

Proof If Y is SSD-dominated, $X \succ_2 Y$ for some $X \in \mathcal{V}$. Take $U = X - Y$. Then $\sum_{i=1}^t X(\omega_{s(i)}) p_{s(i)} \geq \int_0^{\alpha_t} q_X(\beta) d\beta \geq \int_0^{\alpha_t} q_Y(\beta) d\beta = \sum_{i=1}^t y_{s(i)} p_{s(i)}$, where $y_i = Y(\omega_i)$, $i = 1, \dots, T$, and $\alpha_t = \sum_{i=1}^t p_{s(i)}$, $t = 1, \dots, T$. Hence $\sum_{i=1}^t u_{s(i)} p_{s(i)} \geq 0$, $t = 1, \dots, T$.

Conversely, let U and ε be as described. Then $X(\omega_{s(1)}) \leq \dots \leq X(\omega_{s(T)})$, where $X = Y + \varepsilon U$. Indeed, $X(\omega_{s(i)}) \leq X(\omega_{s(i+1)})$ is equivalent to $\varepsilon(u_{s(i)} - u_{s(i+1)}) \leq y_{s(i+1)} - y_{s(i)}$, which holds due to definition of ε if $i \in J$, and due to $\varepsilon(u_{s(i)} - u_{s(i+1)}) \leq 0 \leq y_{s(i+1)} - y_{s(i)}$ if $i \notin J$. Thus,

$$\int_0^{\alpha_t} q_X(\beta) d\beta = \sum_{i=1}^t X(\omega_{s(i)}) p_{s(i)} \geq \sum_{i=1}^t y_{s(i)} p_{s(i)} = \int_0^{\alpha_t} q_Y(\beta) d\beta, \quad t = 1, \dots, T. \quad (4)$$

Because functions $f_X(\alpha) = \int_0^\alpha q_X(\beta) d\beta$ and $f_Y(\alpha) = \int_0^\alpha q_Y(\beta) d\beta$ are piecewise linear with “vertices” at $\alpha = \alpha_t$, $t = 1, \dots, T$, (4) implies (2). Let $t = t_0$ be the smallest index such that $u_{s(t)} \neq 0$. Then $0 \leq \sum_{i=1}^{t_0} u_{s(i)} p_{s(i)} = u_{s(t_0)} p_{s(t_0)}$, hence $u_{s(t_0)} > 0$, and strict inequality holds in (4) for $t = t_0$. Thus, $X \succ_2 Y$.

Finally, because r.v.s Y and $Y + U$ belong to \mathcal{V} , and $\varepsilon \in (0, 1]$, $X = Y + \varepsilon U \in \mathcal{V}$ due to convexity of \mathcal{V} . \square

Proposition 1 cannot be applied directly for constructing a linear programming SSD-efficiency test, because permutation s depends on U and hence unknown in advance.

Let $I_k \subset \{1, \dots, T\}$, $k = 1, \dots, l$ be the sets of indices (of cardinality at least 2) such that $Y(\omega_i) = Y(\omega_j)$ if and only if $i, j \in I_k$ for some k . Let also $J_k = \{i \in \{1, \dots, T\} \mid i < j, \forall j \in I_k, k = 1, \dots, l, \text{ and } I = \{i \in \{1, \dots, T-1\} \mid Y(\omega_i) < Y(\omega_{i+1})\} \cup \{T\}$. Then $\{1, \dots, T\} = I \cup I_1 \cup \dots \cup I_l$.

Proposition 2 For an r.v. U , the following statements are equivalent

- (a) $\sum_{i=1}^t u_{s(i)} p_{s(i)} \geq 0$, $t = 1, \dots, T$;
- (b) $\sum_{i=1}^t u_i p_i \geq 0$, $t \in I$ and $\sum_{i \in J_k} u_i p_i + \sum_{i \in I_k} u_i^- p_i \geq 0$, $k = 1, \dots, l$, where $x^- = \min(x, 0)$.

Proof (a) \Rightarrow (b): Because

$$\sum_{i=1}^t u_{s(i)} p_{s(i)} = \sum_{i=1}^t u_i p_i, \quad t \in I, \quad (5)$$

(a) implies first statement of (b). For $k \in \{1, \dots, l\}$, let t_0 be the largest index $t \in I_k$ with $u_{s(t)} \leq 0$, and let t_0 be the largest index in J_k if $u_{s(t)} > 0$, $\forall t \in I_k$. Then

$$0 \leq \sum_{i=1}^{t_0} u_{s(i)} p_{s(i)} = \sum_{i \in J_k} u_i p_i + \sum_{i \in I_k, s^{-1}(i) \leq s^{-1}(t_0)} u_i p_i = \sum_{i \in J_k} u_i p_i + \sum_{i \in I_k} u_i^- p_i,$$

and (b) follows.

(b) \Rightarrow (a): If $t \in I$, (a) follows from the first statement of (b) and (5). If $t \in I_k$ for some $k \in \{1, \dots, l\}$,

$$\begin{aligned} \sum_{i=1}^t u_{s(i)} p_{s(i)} &= \sum_{i \in J_k} u_i p_i + \sum_{i \in I_k, s^{-1}(i) \leq s^{-1}(t)} u_i p_i \\ &\geq \sum_{i \in J_k} u_i p_i + \sum_{i \in I_k, s^{-1}(i) \leq s^{-1}(t)} u_i^- p_i \geq \sum_{i \in J_k} u_i p_i + \sum_{i \in I_k} u_i^- p_i \geq 0. \end{aligned}$$

□

It follows from Propositions 1 and 2 that the program

$$\begin{aligned} \max \quad & \sum_{t \in I} \left(\sum_{i=1}^t u_i p_i \right) + \sum_{k=1}^l \left(\sum_{i \in J_k} u_i p_i + \sum_{i \in I_k} u_i^- p_i \right), \\ \text{s.t.} \quad & \sum_{i=1}^t u_i p_i \geq 0, \quad t \in I, \quad \sum_{i \in J_k} u_i p_i + \sum_{i \in I_k} u_i^- p_i \geq 0, \quad k = 1, \dots, l \\ & U = (u_1, \dots, u_T) \in \mathcal{V} - Y \end{aligned} \quad (6)$$

has a positive optimal objective value if and only if Y is SSD-dominated, and in this case a dominating portfolio is $X = Y + \varepsilon U$ with $\varepsilon = \min\{1, \min_{i \in J} \frac{y_{s(i+1)} - y_{s(i)}}{u_{s(i)} - u_{s(i+1)}}\}$ as in Proposition 1. The program (6) is not linear because of the presence of u_i^- but can be linearised in a standard way by introducing variables v_i together with constraints $v_i \leq u_i$ and $v_i \leq 0$.

Let

$$\mathcal{V} = \left\{ X \mid X = \sum_{j=1}^n r_j x_j, \quad \sum_{j=1}^n x_j = 1, \quad x_j \geq 0, \quad j = 1, \dots, n \right\}, \quad (7)$$

where r_1, \dots, r_n are the rates of return of n assets, x_j is the fraction of capital invested into asset j , $\sum_{j=1}^n x_j = 1$ is the budget constraint, and $x_j \geq 0$, $j = 1, \dots, n$ are optional no-short-selling constraints. Let $r_{ij} = r_j(\omega_i)$, $i = 1, \dots, T$, $j = 1, \dots, n$ be the return of asset j under scenario ω_i . Then condition $U \in \mathcal{V} - Y$ in (6) becomes $u_i = \sum_{j=1}^n r_{ij} x_j - y_i$, $i = 1, \dots, T$. Because $Y \in \mathcal{V}$, $y_i = \sum_{j=1}^n r_{ij} x_j^0$, $i = 1, \dots, T$, for some x_1^0, \dots, x_n^0 , and the condition $U \in \mathcal{V} - Y$ becomes $u_i = \sum_{j=1}^n r_{ij} (x_j - x_j^0)$, $i = 1, \dots, T$.

Hence, for \mathcal{V} given by (7), program (6) can be written as

$$\begin{aligned} \max_{x_j, u_i, v_i} \quad & \sum_{t \in I} \left(\sum_{i=1}^t p_i u_i \right) + \sum_{k=1}^l \left(\sum_{i \in J_k} p_i u_i + \sum_{i \in I_k} v_i p_i \right), \\ \text{s.t.} \quad & \sum_{i=1}^t p_i u_i \geq 0, \quad t \in I, \quad \sum_{i \in J_k} p_i u_i + \sum_{i \in I_k} v_i p_i \geq 0, \quad k = 1, \dots, l \end{aligned}$$

$$\begin{aligned}
u_i &= \sum_{j=1}^n r_{ij}(x_j - x_j^0), \quad i = 1, \dots, T, \\
v_i &\leq 0, \quad v_i \leq u_i, \quad i \in (I_1 \cup \dots \cup I_l), \\
\sum_{j=1}^n x_j &= 1, \quad x_j \geq 0, \quad j = 1, \dots, n,
\end{aligned} \tag{8}$$

or, after excluding u_i ,

$$\begin{aligned}
\max_{x_j, v_i} & \sum_{t \in I} \left(\sum_{i=1}^t p_i \sum_{j=1}^n r_{ij}(x_j - x_j^0) \right) + \sum_{k=1}^l \left(\sum_{i \in J_k} p_i \sum_{j=1}^n r_{ij}(x_j - x_j^0) \sum_{i \in I_k} v_i p_i \right), \\
\text{s.t.} & \sum_{i=1}^t p_i \sum_{j=1}^n r_{ij}(x_j - x_j^0) \geq 0, \quad t \in I, \quad \sum_{i \in J_k} p_i \sum_{j=1}^n r_{ij}(x_j - x_j^0) \\
& + \sum_{i \in I_k} v_i p_i \geq 0, \quad k = 1, \dots, l \\
& v_i \leq 0, \quad v_i \leq \sum_{j=1}^n r_{ij}(x_j - x_j^0), \quad i \in (I_1 \cup \dots \cup I_l), \\
& \sum_{j=1}^n x_j = 1, \quad x_j \geq 0, \quad j = 1, \dots, n.
\end{aligned} \tag{9}$$

The resulting linear programming test has $T + n$ variables and at most $2T + 1$ constraints.

Example 1 Assume that there are $T = 3$ equiprobable scenarios, $n = 2$ assets with returns $r_1 = (0.24, 0, 0.06)$ and $r_2 = (0.04, 0.12, 0.12)$, and the benchmark portfolio has weights $(0, 1)$. In this case, $Y = (0.04, 0.12, 0.12)$, and assumption (1) does not hold. In the notation introduced before Proposition 2, $l = 1$, $I_1 = \{2, 3\}$, $J_1 = \{1\}$, $I = \{1, 3\}$, and the linear program (9) takes the form

$$\begin{aligned}
\max_{x_1, x_2, v_2, v_3} & \quad 3 \frac{1}{3} (0.24(x_1 - 0) + 0.04(x_2 - 1)) + \frac{1}{3} 0.12(x_2 - 1) \\
& + \frac{1}{3} (0.06(x_1 - 0) + 0.12(x_2 - 1)) + \frac{1}{3} v_2 + \frac{1}{3} v_3, \\
\text{s.t.} & \quad \frac{1}{3} (0.24(x_1 - 0) + 0.04(x_2 - 1)) \geq 0, \\
& \quad \frac{1}{3} (0.24(x_1 - 0) + 0.04(x_2 - 1)) + \frac{1}{3} 0.12(x_2 - 1) + \frac{1}{3} (0.06(x_1 - 0) \\
& + 0.12(x_2 - 1)) \geq 0, \\
& \quad \frac{1}{3} (0.24(x_1 - 0) + 0.04(x_2 - 1)) + \frac{1}{3} v_2 + \frac{1}{3} v_3 \geq 0, \\
& \quad v_2 \leq 0.12(x_2 - 1),
\end{aligned}$$

$$\begin{aligned}
v_3 &\leq 0.06(x_1 - 0) + 0.12(x_2 - 1), \\
v_2 &\leq 0, \quad v_3 \leq 0, \\
x_1 + x_2 &= 1, \quad x_1 \geq 0, \quad x_2 \geq 0,
\end{aligned} \tag{10}$$

which simplifies to

$$\begin{aligned}
\max_{x_1, x_2, v_2, v_3} \quad & 0.26x_1 + 0.12(x_2 - 1) + \frac{1}{3}v_2 + \frac{1}{3}v_3, \\
\text{s.t.} \quad & 6x_1 + (x_2 - 1) \geq 0, \\
& 15x_1 + 14(x_2 - 1) \geq 0, \\
& 0.24x_1 + 0.04(x_2 - 1) + v_2 + v_3 \geq 0, \\
& v_2 \leq 0.12(x_2 - 1), \\
& v_3 \leq 0.06x_1 + 0.12(x_2 - 1), \\
& v_2 \leq 0, \quad v_3 \leq 0, \\
& x_1 + x_2 = 1, \quad x_1 \geq 0, \quad x_2 \geq 0.
\end{aligned} \tag{11}$$

The optimal solution $x_1 = 1$, $x_2 = 0$, $v_2 = -0.12$, $v_3 = -0.06$, with the corresponding objective value $0.26x_1 + 0.12(x_2 - 1) + \frac{1}{3}v_2 + \frac{1}{3}v_3 = 0.08 > 0$, hence the benchmark portfolio is not SSD-efficient. Next, $U = (0.2, -0.12, -0.06)$, $s(i) = i$, $i = 1, 2, 3$, $J = \{1\}$, $\varepsilon = \min\{1, \frac{y_2 - y_1}{u_1 - u_2}\} = 0.25$, and a dominating portfolio $X = Y + 0.25U$ has weights $(0, 1) + 0.25(1, -1) = (0.25, 0.75)$.

Example 1 is a slight modification of Example 4 in [9], which illustrates that their dual test returns the same dominating portfolio and has 14 variables and 10 constraints. In contrast, test (11) has just 4 variables and 6 constraints.

Kopa and Post [9] compared the size of different SSD-efficiency tests in the case $T = 480$, $n = 12$. Table 1 presents their data for Post [15] dual test, Kuosmanen [10] test, Kopa [8] test, Kopa and Post [9] dual test, and Kopa and Post [9] reduced dual test, together with the corresponding data for Post [16] test, Luedtke [12] test

Table 1 SSD-tests comparison

Test	Size (constraints \times variables)	$T = 480, n = 12$	AT	DP	EP
Post [15] dual test, Eq. (12)	$(T + 1) \times (T + n - 1)$	481×491	No	No	No
Kuosmanen [10] test, Th. 6	$(T^2 + T + 1) \times (3T^2 + n)$	$230,881 \times 691,212$	No	Yes	Yes
Kopa [8], Eq. (16)	$(T^2 + T + 1) \times (T^2 + 2T + n)$	$230,881 \times 231,372$	No	Yes	Yes
Post [16] test, Eq. (5)	$(T + 1) \times (T + n - 1)$	481×491	No	Yes	No
Kopa and Post [9] reduced test	$(T + 1) \times (T + n)$	481×492	Yes	No	No
Kopa and Post [9] full test	$(T^2 + 1) \times (T^2 + T + n)$	$230,401 \times 230,892$	Yes	Yes	Yes
Luedtke [12] test, Eq. (cSSD1)	$(3T + n + 1) \times (T^2 + n)$	$1,453 \times 230,412$	Yes	Yes	Yes
Proposed test	$(2T + 1) \times (T + n)$	961×492	Yes	Yes	No

and our proposed test.⁴ It shows that the proposed test has substantially smaller size than the existing tests which allow ties and are able to identify a dominating portfolio.

3 Conclusions and future research

We have constructed a linear program (8)–(9) with $O(T + n)$ variables and constraints, such that its objective value is strictly positive if and only if the evaluated portfolio Y is SSD-dominated within admissible set \mathcal{V} given by (7). If Y is SSD-dominated, the output of the program can be used to construct a dominating portfolio.

The suggested SSD test is relevant for portfolio management: a dominating portfolio may be suggested as an alternative for an investor who is currently holding the benchmark portfolio Y . One may argue that a dominating portfolio may in general be SSD-inefficient, and, even if efficient, it is generally not optimal for the investor who holds the portfolio Y . Indeed, if the exact utility function of the investor is known, then it may be optimized to find an *optimal* portfolio with respect to his/her preferences. However, in practice, an investor rarely knows his/her utility function. In this case, determining the optimal portfolio is impossible, and a risk-averse investor may be advised to buy a dominating portfolio, which, in general, is not optimal, but anyway is better than the portfolio he/she currently holds, no matter what his/her utility function is.

An obvious question for future research is whether there exists a linear programming SSD-efficiency test with $O(T + n)$ variables and constraints returning a dominating portfolio which is in addition *SSD-efficient*. Another interesting research direction would be generalizing the results of this paper to higher order stochastic dominance. We say, that r.v. X dominates r.v. Y by N -th order stochastic dominance, or NSD, and write $X \succ_N Y$, if $Eu(X) \geq Eu(Y)$ for every function $u \in U_N$, with inequality being strict for some u , where U_N is the set of N times differentiable functions $u : \mathbb{R} \rightarrow \mathbb{R}$ such that $(-1)^{n-1}u^{(n)}(x) \geq 0$, $\forall x \in \mathbb{R}$, $n = 1, \dots, N$. An r.v. $Y \in \mathcal{V}$ is called NSD-efficient within set \mathcal{V} , if there are no $X \in \mathcal{V}$ such that $X \succ_N Y$. It would be interesting to obtain a linear programming NSD-efficiency test capable of identifying a dominating portfolio. However, this is not entirely straightforward. Our SSD-efficiency test relies on the quantile characterization (2) of SSD. Theorem 4 in [11] claims (without proof) that a similar representation holds at least for $N = 3$, namely, $X \succ_3 Y$ if and only if

$$\int_0^\alpha \left(\int_0^\beta q_X(\gamma) d\gamma \right) d\beta \geq \int_0^\alpha \left(\int_0^\beta q_Y(\gamma) d\gamma \right) d\beta, \quad \alpha \in (0, 1]. \quad (12)$$

However, Ng [14] provides a counterexample to this statement. To the best of our knowledge, no convenient representation of NSD in terms of quantile functions is

⁴ The columns AT, DP, and EP indicate whether the test allows ties in the return distribution, whether the results of the test can be used to identify a dominating portfolio, and whether the solution portfolio is SSD-efficient, correspondingly.

known for $N \geq 3$. Recently, Post and Kopa [17] derived a representation of the NSD criterion in terms of piecewise polynomials and co-lower partial moments, and used it to develop an efficient linear programming NSD-efficiency test for any N . However, their test cannot identify a dominating portfolio. This issue calls for new ideas.

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