# ON QUANTALOIDS AND QUANTAL CATEGORIES

## R. P. Gylys

#### 0. INTRODUCTION

The notion of quantale, a complete lattice provided with an additional binary operation subject to certain laws, was introduced by C. J. Mulvey [5] as a possible setting for constructive foundations for quantum mechanics (see also [2, 7, 12]). A. M. Pitts [6] and K. I. Rosenthal [8, 9] extended quantales to quantaloids in a way similar to the generalization of groups to groupoids. So a quantaloid can be viewed as a "quantale with many objects" or a quantale (with unit) can be a quantaloid with only one object. It seems that quantaloids already play a new role of great promise as in the theory of Grothendieck toposes [6] as well as in several areas of theoretical computer science [10, 11].

In [3] we proposed a non-commutative version of some notions and constructions of [4] and considered a class of quantales which might possibly assist in understanding the representation theory of non-commutative C\*-algebras. In this paper we extend these quantales to quantaloids.

We begin in Section 1 with a definition and a discussion of quantaloids with which we shall work. Section 2, devoted to the generalization (in our setting) of Girard quantaloids, is not absolutely necessary for what follows. In Section 3 we introduce the notion of quantal category and indicate its relationship with that of K. I. Rosenthal's Q-enriched category. In Section 4 we show that there is a monad on the category of quantal categories, which assigns to every quantal category its associated quantal category of singletons. Finally, in Section 5 we establish that Eilenberg-Moore algebras for this monad include "sheaf structures."

#### 1. QUANTALOIDS

We recall some basic definitions concerning quantaloids [6, 8-11].

Definition 1.1. A quantaloid Q (under the heading "Sl-category" in [6]) is a locally small category whose hom-sets are complete lattices and whose composition preserves sups in each variable separately:

(i) 
$$(\bigvee_j b_j) \circ a = \bigvee_j (b_j \circ a)$$
 and  $b \circ (\bigvee_j a_j) = \bigvee_j (b \circ a_j)$  for all morphisms  $a: u \to v$ ,  $b: v \to w$  and all families of morphisms  $\{a_j: u \to v\}$ ,  $\{b_j: v \to w\}$  in  $Q$ .

On Q there exist two further binary operations ()\() and ()/() left and right (Kan) extensions, respectively: for any morphisms  $a: u \to v$ ,  $b: v \to w$ ,  $c: u \to w$  in Q, the maps  $b \circ (): Q(u, v) \to Q(u, w)$  and ()  $\circ a: Q(v, w) \to Q(u, w)$  preserve suprema and hence have right adjoints:  $b \setminus (): Q(u, w) \to Q(u, v)$  and ()/a:  $Q(u, w) \to Q(v, w)$ . Thus two different "implications," the left extension  $b \setminus c$  of c along b and the right extension c/a of c along a, are given by

- (ii)  $b \setminus c = \bigvee \{d: u \to v \mid b \circ d \le c\}$  and  $c/a = \bigvee \{e: v \to w \mid e \circ a \le c\}$  and satisfy
- (iii)  $d \leqslant b \setminus c \iff b \circ d \leqslant c$  and  $e \leqslant c/a \iff e \circ a \leqslant c$ , respectively.

In this paper the term of quantaloid has one of its slightly wider meanings: Q is with possible lack of two-sided units (morphisms  $id_u$ ,  $id_v$  such that  $id_v \circ a = a = a \circ id_u$  for all  $a: u \to v$  in Mor(Q)), i.e., two-sided

Institute of Mathematics and Informatics, Akademijos 4, 2600 Vilnius, Lithuania. Published in Lietuvos Matematikos Rinkinys, Vol. 35, No. 3, pp. 266–296, July-September, 1995. Original article submitted October 4, 1994.

unitalness of hom-sets is not at all necessary for us because we shall never use it. But we shall work with a quantaloid Q which, in addition, will satisfy

Axioms 1.2. (1) Relations

- (i)  $T(v, v) \circ a \leq a \leq a \circ T(u, u)$ ,
- $(ii) (b \circ T(v, u)) \wedge (a'' \circ T(v', u)) = ((b \circ T(v, v')) \wedge a'') \circ T(v', u)$

$$(T(w,v)\circ a)\wedge (T(w,u')\circ c'')=T(w,v)\circ (a\wedge (T(v,u')\circ c''))$$

for  $a: u \to v$ ,  $a'': v' \to w$ ,  $b: v \to w$ ,  $c'': u \to u'$ ,  $u, u', v, v', w \in Obj(Q)$ , always hold, where T(v'', u'') denotes the top morphism in a hom-set Q(u'', v''). (The reader may find the fact that the reversal of order of objects in T(,) is rather confusing.)

- (2) Given two morphisms  $b: v \to w$ ,  $c: u \to w$  in Q, there exist morphisms  $b': w \to w$ ,  $c': u \to v$  such that
  - (iii)  $(b \circ T(v, u)) \land c \leq b \circ c' \leq c$  and  $(b \circ T(v, u)) \land c \leq b' \circ c \leq b \circ T(v, u)$ .

Recall that a quantaloid with one object is just a quantale which in our setting is not necessarily unital (that is quite unacceptable in [6, 8-11]). In this case the additional Axioms 1.2 coincide with Axioms 1.2 [3] in turn generalizing those for quantales considered in [1, 2, 4]. Let us give an example of a situation in the more general "many objects" case which will be of concern in the motivation of Axioms 1.2.

Let Q be a modular quantaloid, i.e., it is with two-sided units and there is an involution ()\*:  $Q^{op} \rightarrow Q$  satisfying Freyd's law of modularity: given morphisms  $a: u \rightarrow v$ ,  $b: v \rightarrow w$  and  $c: u \rightarrow w$  in Q, then

$$(b \circ a) \land c \leq b \circ (a \land (b^{\#} \circ c))$$

(see [6, 9], where the latter relation is written in a dual form). Let us also suppose that units of Q coincide with top morphisms. Then, obviously, relations 1.2(i) hold. It is easy to check that the remainder of Axioms 1.2 is also satisfied. To see this, put a = T(v, u),  $c = a'' \circ T(v', u)$  in Freyd's law of modularity. Then, the following holds:

$$b \circ T(v, u) \wedge a'' \circ T(v', u) \leqslant b \circ (T(v, u) \wedge b^{\#} \circ a'' \circ T(v, u)) = b \circ b^{\#} \circ a'' \circ T(v', u)$$
$$\leqslant (b \circ T(v, v') \wedge a'') \circ T(v', u) \qquad \text{(by 1.2(i))}$$
$$(\leqslant b \circ T(v, u) \wedge a'' \circ T(v', u) \qquad \text{(obviously)}.$$

We thus obtain the first equality of 1.2(ii). To prove the second one, we let b = T(w, v),  $c = T(w, u') \circ c''$  in Freyd's law of modularity and calculate as follows:

$$T(w, v) \circ a \wedge T(w, u') \circ c'' \leqslant T(w, v) \circ (a \wedge T(w, v)^{\#} \circ T(w, u') \circ c'')$$
$$\leqslant T(w, v) \circ (a \wedge T(v, u') \circ c'') (\leqslant T(w, v) \circ a \wedge T(w, u') \circ c'' \qquad \text{(obviously)}).$$

In order to prove 1.2(iii), we put a = T(v, u) in Freyd's law of modularity. Then we have that

$$b \circ T(v, u) \land c \leq b \circ b^{\#} \circ c \ (= b \circ c')$$
  
  $\leq T(w, w) \circ c \quad \text{(by 1.3(ii) below)}$   
  $\leq c \quad \text{(by 1.2(i))}$ 

and

$$b \circ T(v, u) \wedge c \leqslant b \circ b^{\#} \circ c \ (= b' \circ c) \leqslant b \circ T(v, u)$$
 (by 1.3(ii)),

i.e., relations 1.2(iii) hold. Thus, a modular quantaloid, units of which are top morphisms, provides a useful example of a quantaloid satisfying Axioms 1.2.

Let us now list some properties of quantaloids that will be of interest to us.

LEMMA 1.3. Let Q be a quantaloid satisfying Axioms 1.2. Then for all morphisms  $a, a': u \to v, b: v \to w, c: u \to w, c': u \to u', d: w \to v, d': w \to u, e: v \to u, e': u' \to u$  and for all families  $\{a_j\} \subseteq Q(u, v), \{b_j\} \subseteq Q(v, w), \{c_j\} \subseteq Q(u, w)$  of morphisms in Q, the following relations hold:

(i) 
$$b \circ (b \setminus c) \leq c$$
,  $(c/a) \circ a \leq c$ ;

- (ii)  $a \le a'$  implies  $a \circ e' \le a' \circ e'$ ,  $b \circ a \le b \circ a'$ ;
- (iii)  $T(v, v) \circ a = a$ ,  $\pm(w, v) \circ a = \pm(w, u)$ ,  $a \circ \pm(u, w) = \pm(v, w)$  (here  $\pm(v', u') = \bigvee \emptyset$  is the bottom morphism of Q(u', v'));
- (iv)  $(a \circ e') \setminus d = e' \setminus (a \setminus d)$ ,  $c'/(b \circ a) = (c'/a)/b$ ,  $e' \setminus (e/b) = (e' \setminus e)/b$ ;
- (v)  $(\bigvee_{i}b_{j})\setminus c = \bigwedge_{i}(b_{j}\setminus c), \ c/\bigvee_{i}a_{j} = \bigwedge_{i}(c/a_{j});$
- (vi)  $b \setminus \bigwedge_i c_j = \bigwedge_i (b \setminus c_j), \ (\bigwedge_i c_j)/a = \bigwedge_i (c_j/a);$
- (vii)  $\bigvee_{i} (b \setminus c_{i}) \leq b \setminus \bigvee_{i} c_{i}$ ,  $\bigvee_{i} (c_{i}/a) \leq (\bigvee_{i} c_{i})/a$ ;
- (viii)  $\bigvee_{i} (b_{i} \setminus c) \leq (\bigwedge_{i} b_{j}) \setminus c, \bigvee_{i} (c/a_{i}) \leq c/\bigwedge_{i} a_{i};$
- (ix)  $(d \setminus a) \circ e' \leq d \setminus (a \circ e')$ ,  $b \circ (a/c') \leq (b \circ a)/c'$ ,  $(b/e) \circ T(u, u) = b/e$ ;
- $(x) (e' \setminus e) \circ (e \setminus d') \leqslant e' \setminus d', (c/a) \circ (a/c') \leqslant c/c';$
- (xi)  $d' \setminus e' \leq (a \circ d') \setminus (a \circ e'), c/a \leq (c \circ e')/(a \circ e');$
- (xii)  $(b \circ T(v, u)) \wedge c = b \circ (b \setminus c) = ((b \circ T(v, u))/c) \circ c;$
- (xiii)  $(b \circ T(v, u)) \wedge (T(w, v) \circ a) = b \circ (b \setminus T(w, v) \circ a) = ((b \circ T(v, u))/a) \circ a;$
- $(xiv) [(\bigvee_i b_j) \circ T(v, u)] \wedge c = \bigvee_i [(b_j \circ T(v, u)) \wedge c], [b \circ T(v, u)] \wedge \bigvee_i c_j = \bigvee_i [(b \circ T(v, u)) \wedge c_j].$

*Proof.* Relations (i), (vii), (ix)–(xi) follow immediately from 1.1(iii) and 1.2(i). To prove (ii) note that  $a \le a'$  is equivalent to  $a \lor a' = a'$  which by 1.1(i) yields  $a' \circ e' = (a \lor a') \circ e' = (a \circ e') \lor (a' \circ e')$  and  $b \circ a' = b \circ (a \lor a') = (b \circ a) \lor (b \circ a')$ , i.e., implication (ii) holds. Next, from (ii) and 1.2(i), (iii) we can deduce easily that  $a = T(v, u) \land a \le T(v, u) \land a \le T' \circ a \le T(v, v) \circ a \le a$ , i.e., T(v, v) is the left-sided unit in Q, while the following two equalities of (iii) are just the case " $\{j\}$  is an empty set" in Axiom 1(i). Obviously (ii) and 1.1(iii) imply (iv), (vi), (viii). Further, we verify the left relation of (v) by the following series of implications:

$$\left(\bigvee_{j}b_{j}\right)\circ\left(\left(\bigvee_{j}b_{j}\right)\setminus c\right)\leqslant c \quad \text{(by (i))} \implies \bigvee_{j}\left[b_{j}\circ\left(\left(\bigvee_{i}b_{i}\right)\setminus c\right)\right]\leqslant c$$

$$\implies \left(\forall j\right) \quad b_{j}\circ\left(\left(\bigvee_{i}b_{i}\right)\setminus c\right)\leqslant c \implies \left(\forall j\right) \quad \left(\bigvee_{i}b_{i}\right)\setminus c\leqslant b_{j}\setminus c \quad \text{(by 1.1(iii))}$$

$$\implies \left(\bigvee_{i}b_{i}\right)\setminus c\leqslant \bigwedge_{i}(b_{i}\setminus c), \quad \left(\forall j\right) \quad \bigwedge_{i}(b_{i}\setminus c)\leqslant b_{j}\setminus c$$

$$\implies \left(\forall j\right) \quad b_{j}\circ\bigwedge_{i}(b_{i}\setminus c)\leqslant c \implies \bigvee_{j}\left(b_{j}\circ\bigwedge_{i}(b_{i}\setminus c)\right)\leqslant c$$

$$\implies \left(\bigvee_{j}b_{j}\right)\circ\bigwedge_{i}(b_{i}\setminus c)\leqslant c \implies \bigwedge_{i}(b_{i}\setminus c)\leqslant \left(\bigvee_{i}b_{i}\right)\setminus c,$$

while the verification of the other equality in (v) is similar. Now we prove the left side of (xii). We have

$$b \circ (b \setminus c) = b \circ \left( \bigvee \{ a \in Q(u, v) \mid b \circ a \leqslant c \} \right) = \bigvee \{ b \circ a \in Q(u, w) \mid b \circ a \leqslant c \}$$
 (by 1.1(i))  
 
$$\geqslant (b \circ T(v, u)) \wedge c$$
 (by 1.2(iii)).

Conversely,  $b \circ (b \setminus c) \leq b \circ T(v, u)$  and  $b \circ (b \setminus c) \leq c$ , so that

$$b \circ (b \setminus c) \leq b \circ T(v, u) \wedge c$$
.

The argument for the right side of (xii) and (xiii) is similar. Finally, the first relation in (xiv) holds, since

$$\left[ \left( \bigvee_{i} b_{i} \right) \circ T(v, u) \right] \wedge c = \left( \bigvee_{i} b_{i} \right) \circ \left( \left( \bigvee_{j} b_{j} \right) \setminus c \right) \quad \text{(by (xii))}$$

$$= \bigvee_{i} \left[ b_{i} \circ \left( \left( \bigvee_{j} b_{j} \right) \setminus c \right) \right] \quad \text{(by 1.1 (i))}$$

$$= \bigvee_{i} \left( b_{i} \circ \bigwedge_{i} (b_{j} \setminus c) \right) \quad \text{(by (v))}$$

$$\leq \bigvee_{i} b_{i} \circ (b_{i} \setminus c) = \bigvee_{j} (b_{i} \circ T(v, u) \wedge c) \quad \text{(by (xii))}$$

and

$$\left(\forall j\right) \quad b_j \circ T(v,u) \wedge c \leqslant \left[\left(\bigvee_i b_i\right) \circ T(v,u)\right] \wedge c \Longrightarrow \quad \bigvee_i (b_i \circ T(v,u) \wedge c) \leqslant \left[\left(\bigvee_i b_i\right) \circ T(v,u)\right] \wedge c,$$

while the last relation is verified in a similar way.

Note that by 1.3(xiv) each hom-lattice Q(u, v) is actually a quantale, with the product of a, a':  $u \rightarrow v$ being  $(a \circ T(u, u)) \wedge a'$  and the left-sided unit being T(v, u). (The associativity of this multiplication follows from 1.2(ii). Since  $T(v, u) = T(v, u) \circ T(u, u)$  (by 1.2(i)), it follows that  $(T(v, u) \circ T(u, u)) \wedge a = a$ .)

#### 2. GENERALIZED GIRARD QUANTALOIDS

In this section we shall deal with a class of quantaloids together with "dualizing families" which coincide with the so-called Girard quantaloids (introduced by K. I. Rosenthal in [9], and in the case of quantales - by D. Yetter in [12]) when these dualizing families are "cyclic" and  $T(u, u) = id_u$  ( $\forall u \in Obj(Q)$ ).

Definition 2.1 (cf. Definition 2.6 [9]). Let Q be a quantaloid.  $O = \{O_u: u \to v \mid u \in Q\}$ , a family of morphisms of Q such that  $O_u \circ T(u, u) = O_u$  is a dualizing family iff given a:  $u \to v$  in Mor(Q), then

$$(O_u/a) \setminus O_u = O_v/(a \setminus O_v) = a \circ T(u, u).$$

Q is a generalized Girard quantaloid iff it has a dualizing family O of morphisms.

If Q is a generalized Girard quantaloid and a:  $u \to v$  is a morphism of Q, then we shall use  $a^{\circ}$  and  $a^{\circ}$ to denote two different (in general) "negations"  $a \setminus O_v$  and  $O_u/a$ , respectively. Note that  $({}^{\circ}a)^{\circ} = {}^{\circ}(a^{\circ}) (=:$  $^{\circ}a^{\circ}) = a \circ T(u, u), \ a \leqslant b^{\circ} \text{ iff } b \leqslant ^{\circ}a \text{ (by 1.1(iii))}, \ a^{\circ} \circ T(v, v) = a^{\circ} \text{ (since } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \text{ implies } O_{v} \circ T(v, v) = O_{v} \circ T(v, v) =$  $(a \setminus O_v) \circ T(v, v) \leqslant a \setminus (O_v \circ T(v, v))$  (by 1.3(ix))=  $a \setminus O_v$ ),  $a \circ T(v, v) = a$  (since  $(O_u/a) \circ a = (O_u/a)$ )  $a)\circ T(v,v)\circ a \text{ (by 1.3(iii))}\leqslant O_u \text{ imply } (O_u/a)\circ T(v,v)\leqslant O_u/a \text{ (by 1.1(iii)))}, \ \ ^\circ(a^\circ)^\circ=a^\circ, \ \ ^\circ(^\circ a)^\circ=^\circ a,$  $({}^{\circ}a^{\circ})^{\circ} = a^{\circ} \text{ (since } a \circ ((a \circ T(u, u)) \setminus O_{v}) = a \circ T(u, u) \circ (a \circ T(u, u) \setminus O_{v}) \leqslant O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \setminus O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{v} \text{ imply } (a \circ T(u, u)) \cap O_{v} \leqslant a \setminus O_{$ and since  $a \circ T(u, u) \circ (a \setminus O_v) = a \circ (a \setminus O_v) \leqslant O_v$  imply  $a \setminus O_v \leqslant (a \circ T(u, u)) \setminus O_v$  and  $o(a \circ a) = a$  (since  $(O_u/(a \circ T(u,u))) \circ a \leqslant (O_u/(a \circ T(u,u))) \circ a \circ T(u,u) \leqslant O_u \text{ imply } O_u/(a \circ T(u,u)) \leqslant O_u/a \text{ and since } O_u/(a \circ T(u,u)) \leqslant O_u/(a \circ T($  $(O_u/a) \circ a \circ T(u,u) \leqslant O_u \circ T(u,u) = O_u \text{ imply } O_u/a \leqslant O_u/(a \circ T(u,u)).$ 

The following lemmas list some basic facts about generalized Girard quantaloids.

LEMMA 2.2 (cf. [9, Proposition 2.3] and [4, Lemma 1.4]). Let Q be a generalized Girard quantaloid with a dualizing family  $O = \{O_u: u \to u \mid u \in Obj(Q)\}$  and let  $a, a': u \to v, b: v \to w, c: u \to w, e: v \to u,$ d':  $w \to u$ ,  $\{a_i: u \to v\}$ ,  $\{b_i: u \to v\}$  be, respectively, morphisms and families of morphisms in Q. Then:

- (i)  $(\bigvee_j a_j)^\circ = \bigwedge_j a_j^\circ$ ,  $(\bigvee_j a_j) = \bigwedge_j \circ a_j$ ,  $(\bigwedge_j \circ a_j^\circ)^\circ = \bigvee_j a_j^\circ$ ,  $(\bigwedge_j \circ a_j^\circ) = \bigvee_j \circ a_j$ ; (ii)  $a \setminus b^\circ = (b \circ a)^\circ$ ,  $a \setminus b = (b \circ a)^\circ$ ,  $b \setminus a^\circ = b^\circ \setminus a^\circ$  (=  $(a \circ b)^\circ = (a \circ b)^\circ$ ); (iii)  $a \setminus b^\circ = a \setminus a^\circ$ ,  $a \setminus a^\circ = a \cap a^\circ$ ,  $a \setminus a^\circ = a \cap a^\circ$ ); (iii)  $a \setminus b^\circ = a \setminus a^\circ$ ,  $a \setminus a^\circ = a \cap a^\circ$ ,  $a \setminus a^\circ = a \cap a^\circ$ ,  $a \cap a^\circ = a^$

- (iv)  $a^{\circ} \lor a^{\circ} = (a^{\circ}/a') \lor a^{\circ} = a^{\circ}/(a' \lor a^{\circ});$
- $\begin{array}{l} (\mathsf{v}) \circ a^{\circ} \vee \bigwedge_{j} \circ a_{j}^{\circ} = \bigwedge_{j} (\circ a^{\circ} \vee \circ a_{j}^{\circ}); \\ (\mathsf{v}i) \ b \circ \bigwedge_{j} \circ a_{j}^{\circ} = \bigwedge_{j} b \circ \circ a_{j}^{\circ}, \ (\bigwedge_{j} \circ b_{j}^{\circ}) \circ \circ a^{\circ} = \bigwedge_{j} (b_{j} \circ \circ a^{\circ}). \end{array}$

*Proof.* (i) The first two equalities are none other than 1.3(v). Since  $\bigwedge_j {}^{\circ}a_j^{\circ} = {}^{\circ}(\bigvee_j a_j^{\circ}) = (\bigvee_j {}^{\circ}a_j)^{\circ}$ , it follows that  $(\bigwedge_j {}^{\circ}a_j^{\circ})^{\circ} = {}^{\circ}(\bigvee_j a_j^{\circ})^{\circ} = (\bigvee_j a_j^{\circ}) \circ T(v,v) = \bigvee_j (a_j^{\circ} \circ T(v,v)) = \bigvee_j a_j^{\circ}$  and  ${}^{\circ}(\bigwedge_j {}^{\circ}a_j^{\circ}) = {}^{\circ}(\bigvee_j {}^{\circ}a_j)^{\circ} = (\bigvee_j {}^{\circ}a_j) \circ T(v,v) = \bigvee_j ({}^{\circ}a_j \circ T(v,v)) = \bigvee_j {}^{\circ}a_j$ .

(ii) The first two equalities follow immediately from 1.3(iv). Next, we have that  $({}^{\circ}b \setminus {}^{\circ}a^{\circ}) \circ a^{\circ} = ({}^{\circ}b \setminus {}^{\circ}a^{\circ}) \circ ({}^{\circ}a^{\circ}) \circ ({}^{\circ}a^{\circ}) \circ ({}^{\circ}a^{\circ}) \circ ({}^{\circ}a^{\circ}) \circ ({}^{\circ}a^{\circ}) \circ ({}^{\circ}a^{\circ}) \circ ({}^{\circ}b^{\circ}) \circ ({}^{\circ}b^{\circ})$ 

(iii) In view of 1.3(x) and 1.1(iii) we obtain

$$c \circ (c^{\circ}/b^{\circ}) = c \circ T(u, u) \circ (c^{\circ}/b^{\circ}) = (O_{w}/c^{\circ}) \circ (c^{\circ}/b^{\circ}) \leqslant {^{\circ}}b^{\circ} \implies c^{\circ}/b^{\circ} \leqslant c \setminus {^{\circ}}b^{\circ},$$

$$(c \setminus {^{\circ}}b^{\circ}) \circ b^{\circ} = (c \setminus {^{\circ}}b^{\circ}) \circ ({^{\circ}}b^{\circ})^{\circ} \leqslant c^{\circ} \implies c \setminus {^{\circ}}b^{\circ} \leqslant c^{\circ}/b^{\circ},$$

$$({^{\circ}}c \setminus {^{\circ}}a) \circ a \leqslant ({^{\circ}}c \setminus {^{\circ}}a) \circ ({^{\circ}}a \setminus O_{u}) \leqslant {^{\circ}}c^{\circ} \implies {^{\circ}}c \setminus {^{\circ}}a \leqslant {^{\circ}}c^{\circ}/a,$$

$${^{\circ}}c \circ ({^{\circ}}c^{\circ}/a) = {^{\circ}}({^{\circ}}c^{\circ}) \circ ({^{\circ}}c^{\circ}/a) \leqslant {^{\circ}}a \implies {^{\circ}}c^{\circ}/a \leqslant {^{\circ}}c \setminus {^{\circ}}a.$$

(iv) We have

$$a^{\circ} \vee a'^{\circ} = ({}^{\circ}a^{\circ} \wedge {}^{\circ}a'^{\circ})^{\circ} \qquad \text{(by (i))}$$

$$= (a \circ T(u, u) \wedge {}^{\circ}a'^{\circ})^{\circ} = (a \circ (a \setminus {}^{\circ}a'^{\circ}))^{\circ} \qquad \text{(by 1.3(xi))}$$

$$= (a \setminus {}^{\circ}a'^{\circ}) \setminus a^{\circ} \qquad \text{(by (iii))}$$

$$= (a^{\circ}/a'^{\circ}) \setminus a^{\circ} \qquad \text{(by (iii))},$$

while the remainder follows by a similar calculation.

(v) Because of (i) and 1.3(xiv) we obtain that

$${}^{\circ}a^{\circ} \vee \bigwedge_{j} {}^{\circ}a_{j}^{\circ} = ({}^{\circ}a)^{\circ} \vee \bigwedge_{j} ({}^{\circ}a_{j})^{\circ} = ({}^{\circ}a)^{\circ} \vee \left(\bigvee_{j} {}^{\circ}a_{j}\right)^{\circ}$$

$$= \left({}^{\circ}({}^{\circ}a)^{\circ} \wedge {}^{\circ}\left(\bigvee_{j} {}^{\circ}a_{j}\right)^{\circ}\right)^{\circ} = \left({}^{\circ}({}^{\circ}a)^{\circ} \wedge \bigvee_{j} {}^{\circ}({}^{\circ}a_{j})^{\circ}\right)^{\circ}$$

$$= \left(\bigvee_{j} ({}^{\circ}({}^{\circ}a)^{\circ} \wedge {}^{\circ}({}^{\circ}a_{j})^{\circ})\right)^{\circ} = \bigwedge_{j} ({}^{\circ}({}^{\circ}a)^{\circ} \wedge {}^{\circ}({}^{\circ}a_{j})^{\circ})^{\circ} = \bigwedge_{j} ({}^{\circ}a^{\circ} \vee {}^{\circ}a_{j}^{\circ}).$$

(vi) In view of

$$b \setminus \bigwedge_{j} b \circ {}^{\circ}a_{j}^{\circ} \left[ = \bigwedge_{j} (b \setminus b \circ {}^{\circ}a_{j}^{\circ}) \quad (by \ 1.3(vi)) \right]$$

$$= \bigwedge_{j} (b \setminus {}^{\circ}(b \circ a_{j})^{\circ}) = \bigwedge_{j} (b^{\circ}/(b \circ a_{j})^{\circ}) = \bigwedge_{j} (b^{\circ}/(a_{j} \setminus b^{\circ}))$$

$$= \bigwedge_{j} ({}^{\circ}(b^{\circ})^{\circ}/(a_{j} \setminus {}^{\circ}(b^{\circ})^{\circ})) = \bigwedge_{j} ({}^{\circ}(b^{\circ})^{\circ} \vee {}^{\circ}a_{j}^{\circ})$$

$$= {}^{\circ}(b^{\circ})^{\circ} \vee \bigwedge_{j} {}^{\circ}a_{j}^{\circ} \quad (by \ (v)) = {}^{\circ}(b^{\circ})^{\circ} \vee \left(\bigwedge_{j} {}^{\circ}b_{j}^{\circ}\right)^{\circ}$$

$$= {}^{\circ}(b^{\circ})^{\circ} / \left(\bigwedge_{j} {}^{\circ}a_{j}^{\circ} \setminus {}^{\circ}(b^{\circ})^{\circ}\right) = b^{\circ} / \left(\bigwedge_{j} {}^{\circ}a_{j}^{\circ} \setminus b^{\circ}\right)$$

$$= b^{\circ} / \left(b \circ \bigwedge_{j} {}^{\circ}a_{j}^{\circ}\right)^{\circ} = b / \left(b \circ \bigwedge_{j} {}^{\circ}a_{j}^{\circ}\right)^{\circ} = b \setminus b \circ \left(\bigwedge_{j} {}^{\circ}a_{j}^{\circ}\right)^{\circ}\right]$$

$$= b \setminus \left(b \circ \bigwedge_{j} {}^{\circ}a_{j}^{\circ}\right)$$

we obtain that

$$\bigwedge_{j} b \circ {}^{\circ}a_{j}^{\circ} = b \circ T(v, v) \wedge \bigwedge_{j} b \circ {}^{\circ}a_{j}^{\circ} = b \circ \left( b \setminus \bigwedge_{j} b \circ {}^{\circ}a_{j}^{\circ} \right) = b \circ \left( b \setminus b \circ \bigwedge_{j} {}^{\circ}a_{j}^{\circ} \right)$$

$$= b \circ T(v, v) \wedge \left( b \circ \bigwedge_{j} {}^{\circ}a_{j}^{\circ} \right) = b \circ \bigwedge_{j} {}^{\circ}a_{j}^{\circ}.$$

Now interchanging the role of pairs  $(b, a_i)$  and  $(b_i, a)$  we can verify the last equality.

LEMMA 2.3. Let  $(Q, \{O_v\})$  be a generalized Girard quantaloid. Then it has a second composition, () + (), defined by

$$b + a = ({}^{\circ}a \circ {}^{\circ}b)^{\circ} [= {}^{\circ}(a^{\circ} \circ b^{\circ}) = {}^{\circ}b \setminus {}^{\circ}a^{\circ} = {}^{\circ}b^{\circ}/a^{\circ} \quad (by \ 2.2(ii))]$$

and satisfying

- (i) (b+a)+e'=b+(a+e');
- (ii)  $O_v + a = a + O_u = a \circ T(u, u)$ ;
- (iii)  $({}^{\circ}a + {}^{\circ}b)^{\circ} = {}^{\circ}(a^{\circ} + b^{\circ}) = b \circ a \circ T(u, u);$
- (iv)  $b + \bigvee_{i} a_{j} = \bigvee_{i} (b + a_{j}), \ (\bigvee_{i} b_{j}) + a = \bigvee_{i} (b_{j} + a);$
- (v)  $b + \bigwedge_i {}^{\circ}a_i^{\circ} = \bigwedge_i (b + a_i), \ (\bigwedge_i {}^{\circ}b_i^{\circ}) + a = \bigwedge_i (b_i + a);$
- (vi)  $a^{\circ} \lor a'^{\circ} = (a' \circ a^{\circ}) + a = a + (a' \circ a');$
- (vii)  $b \circ (a + e') \leq (b \circ a) + e'$ ,  $(b + a) \circ e' \leq b + (a \circ e')$  (o+ associativity);
- (viii)  $e' \circ (e + d) \leq (e' \setminus e) + d$ ,

where  $a, a': u \to v, b: v \to w, e': u' \to u, d: w \to v, \{a_j\} \subseteq Q(u, v), \{b_j\} \subseteq Q(v, w)$  are, respectively, morphisms and families of morphisms in Q.

Proof. To prove (i) note that

$$(b+a) + e' = {}^{\circ}(b+a) \setminus {}^{\circ}e' {}^{\circ} = ({}^{\circ}a \circ {}^{\circ}b) \setminus {}^{\circ}e' {}^{\circ} = {}^{\circ}b \setminus ({}^{\circ}a \setminus {}^{\circ}e' {}^{\circ})$$
 (by 1.3(iv))  
=  ${}^{\circ}b \setminus (a+e') = {}^{\circ}b \setminus {}^{\circ}(a+e') {}^{\circ} = b + (a+e'),$ 

while relations (ii), (iii) follow immediately from definition of +. Next, we have that

$$b + \bigvee_{j} a_{j} = {\circ} \left( \left( \bigvee_{j} a_{i} \right)^{\circ} \circ b^{\circ} \right) = {\circ} \left( \left( \bigwedge_{j} a_{j}^{\circ} \right) \circ b^{\circ} \right)$$

$$= {\circ} \left( \left( \bigwedge_{j} {\circ} (a_{j}^{\circ})^{\circ} \right) \circ {\circ} (b^{\circ})^{\circ} \right) = {\circ} \left( \bigwedge_{j} a_{j}^{\circ} \circ {\circ} (b^{\circ})^{\circ} \right) \qquad \text{(by 2.2(vi))}$$

$$= \bigvee_{j} {\circ} (a_{j}^{\circ} \circ b^{\circ}) = \bigvee_{j} (b + a_{j})$$

and

$$\left(\bigvee_{j} b_{j}\right) + a = {\circ} \left(a^{\circ} \circ \bigwedge_{j} b_{j}^{\circ}\right) = {\circ} \left(a^{\circ} \circ \bigwedge_{j} {\circ} (b_{j}^{\circ})^{\circ}\right) = {\circ} \left(\bigwedge_{j} a^{\circ} \circ {\circ} (b_{j}^{\circ})^{\circ}\right)$$

$$= \bigvee_{j} {\circ} (a^{\circ} \circ b_{j}^{\circ}) = \bigvee_{j} (b_{j} + a).$$
(by 2.2(vi))

To verify (v) we proceed as follows:

$$b + \bigwedge_{j}^{\circ} a_{j}^{\circ} = {\circ} \left( \left( \bigwedge_{j}^{\circ} a_{j}^{\circ} \right)^{\circ} \circ b^{\circ} \right) = {\circ} \left( \left( \bigvee_{j}^{\circ} a_{j}^{\circ} \right) \circ b^{\circ} \right)$$
 (by 2.2(i))  
$$= {\circ} \left( \bigvee_{j}^{\circ} (a_{j}^{\circ} \circ b^{\circ}) \right)$$
 (by 1.1(i))  
$$= \bigwedge_{j}^{\circ} (a_{j}^{\circ} \circ b^{\circ}) = \bigwedge_{j}^{\circ} (b + a_{j}),$$

$$\left(\bigwedge_{j}^{\circ}b_{j}^{\circ}\right)+a={\circ}\left(a^{\circ}\circ\left(\bigwedge_{j}^{\circ}b_{j}^{\circ}\right)^{\circ}\right)={\circ}\left(a^{\circ}\circ\bigvee_{j}b_{j}^{\circ}\right)={\circ}\left(\bigvee_{j}(a^{\circ}\circ b_{j}^{\circ})\right)=\bigwedge_{j}^{\circ}(a^{\circ}\circ b_{j}^{\circ})=\bigwedge_{j}^{\circ}(b_{j}+a).$$

Finally, (vi) follows from 2.2(iv), while (vii) and (viii) follow from 1.1(iii).

LEMMA 2.4. Let Q be a generalized Girard quantaloid. Then Q has two "coimplications," ()  $\rightarrow$  () and ()  $\leftarrow$  (), such that  $^{\circ}(b \rightarrow c)^{\circ} = b \rightarrow c$ ,  $^{\circ}(c \leftarrow a)^{\circ} = c \leftarrow a$  with

- (i)  $b + a \geqslant c \iff {}^{\circ}a^{\circ} \geqslant b \rightarrow c \quad (+ \rightarrow residuation),$
- (ii)  $b+a\geqslant c\iff {}^{\circ}b^{\circ}\geqslant c \leftarrow a \quad (+\leftarrow residuation),$  and it holds true that

(iii)  $b \rightarrow c = \bigwedge \{ a^{\circ} \mid b + a \geqslant c \},$ 

- (iv)  $c \leftarrow a = \bigwedge \{ a \mid b + a \geqslant c \},\$
- where  $a: u \to v, b: v \to w, c: u \to w$  are morphisms in Q.

*Proof.* Assume  $b \to c$  to be given. By (i)  $(b \to c)^\circ = b \to c$  implies  $b + (b \to c) \ge c$ . So it is clear that the infimum in (iii) does not exceed  $(b \to c)^\circ = b \to c$ . But if  $\bigwedge \{a^\circ \mid b + a \ge c\} < b \to c$ , then there exists some  $d \in \text{Mor}(Q)$  with  $d^\circ < b \to c$  and  $d^\circ < b \to c$  and  $d^\circ < b \to c$ . For such  $d^\circ = b$  we get a contradiction by  $d^\circ = b \to c > d^\circ$ . Let the coimplication  $d^\circ = b \to c > d^\circ$ . Let the coimplication  $d^\circ = b \to c > d^\circ$ .

$$\overset{\circ}{\left(\bigwedge\{{}^{\circ}a^{\circ}\mid b+a\geqslant c\}\right)}^{\circ} = \overset{\circ}{\left(\bigvee\{a^{\circ}\mid b+a\geqslant c\}\right)} = \bigwedge\{{}^{\circ}a^{\circ}\mid b+a\geqslant c\},$$

$$b+a\geqslant c \implies {}^{\circ}a^{\circ}\geqslant \bigwedge\{{}^{\circ}a'^{\circ}\mid b+a'\geqslant c\} = b \to c,$$

$$\overset{\circ}{a}^{\circ}\geqslant b \to c \implies b+a=b+{}^{\circ}a^{\circ}\geqslant b+(b\to c)=b+\bigwedge\{{}^{\circ}a'^{\circ}\mid b+a'\geqslant c\}$$

$$= \bigwedge\{b+a'\mid b+a'\geqslant c\}\geqslant c,$$

i.e., + - residuation holds, while the argument for the other coimplication is similar.

PROPOSITION 2.5. Let Q be a generalized Girard quantaloid. For all a, a':  $u \rightarrow v$ , b:  $v \rightarrow w$ , c:  $u \rightarrow w$ , c':  $u \rightarrow u'$ , d:  $w \rightarrow v$ , d':  $w \rightarrow u$ , e:  $v \rightarrow u$ , e':  $u' \rightarrow u$ , l, m, n:  $u \rightarrow u$ ,  $\{a_j\} \subseteq Q(u,v)$ ,  $\{b_j\} \subseteq Q(v,w)$ ,  $\{c_j\} \subseteq Q(u,w)$ , the following relations hold:

- (i)  $b \leftarrow e = b \circ e^{\circ}, c \rightarrow b = {}^{\circ}c \circ b \circ T(v, v);$
- (ii)  $b + (b \rightarrow c) \ge c$ ,  $(c \leftarrow a) + a \ge c$ ;
- (iii)  $a \geqslant a'$  implies  $a + e' \geqslant a' + e'$ ,  $b + a \geqslant b + a'$ ;
- (iv)  $c' \leftarrow (b+a) = (c' \leftarrow a) \leftarrow b$ ,  $(a+e') \rightarrow c = e' \rightarrow (a \rightarrow c)$ ,  $e' \rightarrow (e \leftarrow b) = (e' \rightarrow e) \leftarrow b$ ;
- (v)  $c \leftarrow \bigwedge_{j} {}^{\circ}a_{j}^{\circ} = \bigvee_{j} (c \leftarrow a_{j}), (\bigwedge_{j} {}^{\circ}b_{j}^{\circ}) \rightarrow c = \bigvee_{j} (b_{j} \rightarrow c);$
- (vi)  $(\bigvee_j c_j) a = \bigvee_j (c_j a), b \bigvee_j c_j = \bigvee_j (b c_j);$
- (vii)  $\bigwedge_{j} (c_{j} a) \ge (\bigwedge_{j} c_{j}) a, \bigwedge_{j} (b c_{j}) \ge b \bigwedge_{j} c_{j};$
- (viii)  $\bigwedge_{j} (c \leftarrow a_{j}) \geqslant c \leftarrow (\bigvee_{j} a_{j}), \bigwedge_{j} (b_{j} \rightarrow c) \geqslant (\bigvee_{j} b_{j}) \rightarrow c;$
- (ix)  $b + (a \leftarrow c) \ge (b+a) \leftarrow c$ ,  $(d \rightarrow b) + e' \ge d \rightarrow (b+e')$ ;
- $(x) (c \leftarrow a) + (a \leftarrow c') \geqslant c \leftarrow c', (e' \rightarrow e) + (e \rightarrow d') \geqslant e' \rightarrow d';$

(xi) 
$$c \leftarrow a \geqslant (c + e') \leftarrow (a + e'), d' \rightarrow e' \geqslant (a + d') \rightarrow (a + e');$$

- (xii)  $a^{\circ} \lor a^{\circ} = a + (a \rightarrow a') = (a' \leftarrow a) + a;$
- (xiii)  $a \leftarrow m \leq {}^{\circ}a^{\circ}, m \rightarrow e \leq {}^{\circ}e^{\circ};$
- (xiv)  $b \geqslant b'$  implies  $b \leftarrow e \geqslant b' \leftarrow e$  and  $c \rightarrow b \geqslant c \rightarrow b'$ ,  $a \geqslant a'$  implies  $c \leftarrow a \leqslant c \leftarrow a'$  and  $a \rightarrow d \leqslant a' \rightarrow d$ ;
- $(xv)^{\circ}(b \leftarrow e) = {}^{\circ}e \setminus {}^{\circ}b, (c \rightarrow b)^{\circ} = b^{\circ}/c^{\circ};$
- (xvi)  $(m \leftarrow n) \rightarrow m = m \leftarrow (n \rightarrow m) = {}^{\circ}m^{\circ} \wedge {}^{\circ}n^{\circ};$
- (xvii)  $(m \leftarrow n) \rightarrow (m \leftarrow c) = ({}^{\circ}m^{\circ} \wedge n) \leftarrow c,$  $(e \rightarrow m) \leftarrow (n \rightarrow m) = e \rightarrow ({}^{\circ}m^{\circ} \wedge n);$
- (xviii)  $(m \leftarrow l) \leftarrow (({}^{\circ}m^{\circ} \wedge n) \leftarrow l) = (m \leftarrow l) \wedge (m \leftarrow n),$  $(l \rightarrow ({}^{\circ}m^{\circ} \wedge n)) \rightarrow (l \rightarrow m) = (l \rightarrow m) \wedge (n \rightarrow m);$
- (xix)  $m \leftarrow \perp (u, u) = \perp (u, u) \rightarrow m = {}^{\circ}m^{\circ}$  (recall that  $\perp (u, u)$  is the bottom morphism of Q(u, u));
- $(xx) m \leftarrow m = m \rightarrow m = \bot(u, u);$
- (xxi) If  $m \leftarrow n = \bot(u, u)$  or  $n \rightarrow m = \bot(u, u)$ , then  ${}^{\circ}m^{\circ} \leqslant {}^{\circ}n^{\circ}$ ;
- (xxii)  $m \leftarrow n = {}^{\circ}m^{\circ} = n \rightarrow m \iff {}^{\circ}m^{\circ} \wedge {}^{\circ}n^{\circ} = \bot(u, u).$

*Proof.* (i) Since  ${}^{\circ}(b \circ e^{\circ}){}^{\circ}/e^{\circ} \ge b$  always, it follows that  $b \circ e^{\circ} + e \ge b$ , i.e.,  $b \circ e^{\circ} \ge b \leftarrow e$ . Further, by 2.4(iv) we have

$$b \leftarrow e = \bigwedge \{ c^{\circ} : u \rightarrow w \mid c + e \geqslant b \}.$$

But  $c + e \ge b$  iff  $b \circ e^{\circ} \le {}^{\circ}c^{\circ}$ . So

$$b \circ e^{\circ} \leq \bigwedge \{ {}^{\circ}c^{\circ} \mid c + e \geq b \} = b \leftarrow e.$$

The next equality in (i) can be checked similarly.

Relations (ii)—(xv) follow immediately from (i). (They may be considered as a dual case of 1.3.) (xvi) Using (i), 2.2(ii), 1.3(xii), and 1.2(ii), we get

$$(m \leftarrow n) \rightarrow m = {}^{\circ}(m \circ n^{\circ}) \circ m \circ T(u, u) = (n \circ T(u, u)/m) \circ m \circ T(u, u)$$

$$= (n \circ T(u, u) \wedge m) \circ T(u, u) = n \circ T(u, u) \wedge m \circ T(u, u)$$

$$= m \circ T(u, u) \wedge {}^{\circ}n^{\circ} = m \circ (m \setminus {}^{\circ}n^{\circ}) = m \circ ({}^{\circ}n \circ m)^{\circ} = m \leftarrow (n \rightarrow m).$$

(xvii) In view of (iii), (xvi), 1.2(ii), and 1.3(iii) we obtain

$$(m \leftarrow n) \rightarrow (m \leftarrow c) = ((m \leftarrow n) \rightarrow m) \leftarrow c = (m \circ T(u, u) \wedge n \circ T(u, u)) \leftarrow c$$

$$= (m \circ T(u, u) \wedge n) \circ T(u, u) \circ c^{\circ} = (m \circ T(u, u) \wedge n) \circ c^{\circ}$$

$$= (m \circ T(u, u) \wedge n) \leftarrow c,$$

$$(e \rightarrow m) \leftarrow (n \rightarrow m) = e \rightarrow (m \leftarrow (n \rightarrow m)) = e \rightarrow (m \circ T(u, u) \wedge n \circ T(u, u))$$

$$= {\circ}e \circ (m \circ T(u, u) \wedge n) \circ T(u, u) \circ T(u, u) = e \rightarrow (m \circ T(u, u) \wedge n).$$

(xviii) Using (xvii), (xvi) we get

$$(m \leftarrow l) \leftarrow (({}^{\circ}m^{\circ} \wedge n) \leftarrow l) = (m \leftarrow l) \leftarrow ((m \leftarrow n) \rightarrow (m \leftarrow l))$$

$$= (m \leftarrow l) \wedge (m \leftarrow n),$$

$$(l \rightarrow ({}^{\circ}m^{\circ} \wedge n)) \rightarrow (l \rightarrow m) = ((l \rightarrow m) \leftarrow (n \rightarrow m)) \rightarrow (l \rightarrow m)$$

$$= (l \rightarrow m) \wedge (n \rightarrow m).$$

(xix) From (xvi), (xiv), (xiii) we have

$$m \circ T(u, u) = (m \leftarrow m) \rightarrow m \leqslant \bot(u, u) \rightarrow m \leqslant m \circ T(u, u),$$

$$m \circ T(u, u) = m \leftarrow (m \rightarrow m) \leqslant m \leftarrow \bot(u, u) \leqslant m \circ T(u, u),$$

which imply the assertion.

(xx) From (i), (xvi), (xix) we get

$$m \leftarrow m = m \circ m^{\circ} = m \circ (m \circ T(u, u))^{\circ} = m \leftarrow m \circ T(u, u) = m \leftarrow (\bot(u, u) \rightarrow m)$$
$$= m \circ T(u, u) \land \bot(u, u) \circ T(u, u) = \bot(u, u)$$

and

$$m \rightarrow m = {}^{\circ}m \circ m = {}^{\circ}(m \circ T(u, u)) \circ m = m \circ T(u, u) \rightarrow m = (m \leftarrow \bot(u, u)) \rightarrow m = \bot(u, u).$$

(xxi) From (xix), (xvi) we get

$${}^{\circ}m^{\circ} = m \leftarrow \bot(u, u) = m \leftarrow (n \rightarrow m) = {}^{\circ}m^{\circ} \wedge {}^{\circ}n^{\circ}$$
 or  ${}^{\circ}m^{\circ} = \bot(u, u) \rightarrow m = (m \leftarrow n) \rightarrow m = {}^{\circ}m^{\circ} \wedge {}^{\circ}n^{\circ}$ , i.e.,  ${}^{\circ}m^{\circ} \leq {}^{\circ}n^{\circ}$ .

(xxii) Let  $m \leftarrow n = {}^{\circ}m^{\circ}$ . Then  ${}^{\circ}m^{\circ} \wedge {}^{\circ}n^{\circ} = (m \leftarrow n) \rightarrow m = {}^{\circ}m^{\circ} \rightarrow m = m \rightarrow m = \bot(u, u)$ . Let  ${}^{\circ}m^{\circ} \wedge {}^{\circ}n^{\circ} = \bot(u, u)$ . Then  $(m \leftarrow n) \rightarrow m = m \leftarrow (n \rightarrow m) = \bot(u, u)$ , and by (xxi) we have  ${}^{\circ}m^{\circ} \leq m \leftarrow n$ ,  ${}^{\circ}m^{\circ} = n \rightarrow m$ , which imply, by (xiii),  $m \leftarrow n = {}^{\circ}m^{\circ} = n \rightarrow m$ .

We finish this section with an "example" of a generalized Girard quantaloid. Recall that every hom-lattice Q(u, v) of a quantaloid Q obeying 1.2 is a quantale with the product  $(a, b) \mapsto (b \circ T(u, u)) \wedge a$  and with the left-sided unit T(v, u). Let us consider the case of quantaloids where, in addition, every Q(u, v) is such that for all a in Q(u, v),

$$\vec{a} = \vec{a} = a \circ T(u, u), \tag{1}$$

$$a \setminus \bot(v, u) \leqslant T(u, v) \circ \bar{a}^{\dagger}, \quad \bot(v, u)/a \leqslant \bar{a} \circ T(u, v),$$
 (2)

where

$$\bar{a}:=\bigvee\{b\in Q(u,v)\mid (a\circ T(u,u))\wedge b=\bot(v,u)\}$$

and

$$\overline{a}$$
: =  $\bigvee \{b \in Q(u, v) \mid (b \circ T(u, u)) \land a = \bot(v, u)\}$ 

are "complements" of a in Q(u, v). (We still keep the notation  $\bot(v, u)$  for the bottom morphism of Q(u, v).) Moreover, let the equality

$$T(u, v) \circ T(v, u) = T(u, u) \tag{3}$$

always hold.

Notice some analogy between (1) and the Boolean law of "double negation":  $\overline{a}^{\dagger}$  (:=  $\neg \neg a$ ) = a. By 1.3(xiv), Q(u, v) is a Heyting algebra with the operation of "pseudocomplementation"  $\overline{()}$  (=  $\overline{()}$ ), assuming  $T(u, u) = id_u$ . Thanks to (1), this operation is precisely what can be kept as an "authentic complementation," making Q(u, v) into a Boolean algebra. Observe also that conditions (2) are satisfied if, for example, for all  $a, b \in Q(u, v)$ ,  $b \circ T(u, v) \circ a = \bot(v, u)$  implies  $b \circ T(u, u) \wedge a = \bot(v, u)$ .

PROPOSITION 2.6 (cf. [9, Theorem 2.3]). Let Q be a quantaloid satisfying additional Axioms (1)–(3). Then Q is a generalized Girard quantaloid with the dualizing family  $\bot = \{\bot(u,u) \mid u \in Q\}$ .

*Proof*. We first need to verify the following inequalities:

$$a \setminus \bot(v,v) \leqslant T(u,v) \circ \bar{a} \circ T(u,v), \qquad \bot(u,u)/a \leqslant T(u,v) \circ \bar{a} \circ T(u,v). \tag{i}$$

From (2) we have

$$(a \setminus \bot(v,u))/T(v,u) \leq (T(u,v) \circ \overline{a})/T(v,u).$$

But

$$\pm (v, u)/T(v, u) \leqslant (\pm (v, u) \circ T(u, v))/(T(v, u) \circ T(u, v)) \qquad \text{(by 1.3(xi))}$$

$$= \pm (v, v)/T(v, v) \qquad \text{(by (3))}$$

$$\leqslant (\pm (v, v)/T(v, v)) \circ T(v, v) \qquad \text{(by 1.2(i))}$$

$$\leqslant \pm (v, v),$$

so that, by 1.3(iv), 1.3(xi), 1.2(i),

$$a \setminus \bot(v, v) = (a \setminus \bot(v, u))/T(v, u) \leqslant (T(u, v) \circ \bar{a})/T(v, u)$$

$$\leqslant (T(u, v) \circ \bar{a} \circ T(u, v))/(T(v, u) \circ T(u, v))$$

$$= (T(u, v) \circ \bar{a} \circ T(u, v))/T(v, v) \quad \text{(by (3))}$$

$$\leqslant (T(u, v) \circ \bar{a} \circ T(u, v)/T(v, v)) \circ T(v, v) \leqslant T(u, v) \circ \bar{a} \circ T(u, v),$$

while the argument for the other inequality of (i) is similar. Next,  $a \circ T(u, v) \circ \bar{a} \leqslant \bar{a}$ ,  $\bar{a} \circ T(u, v) \circ a \leqslant a$  imply, by 1.1(iii),  $T(u, v) \circ \bar{a} \leqslant a \setminus \bar{a}$  and  $T(u, v) \circ a \leqslant \bar{a} \setminus a$ . Using the latter pair of inequalities, we proceed as follows:

$$a \circ T(u, v) \circ \vec{a} \leq a \circ (a \setminus \vec{a}) = a \circ T(u, u) \wedge \vec{a} = \bot(v, u)$$
 (by 1.3(xii)),  

$$\vec{a} \circ T(u, v) \circ a \leq \vec{a} \circ (\vec{a} \setminus a) = (\vec{a} \circ T(u, u)) \wedge a = \bot(v, u),$$

$$a \circ T(u, v) \circ \vec{a} \circ T(u, v) \leq \bot(v, u) \circ T(u, v) = \bot(v, v),$$

$$T(u, v) \circ \vec{a} \circ T(u, v) \circ a \leq T(u, v) \circ \bot(v, u) = \bot(u, u).$$

So in view of 1.1(iii), (2), and (i) we obtain that for every  $a \in Q(u, v)$ ,

$$a \setminus \bot(v, u) = T(u, v) \circ \vec{a}, \qquad \bot(v, u)/a = \overline{a} \circ T(u, v),$$

$$a \setminus \bot(v, v) = T(u, v) \circ \vec{a} \circ T(u, v), \qquad \bot(u, u)/a = T(u, v) \circ \overline{a} \circ T(u, v). \tag{ii}$$

Now we are ready to verify that ⊥ is a dualizing family. We have

But

Hence

But

$$a \circ T(u, u) \leq (a \circ T(u, v))/T(u, v) \leq (a \circ T(u, v) \circ T(v, u))/(T(u, v) \circ T(v, u))$$

$$= (a \circ T(u, u))/T(u, u) \quad \text{(by (3))}$$

$$\leq (a \circ T(u, u)/T(u, u)) \circ T(u, u) \leq a \circ T(u, u).$$

Therefore

$$\perp (v, v)/(a \setminus \perp (v, v)) = a \circ T(u, u).$$

Further.

$$(\pm(u,u)/a) \setminus \pm(u,u) = (T(u,v) \circ \overline{a} \circ T(u,v)) \setminus \pm(u,u) \quad \text{(by (ii))}$$

$$= (\overline{a} \circ T(u,v)) \setminus (T(u,v) \setminus \pm(u,u)) \quad \text{(by 1.3(iv))}$$

$$= (\overline{a} \circ T(u,v)) \setminus \pm(v,u) = T(u,v) \setminus (\overline{a} \setminus \pm(v,u))$$

$$= T(u,v) \setminus (T(u,v) \circ \overline{a}) \quad \text{(by (ii))}$$

$$= T(u,v) \setminus (T(u,v) \circ a \circ T(u,u)) \quad \text{(by (1))}.$$

Since

$$a \circ T(u, u) \leqslant T(u, v) \setminus (T(u, v) \circ a \circ T(u, u)) \leqslant (T(v, u) \circ T(u, v)) \setminus (T(v, u) \circ T(u, v) \circ a \circ T(u, u))$$

$$= T(v, v) \setminus (T(v, v) \circ a \circ T(u, u)) \qquad \text{(by (3))}$$

$$= T(v, v) \circ (T(v, v) \setminus a \circ T(u, u)) \qquad \text{(by 1.3(iii))}$$

$$\leqslant a \circ T(u, u),$$

we conclude that

$$(\bot(u,u)/a)\setminus\bot(u,u)=a\circ T(u,u),$$

as claimed.

#### 3. QUANTAL CATEGORIES

A quantaloid Q can also be viewed as a locally ordered base bicategory in the sense of R. H. Street, so it is natural to consider categories with hom-sets enriched in Q (as K. I. Rosenthal [9, 10] does). However, we shall go a little bit outside the usual context of enriched category theory, since our quantaloid is not very ordinary and convenient in the categorical sense.

We thus turn to a "non-commutative many-object" version of U. Höhle's [4] concept of set structured with an "existence predicate."

Definition 3.1. Let Q be a quantaloid defined in 1.2. A Q-category (or quantal category) is a triple  $(X, \widetilde{()}, E)$  (below we shall suppress writing  $\widetilde{()}$  explicitly and these data will take the form (X, E)), where X is a (small) set,  $\widetilde{()}$ :  $X \to \operatorname{Obj}(Q)$  is a map assigning to  $x \in X$  an object  $\widetilde{x} \in Q$ ,  $E: X \times X \to \operatorname{Mor}(Q)$  is a map ("enrichment"), which assigns to every pair  $x, x' \in X$  a morphism E(x, x'):  $\widetilde{x'} \to \widetilde{x}$  in Q such that

- (i)  $E(x, x') \leq (E(x, x) \circ T(\tilde{x}, \tilde{x'})) \wedge (T(\tilde{x}, \tilde{x'}) \circ E(x', x'))$  (strictness),
- (ii)  $E(x, x') \circ (E(x', x') \setminus E(x', x'')) \leqslant E(x', x'')$  (transitivity)

for all x, x' x'' in X. On any Q-category (X, E) we define extents of existence and (ordinary) equivalence by (iii) Ex = E(x, x), and

(iv)  $x \approx x' \iff \tilde{x} = \tilde{x'}, \ Ex \vee Ex' \leqslant E(x, x') \wedge E(x', x).$ 

A Q-category (X, E) is separated iff (X, E) satisfies the additional axiom:

(v)  $x \approx x'$  implies x = x' (separation).

We can associate with every Q-category (X, E) a separated Q-category  $(\hat{X}, \hat{E})$  in usual way:

(vi)  $\hat{X} = X/\approx$ ,  $\hat{E}(\hat{x}, \hat{x'}) = E(y, y')$ , where  $y \in \hat{x}$ ,  $y' \in \hat{x'}$ .

Axioms 3.1 generalize to the "many-object" case those for Q-sets (cf. [3, 4]). Here are some easy consequences of the definitions from 1.3(iii), (xiii), 1.2(ii).

PROPOSITION 3.2. Let Q be a quantaloid and (X, E) be a Q-category. For all  $x, x', x'' \in X$ ,

- (i)  $E(x, x') \circ (Ex' \setminus E(x', x'')) = (E(x, x') \circ T(\tilde{x'}, \tilde{x'})/Ex') \circ E(x', x''),$
- (ii)  $(E(x, x') \circ T(\tilde{x'}, \tilde{x'})) \wedge (T(\tilde{x}, \tilde{x'}) \circ Ex') = E(x, x'),$
- (iii)  $E(x, x') = Ex \circ (Ex \setminus E(x, x')) = (E(x, x') \circ T(x', x')/Ex') \circ Ex'$ .

Let us compare quantal categories over Q obeying Axioms 3.1 with quantal categories in the sense of [9].

Definition 3.3 ([9, Definition 3.1]). Let Q be a quantaloid containing two-sided units but in general unsatisfying Axioms 1.2. A set X is a Q-category iff it comes equipped with the following data:

(1) a map () assigning to  $x \in X$  an object  $\tilde{x} \in Q$ ;

- (2) an enrichment which assigns to every pair  $x, x' \in X$  a morphism E(x, x'):  $\tilde{x'} \to \tilde{x}$  in Q such that
  - (i)  $id_{\bar{x}} \leq E(x, x)$  for all  $x \in X$ ,
  - (ii)  $E(x, x') \circ E(x', x'') \leqslant E(x, x'')$  for all  $x, x', x'' \in X$ .

PROPOSITION 3.4. Let Q be a quantaloid having two-sided units and satisfying Axioms 1.2. Then every Q-category in the sense of 3.3 is also a Q-category in the sense of 3.1.

*Proof.* Let (X, E) be a Q-category in the sense of 3.3. Then 3.3(i) implies  $E(x, x) = T(\tilde{x}, \tilde{x})$ , since, by 1.3(iii),  $id_{\tilde{x}} = T(\tilde{x}, \tilde{x})$ . Consequently, in view of 1.3(iii), (xii) we have

$$E(x',x') \setminus E(x',x'') = T(\tilde{x'},\tilde{x'}) \setminus E(x',x'') = T(\tilde{x'},\tilde{x'}) \circ (T(\tilde{x'},\tilde{x'}) \setminus E(x',x''))$$

$$= T(\tilde{x'},\tilde{x'}) \circ T(\tilde{x'},\tilde{x''}) \wedge E(x',x'') = T(\tilde{x'},\tilde{x''}) \wedge E(x',x'') = E(x',x'').$$

Hence (X, E) is a Q-category in the sense of 3.1.

It is interesting to see that in a certain sense these definitions lead to the same things.

PROPOSITION 3.5. Let (X, E) be a Q-category in the sense of 3.1, in the setting of the preceding proposition. Then (X, E') is a Q-category in the sense of 3.3, where E' is defined by  $E'(x, x') = Ex \setminus E(x, x')$ .

Proof. Because of 3.2(iii) and 1.1(iii), we have that inequality 3.1(ii) takes the form

$$(Ex \setminus E(x, x')) \circ (Ex' \setminus E(x', x'')) \leqslant Ex \setminus E(x, x''),$$

i.e., E' obeys 3.3(ii), while 3.3(i) follows immediately from 1.1(iii):

$$Ex \circ id_{\tilde{r}} \leqslant Ex \implies id_{\tilde{r}} \leqslant Ex \setminus Ex = E'(x,x).$$

Note that below we shall not consider Q-categories in the sense of 3.3, rather Q-category will always refer to 3.1.

Definition 3.6. If (X, E), (Y, F) are separated Q-categories (in the sense of 3.1), a Q-functor  $f: (X, E) \rightarrow (Y, F)$  of Q-categories is a map  $f: X \rightarrow Y$  satisfying

- (i)  $\tilde{x} = f(x)$ , Ex = Ff(x) (strictness),
- (ii)  $E(x, x') \leq F(f(x), f(x'))$  (preservation of structure)

for all x, x' in X.

Obviously separated Q-categories and Q-functors together with the usual composition of maps constitute some category Q-CAT.

#### 4. A MONAD ON Q-CAT

In this section we construct a monad on the category Q-CAT. This can be done by extending Höhle's [4] notion of singletons.

Definition 4.1. For a Q-category (X, E), a singleton is a pair s of maps  $\langle s, s \rangle$ ,  $\langle s, s \rangle$ :  $X \to \text{Mor}(Q)$  assigning to every  $x \in X$  morphisms  $\langle x, s \rangle$ :  $\tilde{s} \to \tilde{x}$  and  $\langle s, x \rangle$ :  $\tilde{x} \to \tilde{s}$  (for some object  $\tilde{s} \in Q$ ) and satisfying the following conditions:

- (i)  $\llbracket s \rrbracket \circ T(\tilde{s}, \tilde{s}) = \bigvee_{x \in X} \langle s, x \rangle \circ T(\tilde{x}, \tilde{s}),$
- (ii)  $E(x, x') \circ (Ex' \setminus \langle x', s \rangle) \leq \langle x, s \rangle$ ,  $\langle s, x \rangle \circ (Ex \setminus E(x, x')) \leq \langle s, x' \rangle$  (extensionality),
- (iii)  $\langle x, s \rangle \circ (\llbracket s \rrbracket \setminus \langle s, x' \rangle) \leqslant E(x, x')$  (singleton condition),
- $(iv) \langle x, s \rangle \leqslant T(\tilde{x}, \tilde{s}) \circ \llbracket s \rrbracket, \langle s, x \rangle \leqslant \llbracket s \rrbracket \circ T(\tilde{s}, \tilde{x}), \llbracket s \rrbracket \setminus \llbracket s \rrbracket \leqslant \langle x, s \rangle \setminus \langle x, s \rangle$

for all x, x' in X, where  $[s]: = \bigvee_{x \in X} T(\bar{s}, \tilde{x}) \circ \langle x, s \rangle$ .

Here are some consequences of the definition.

PROPOSITION 4.2. Let s be a singleton of (X, E). For all  $x, x', x'' \in X$ , (i)  $\langle x, s \rangle \circ (\llbracket s \rrbracket \setminus \langle s, x' \rangle) = (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \langle s, x' \rangle$ ,

(ii)  $\langle x, s \rangle \leq Ex \circ T(\tilde{x}, \tilde{s}), \langle s, x \rangle \leq T(\tilde{s}, \tilde{x}) \circ Ex$  (strictness),

(iii)  $E(x, x') \circ (Ex' \setminus \langle x', s \rangle) = (E(x, x') \circ T(x', x') / Ex') \circ \langle x', s \rangle$ ,

(iv)  $\langle s, x \rangle \circ (Ex \setminus E(x, x')) = (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x})/Ex) \circ E(x, x'),$ 

 $(\mathsf{v})\ (\langle s,x\rangle \circ T(\tilde{x},\tilde{x})) \wedge (T(\tilde{s},\tilde{x})\circ Ex) = \langle s,x\rangle,$ 

(vi)  $(\langle x, s \rangle \circ T(\tilde{s}, \tilde{s})) \wedge (T(\tilde{x}, \tilde{s}) \circ \llbracket s \rrbracket) = \langle x, s \rangle$ .

Proof. (vi) We have

$$\langle x, s \rangle \leqslant \langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) \wedge T(\tilde{x}, \tilde{s}) \circ \llbracket s \rrbracket$$
 (by 4.1(iv), 1.2(i))  

$$= (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \llbracket s \rrbracket$$
 (by 1.3(xiii))  

$$\leqslant (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \llbracket s \rrbracket \circ (\langle x, s \rangle \setminus \langle x, s \rangle)$$
 (by 4.1(iv))  

$$\leqslant \langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) \circ (\langle x, s \rangle \setminus \langle x, s \rangle)$$
  

$$= \langle x, s \rangle \circ (\langle x, s \rangle \setminus \langle x, s \rangle) = \langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) \wedge \langle x, s \rangle = \langle x, s \rangle,$$

which implies (vi).

(i) The following holds:

$$\langle x, s \rangle \circ (\llbracket s \rrbracket \setminus \langle s, x' \rangle) = (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) \wedge T(\tilde{x}, \tilde{s}) \circ \llbracket s \rrbracket) \circ (\llbracket s \rrbracket \setminus \langle s, x' \rangle)$$

$$= (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ (\llbracket s \rrbracket \circ (\llbracket s \rrbracket \setminus \langle s, x' \rangle)$$

$$= (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ (\llbracket s \rrbracket \circ T(\tilde{s}, \tilde{x'}) \wedge \langle s, x' \rangle)$$

$$= (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \langle s, x' \rangle$$

$$(\text{by } 1.3(\text{xiii}))$$

$$= (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \langle s, x' \rangle$$

$$(\text{by } 4.1(\text{iv})).$$

(ii) In view of (vi), 1.2(i), 1.3(xiii), 4.1(i), (iii), (iv) and 3.1(i) we obtain

$$\begin{aligned} \langle x, s \rangle &= \langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) \wedge T(\tilde{x}, \tilde{s}) \circ \llbracket s \rrbracket \leqslant (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \llbracket s \rrbracket \circ T(\tilde{s}, \tilde{s}) \\ &= (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \bigvee_{x' \in X} \langle s, x' \rangle \circ T(\tilde{x'}, \tilde{s}) \\ &\leqslant \bigvee_{x' \in X} E(x, x') \circ T(\tilde{x'}, \tilde{s}) \leqslant \bigvee_{x' \in X} Ex \circ T(\tilde{x}, \tilde{x'}) \circ T(\tilde{x'}, \tilde{s}) = Ex \circ T(\tilde{x}, \tilde{s}) \end{aligned}$$

and

$$\begin{aligned} \langle s, x \rangle &= \llbracket s \rrbracket \circ T(\tilde{s}, \tilde{x}) \wedge \langle s, x \rangle = \llbracket s \rrbracket \circ (\llbracket s \rrbracket \setminus \langle s, x \rangle) \\ &= \bigvee_{x' \in X} T(\tilde{s}, \tilde{x'}) \circ \langle x', s \rangle \circ (\llbracket s \rrbracket \setminus \langle s, x \rangle) \leqslant \bigvee_{x' \in X} T(\tilde{s}, \tilde{x'}) \circ E(x', x) \\ &\leqslant \bigvee_{x' \in X} T(\tilde{s}, \tilde{x'}) \circ T(\tilde{x'}, \tilde{x}) \circ Ex = T(\tilde{s}, \tilde{x}) \circ Ex. \end{aligned}$$

(iii) The proof is based on (ii), 3.2(iii) and 1.3(xiii):

$$E(x, x') \circ (Ex' \setminus \langle x', s \rangle) = (E(x, x') \circ T(\tilde{x'}, \tilde{x'})/Ex') \circ Ex' \circ (Ex' \setminus \langle x', s \rangle)$$

$$= (E(x, x') \circ T(\tilde{x'}, \tilde{x'})/Ex') \circ (Ex' \circ T(\tilde{x'}, \tilde{s}) \wedge \langle x', s \rangle)$$

$$= (E(x, x') \circ T(\tilde{x'}, \tilde{x'})/Ex') \circ \langle x', s \rangle).$$

(iv) By virtue of (ii), 1.2(ii), 1.3(xiii) and 3.2(iii) we have that

$$\langle s, x \rangle \circ (Ex \setminus E(x, x')) = \langle s, x \rangle \circ T(\tilde{x}, \tilde{x}) \circ (Ex \setminus E(x, x'))$$

$$= (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x}) \wedge T(\tilde{s}, \tilde{x}) \circ Ex \circ T(\tilde{x}, \tilde{x})) \circ (Ex \setminus E(x, x'))$$

$$= (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x}) \wedge T(\tilde{s}, \tilde{x}) \circ Ex) \circ T(\tilde{x}, \tilde{x}) \circ (Ex \setminus E(x, x'))$$

$$= (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x})/Ex) \circ Ex \circ (Ex \setminus E(x, x'))$$

$$= (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x})/Ex) \circ E(x, x').$$

### (v) The equality follows from

$$\langle s, x \rangle = \langle s, x \rangle \wedge T(\tilde{s}, \tilde{x}) \circ Ex \qquad \text{(by (ii))}$$

$$\leq \langle s, x \rangle \circ T(\tilde{x}, \tilde{x}) \wedge T(\tilde{s}, \tilde{x}) \circ Ex$$

$$= (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x})/Ex) \circ Ex \qquad \text{(by 1.3(xiii))}$$

$$\leq \langle s, x \rangle \qquad \text{(by (iv), 4.1(ii))}.$$

PROPOSITION 4.3. Given a quantaloid Q, a Q-category (X, E), an element  $x \in X$ , an object  $u \in \text{Obj}(Q)$  with  $T(\tilde{x}, u) \circ T(u, \tilde{x}) = T(\tilde{x}, \tilde{x})$ , and a morphism  $a: u \to \tilde{x}$  in  $Q(u, \tilde{x})$ , we produce singletons  $E^x$  and  $E^{(x,a)}$  (of (X, E)) by putting

(i)  $\langle x', E^x \rangle = E(x', x), \langle E^x, x' \rangle = E(x, x')$  (with  $\llbracket E^x \rrbracket = Ex$ );

(ii)  $\langle x', E^{(x,a)} \rangle = E(x', x) \circ (Ex \setminus a) = (E(x', x) \circ T(\tilde{x}, \tilde{x})/Ex) \circ ((Ex \circ T(\tilde{x}, u)) \wedge a) \text{ (by 3.2(iii), 1.3(xii))}$ =  $(E(x', x) \circ T(\tilde{x}, \tilde{x})/Ex) \circ T(\tilde{x}, u) \circ [E^{(x,a)}]$ ,

$$\begin{split} \left\langle E^{(x,a)}, x' \right\rangle &= T(u, \tilde{x}) \circ (a \circ T(u, \tilde{x})) / Ex) \circ E(x, x') \\ &\quad (= T(u, \tilde{x}) \circ ((a \circ T(u, \tilde{x})) \wedge Ex) \circ (Ex \setminus E(x, x')) \quad \text{(by 3.2(iii), 1.3(xiii))} \\ &= T(u, \tilde{x}) \circ ((Ex \circ T(\tilde{x}, u)) \wedge a) \circ T(u, \tilde{x}) \circ (Ex \setminus E(x, x')) \quad \text{(by 1.2(ii))} \\ &= \llbracket E^{(x,a)} \rrbracket \circ T(u, \tilde{x}) \circ (Ex \setminus E(x, x')) \end{split}$$

(with  $\llbracket E^{(x,a)} \rrbracket = T(u, \tilde{x}) \circ ((Ex \circ T(\tilde{x}, u)) \wedge a)$ .

PROPOSITION 4.4. On the set S(X, E) of all singletons of (X, E), there is a Q-category structure given by the map  $\widetilde{()}$ :  $S(X, E) \to \operatorname{Obj}(Q)$  assigning to every  $s \in S(X, E)$  the object  $\widetilde{s} \in Q$  (defined in 4.1) and by the enrichment  $[\![ , ]\!]$ :  $S(X, E) \times S(X, E) \to \operatorname{Mor}(Q)$ :

$$\llbracket s, s' \rrbracket := \left( \llbracket s \rrbracket \circ \bigwedge_{s \in X} \langle x, s \rangle \setminus \langle x, s' \rangle \right) \wedge \left( \left( \bigwedge_{x \in X} \langle s, x \rangle / \langle s', x \rangle \right) \circ \llbracket s' \rrbracket \right).$$

Moreover, [s, s] = [s].

*Proof.* The last equality follows immediately from 4.1(iv), 1.3(iii) and the fact that  $\langle s, x \rangle / \langle s, x \rangle = T(\tilde{s}, \tilde{s})$ . Now the strictness 3.1(i) of [, ] is evident. Let us prove the transitivity of [, ]. We have that

$$\begin{bmatrix} [s,s'] \circ ([s',s'] \setminus [s',s'']) = \left( \left( [s] \circ \bigwedge_{x \in X} \langle x,s \rangle \setminus \langle x,s' \rangle \right) \\
\wedge \left( \bigwedge_{x \in X} \langle s,x \rangle / \langle s',x \rangle \right) \circ [s'] \right) \circ \left( [s'] \setminus \left( [s'] \circ \bigwedge_{x \in X} \langle x,s' \rangle \mid \langle x,s'' \rangle \right) \wedge \left( \bigwedge_{x \in X} \langle s',x \rangle / \langle s'',x \rangle \right) \circ [s'] \right) \\
\leq \left( \left( [s] \circ \bigwedge_{x \in X} \langle x,s \rangle \setminus \langle x,s' \rangle \right) \circ T(\tilde{s},\tilde{s}') \wedge T(\tilde{s},\tilde{s}') \circ [s'] \right) \circ \left( [s'] \setminus [s'] \circ \bigwedge_{x \in X} \langle x,s' \rangle \setminus \langle x,s'' \rangle \right) \\
= \left( \left( [s] \circ \bigwedge_{x \in X} \langle x,s \rangle \setminus \langle x,s' \rangle \right) \circ T(\tilde{s},\tilde{s}') / [s'] \right) \circ [s'] \circ \left( [s'] \setminus [s'] \circ \bigwedge_{x \in X} \langle x,s' \rangle \setminus \langle x,s'' \rangle \right) \quad \text{(by 1.3(xiii))} \\
\leq [s] \circ \left( \bigwedge_{x \in X} \langle x,s \rangle \setminus \langle x,s' \rangle \right) \circ \left( \langle x,s' \rangle \setminus \langle x,s'' \rangle \right) \leq [s] \circ \bigwedge_{x \in X} \langle x,s \rangle \setminus \langle x,s'' \rangle \quad \text{(by 1.3(x))}$$

and that

$$\begin{aligned}
& [s, s'] \circ ([s', s'] \setminus [s', s'']) \leqslant \left( \bigwedge_{x \in X} \langle s, x \rangle / \langle s', x \rangle \right) \circ [s'] \circ \left( [s'] \setminus \left( \bigwedge_{x \in X} \langle s', x \rangle / \langle s'', x \rangle \right) \circ [s''] \right) \\
& \leqslant \left( \bigwedge_{x \in X} \langle s, x \rangle / \langle s', x \rangle \right) \circ \left( \bigwedge_{x \in X} \langle s', x \rangle / \langle s'', x \rangle \right) \circ [s''] \quad \text{(by 1.3(i))} \\
& \leqslant \left( \bigwedge_{x \in X} \langle \langle s, x \rangle / \langle s'', x \rangle \right) \circ [s''] \quad \text{(by 1.3(x))}.
\end{aligned}$$

These imply the transitivity of [, ].

Proposition 4.5.

(i)  $\Sigma(X, E) := (S(X, E), [[,]])$  is a separated Q-category;

(ii) 
$$\llbracket E^x, s \rrbracket = \langle x, s \rangle$$
 and  $\llbracket s, E^x \rrbracket = \langle s, x \rangle$  for each  $(x, s)$  in  $X \times S(X, E)$ .

*Proof.* (i) We need to verify the following implication:

$$\tilde{s} = \tilde{s'}, \quad [s] \lor [s'] \leqslant [s, s'] \land [s', s] \implies s = s'.$$

Thus, to establish the separation it suffices to show that if

$$\tilde{s} = \tilde{s'}, \qquad [\![s]\!] \leqslant [\![s]\!] \circ \bigwedge_{x \in X} \langle x, s \rangle \setminus \langle x, s' \rangle, 
[\![s']\!] \leqslant \left( \bigwedge_{x \in X} \langle s', x \rangle / \langle s, x \rangle \right) \circ [\![s]\!], 
[\![s']\!] \leqslant [\![s']\!] \circ \bigwedge_{x \in X} \langle x, s' \rangle \setminus \langle x, s \rangle, 
[\![s]\!] \leqslant \left( \bigwedge_{x \in X} \langle s, x \rangle / \langle s', x \rangle \right) \circ [\![s']\!],$$

then s = s', and this is so since these datum lead to the following relations:

$$\langle x, s \rangle = (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \llbracket s \rrbracket \qquad \text{(by 4.2(vi), 1.3(xiii))}$$

$$\leqslant (\langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \llbracket s \rrbracket \circ \bigwedge_{x \in X} \langle x, s \rangle \setminus \langle x, s' \rangle$$

$$\leqslant \langle x, s \rangle \circ T(\tilde{s}, \tilde{s}) \circ \bigwedge_{x \in X} \langle x, s \rangle \setminus \langle x, s' \rangle$$

$$\leqslant \langle x, s \rangle \circ (\langle x, s \rangle \setminus \langle x, s' \rangle) \leqslant \langle x, s' \rangle,$$

$$\langle s, x \rangle = \llbracket s \rrbracket \circ T(\tilde{s}, \tilde{x}) \wedge \langle s, x \rangle \qquad \text{(by 4.1(iv))}$$

$$= \llbracket s \rrbracket \circ (\llbracket s \rrbracket \setminus \langle s, x \rangle) \qquad \text{(by 1.3(xiii))}$$

$$\leqslant \left( \bigwedge_{x \in X} \langle s', x \rangle / \langle s, x \rangle \right) \circ \llbracket s \rrbracket \circ (\llbracket s \rrbracket \setminus \langle s, x \rangle)$$

$$\leqslant (\langle s', x \rangle / \langle s, x \rangle) \circ \langle s, x \rangle \leqslant \langle s', x \rangle \qquad \text{(by 1.3(i))}.$$

 $\langle x, s' \rangle \leqslant \langle x, s \rangle$  and  $\langle s', x \rangle \leqslant \langle s, x \rangle$  (interchanging the role of s and s').

(ii) Let s be in  $\Sigma(X, E)$  and let  $x \in X$ . Then

$$\llbracket E^x, s \rrbracket = \left( Ex \circ \bigwedge_{x' \in X} E(x', x) \setminus \langle x', s \rangle \right) \wedge \left( \bigwedge_{x' \in X} E(x, x') / \langle s, x' \rangle \right) \circ \llbracket s \rrbracket.$$

Observe that

$$\langle x, s \rangle \leqslant Ex \circ (Ex \setminus \langle x, s \rangle) \leqslant Ex \circ \bigwedge_{x' \in X} E(x', x) \setminus \langle x', s \rangle \qquad \text{(by 4.1(iii), 1.1(iii))}$$
$$\leqslant Ex \circ (Ex \setminus \langle x, s \rangle) \leqslant \langle x, s \rangle,$$

which imply

$$Ex \circ \bigwedge_{x' \in X} E(x', x) \setminus \langle x', s \rangle = \langle x, s \rangle.$$

Noting still more that

$$\left(\bigwedge_{x'\in X} E(x,x')/\langle s,x'\rangle\right) \circ \llbracket s\rrbracket \geqslant ((\langle x,s\rangle \circ T(\tilde{s},\tilde{s})/\llbracket s\rrbracket) \circ \llbracket s\rrbracket \qquad \text{(by 4.2(i), 4.1(iii), 1.1(iii))}$$
$$\geqslant \langle x,s\rangle \circ T(\tilde{s},\tilde{s}) \geqslant (\langle x,s\rangle, s\rangle, s\rangle$$

we prove the left equality of (ii). Similarly,

$$\llbracket s, E^x \rrbracket = \left( \llbracket s \rrbracket \circ \bigwedge_{x' \in X} \langle x', s \rangle \setminus E(x', x) \right) \wedge \left( \bigwedge_{x' \in X} \langle s, x' \rangle / E(x, x') \right) \circ Ex = \langle s, x \rangle,$$

since 4.1(iii), 1.1(iii), 1.3(i) imply

$$\langle s, x \rangle \leqslant \llbracket s \rrbracket \circ \bigwedge_{r' \in X} \langle x', s \rangle \setminus E(x', x)$$

and since the following relations hold:

$$\langle s, x \rangle \leqslant (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x})/Ex) \circ Ex \leqslant \left( \bigwedge_{x' \in X} \langle s, x' \rangle / E(x, x') \right) \circ Ex \qquad \text{(by 4.2(iv), 4.1(ii), 1.1(iii))}$$

$$\leqslant (\langle s, x \rangle / Ex) \circ Ex \leqslant \langle s, x \rangle.$$

Now we are going to define a monad on the category Q-CAT. We begin with a series of lemmas, the straightforward proofs of which will be omitted for reasons of space. (Those proofs parallel the respective proofs given in [3].)

LEMMA 4.6. Let  $f: (X, E) \to (Y, F)$  be a morphism in Q-CAT. Then the map  $\Sigma(f): \Sigma(X, E) \to \Sigma(Y, F)$  assigning to every  $s \in S(X, E)$  the singleton  $(\Sigma(f))(s)$  in S(Y, F) defined by

$$\langle y, (\Sigma(f))(s) \rangle := \bigvee_{x \in X} F(y, f(x)) \circ (Ex \setminus \langle x, s \rangle) \Big( = \bigvee_{x \in X} (F(y, f(x)) \circ T(\tilde{x}, \tilde{x}) / Ex) \circ \langle x, s \rangle$$

$$(by 3.2(iii), 3.6(i), 1.3(xiii), 4.2(ii)) \Big),$$

$$\langle (\Sigma(f))(s), y \rangle := \bigvee_{x \in X} \langle s, x \rangle \circ (Ex \setminus F(f(x), y)) (= \bigvee_{x \in X} (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x}) / Ex) \circ F(f(x), y)$$

$$(by 4.2(v), 1.3(xiii), 3.6(i), 3.2(iii)))$$

is a morphism in Q-CAT.

LEMMA 4.7. Let  $\Sigma$  be an object map which associates with every separate Q-category (X, E) the singleton space  $\Sigma(X, E)$  and a morphism map which associates with every morphism  $f: (X, E) \to (Y, F)$  in Q-CAT the morphism  $\Sigma(f): \Sigma(X, E) \to \Sigma(Y, F)$  defined in 4.6. Then  $\Sigma$  is a (covariant) functor in Q-CAT.

LEMMA 4.8. If  $\sigma(s) = (\langle s, \sigma \rangle: \tilde{\sigma} \to \tilde{s}, \langle \sigma, s \rangle: \tilde{s} \to \tilde{\sigma}), s \in S(X, E), is a singleton of <math>\Sigma(X, E),$  then  $\sigma(E^x) = (\langle E^x, \sigma \rangle: \tilde{\sigma} \to \tilde{x}, \langle \sigma, E^x \rangle: \tilde{x} \to \tilde{\sigma}), x \in X,$  is a singleton of (X, E) such that

$$\llbracket \sigma(E^{()}) \rrbracket_{S(X,E)} = \llbracket \sigma \rrbracket_{S(\Sigma(X,E))}.$$

(Recall that  $E^x = (E(\cdot, x), E(x, \cdot))$ .)

The following lemma completes the list of preparatory results.

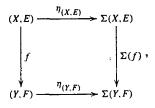
LEMMA 4.9. There are natural transforms  $\eta$ : Id  $\to \Sigma$  and  $\mu$ :  $\Sigma^2 \to \Sigma$  given by

$$\eta_{(X,E)}(x) = E^x, \qquad \mu_{(X,E)}(\sigma) = \sigma(E^{()})$$

for all objects (X, E) of Q-CAT, where Id is the identity functor on Q-CAT, and  $\Sigma$  is the functor produced in 4.6.

*Proof.* First, we must be sure that the maps  $\eta_{(X,E)}$  and  $\mu_{(X,E)}$  are morphisms of Q-CAT for any separated Q-category (X, E). Axioms 3.6(i), (ii) for  $\eta_{(X,E)}$  follow immediately from 4.5(ii), while the strictness of  $\mu_{(X,E)}$  is already met in 4.8(i). In order to verify 3.6(ii) for  $\mu_{(X,E)}$  we proceed as follows:

Next, the diagrams



$$\Sigma^{2}(X,E) \xrightarrow{\mu_{(X,E)}} \Sigma(X,E)$$

$$\downarrow \Sigma^{2}(f) \qquad \qquad \downarrow \Sigma(f)$$

$$\Sigma^{2}(Y,F) \xrightarrow{\mu_{(Y,F)}} \Sigma(Y,F)$$

always commute: we verify the commutativity of the first diagram as follows:

$$\left\langle y, \left( \Sigma(f) \right) \left( \eta_{(X,E)}(x) \right) \right\rangle = \bigvee_{x' \in X} \left( F(y, f(x')) \circ T(\tilde{x'}, \tilde{x'}) / Ex' \right) \circ E(x', x)$$

$$\leqslant \bigvee_{x' \in X} \left( F(y, f(x')) \circ T(\tilde{x'}, \tilde{x'}) / Ff(x') \right) \circ F(f(x'), f(x))$$
 (by 3.6)
$$\leqslant F(y, f(x)) \quad \text{(by 3.2(i), 3.1(ii))}$$

$$\left( = \left\langle y, \eta_{(X,F)}(f(x)) \right\rangle \right)$$

$$= \left( F(y, f(x)) \circ T(\widetilde{f(x)}, \widetilde{f(x)}) / Ff(x) \right) \circ Ff(x) \quad \text{(by 3.2(iii))}$$

$$= \left( F(y, f(x)) \circ T(\tilde{x}, \tilde{x}) / Ex \right) \circ Ex \quad \text{(3.6(i))}$$

$$\leqslant \bigvee_{x' \in X} \left( F(y, f(x')) \circ T(\tilde{x'}, \tilde{x'}) / Ex' \right) \circ E(x', x)$$

$$= \left\langle y, \left( \Sigma(f) \right) (\eta_{(X,E)}(x)) \right\rangle,$$

$$\left\langle \left( \Sigma(f) \right) (\eta_{(X,E)}(x)), y \right\rangle = \left\langle \eta_{(Y,F)}(f(x)), y \right\rangle \quad \text{(similarly)}.$$

Now we establish the commutativity of the second one: if  $\sigma$  is in  $S(\Sigma(X, E))$  and y in Y, then

$$\left\langle y, \mu_{(Y,F)} \Big( \Big( \Sigma \big( \Sigma \big( f \big) \Big) \Big) (\sigma) \Big) \right\rangle = \left\langle F^{Y}, \Big( \Sigma \big( \Sigma \big( f \big) \Big) \Big) (\sigma) \right\rangle \qquad \text{(by definition of } \mu)$$

$$= \bigvee_{s \in S(X,E)} \left[ F^{Y}, \Big( \Sigma \big( f \big) \Big) (s) \right]_{S(Y,F)} \circ (\llbracket s \rrbracket \setminus \langle s, \sigma \rangle) \qquad \text{(by 4.6)}$$

$$= \bigvee_{s \in S(X,E)} \left\langle y, \Big( \Sigma \big( f \big) \Big) (s) \right\rangle \circ (\llbracket s \rrbracket \setminus \langle s, \sigma \rangle) \qquad \text{(by 4.5(ii))}$$

$$= \bigvee_{s \in S(X,E)} \left( \bigvee_{x \in X} (F(y,f(x)) \circ T(\bar{x},\bar{x})/Ex) \circ \langle x,s \rangle \right) \circ (\llbracket s \rrbracket \setminus \langle s,\sigma \rangle)$$

$$= \bigvee_{x \in X} (F(y,f(x)) \circ T(\bar{x},\bar{x})/Ex) \circ \bigvee_{s \in S(X,E)} \left[ \llbracket E^{x},s \rrbracket \right] \circ (\llbracket s \rrbracket \setminus \langle s,\sigma \rangle)$$

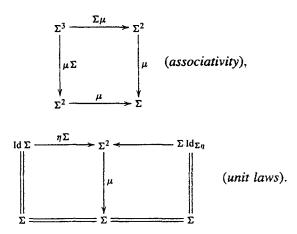
$$= \bigvee_{x \in X} (F(y,f(x)) \circ T(\bar{x},\bar{x})/Ex) \circ \left\langle E^{x},\sigma \right\rangle \qquad \text{(by 4.1(ii), 1.3(xii), 4.2(ii))}$$

$$= \left\langle y, \Big( \Sigma \big( f \big) \Big) (\sigma(E^{(1)}) \Big), \qquad (\text{analogously)}.$$

Finally, we arrive at

THEOREM 4.10. The triple  $(\Sigma, \eta, \mu)$  consisting of the functor  $\Sigma$ : Q-CAT  $\to$  Q-CAT and the two natural transforms  $\eta$ : Id  $\to \Sigma$ ,  $\mu$ :  $\Sigma^2 \to \Sigma$  (described in 4.6, 4.9) is a monad on the category Q-CAT, i.e., the

following diagrams always commute:



*Proof.* A verification of the unit laws for  $\eta$  is easy. Let us prove the associativity of  $\mu$ . If  $\tau$  is in  $S(\Sigma(\Sigma(X, E)))$  and x is in X, then we have

$$\left\langle x, \mu_{(X,E)} \Big( \Big( \Sigma \Big( \mu_{(X,E)} \Big) \Big) (\tau) \Big) \right\rangle = \left\langle E^x, \Big( \Sigma \Big( \mu_{(X,E)} \Big) \Big) (\tau) \right\rangle$$
 (by definition of  $\mu$ ) 
$$= \bigvee_{\sigma \in S \Big( \Sigma(X,E) \Big)} \left[ E^x, \mu_{(X,E)} (\sigma) \right]_{S(X,E)} \circ (\llbracket \sigma \rrbracket \setminus \langle \sigma, \tau \rangle)$$
 (by 4.6) 
$$= \bigvee_{\sigma \in S \Big( \Sigma(X,E) \Big)} \left\langle x, \mu_{(X,E)} \right\rangle \circ (\llbracket \sigma \rrbracket \setminus \langle \sigma, \tau \rangle)$$
 (by 4.5(ii)) 
$$= \bigvee_{\sigma \in S \Big( \Sigma(X,E) \Big)} \left\{ E^x, \sigma \right\rangle \circ (\llbracket \sigma \rrbracket \setminus \langle \sigma, \tau \rangle)$$
 
$$= \bigvee_{\sigma \in S \Big( \Sigma(X,E) \Big)} \left[ \left( \Big[ , E^x \Big], \Big[ E^x, \Big] \right), \sigma \Big]_{S \Big( \Sigma(X,E) \Big)} \circ (\llbracket \sigma \rrbracket \setminus \langle \sigma, \tau \rangle)$$
 (by 4.5(ii)) 
$$= \left\langle (\Big[ , E^x \Big], \Big[ E^x, \Big] \right), \tau \right\rangle$$
 (by 4.1(ii), 1.3(xii), 4.2(ii)) 
$$= \left\langle E^x, \mu_{\Sigma(X,E)} (\tau) \right\rangle$$
 
$$= \left\langle x, \mu_{(X,E)} \Big( \mu_{\Sigma(X,E)} (\tau) \right\rangle$$
 (by definition of  $\mu$ ), 
$$\left\langle \mu_{(X,E)} \Big( \Big( \Sigma \Big( \mu_{(X,E)} \Big) \Big) (\tau) \right\rangle, x \right\rangle = \left\langle \mu_{(X,E)} \Big( \mu_{\Sigma(X,E)} (\tau) \right), x \right\rangle$$
 (analogously).

### 5. SHEAVES OVER QUANTALOIDS

Having defined the monad  $(\Sigma, \eta, \mu)$  on the category Q-CAT, it is now time to extend Höhle's concept of sheaves.

Definition 5.1. Let (X, E) be a separated Q-category and u be an object of Q. (i) A subset B of X is called *compatible* when  $\{(Ex \circ T(\bar{x}, u), T(u, \bar{x}) \circ Ex) \mid x \in B\}$  constitute a singleton of the Q-category  $(B, E \mid_{B \times B})$ , where  $E \mid_{B \times B}$  stands for the restriction of E(, ) to  $B \times B$ . (ii) An element  $x_B$  of X is said to be a join of a compatible subset B of (X, E) when  $\tilde{x}_B = u$ ,

$$E(x,x_B) = \bigvee_{x' \in B} E(x,x') \circ T(\tilde{x'},u) \quad \text{and} \quad E(x_B,x) = \bigvee_{x' \in B} T(u,\tilde{x'}) \circ E(x',x).$$

PROPOSITION 5.2. Let (X, E) be a separated Q-category, u an object of Q, B a subset of X, and s a pair of maps  $\langle s, s \rangle$ ,  $\langle s, s \rangle$ :  $B \to \text{Mor}(Q)$  assigning to every  $x \in B$  morphisms  $\langle x, s \rangle$ :  $u \to \tilde{x}$  and  $\langle s, x \rangle$ :  $\tilde{x} \to u$ , respectively. Then the following assertions are equivalent:

- (i) s is a singleton of  $(B, E|_{B\times B})$ ;
- (ii) The pair  $\hat{s}$  of maps  $\langle , \hat{s} \rangle$  and  $\langle \hat{s}, \rangle$  from X to Mor(Q) defined by

$$\langle x, \hat{s} \rangle = \bigvee_{x' \in B} E(x, x') \circ (Ex' \setminus \langle x', s \rangle) \qquad \left( = \bigvee_{x' \in B} (E(x, x') \circ T(\tilde{x'}, \tilde{x'}) / Ex') \circ \langle x', s \rangle \right)$$

and

$$\langle \hat{s}, x \rangle = \bigvee_{x' \in B} (\langle s, x' \rangle \circ T(\tilde{x'}, \tilde{x'}) / Ex') \circ E(x, x') \qquad \left( = \bigvee_{x' \in B} \langle s, x' \rangle \circ (Ex' \setminus E(x', x)) \right)$$

is a singleton of (X, E) such that  $\hat{s}(x) = s(x)$  for any x in B and

$$[\![\hat{s}]\!] = \bigvee_{x \in B} T(u, \tilde{x}) \circ \langle x, s \rangle.$$

Moreover, for two such extensions  $\hat{s}$  and  $\hat{s'}$  of  $(B, E|_{B\times B})$ -singletons s and s', respectively, the following equality:

$$[\hat{s}, \hat{s'}]_{S(X,E)} = [s, s']_{S(B,E|_{B\times B})}$$

always holds.

*Proof.* The implication (ii)  $\Rightarrow$  (i) is obvious. In order to verify (i)  $\Rightarrow$  (ii) we first observe that if  $x \in B$ , then

$$\langle x, s \rangle = Ex \circ (Ex \setminus \langle x, s \rangle) \leqslant \bigvee_{x' \in B} E(x, x') \circ (Ex' \setminus \langle x', s \rangle) \qquad (= \langle x, \hat{s} \rangle)$$

$$\leqslant \langle x, s \rangle \qquad \text{(by 4.1(ii))}$$

$$\langle s, x \rangle = (\langle s, x \rangle \circ T(\tilde{x}, \tilde{x})/Ex) \circ Ex \leqslant \bigvee_{x' \in B} (\langle s, x' \rangle \circ T(\tilde{x'}, \tilde{x'})/Ex') \circ E(x', x) \qquad (= \langle \hat{s}, x \rangle)$$

$$\leqslant \langle s, x \rangle,$$

i.e.,  $\langle x, \hat{s} \rangle = \langle x, s \rangle$  and  $\langle \hat{s}, x \rangle = \langle s, x \rangle$ . Let us now prove that  $\hat{s}$  is a singleton of (X, E). But this is an immediate consequence of 4.6. Moreover, by 4.6 the map  $\hat{s} \mapsto \hat{s}$  is a Q-CAT-morphism from  $(B, E|_{B \times B})$  to (X, E), i.e., for every s in  $S(B, E|_{B \times B})$ ,  $\hat{s}$  is in S(X, E),  $[[\hat{s}]] = [[s]]$ , and  $[[s, s']] \leq [[\hat{s}, \hat{s}']]$ . Thus, it only remains to verify the opposite inequality:  $[[\hat{s}, \hat{s}']] \leq [[s, s']]$ , but this is true, since

COROLLARY 5.3 (cf. [4, Lemma 4.2]). Let (X, E) be a separated Q-category, u an object of Q, and B a subset of X. Then the following assertions are equivalent:

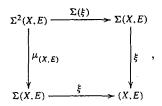
- (i) B is compatible;
- (ii) The pair  $\hat{s}$  of maps  $\langle , \hat{s} \rangle$ ,  $\langle \hat{s}, \rangle$  from X to Mor(Q) defined by

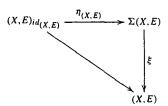
$$\langle x, \hat{s} \rangle = \bigvee_{x' \in B} E(x, x') \circ T(\tilde{x'}, u) \quad and \quad \langle \hat{s}, x \rangle = \bigvee_{x' \in B} T(u, \tilde{x'}) \circ E(x', x)$$

is a singleton of (X, E) with  $[\hat{s}] = \bigvee_{x \in B} T(u, \tilde{x}) \circ Ex \circ T(\tilde{x}, u)$ .

Before going further, we recall the definition of (Eilenberg-Moore) algebras for a monad.

Definition 5.4. Let  $M = (\Sigma, \eta, \mu)$  be the monad specified in 4.10. An *M-algebra* is a pair  $((X, E), \xi)$  of a separated Q-category (X, E) and a Q-CAT-morphism  $\xi: \Sigma(X, E) \to (X, E)$  which makes both diagrams





commute.

PROPOSITION 5.5 (cf. [4, Proposition 4.3]). Let  $((X, E), \xi)$  be an M-algebra, u an object of Q and B a compatible subset of (X, E). Then B has a (unique) join.

*Proof.* Consider the maps  $\langle , s \rangle$ ,  $\langle s, \rangle$ :  $X \to \text{Mor}(Q)$  defined by

$$\langle x, s \rangle = \bigvee_{x' \in B} E(x, x') \circ T(\tilde{x'}, u) \quad \text{and} \quad \langle s, x \rangle = \bigvee_{x' \in B} T(u, \tilde{x'}) \circ E(x', x).$$

Since  $s = (\langle , s \rangle, \langle s, \rangle)$  is a singleton of (X, E) (by 5.3) and  $\xi$  is a Q-CAT-morphism, we obtain  $\langle x, s \rangle = [E^x, s] \in E(x, \xi(s))$  and  $\langle s, x \rangle = [s, E^x] \in E(\xi(s), x)$  for all x in X. To establish the opposite inequalities,

we proceed as follows:

$$E(x,\xi(s)) = (E(x,\xi(s)) \circ T(u,u)/E\xi(s)) \circ E\xi(s) \qquad \text{(by 3.2(iii))}$$

$$= (E(x,\xi(s)) \circ T(u,u)/E\xi(s)) \circ \llbracket s \rrbracket \qquad \text{(by 3.6(i))}$$

$$= (E(x,\xi(s)) \circ T(u,u)/E\xi(s)) \circ \bigvee_{x' \in B} T(u,\tilde{x'}) \circ Ex' \circ T(\tilde{x'},u) \qquad \text{(by 5.3)}$$

$$= \bigvee_{x' \in B} (E(x,\xi(s)) \circ T(u,u)/E\xi(s)) \circ \langle s,x' \rangle \circ T(\tilde{x'},u)$$

$$(\operatorname{since}\langle s,x' \rangle = T(u,\tilde{x'}) \circ Ex', \text{ when } x' \text{ is in } B)$$

$$= \bigvee_{x' \in B} (E(x,\xi(s)) \circ T(u,u)/E\xi(s)) \circ \llbracket s,E^{x'} \rrbracket \circ T(\tilde{x'},u) \qquad \text{(by 4.5(ii))}$$

$$\leqslant \bigvee_{x' \in B} (E(x,\xi(s)) \circ T(u,u)/E\xi(s)) \circ E(\xi(s),x') \circ T(\tilde{x'},u) \qquad \text{(by 3.6(ii))}$$

$$\leqslant \bigvee_{x' \in B} E(x,x') \circ T(\tilde{x'},u) \qquad \text{(by 3.2(i), 3.1(ii))}$$

$$= \langle s,x \rangle.$$

 $E(\xi(s), x) = \langle s, x \rangle$  (analogously). Hence  $\xi(s)$  is the join of B. Since we consider only separated Q-categories, this join is unique.

PROPOSITION 5.6 (cf. [4, Proposition 4.4] and [3, Proposition 4.6]). Every M-algebra  $((X, E), \xi)$  induces a "preasheaf structure"  $(X, \neg, E)$  on X over Q by

$$x \, \exists a = \xi \left( E^{(x,a)} \right), \tag{i}$$

with the following properties: for all x in X, u, v in Obj(Q) with the relations:  $T(\tilde{x}, u) \circ T(u, \tilde{x}) = T(\tilde{x}, \tilde{x}) = T(\tilde{x}, v) \circ T(v, \tilde{x})$ ,  $T(u, v) \circ T(v, u) = T(u, u)$ , and for all  $a: u \to \tilde{x}$ ,  $b: v \to u$ ,

$$x \exists Ex = x, \qquad E(x \exists a) = T(u, \tilde{x}) \circ ((Ex \circ T(\tilde{x}, u)) \land a) \qquad \Big( = \Big[ E^{(x, a)} \Big] \Big),$$

$$(x \exists a) \exists b = x \exists ((a \circ T(u, v)) \land (T(\tilde{x}, u) \circ b)), \qquad (ii)$$

where  $E^{(x,a)}$  is determined by 4.3(ii), and E by 3.1(iii). In particular,  $(X, \neg, E)$  is compatible with the underlying Q-category structure by

$$E(x \mid a, x \mid (T(\tilde{x}, u) \circ b)) = T(u, \tilde{x}) \circ ((Ex \circ T(\tilde{x}, v)) \wedge (a \circ T(u, v))) \wedge (T(\tilde{x}, u) \circ b)). \tag{iii}$$

*Proof.* The strictness 3.6(i) of  $\xi$  immediately implies the first two equalities of (ii). In order to prove the last equality of (ii), for given  $x \in X$ , u,  $v \in \mathrm{Obj}(Q)$  (with  $T(\tilde{x}, u) \circ T(u, \tilde{x}) = T(\tilde{x}, \tilde{x}) = T(\tilde{x}, v) \circ T(v, \tilde{x})$  and  $T(u, v) \circ T(v, u) = T(u, u)$ ),  $a: u \to \tilde{x}$  and  $b: v \to u$ , let us consider a singleton  $\sigma = (\langle \ , \sigma \rangle, \langle \sigma, \ \rangle)$  of  $\Sigma(X, E)$  defined by  $\langle s, \sigma \rangle = [\![s, E^{(x,a)}]\!] \circ ([\![E^{(x,a)}]\!] \setminus b)$  and  $\langle \sigma, s \rangle = T(v, u) \circ (b \circ T(v, u) / [\![E^{(x,a)}]\!]) \circ [\![E^{(x,a)}, s]\!]$  (cf.

4.3(ii)). Put  $s = E^{x'}$ , then

$$\begin{split} \left\langle E^{x'}, \sigma \right\rangle &= \left[ E^{x'}, E^{(x,a)} \right] \circ \left( \left[ E^{(x,a)} \right] \setminus b \right) = \left\langle x', E^{(x,a)} \right\rangle \circ \left( \left[ E^{(x,a)} \right] \setminus b \right) & \text{(by 4.5(ii))} \\ &= \left( E(x', x) \circ T(\tilde{x}, \tilde{x}) / Ex) \circ T(\tilde{x}, u) \circ \left[ E^{(x,a)} \right] \circ \left( \left[ E^{(x,a)} \right] \setminus b \right) & \text{(by 4.3(ii))} \\ &= \left( E(x', x) \circ T(\tilde{x}, \tilde{x}) / Ex) \circ T(\tilde{x}, u) \circ \left( \left[ E^{(x,a)} \right] \circ T(u, v) \wedge b \right) & \text{(by 1.3(xii))} \\ &= \left( E(x', x) \circ T(\tilde{x}, \tilde{x}) / Ex) \circ T(\tilde{x}, u) \circ \left( \left( T(u, \tilde{x}) \circ (Ex \circ T(\tilde{x}, u) \wedge a) \circ T(u, v) \right) \wedge b \right) \\ &= \left( E(x', x) \circ T(\tilde{x}, \tilde{x}) / Ex) \circ (Ex \circ T(\tilde{x}, v) \wedge (a \circ T(u, v) \wedge T(\tilde{x}, u) \circ b) \right) & \text{(by 1.2(ii))} \\ &= \left\langle x', E^{(x,a\circ T(u,v) \wedge T(x,u) \circ b)} \right\rangle, \end{split}$$

 $\langle \sigma, E^{x'} \rangle = \langle E^{(x,a\circ T(u,v)\wedge T(\tilde{x},u)\circ b)}, x' \rangle$  (similarly). Hence  $\mu_{(X,E)}(\sigma) (= \sigma(E^{()}))$  (by 4.9)  $= E^{(x,a\circ T(u,v)\wedge T(\tilde{x},u)\circ b)}$ . We further have

$$\langle x', (\Sigma(\xi))(\sigma) \rangle = \bigvee_{s \in S(X,E)} (E(x', \xi(s)) \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \langle s, \sigma \rangle \qquad \text{(by 4.6)}$$

$$= \bigvee_{s \in S(X,E)} (E(x', \xi(s)) \circ T(\tilde{s}, \tilde{s}) / \llbracket s \rrbracket) \circ \llbracket s, E^{(x,a)} \rrbracket \circ \left( \llbracket E^{(x,a)} \rrbracket \setminus b \right)$$

$$= \left( E(x', \xi(E^{(x,a)}) \right) \circ \left( \llbracket E^{(x,a)} \rrbracket \setminus b \right) \qquad \text{(by 3.6, 3.1(ii))}$$

$$= E(x', x \exists a) \circ (E(x \exists a) \setminus b) = \langle x', E^{(x \exists a,b)} \rangle,$$

 $\langle (\Sigma(\xi))(\sigma), x' \rangle = \langle E^{(x \neg a,b)}, x' \rangle$  (analogously). Hence

$$x \, \exists (a \circ T(u, v) \land T(\tilde{x}, u) \circ b) = \xi \left( E^{(x, a \circ T(u, v) \land T(\tilde{x}, u) \circ b)} \right) = \xi \left( \mu_{(X, E)}(\sigma) \right)$$

$$= \xi \left( \left( \Sigma(\xi) \right) (\sigma) \right) \quad \text{(by 5.4)}$$

$$= \xi \left( E^{(x \, \exists a, b)} \right) = (x \, \exists a) \, \exists b.$$

To prove (iii) note that

$$\begin{split} \left[E^{(x,a)}, E^{(x,T(\tilde{x},u)\circ b)}\right] \leqslant \left[E^{(x,a)}\right] \circ \bigwedge_{x'\in X} (E(x',x)\circ (Ex\setminus a)\setminus E(x',x)\circ (Ex\setminus T(\tilde{x},u)\circ b)) & \text{(by 4.4, 4.3(ii))} \\ \leqslant \left[E^{(x,a)}\right] \circ (Ex\circ (Ex\setminus a)\setminus Ex\circ (Ex\setminus T(\tilde{x},u)\circ b)) & \\ = T(u,\tilde{x})\circ ((Ex\circ T(\tilde{x},u)\wedge a) & \\ \circ T(u,v)\wedge (Ex\circ T(\tilde{x},v)\wedge T(\tilde{x},u)\circ b)) & \text{(by 1.3(xii) twice)} \\ = T(u,\tilde{x})\circ (Ex\circ T(\tilde{x},v)\wedge a\circ T(u,v)\wedge T(\tilde{x},u)\circ b) & \text{(by 1.2(ii))} \end{split}$$

and

$$\begin{split} \left[E^{(x,a)}, E^{(x,T(\bar{x},u)\circ b)}\right] \geqslant & \left[\left[E^{(x,a)}\right] \circ \bigwedge_{x' \in X} \left(\left(E(x',x) \circ T(\bar{x},\bar{x})/Ex\right) \circ T(\bar{x},u)\right) \\ & \circ \left[E^{(x,a)}\right] \setminus \left(E(x',x) \circ T(\bar{x},\bar{x})/Ex\right) \circ T(\bar{x},u) \circ T(u,v) \circ \left[E^{(x,T(\bar{x},u)\circ b)}\right]\right) \right] \\ & \wedge \left[\left(\bigwedge_{x' \in X} \left[E^{(x,a)}\right] \circ T(u,v) \circ T(v,\bar{x}) \circ \left(Ex \setminus E(x,x')\right)/\left[E^{(x,T(\bar{x},u)\circ b)}\right] \circ T(v,\bar{x})\right] \\ & \circ \left(Ex \setminus E(x,x')\right)\right) \circ \left[E^{(x,T(\bar{x},u)\circ b)}\right] \end{split}$$
 (by 4.4, 4.3(ii))

$$\geqslant \left( \left[ E^{(x,a)} \right] \circ \left( \left[ E^{(x,a)} \right] \setminus T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\land \left( \left( \left[ E^{(x,a)} \right] \circ T(u,v) / \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\land \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(v,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,v) \wedge T(u,\tilde{x}) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \right)$$

$$\Rightarrow \left( \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \wedge$$

i.e.,  $\left[E^{(x,a)}, E^{(x,T(\tilde{x},u)\circ b)}\right] = T(u,\tilde{x})\circ (Ex\circ T(\tilde{x},v)\wedge a\circ T(u,v)\wedge T(\tilde{x},u)\circ b).$  Then

$$T(u,\tilde{x}) \circ (Ex \circ T(\tilde{x},v) \wedge a \circ T(u,v) \wedge T(\tilde{x},u) \circ b)$$

$$= \left[ E^{(x,a)}, E^{(x,T(\tilde{x},u)\circ b)} \right] \leqslant \left[ E^{(x,a)} \right] \circ T(u,v) \wedge T(u,v) \circ \left[ E^{(x,T(\tilde{x},u)\circ b)} \right] \qquad \text{(by 3.1(i))}$$

$$= T(u,\tilde{x}) \circ (Ex \circ T(\tilde{x},u) \wedge a) \circ T(u,v) \wedge T(u,v) \circ T(v,\tilde{x}) \circ (Ex \circ T(\tilde{x},v) \wedge T(\tilde{x},u) \circ b)$$

$$= T(u,\tilde{x}) \circ Ex \circ T(\tilde{x},v) \wedge T(u,\tilde{x}) \circ a \circ T(u,v) \wedge T(u,\tilde{x}) \circ Ex \circ T(\tilde{x},v) \wedge T(u,u) \circ b \qquad \text{(by 1.2(ii))}$$

$$= T(u,\tilde{x}) \circ (Ex \circ T(\tilde{x},v) \wedge a \circ T(u,v) \wedge T(\tilde{x},u) \circ b) \qquad \text{(by 1.2(ii))},$$

which proves (iii).

Note that in the case where Q is a commutative quantale U. Höhle [4] calls M-algebras also sheaves over Q.

#### REFERENCES

- 1. U. Berni-Canani, F. Borceux, and R. Succi Cruciani, A theory of quantal sets, J. Pure Appl. Algebra, 62, 123-136 (1989).
- 2. F. Borceux and G. van den Bossche, Quantales and their sheaves, Order, 3, 61-87 (1986).
- 3. R. P. Gylys, Quantal sets and sheaves over quantales, Lith. Math. J., 34, 8-29 (1994).
- 4. U. Höhle, M-valued sets and sheaves over integral commutative cl-monoids, in: Applications of Category Theory to Fuzzy Subsets, S. E. Rodabaug et al. (eds), Kluwer Academic Publishers, The Netherlands (1992), pp. 33-72.
- 5. C. J. Mulvey, &, Rend. Circ. Mat. Palermo, 12, 99-104 (1986).
- 6. A. M. Pitts, Applications of sup-lattice enriched category theory to sheaf theory, *Proc. London Math. Soc.*, **57**, 433–480 (1988).
- 7. K. I. Rosenthal, Quantales and Their Applications, Pitman Research Notes in Math. 234, Longman, Scientific and Technical, Harlow (1990).
- 8. K. I. Rosenthal, Free quantaloids, J. Pure Appl. Algebra, 72, 67-82 (1991).
- 9. K. I. Rosenthal, Girard quantaloids, Math. Struct. Comp. Science, 2, 93-108 (1992).
- 10. K. I. Rosenthal, Quantaloidal nuclei, the syntactic congruence and tree automata, J. Pure Appl. Algebra, 77, 189-205 (1992).
- 11. K. I. Rosenthal, A categorical look at tree automata and context-free languages, to appear in Math. Struct. Comp. Science.
- D. Yetter, Quantales and (non-commutative) linear logic, J. Symbolic Logic, 55, 41–64 (1990).