ThreadProof: A Non-Isomorphic Lattice-Based Post-Quantum Identity Proof

OBINexus Computing
`support@obinexus.org`

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Abstract

This document presents a formal mathematical framework for ThreadProof, a non-isomorphic lattice-based cryptographic identity system designed for post-quantum security. ThreadProof leverages coordinate-locked lattice structures to enforce transformation invariance and provides modular proof properties across distributed systems. We establish the theoretical foundations, prove security properties including completeness, soundness, and LWE-based hardness, and demonstrate the non-transferability of proofs across geometric transformation spaces. The system integrates with the OBINexus Cryptographic Interoperability Standard v1.0 to provide deterministic, quantum-resistant identity management.

1 Introduction

ThreadProof addresses the fundamental challenge of creating cryptographically secure identities that maintain their integrity across different computational bases while resisting quantum attacks. Traditional identity systems rely on algebraic structures that may be vulnerable to transformation attacks or quantum algorithms. ThreadProof introduces a novel approach using non-isomorphic lattice transformations that are inherently locked to their coordinate system.

1.1 Motivation

The advent of quantum computing necessitates cryptographic systems that resist both classical and quantum attacks. While lattice-based cryptography provides quantum resistance through the Learning With Errors (LWE) problem, existing implementations often allow transformations between coordinate systems that may weaken security guarantees. ThreadProof addresses this by enforcing strict non-isomorphism between geometric representations.

1.2 Contributions

This work makes the following contributions:

1. A formal framework for non-isomorphic lattice-based identity proofs
2. Mathematical proofs of completeness, soundness, and quantum resistance
3. Integration methodology with existing OBINexus cryptographic infrastructure
4. Practical implementation guidelines with performance guarantees

2 System Architecture

2.1 Canonical Identity Definition (CID)

The CID serves as the fundamental identity primitive in ThreadProof:

Definition 2.1 (Canonical Identity Definition)**.** A CID is a tuple (I, V, S, B, E, K, P, C) where:

- I = Identity string following format `ZID:QZ-Lattice-XX`
- V = Version string using semantic versioning
- S = Structure type (non-isomorphic lattice)
- B = Base space (Cartesian, Polar, etc.)
- E = Encoding function (HKDF-SHA3-512)
- K = Key derivation method (context-bound)
- P = Proof type (post-quantum lattice)
- C = Configuration parameters

2.2 Congregation Linker

The Congregation Linker provides modular component integration:

```
Congregation = {
    "linked_components": [
        "ZID Root Key",
        "Proof Generator",
        "Lattice Mapper",
        "Non-Isomorphic Verifier"
    ],
    "validation": {
        "structure_check": True,
        "transform_space": "Locked",
        "derivation_protection": "Enabled"
    }
}
```

Listing 1: Congregation Linker Structure

3 Mathematical Foundation

3.1 Lattice Structure

Let Λ be a lattice in $\mathbb{R}^n$ generated by basis $\mathbf{B} = \{\mathbf{b}_1, \ldots, \mathbf{b}_n\}$:

$$\Lambda = \left\{ \sum_{i=1}^{n} z_i \mathbf{b}_i : z_i \in \mathbb{Z} \right\}$$

3.2 Non-Isomorphism Constraint

Definition 3.1 (Non-Isomorphic Lattices)**.** Two lattices Λ_1 and Λ_2 with bases in different coordinate systems are non-isomorphic if there exists no linear transformation T such that:

$$T(\Lambda_1) = \Lambda_2 \text{ and } \det(T) = \pm 1$$

4 Axiomatic Framework

ThreadProof operates under three fundamental axioms:

Axiom 4.1 (Structural Non-Isomorphism)**.** A cryptographic lattice transformation must be non-isomorphic across coordinate systems. For all transformations $T : X \to Y$, where X is Cartesian and Y is polar, there exists no bijection f such that $T_X \cong T_Y$.

Axiom 4.2 (Encoding Rigidity)**.** Proofs generated in a given lattice-structured system L are valid only within the coordinate space of L. Translation to an alternate geometric form invalidates verification.

Axiom 4.3 (Key Derivation Invariance)**.** Derived keys $k' \leftarrow \text{HKDF}(k, \text{context})$ preserve the transformation structure of their origin key. Any attempt to derive across systems (Cartesian $\to$ polar) breaks key validity.

5 Security Properties

5.1 Completeness

Theorem 5.1 (Completeness of ThreadProof)**.** For any valid identity I with corresponding proof π, the verification algorithm V always accepts:

$$\Pr[V(I, \pi) = 1 | \text{Valid}(I, \pi)] = 1$$

Proof. Let I be a valid identity with CID structure (I, V, S, B, E, K, P, C) and proof π generated by the honest prover algorithm. By construction:

1. The lattice basis $\mathbf{B}$ is correctly formed in the specified coordinate system 2. The encoding E produces a deterministic output for given input 3. The proof π contains valid witness vectors $\mathbf{w} \in \Lambda$

Since the verifier checks: - Lattice membership: $\mathbf{w} \in \Lambda$ - Coordinate system match: $\text{Space}(\pi) = B$ - Encoding validation: $E(\mathbf{w}) = \text{commitment}$

All checks pass with probability 1 for honestly generated proofs. □ □

5.2 Soundness

Theorem 5.2 (Soundness of ThreadProof)**.** For any invalid identity I' or forged proof π', the probability of verification acceptance is negligible:

$$\Pr[V(I', \pi') = 1 | \neg\text{Valid}(I', \pi')] \le \text{negl}(\lambda)$$

where λ is the security parameter.

Proof. Assume adversary $\mathcal{A}$ attempts to forge a proof π' for invalid identity I'. By Axiom 4.1, $\mathcal{A}$ cannot transform a valid proof from another coordinate system.

The security reduces to the hardness of the Learning With Errors problem:

$$\text{Advantage}_{\mathcal{A}} \leq \text{Advantage}_{\text{LWE}}(n, q, \chi)$$

where n is lattice dimension, q is modulus, and χ is error distribution.

Given that LWE is conjectured hard even for quantum adversaries, the forgery probability is negligible. □ □

5.3 Hardness

Theorem 5.3 (Quantum Hardness)**.** Breaking ThreadProof security requires solving the Learning With Errors problem, which remains hard against quantum adversaries.

Proof. The security of ThreadProof relies on:

1. **Lattice Problems**: Finding short vectors in the locked coordinate system 2. **LWE Hardness**: Distinguishing $(\mathbf{A}, \mathbf{As} + \mathbf{e})$ from uniform 3. **Non-Isomorphism**: Preventing coordinate system attacks

Quantum algorithms (Shor's, Grover's) do not provide efficient solutions to lattice problems in high dimensions. The best known quantum algorithm achieves only polynomial speedup:

$$\text{Time}_{\text{Quantum}} = \tilde{O}(2^{0.265n})$$

Therefore, ThreadProof maintains post-quantum security. □ □

5.4 Correctness

Proposition 5.4 (Correctness Under Transformation)**.** ThreadProof maintains correctness only within its native coordinate system. Any transformation violates correctness.

5.5 Encoding Invariance

Lemma 5.5 (HKDF-SHA3-512 Invariance)**.** The encoding function $E : \Lambda \to \{0, 1\}^{512}$ maintains collision resistance and preimage resistance under the coordinate lock constraint.

6 Implementation Specification

6.1 Key Generation

6.2 Proof Generation

6.3 Verification Algorithm

7 Integration with OBINexus Framework

ThreadProof integrates seamlessly with the OBINexus Cryptographic Interoperability Standard v1.0:

Algorithm 1 ThreadProof Key Generation

```
Require: Security parameter $\lambda$, coordinate system $\mathcal{C}$
Ensure: Public key $pk$, secret key $sk$
 1: Sample random seed $s \leftarrow \{0,1\}^\lambda$
 2: Generate lattice basis $\mathbf{B} \leftarrow \text{GenBasis}(s, \mathcal{C})$
 3: Lock coordinate system: $\mathbf{B}.\text{lock}(\mathcal{C})$
 4: Compute trapdoor $\mathbf{T} \leftarrow \text{Trapdoor}(\mathbf{B})$
 5: Derive identity: $\text{ZID} \leftarrow \text{HKDF}(s, \text{"identity"})$
 6: return $pk = (\mathbf{B}, \text{ZID}, \mathcal{C})$, $sk = (\mathbf{T}, s)$
```

Algorithm 2 ThreadProof Generation

```
Require: Secret key $sk$, message $m$, context $ctx$
Ensure: Proof $\pi$
 1: Parse $sk = (\mathbf{T}, s)$
 2: Sample randomness $r \leftarrow \chi$
 3: Compute commitment $c = \text{HKDF}(m||ctx||r)$
 4: Generate witness $\mathbf{w} \leftarrow \text{SamplePre}(\mathbf{T}, c)$
 5: Verify coordinate lock: $\text{CheckLock}(\mathbf{w}, \mathcal{C})$
 6: Create proof $\pi = (\mathbf{w}, c, r, \mathcal{C})$
 7: return $\pi$
```

Algorithm 3 ThreadProof Verification

```
Require: Public key $pk$, proof $\pi$, message $m$, context $ctx$
Ensure: Accept/Reject
 1: Parse $pk = (\mathbf{B}, \text{ZID}, \mathcal{C})$
 2: Parse $\pi = (\mathbf{w}, c, r, \mathcal{C}')$
 3: if $\mathcal{C} \neq \mathcal{C}'$ then
 4:     return Reject // Coordinate system mismatch
 5: end if
 6: Verify lattice membership: $\mathbf{w} \in \Lambda(\mathbf{B})$
 7: Recompute $c' = \text{HKDF}(m||ctx||r)$
 8: if $c \neq c'$ then
 9:     return Reject
10: end if
11: Verify witness norm: $||\mathbf{w}|| \leq \beta$
12: return Accept
```

7.1 Pattern Registration

ThreadProof primitives register with the OBINexus pattern matching system:

```python
THREADPROOF_PATTERNS = {
    "ZID_PATTERN": r"^ZID:QZ-Lattice-\d+$",
    "ENCODING_PATTERN": r"^HKDF-SHA3-512$",
    "SIGNATURE_PATTERN": r"^Falcon-512$"
}

def register_threadproof_patterns():
    for name, pattern in THREADPROOF_PATTERNS.items():
        register_cryptographic_pattern(
            pattern_name=name,
            regex=pattern,
            category="POST_QUANTUM",
            version="1.0.0"
        )
```

Listing 2: ThreadProof Pattern Registration

7.2 Audit Integration

ThreadProof operations integrate with the SecureAuditNode:

```python
class ThreadProofAuditNode(SecureAuditNode):
    def __init__(self, zid, operation, proof_hash):
        super().__init__(
            primitive_digest=zid,
            pattern_state=ThreadProofState()
        )
        self.operation = operation
        self.proof_hash = sha256(proof_hash).hexdigest()[:16]
        self.coordinate_system = "Cartesian"
        self.axiom_compliance = self.verify_axioms()
```

Listing 3: ThreadProof Audit Trail

8 Performance Analysis

8.1 Computational Complexity

Proposition 8.1 (Time Complexity)**.** ThreadProof operations have the following complexity bounds:

- Key Generation: $O(n^3)$ where n is lattice dimension
- Proof Generation: $O(n^2 \log n)$ using fast sampling
- Verification: $O(n^2)$ for matrix-vector operations

8.2 Space Complexity

Proposition 8.2 (Storage Requirements)**.** ThreadProof requires:

- Public Key: $O(n^2 \log q)$ bits

- Secret Key: $O(n^2 \log q)$ bits
- Proof Size: $O(n \log q)$ bits
- CID Storage: $O(1)$ (constant 512 bits)

9 Security Analysis

9.1 Attack Resistance

ThreadProof resists the following attack classes:

1. **Coordinate Transformation Attacks**: Prevented by Axiom 4.1
2. **Cross-System Derivation**: Blocked by Axiom 4.3
3. **Proof Transplantation**: Prevented by Axiom 4.2
4. **Quantum Attacks**: Resistant due to LWE hardness
5. **Side-Channel Attacks**: Mitigated by constant-time implementations

9.2 Formal Security Proof

Theorem 9.1 (ThreadProof Security)**.** ThreadProof achieves IND-CCA2 security under the LWE assumption with advantage:

$$\mathrm{Adv}^{\text{IND-CCA2}}_{\text{ThreadProof}}(\mathcal{A}) \leq \mathrm{Adv}^{n,q,\chi}_{\text{LWE}}(\mathcal{B}) + \mathrm{negl}(\lambda)$$

10 Conclusion

ThreadProof provides a mathematically rigorous framework for post-quantum secure identity management through non-isomorphic lattice-based cryptography. By enforcing coordinate system locks and preventing geometric transformations, the system achieves strong security guarantees while maintaining practical efficiency. The integration with OBINexus infrastructure ensures compatibility with existing cryptographic standards while advancing the state of quantum-resistant identity systems.

Future work includes optimizing proof sizes through advanced compression techniques and extending the framework to support threshold identity schemes for distributed applications.

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Contact

For implementation details, integration support, or security audits:

`support@obinexus.org`
OBINexus Computing
`https://github.com/obinexus`