# Polynomial semantics of probabilistic circuits

Oliver Broadrick, UCLA

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Based on joint work with Honghua Zhang and Guy Van den Broeck

▶ Tractable probabilistic models

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- ▶ Several circuit semantics in the literature ...
- ▶ are equivalent! (for binary random variables)
- ► And, don't all extend to non-binary variables

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AI research

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▶ Deep learning and formal methods: "neuro-symbolic AI"

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#### AI research

- ▶ Deep learning and formal methods: "neuro-symbolic AI"
- ▶ Applications: images, language, audio, medicine, science, economics, etc.

$X_1$	$X_2$	Pr
0	0	.1
0	1	.2
1	0	.3
1	1	.4

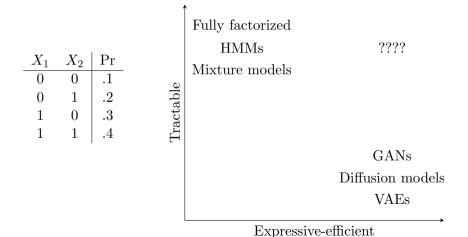
► Expressive-efficient representation

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- ► Tractable inference

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/19

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If  $X = Y \sqcup Z$ , then what is Pr[Y = y]?

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$$\Pr[\boldsymbol{Y} = \boldsymbol{y}] = \sum_{\boldsymbol{z}} \Pr[\boldsymbol{Y} = \boldsymbol{y}, \boldsymbol{Z} = \boldsymbol{z}]$$

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In general:

$$egin{array}{c|c|c|c} X_1 & X_2 & \operatorname{Pr} \\ \hline 0 & 0 & .1 \\ 0 & 1 & .2 \\ 1 & 0 & .3 \\ 1 & 1 & .4 \\ \hline \end{array}$$

For example:

$$Pr[X_1 = 1] = Pr[X_1 = 1, X_2 = 0] + Pr[X_1 = 1, X_2 = 1]$$

$$= 0.3 + 0.4$$

$$= 0.7$$

 $\Pr[oldsymbol{Y} = oldsymbol{y}] = \sum \Pr[oldsymbol{Y} = oldsymbol{y}, oldsymbol{Z} = oldsymbol{z}]$ 

If  $X = Y \sqcup Z$ , then what is Pr[Y = y]?

 $D_n[\mathbf{V} = \alpha] = \sum D_n[\mathbf{V} = \alpha, \mathbf{Z} = \alpha]$ 

In general:

$X_1$	$X_2$	Pr	$\Pr[\mathbf{Y} = \mathbf{y}] = \sum_{\mathbf{z}} \Pr[\mathbf{Y} = \mathbf{y}, \mathbf{Z} = \mathbf{z}]$
0	0	.1	~
0	1 0 1	.2	For example:
1	0	.3	Tor example.
1	1	.4	$\Pr[X_1 = 1] = \Pr[X_1 = 1, X_2 = 0] + \Pr[X_1 = 1, X_2 = 1]$
			=0.3+0.4
			=0.7

Goal: Find a model of polysize that supports marginal inference in polytime, for as large a set of probability distributions as possible.

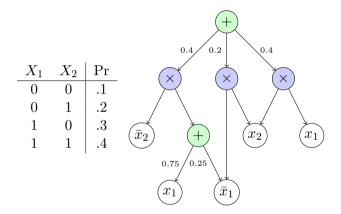
▶ Bayesian Networks (of bounded treewidth) (BNs)

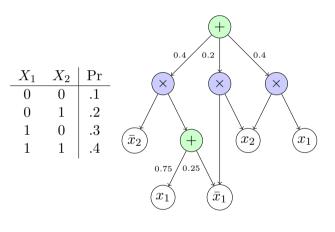
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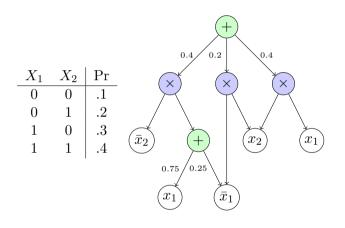
- ▶ Bayesian Networks (of bounded treewidth) (BNs)
- ▶ Determinantal Point Processes (DPPs)
- ▶ Probabilistic Sentential Decision Diagrams (PSDDs)
- ► Probabilistic Circuits!

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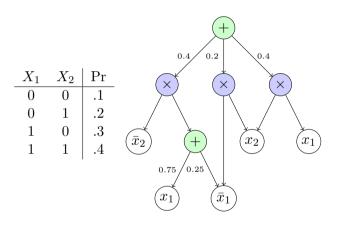


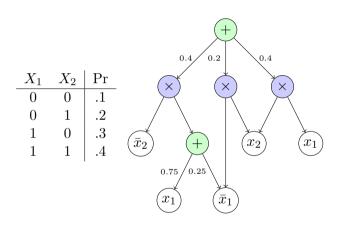


Marginal Inference:

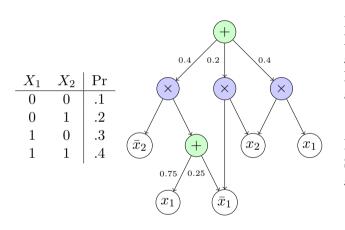


Marginal Inference: If  $X_i = b$  for  $b \in \{0, 1\}$ , set  $x_i = b$  and  $\bar{x}_i = 1 - b$ .



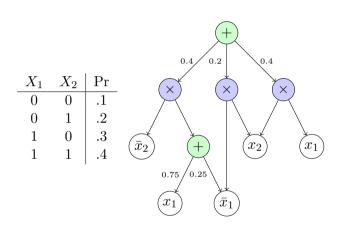


$$\begin{aligned} &\Pr[X_1 = 1]?\\ &\text{Set } x_1 = 1, \ \bar{x}_1 = 0, \ x_2 = 1,\\ &\bar{x}_2 = 1 \end{aligned}$$



$$Pr[X_1 = 1]?$$
  
Set  $x_1 = 1$ ,  $\bar{x}_1 = 0$ ,  $x_2 = 1$ ,  $\bar{x}_2 = 1$ 

$$p(x_1, x_2, \dots, x_n, \bar{x}_1, \bar{x}_2, \dots, \bar{x}_n) = .1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2$$



$$\Pr[X_1 = 1]$$
?  
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 "Network polynomial"

## **Circuit Semantics**

Polynomial	Notation	Inference	Citation
Network polynomial	$p(x_1,\ldots,x_n,\bar{x}_1,\ldots,\bar{x}_n)$	<b>✓</b>	Darwiche [2003]

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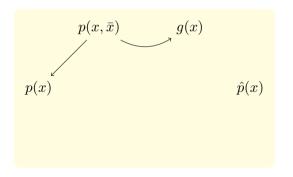
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Fourier transform	$\hat{p}(x_1,\ldots,x_n)$	<b>✓</b>	Yu et al. [2023]

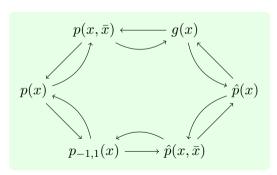
## How do they relate?

```
p(x,\bar{x})
           Network polynomial
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 p(x)
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p(x)
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## How do they relate?

 $p(x, \bar{x})$  Network polynomial p(x) Likelihood polynomial g(x) Generating function  $\hat{p}(x)$  Fourier transform





$$p(x_1, x_2) = .2x_1 + .1x_2 + .1$$

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Can we do marginal inference?

	$X_2$	
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0	1	.2
1	0	.3
1	0 1 0 1	.4

$$\begin{array}{c|cccc} X_1 & X_2 & \Pr \\ \hline 0 & 0 & .1 \\ 0 & 1 & .2 \\ 1 & 0 & .3 \\ 1 & 1 & .4 \\ \end{array}$$

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Can we do marginal inference? Relation to network polynomial?

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2	
Can we do marginal inference?	
Relation to network polynomial?	
Transformation from network to a likelihood:	

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 $p(x_1, x_2) = .2x_1 + .1x_2 + .1$ Can we do marginal inference?
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Transformation from network to a likelihood:  $p(x_1, x_2, \bar{x}_1, \bar{x}_2) = .1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2$ 

 $p(x_1, x_2, 1 - x_1, 1 - x_2)$ 

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 $p(x_1, x_2, 1 - x_1, 1 - x_2)$ 

 $= .2x_1 + .1x_2 + .1$ 

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$$= .1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2 = p(x_1, x_2, \bar{x}_1, \bar{x}_2)$$

Wait a second... 
$$\left(\prod_{i=1}^{n} (x_i + \bar{x}_i)\right) p\left(\frac{x_1}{x_1 + \bar{x}_1}, \dots, \frac{x_n}{x_n + \bar{x}_n}\right)$$

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Theorem 2 (Strassen). You can remove the divisions in polynomial time!

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Lemma 3. For a circuit computing f of degree d, we can obtain circuits computing  $H_0[f], H_1[f], \ldots, H_d[f]$  the homogeneous parts of f, i.e.  $H_i[f]$  has degree i and  $f = \sum_i H_i[f]$ .

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Idea: Move divisions to the root using a/b + c/d = (ad + bc)/bd and a/b \* c/d = ac/bd.

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Idea: Move divisions to the root using a/b + c/d = (ad + bc)/bd and a/b \* c/d = ac/bd.

Then for circuit a/b computing polynomial f = a/b of degree d, assume b(0) = 1, and we have

$$H_i[f] = H_i[a(1+(1-b)+(1-b)^2+\ldots+(1-b)^d].$$

### Progress update

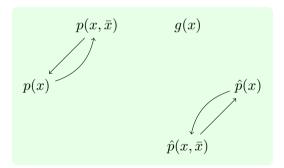
```
p(x,\bar{x})
           Network polynomial
 p(x)
           Likelihood polynomial
           Generating function
 g(x)
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         p(x,\bar{x})
                           g(x)
                                      \hat{p}(x)
p(x)
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## Progress update

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        p(x,\bar{x})
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## Progress update

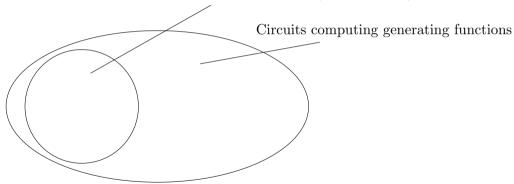
 $p(x, \bar{x})$  Network polynomial p(x) Likelihood polynomial g(x) Generating function  $\hat{p}(x)$  Fourier transform



Generating functions: Why?

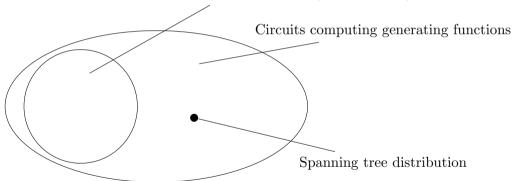
## Generating functions: Why?

Circuits computing network polynomials (S, D, positive)



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Circuits computing network polynomials (S, D, positive)



$$g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$$

Can we do marginal inference?

$X_1$	$X_2$	Pr
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$$g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$$

Can we do marginal inference? For  $X_i = 1$ , set  $x_i = t$ 

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Can we do marginal inference?

For  $X_i = 1$ , set  $x_i = t$ 

For  $X_i = 0$ , set  $x_i = 0$ 

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Can we do marginal inference?

For  $X_i = 1$ , set  $x_i = t$ 

For  $X_i = 0$ , set  $x_i = 0$ 

For  $X_i = ?$ , set  $x_i = 1$ 

$X_1$	$X_2$	Pr
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$$g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$$

Can we do marginal inference?

For  $X_i = 1$ , set  $x_i = t$ 

For  $X_i = 0$ , set  $x_i = 0$ 

For  $X_i = ?$ , set  $x_i = 1$ 

$X_1$	$X_2$	Pr
0	0	.1
0	1	.2
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$$Pr[X_1 = 1]?$$

$$g(t, 1) = .1 + .2 + .3t + .4t = .3 + .7t$$

## Generating functions

$$g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$$

Can we do marginal inference?

For  $X_i = 1$ , set  $x_i = t$ 

For  $X_i = 0$ , set  $x_i = 0$ 

For  $X_i = ?$ , set  $x_i = 1$ 

 $Pr[X_1 = 1]?$  g(t, 1) = .1 + .2 + .3t + .4t = .3 + .7t

Relation to network polynomial?

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$$\begin{array}{c|cccc}
0 & 1 & .2 \\
1 & 0 & .3 \\
1 & 1 & .4
\end{array}$$

Relation to network polynomial?

Transformation from network to a generating:  $p(x_1, x_2, \bar{x}_1, \bar{x}_2) = .1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2$ 

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$$X_i = ?$$
, set  $x_i = 1$ 
 $X_1 \quad X_2 \mid \Pr$ 
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Theorem 4. Let Pr be a probability distribution on n binary random variables. Then a circuit of size s computing the probability generating function for Pr can be transformed to a circuit of size  $O(sn^2)$  computing the network polynomial for Pr.

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Example: Starting with  $g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$ , we form

$$(\bar{x}_1\bar{x}_2)\left(.1 + .2\frac{x_2}{\bar{x}_2} + .3\frac{x_1}{\bar{x}_1} + .4\frac{x_1}{\bar{x}_1}\frac{x_2}{\bar{x}_2}\right)$$

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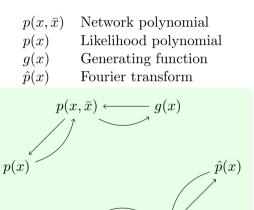
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Strassen to the rescue!

```
p(x,\bar{x})
          Network polynomial
p(x)
           Likelihood polynomial
g(x)
           Generating function
\hat{p}(x)
           Fourier transform
      p(x,\bar{x})
                         g(x)
                       \hat{p}(x,\bar{x})
```

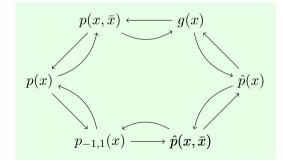
```
p(x,\bar{x})
           Network polynomial
p(x)
           Likelihood polynomial
g(x)
           Generating function
\hat{p}(x)
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      p(x,\bar{x}) \longleftarrow g(x)
                        \hat{p}(x,\bar{x})
```

```
p(x,\bar{x})
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\hat{p}(x)
```



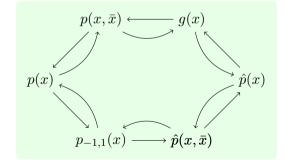
Proposition 1. For binary random variables, probability generating functions g(x) and Fourier polynomials  $\hat{p}(x)$  are the same function(!), on respective domains  $\{-1,1\}^n$  and  $\{0,1\}^n$ , up to the bijection  $\phi: \{0,1\} \to \{-1,1\}$  given by  $\phi(b) = (-1)^b$  applied bitwise.

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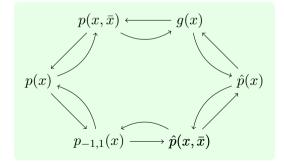
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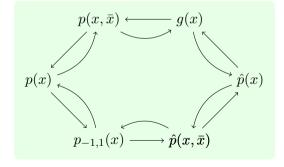


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- several distinct circuit-based models are equally succinct
- distinct inference algorithms in a common framework

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Non-binary distributions?

## Non-binary distributions?

Let  $\Pr: K^n \to \mathbb{R}$  be a probability mass function with  $K = \{0, 1, 2, \dots, k-1\}$ . Then the probability generating polynomial of  $\Pr$  is

$$g(x) = \sum_{(d_1, d_2, \dots, d_n) \in K^n} \Pr(d_1, \dots, d_n) x_1^{d_1} x_2^{d_2} \cdots x_n^{d_n}.$$
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 (1)

Theorem 5. For  $k \geq 4$ , computing likelihoods on a circuit for g(x) is #P-hard. Proof idea: Reduce from 0, 1-permanent.

#### What we did:

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#### What's next?

- ▶ Are there more succinct tractable representations? e.g., do we need multilinearity?
- ► Can we characterize *all* tractable marginal inference?
- ▶ How can theoretically more expressive models be learned/constructed in practice?

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