Fall 2025-2026 CSSE 220

CSSE 220 – Object-Oriented Software Development Rose-Hulman Institute of Technology

Worksheet 19

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 ${\rm CSSE}~220$ Fall 2025-2026

8.	Review the code for Selection Sort: 1). In selection sort, for an array of length n, how many times does the OUTER LOOP execute, in terms of n?
	2). In selection sort, for an array of length n, how many times is Comparison called during the FIRST iteration of the outer loop?
	3) during the SECOND iteration of the outer loop?
	4) during the SECOND-TO-LAST iteration of the outer loop?
	5) during the LAST iteration of the outer loop?
9.	Formula for the number of times selectionSort() calls Comparison for an array size n:
10	Summary Table for Selection Sort:
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Case	Number of Comparisons	BigO
Worst		
Average		
Best		

11. Suppose the insertion sort algorithm is applied to the initial array above. Show the state of the array immediately following each of the first three iterations of the outer loop. Clearly mark (as in 0th is alreday marked for you) the sorted part of the array after each iteration.

0 st itera	tion:								
9	5	7	10	18	1	12	8	16	4
1 st itera	tion:								
2 nd itera	ation:	•				'		,	
3 rd itera	ition:			ı			ı		

12. Summary Table for InsertionSort:

Case	BigO
Worst (reversed)	
Average	
Best (sorted)	

13. True/False Binary search must be sorted



Fall 2025-2026 CSSE 220

14. Summary Table for Binary Search:

Case	BigO
Worst	
Best	
(ony one comparison is needed)	

15. What is run-time in terms of Big-O?

```
public static void function3(int[] array) {
2
        for(int i = 1; i <= array.length; i++) {</pre>
3
            for(int j = array.length; j >= 1; j = j / 2) {
                  if(array[i-1] <= array[j-1]) {</pre>
4
                      array[j-1] = array[i-1];
5
6
             }
          }
8
9 }
10
  //Runtime:
11
```

16. Summary Table for MergeSort:

Case	BigO
Worst	
Best	

17. Suppose the merge sort algorithm from class is applied to the initial array.

9	5	7	10	18	1	12	8	16	4	
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Show the state of the two sub-arrays immediately before the final merge.

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28 October 2025