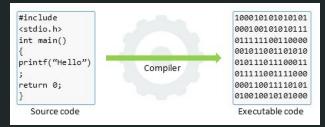


# Intro to Reverse Engineering

with a not so good reverse-engineer

## What is Reverse Engineering?

 A typical program is compiled into a bytecode / assembly language from source code



- Reverse engineering is the opposite of this process.
- If we ever wanted to get an idea of exactly what the binary does, we cannot be sure without peeking inside
- It's done to find security exploits, audit software, analyze malware, perform IP theft, create open-source drivers, crack software and create game hacks

## Choose your weapon

- Our target applications are going to be running on Windows on the x86 architecture
- How are we going to reverse engineer them?
- There are a LOT of tools, all with pros and cons:
  - IDA Pro + Hex Rays decompiler
  - Ghidra
  - Ollydba
  - Binary Ninja
  - Radare2 + Cutter UI
  - Hopper 0



- IDA Pro has always been the most mature disassembler and was king for 15+ years.
- However, it costs ~\$7000, \$1879 for the disassembler & \$5258 for the Hex-Rays decompiler
- The freeware edition is very powerful but lacks Hex-Rays, you may need to source it somewhere else
- The alternative is Ghidra, released by the NSA for... free. And it includes a decompiler
- But hey, we're going to use IDA cause it makes us feel rich



## Disassembler vs Decompiler

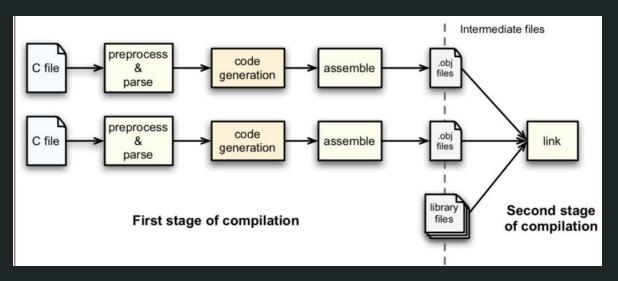
- A disassembler simply reads the machine code and translate it back into its relevant instructions.
- This is trivial in comparison to a decompiler
- A decompiler turns disassembled code back into (somewhat pseudo) source code

```
sub_7FFA38B204F0 proc near
   var 28= gword ptr -28h
   var 20= gword ptr -20h
   var 18= gword ptr -18h
   var 10= dword ptr -10h
    arg 18= dword ptr 20h
000 sub
048 and
            [rsp+48h+arg_18], 0
048 mov
            r9b, r8b
048 movzx
            r8d. dx
048 mov
048 lea
            rcx, [rsp+48h+arg 18]
048 call
            sub 7FFA38B28C24
048 mov
            eax, [rsp+48h+arg 18]
048 lea
           r9, aMilanreadrawda : "MilanReadRawData"
048 mov
            rcx, cs:pDebugStruct
048 lea
            r8, aEGitGenevaRemo : "e:\\git\\geneva remote\'
048 mov
            [rsp+48h+var 10], eax
048 mov
            edx. 8
048 lea
            rax, aExitX
                          ; "Exit:%x"
048 mov
            [rsp+48h+var_18], rax
            [rsp+48h+var_20], 0
048 and
048 mov
            dword ptr [rsp+48h+var_28], 463h
048 call
                            ; check if this is nullptr
048 mov
            eax, [rsp+48h+arg 18]
048 add
            rsp, 48h
000 retn
    sub 7FFA38B204F0 endp
```

```
fastcall sub 7FFA38B204F0( int64 al, unsigned int16 a2, char a3,
     _int64 v5; // [rsp+20h] [rbp-28h]
    unsigned int v6; // [rsp+38h] [rbp-10h]
    unsigned int v7: // [rsp+68h] [rbp+20h]
     v7 = 0:
     LOBYTE(a4) = a3:
     sub 7FFA38B28C24(&v7, a1, a2, a4):
     v6 = v7;
    LODWORD(v5) = 1123;
12
     DebugPrint(
13
       ( int64)pDebugStruct.
14
       8u.
      L"e:\\git\\geneva_remote\\winfpcode\\driver\\common\\chipoperation\\milanfseries\\md
15
       L"MilanReadRawData",
16
17
       v5.
18
19
       L"Exit: %x".
20
       v6);
    return v7;
223
```

# Compilation 101

• We're gonna need to know how a C and some basic C++ files end up as an executable, how are "dependencies" managed?



## Compilation 101

- There exists two types of linkage, static linkage & dynamic linkage.
- Static linking a library means the library binary is copied into our application executable. When the application calls a library function, the address of the function to be called is simply hardcoded in
- Dynamic linking allows for a different library file to be loaded at runtime. This could mean we could update our library and as long as we kept the same **Application Binary Interface**, everything will still work.
- Typically, the OS libraries and C runtime (i.e the library that contains the implementations of C functions) are dynamically linked.
- This is where the term **Dynamic Link Library (.dll)** on Windows comes from. The Linux equivalent would be the **.so, shared object**
- These libraries are built to be relocatable, so they can be loaded at any memory address
- Should also be noted that we ourselves could dynamically load a library. More on this later

## Portable Executables and Dynamic Linking

- How does the OS know what DLLs to load when a program boots?
- Portable Executables (.exe & .dll) have an import table
- An .exe which would import a function from a .dll would list the DLL name on it's Import table
- Whilst that DLL would list it's exported functions on it's export address table (EAT) of it's PE
- When we launch our program:
  - The OS reads our Import Table
  - Maps the requested library into the memory space
  - Adds a record to the Import Address Table of our application that points to the function
- When the application wants to call that function, it simply looks at it's IAT, follows the record to the loaded library and calls the function
- IDA, CFF Explorer Demo + Limits of Static Analysis

## **API** Monitoring

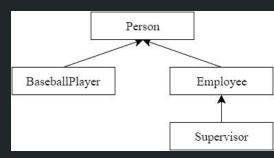
- The nature of exports & imports means the symbol names must generally be human-readable.
- This mean it's inevitable that an application will call a known function to us
- This makes it very easy to see the behaviour of the application, as we can see it use calls to the OS
  and public, known libraries
- These calls can easily be monitored with debugging breakpoints & other techniques
- API Monitor Demo + IDA Pro

#### Classes

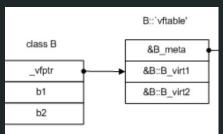
- Classes and structs members will be stored linearly in memory in the order they were declared
- There may be padding bytes in between items to align them to power of 2 boundaries
- We will investigate a way of modeling such items later
- When one class inherits another, the memory profile is simply their members concatenated together

# Function run-time type information (RTTI) & vtables

- A byproduct of function inheritance in C++, specifically in MSVC
- RTTI is used to implement dynamic\_cast and it's downcasting ability
- An instance of a class holds a pointer to a vtable as it's first member
- A vtable is a list of function pointers for each overridable method on that class



- When a class wants to override a method on a superclass, it uses a new vtable (and new pointer) with the relevant function pointer in the table changed to the overridden method
- https://godbolt.org/z/QFFjX6
- https://bit.ly/2NPNEdJ
- VTables store the name of the class types, giving us effective hints at what the class does

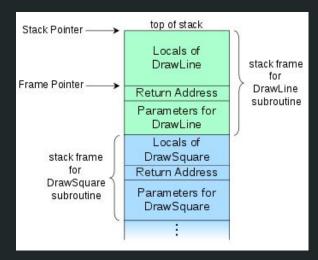


# Strings

- Remember that IDA isn't perfect and it may not find every string
- Make sure you set the string type in Compiler options correctly, Win32 often favors wchar\_t 16bit characters more than char 8bits.
- Don't underestimate the power of the xref feature

#### The Hard Work

- Now comes the most difficult part of reverse engineering: figuring out what the hell is going on
- The bad part: you're going to need to learn assembly
- The even worse part: you're going need to learn all the calling conventions
- https://en.wikipedia.org/wiki/X86 calling conventions#List of x86 calling conventions
- You're going to have to model structs as you examine function calls
- You're going to have to alias everything
- Demo?



## Breaking It

- Now that we know how something works, how do we break it?
- Let's do something fun: game hacks
- Generally game hacks can be categorized as two types: internal & external
- An external hack is one that lives outside the games process and modifies memory of the game using OS calls, namely ReadProcessMemory & WriteProcessMemory
- Some people believe this to be more "secure" against anti-cheats, but in general it tends to be pretty slow
- Internal game hacks inject code into the games process and runs it. An existing thread is either hijacked or we start a new one
- These perform just as fast as the game does, and you can even use the game's internal APIs if you wish

# External Hack Demo: ReClass & SigScans

- ReClass is a tool that allows you to model structs inside a process, you can then export them as C/C++ header files. We could then import them to IDA or in our hack code
- Sig Scanning is a method where we can scan the game process for a set of bytes.
- A signature can be made up of any number of bytes with "mask" bytes whose values are ignored during the search.
- You can generally make these signatures by hand
- CSGO Demo w/ CDebugOverlay, IVEntityList & s\_pTextOverlays

## Internal Hacks: DLL Injection & Hooks

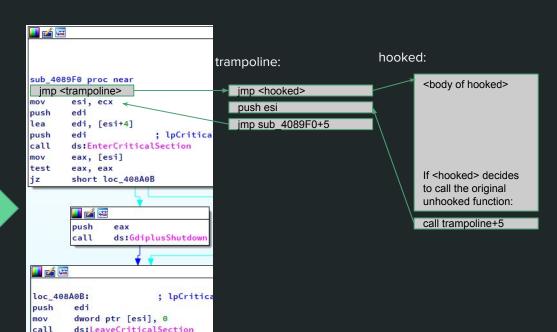
- We can build our game hack as a DLL and force the game to load it
- There are a few methods to do this:
  - Call CreateRemoteThread & use it to run LoadLibrary("uwu.dll")
  - Pause all threads, use NtSetContextThread
  - Manual mapping

- Once we're in, what can we do?
  - Our main goal would be to "hook" functions
  - Function hooking is similar to writing an event handler
  - Whenever the hooked function gets called, our function gets called
  - We can modify any parameters that get passed, completely block the call, extend functionality, etc...

## JMP Hooking

- JMP hooking involves modifying the contents of a function directly. This means that memory permissions of the page containing the function will need to be modified.
- We will need to replace the first 5 bytes of the function with a jmp to our "trampoline"
- We will need to allocate memory for the "trampoline" stub, the trampoline stub contains a jmp to our hooked function, and the instruction we removed above followed by a jmp back to the original function

```
sub_4089F0 proc near
       esi
push
       esi, ecx
mov
       edi
push
lea
       edi, [esi+4]
push
       edi
                       ; lpCriticalSection
       ds:EnterCriticalSection
call
mov
       eax, [esi]
       eax, eax
test
jz
       short loc_408A0B
       push
               eax
               ds:GdiplusShutdown
        call
II 🚄
loc_408A0B:
                       ; lpCriticalSection
        edi
push
        dword ptr [esi], 0
mov
call
        ds:LeaveCriticalSection
pop
        edi
pop
        esi
retn
sub_4089F0 endp
```



edi

esi

sub\_4089F0 endp

pop

pop

retn

```
void Attached();
 void APIENTRY h glBegin(GLenum mode);
□void Attached()
     hOpenGL = GetModuleHandle("opengl32.dl1");
     if (hOpenGL == INVALID HANDLE VALUE) return;
     o_glBegin = (glBegin_t)DetourFunction((LPBYTE)GetProcAddress(hOpenGL, "glBegin"), (LPBYTE)&h_glBegin);

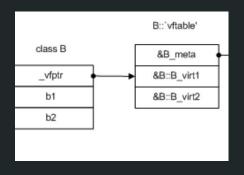
    □void APIENTRY h glBegin(GLenum mode)

     float curcol[4];
     glGetFloatv(GL_CURRENT_COLOR, curcol);
     glDisable(GL_DEPTH_TEST);
     glEnable(GL_BLEND);
     glBlendFunc(GL SRC ALPHA, GL ONE MINUS SRC ALPHA);
     glColor4f(curcol[0], curcol[1], curcol[2], 0.5f);
     (*o_glBegin)(mode);
■BOOL WINAPI DllMain(HINSTANCE hDllInstance, DWORD nReason, LPVOID Reserved)
     switch (nReason)
     case DLL PROCESS ATTACH:
         DisableThreadLibraryCalls(hDllInstance);
         hThisDll = hDllInstance;
         hThisThread = CreateThread(0, 0, (LPTHREAD START ROUTINE)Attached, 0, 0, NULL);
         return (hThisThread != INVALID_HANDLE_VALUE);
         break;
     return 0;
```

Detnuct Chart t

# VTable/VMT Hooking

- Remember those vtables from earlier? The vfptr is the first member variable of a class
- We could create a copy of the vtable and modify a function pointer from its original address to the address of our hooked function
- Then all we'll need to do is change the vfptr to our copy
- We can now hook any method that has been declared as virtual
- This method is incredibly popular with game overlays / screen recorders
- One goal is to hook IDXGISwapChain::Present. Our hooked function will be called whenever the screen is drawn by DirectX. i.e from there we can draw what we wish on the game screen, even if fullscreen



## VEH / PAGE\_GUARD Hooking

- Vectored Exception Handling can be used as a hooking mechanism by triggering exceptions.
- We can set the PAGE\_GUARD bit on a memory page, whenever the CPU tries
  to access any memory in that page it triggers the
  STATUS\_GUARD\_PAGE\_VIOLATION exception, and removes the
  PAGE\_GUARD bit
- Think of it like a bear trap
- We can register an VEH exception handler and catch that exception
- By checking the exception context, we can check if the instruction pointer is pointed at the address of a function we want to hook (i.e the CPU tried to read the first instruction of the function). If so, enable single stepping or not, set PAGE\_GUARD enabled again.
- If it is the address, we can set the instruction pointer to our detoured function

# IAT (Import Address Table) Hooking

- All that needs to be done is to change the pointer to the function on the IAT
- Save the original so you can call it if needed

#### Conclusion & trivia

- ReactOS (they're going to be at FOSDEM!)
- Wine
- Phoenix Technologies / clean-room reverse engineering
- Movfuscator

# Questions