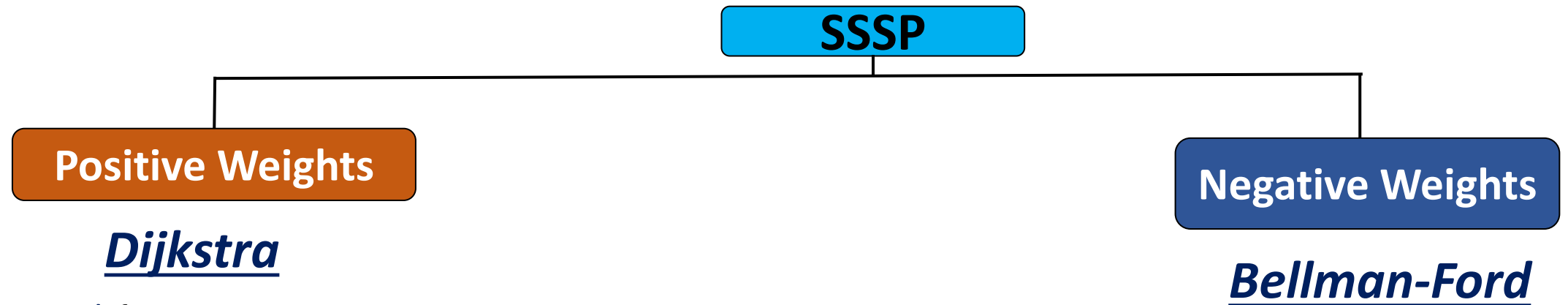


# CS2x1:Data Structures and Algorithms

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# Recap: SSSP



*DIJKSTRA* ( $G, w, s$ ) {

1 **INITIALIZE-SINGLE-SOURCE** ( $G, s$ )

2  $S = \emptyset$

3  $Q = G.V$

4 **while**  $Q \neq \emptyset$  ;

5    $u = \text{EXTRACT-MIN}(Q)$

6    $S = S \cup \{u\}$

7   **for** each vertex  $v \in Q$ .  $\text{Adj}[u]$

8     **RELAX** ( $u, v, w$ )

}

**INITIALIZE-SINGLE-SOURCE** ( $G, s$ )

{

1 **for** each  $v \in G.V$

2    $v.d = \infty$

3    $v.\pi = \text{NIL}$

4  $s.d = 0$

**RELAX** ( $u, v, w$ ) {

1 **if**  $v.d > u.d + w(u, v)$

2    $v.d = u.d + w(u, v)$

3    $v.\pi = u$

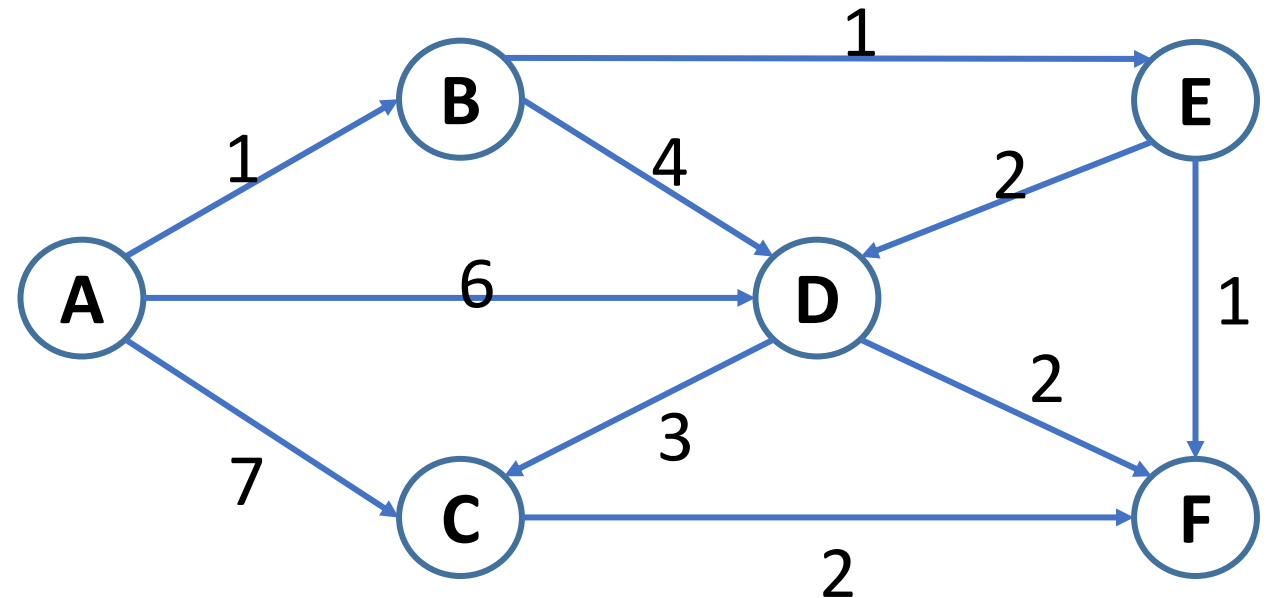
Total time complexity:  $\underline{O(V) + O(V \log V) + O(V) + O(E \log V)}$   
:  $O(E \log V)$

# Exercise: Dijkstra

Consider the following digraph starting at **vertex A** and apply Dijkstra's single source shortest path algorithm on it.

Select the correct order from the following in which vertices are removed from the Priority Queue?

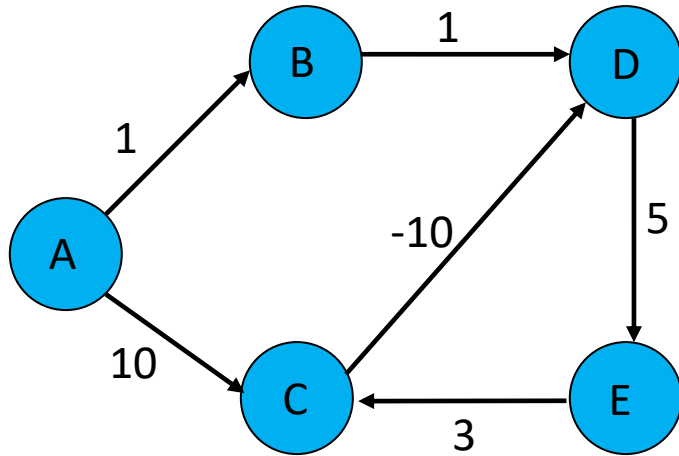
- a. A, B, C, D, E, F
- b. A, B, F, E, C, D
- c. A, B, F, E, D, C
- d. A, B, E, F, D, C



# Exercise: Dijkstra (1)

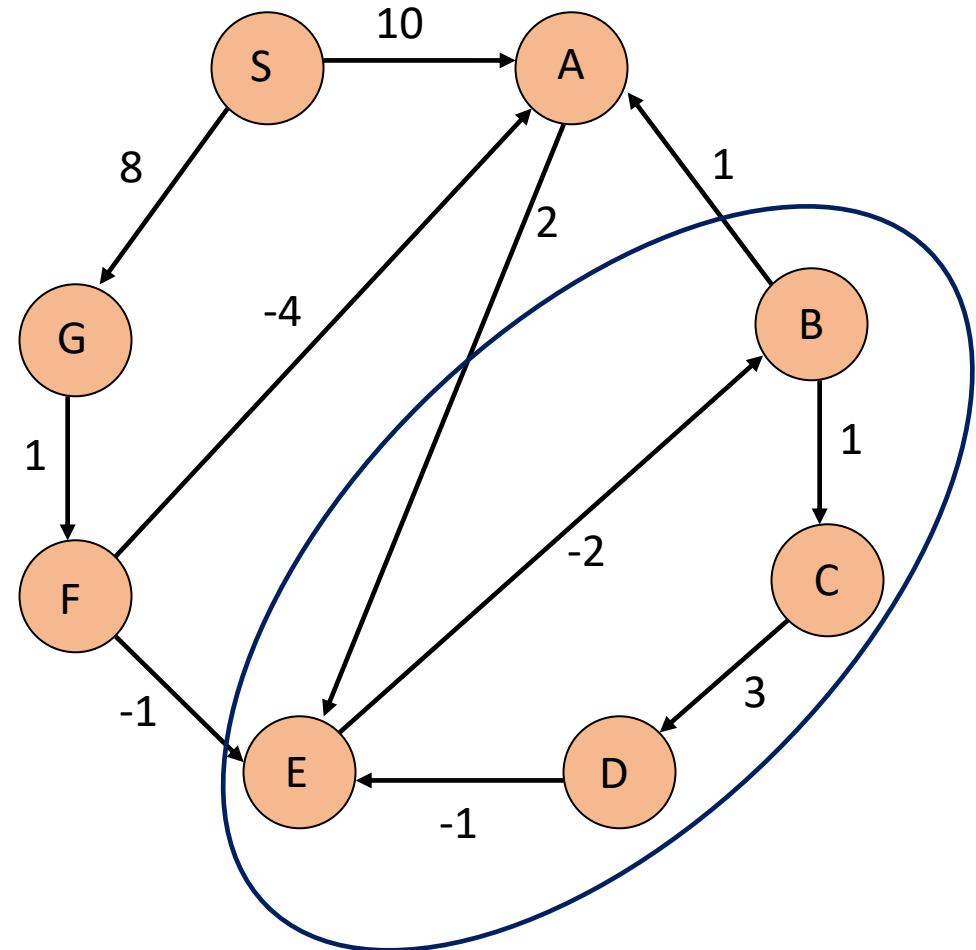
## Negative cycles

What is the shortest path from A to E?

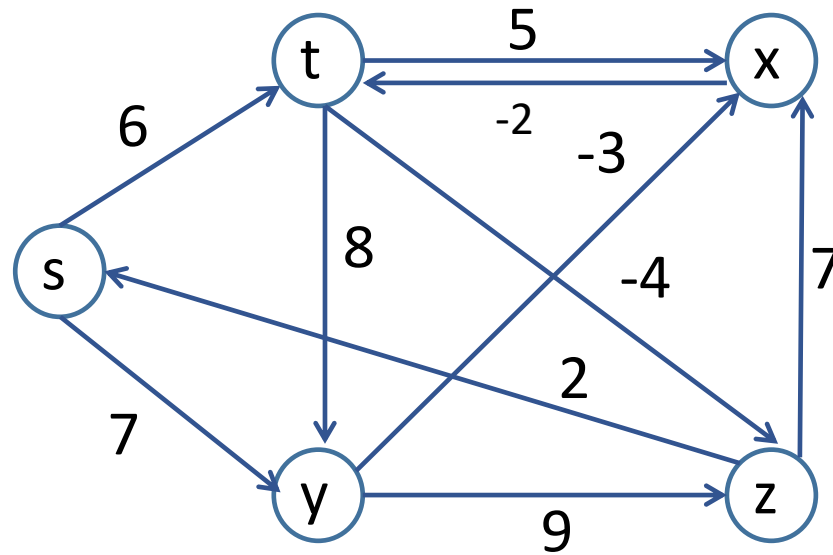
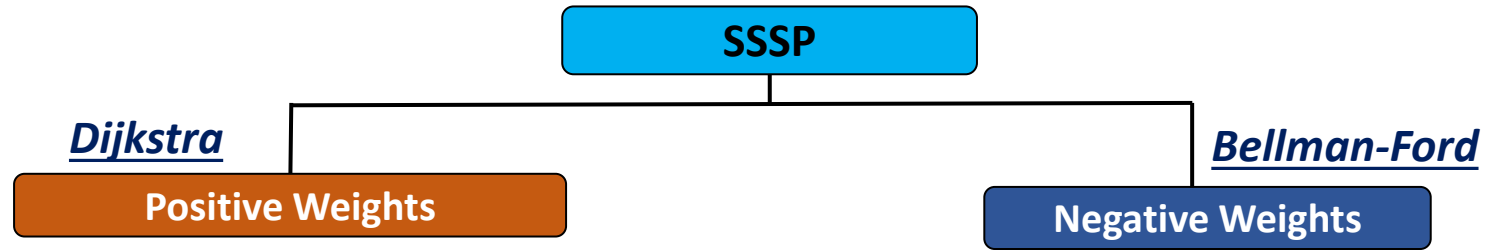


- ❖ Dijkstra doesn't work for **Graphs with negative weight cycle**, Bellman-Ford works for such graphs.
- ❖ Bellman-Ford is also simpler than Dijkstra and suites well for distributed systems.
- ❖ It does not use Priority Queue

What is the shortest path from S to D?



# SSSP → Dijkstra



# SSSP: Bellman-Ford

1

```

BELLMAN-FORD (G, w, s) {
1  INITIALIZE-SINGLE-SOURCE (G, s)
2  for i = 1 to |G.V| - 1
3    for each edge (u, v) ∈ G.E
4      RELAX (u, v, w)
5  for each edge (u, v) ∈ G.E
6    if v. d > u. d + w(u, v)
7      return False
8  return True
    
```

2

```

INITIALIZE-SINGLE-SOURCE (G, s) {
1  for each v ∈ G.V
2    v. d = ∞
3    v. π = NIL
4  s. d = 0
    
```

3

```

RELAX (u, v, w) {
1  if v. d > u. d + w(u, v)
2    v. d = u. d + w(u, v)
3    v. π = u
    
```

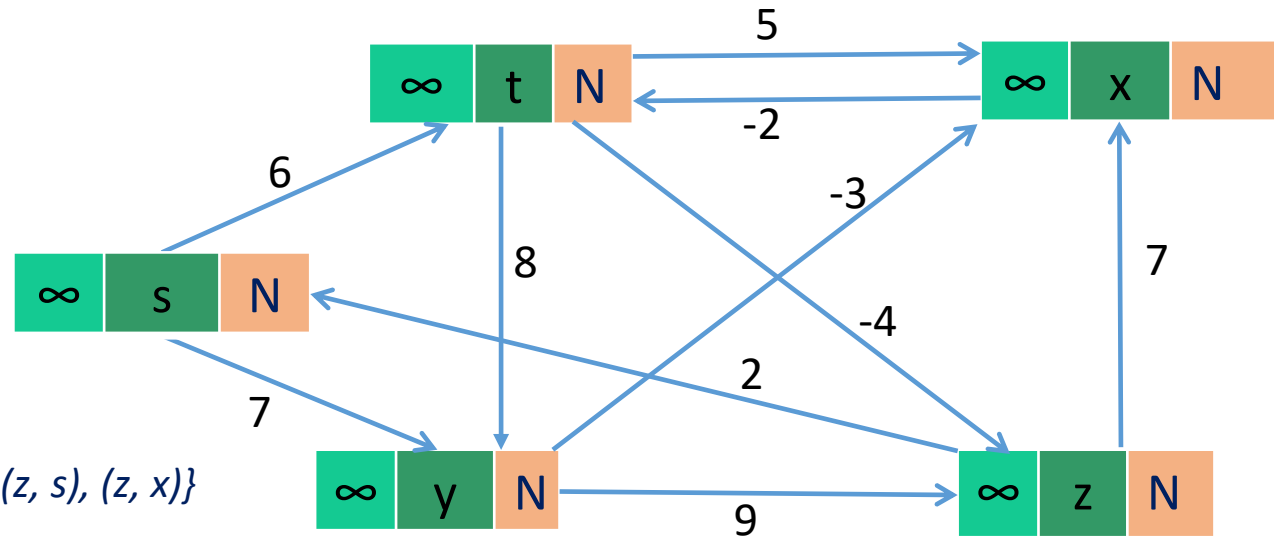
Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 - 4: for i = 1 to 5 - 1

Step 3 : # of edges =

{ (s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x) }

Step 4 : RELAX (s, t, w); RELAX (s, y, w)



# SSSP: Bellman-Ford (1)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 1$  to  $5 - 1$

Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $s, t, w$ );

$\infty > 0 + 6$ ;

$t.d = 6$ ;

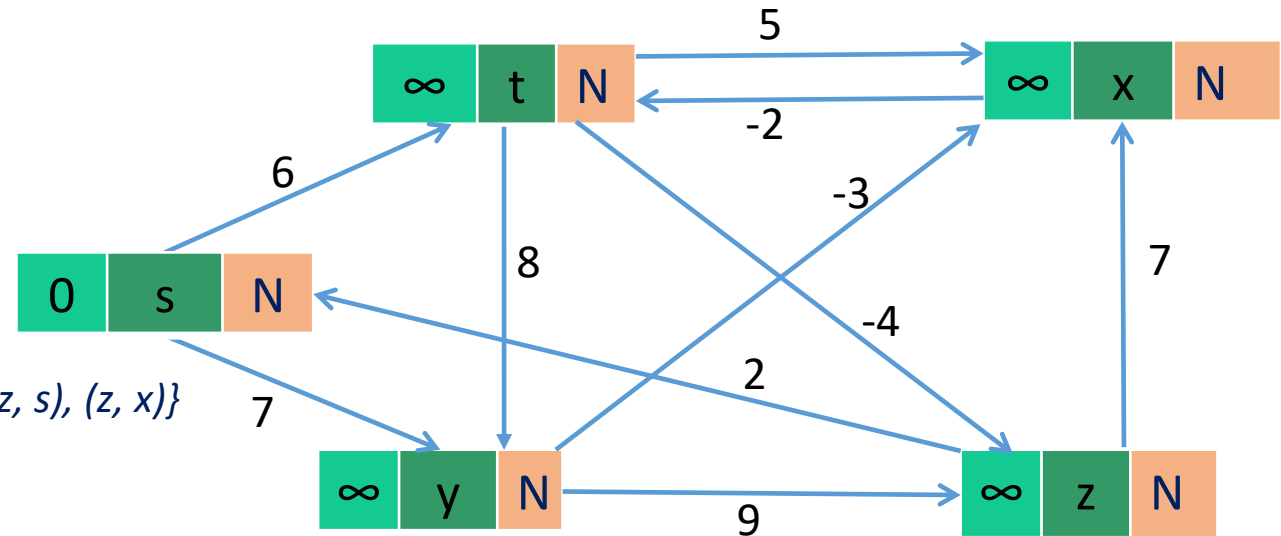
$t.\pi = s$ ;

RELAX ( $s, y, w$ )

$\infty > 0 + 7$ ;

$y.d = 7$ ;

$y.\pi = s$ ;



# SSSP: Bellman-Ford (2)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 1$  to  $5 - 1$

Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $t, x, w$ );

RELAX ( $t, y, w$ );

RELAX ( $t, z, w$ )

$\infty > 6 + 5$ ;

$x.d = 11$ ;

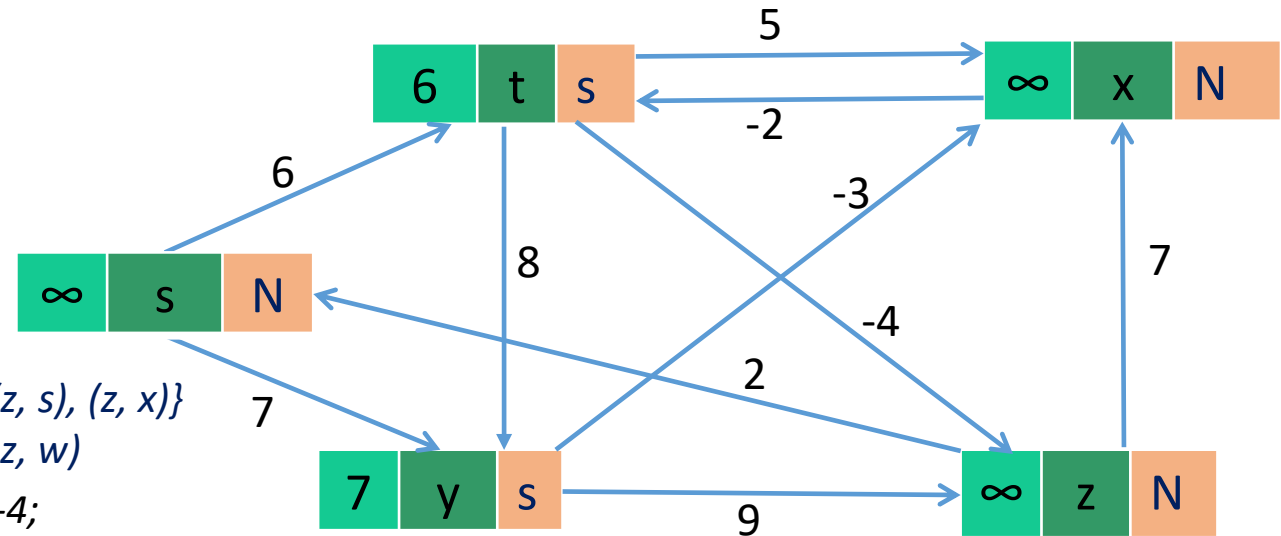
$x.\pi = t$ ;

$7 > 6 + 8$ ;

$\infty > 6 + -4$ ;

$z.d = 2$ ;

$z.\pi = t$ ;





# SSSP: Bellman-Ford (3)

```

BELLMAN-FORD (G, w, s) {
1 INITIALIZE-SINGLE-SOURCE (G, s)
2 for i = 1 to |G.V| - 1
3   for each edge (u, v) ∈ G.E
4     RELAX (u, v, w)
5 for each edge (u, v) ∈ G.E
6   if v. d > u. d + w(u, v)
7     return False
8 return True
    
```

```

INITIALIZE-SINGLE-SOURCE (G, s) {
1 for each v ∈ G.V
2   v. d = ∞
3   v. π = NIL
4 s. d = 0
    
```

```

RELAX (u, v, w) {
1 if v. d > u. d + w(u, v)
2   v. d = u. d + w(u, v)
3   v. π = u
    
```

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for i = 1 to 5 - 1

Step 3 : # of edges =

{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)}

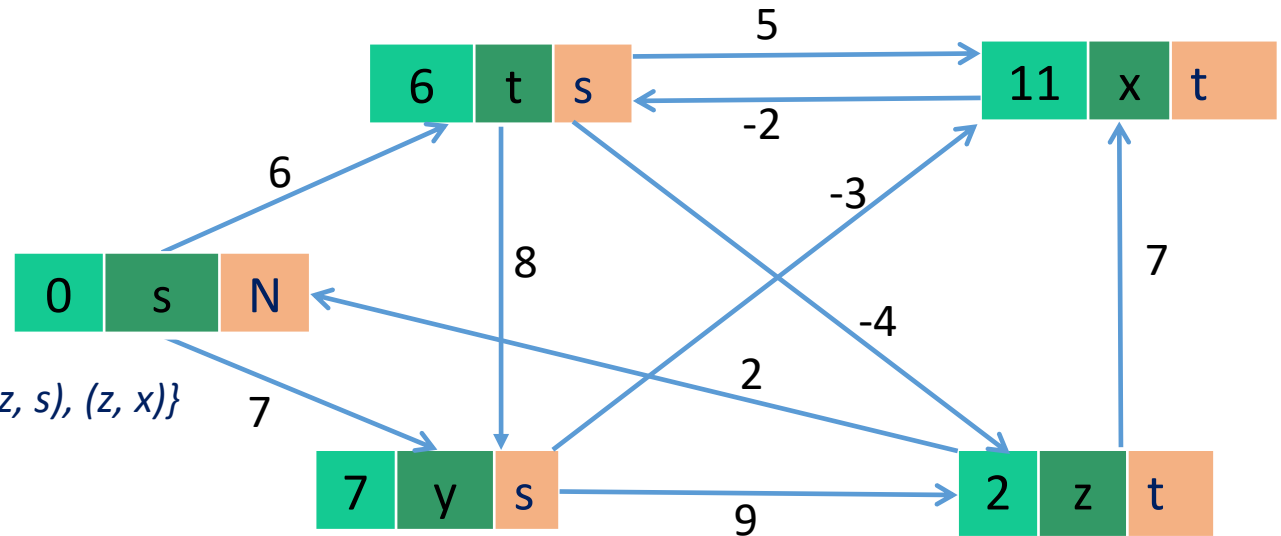
Step 4 : RELAX (y, x, w); RELAX (y, z, w);

11 > 7 + -3;

x. d = 4;

x. π = y;

2 > 7 + 9;



# SSSP: Bellman-Ford (4)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 1$  to  $5 - 1$

Step 3 : # of edges =

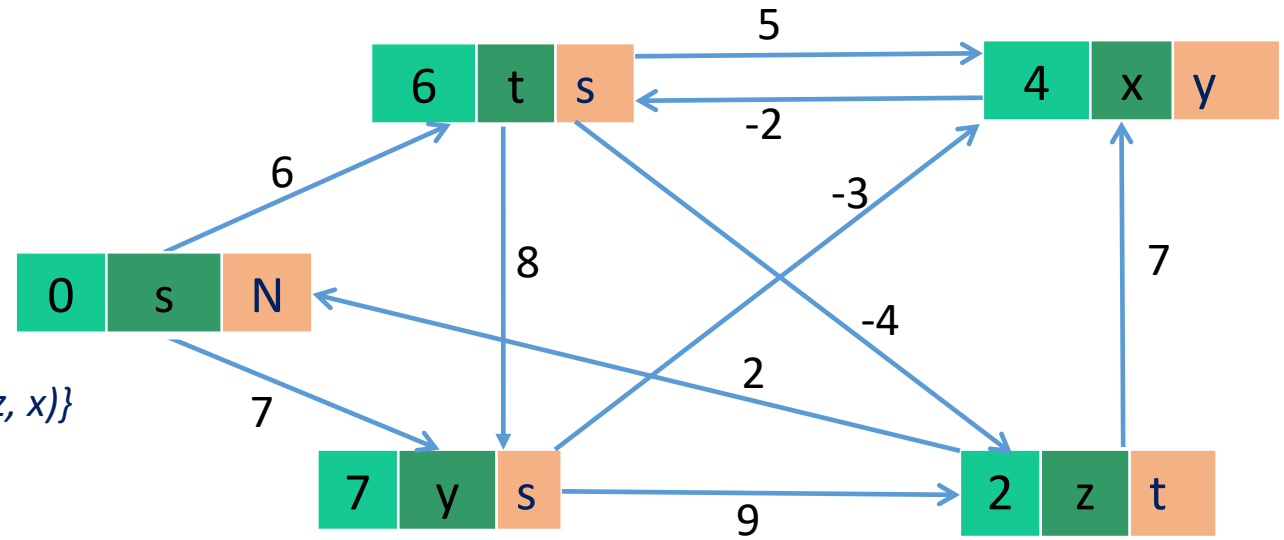
$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $x, t, w$ );

$6 > 4 + -2$ ;

$t.d = 2$ ;

$t.\pi = x$ ;



# SSSP: Bellman-Ford (5)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 1$  to  $5 - 1$

Step 3 : # of edges =

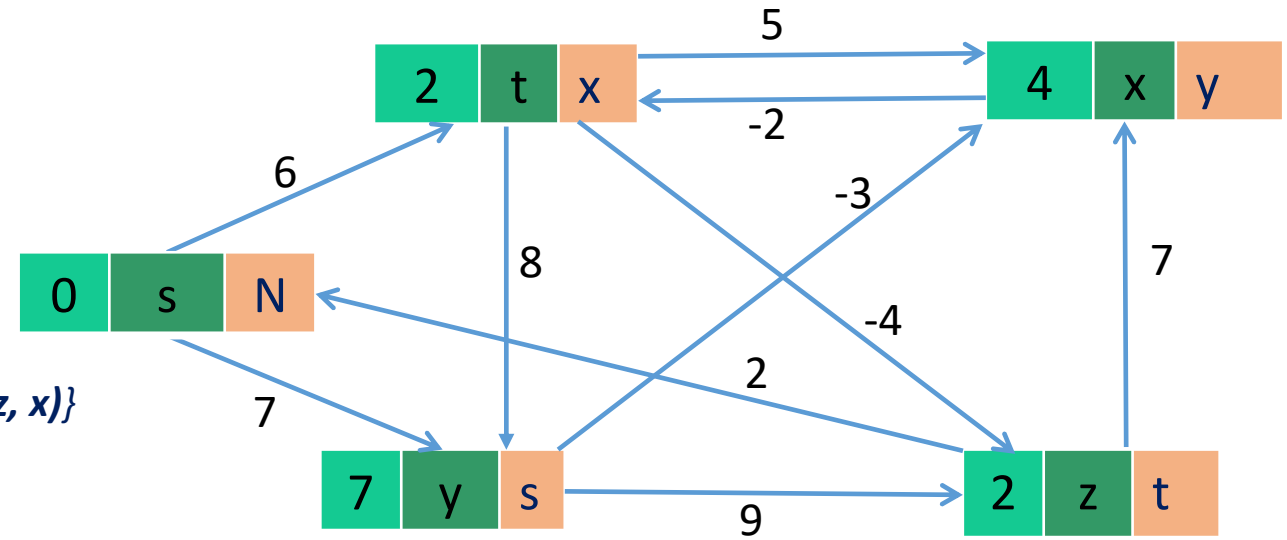
$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $z, s, w$ );

$0 > 2 + 2;$

RELAX ( $z, x, w$ )

$4 > 2 + 7;$



# SSSP: Bellman-Ford (6)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 2$

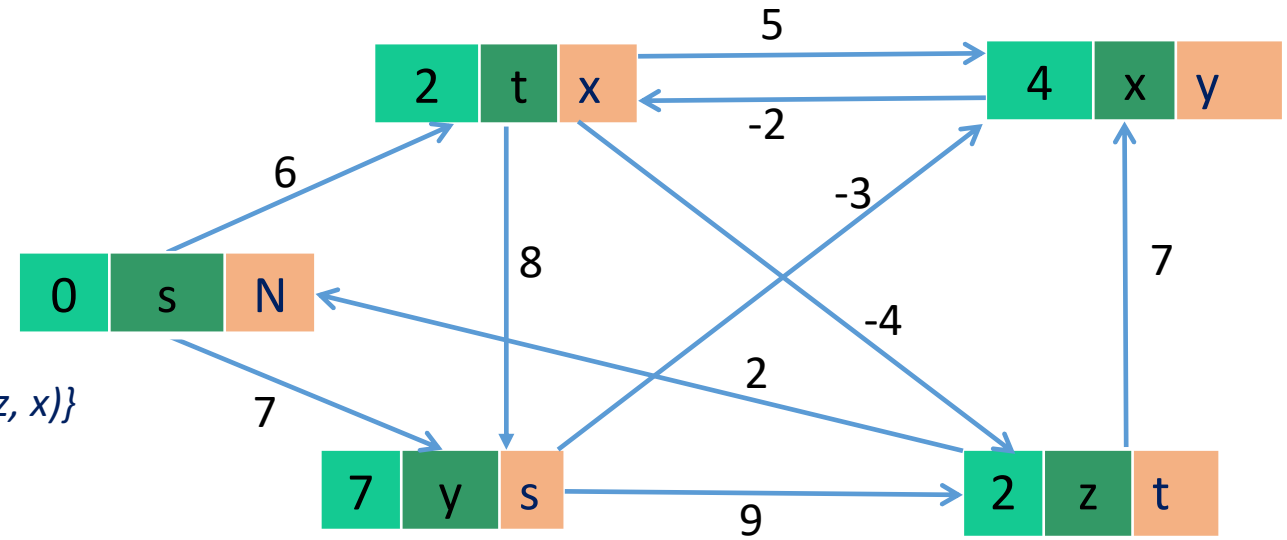
Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $s, t, w$ );     RELAX ( $s, y, w$ )

$2 > 0 + 6;$

$7 > 0 + 7;$



# SSSP: Bellman-Ford (7)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 2$

Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $t, x, w$ );

RELAX ( $t, y, w$ );

RELAX ( $t, z, w$ )

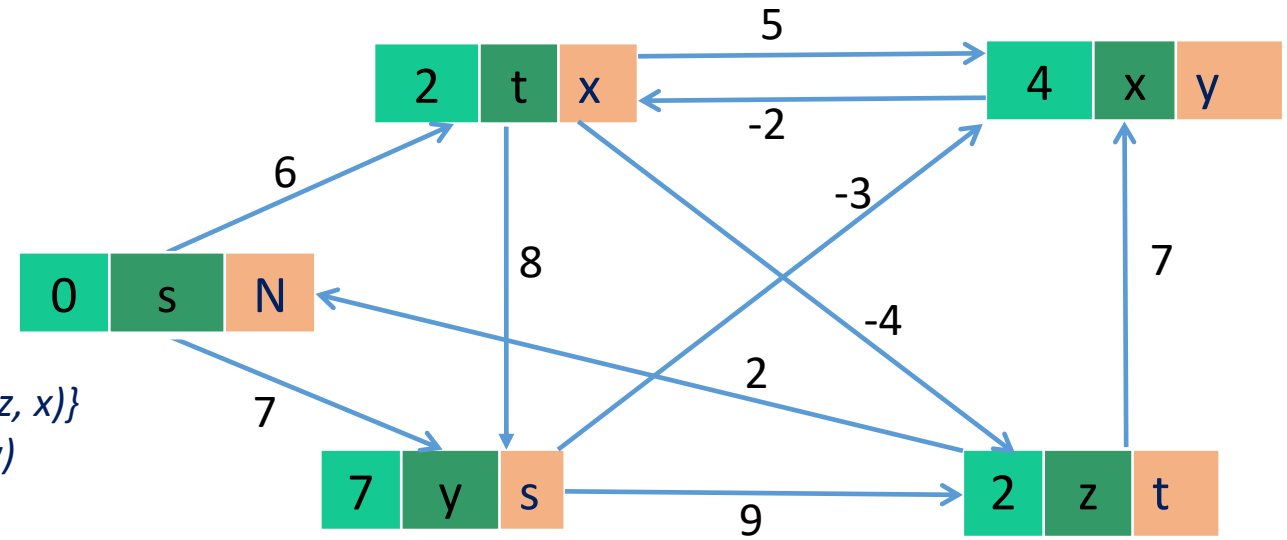
$4 > 2 + 5;$

$7 > 2 + 8;$

$2 > 2 + -4;$

$z.d = -2$

$z.\pi = t$



# SSSP: Bellman-Ford (8)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 2$

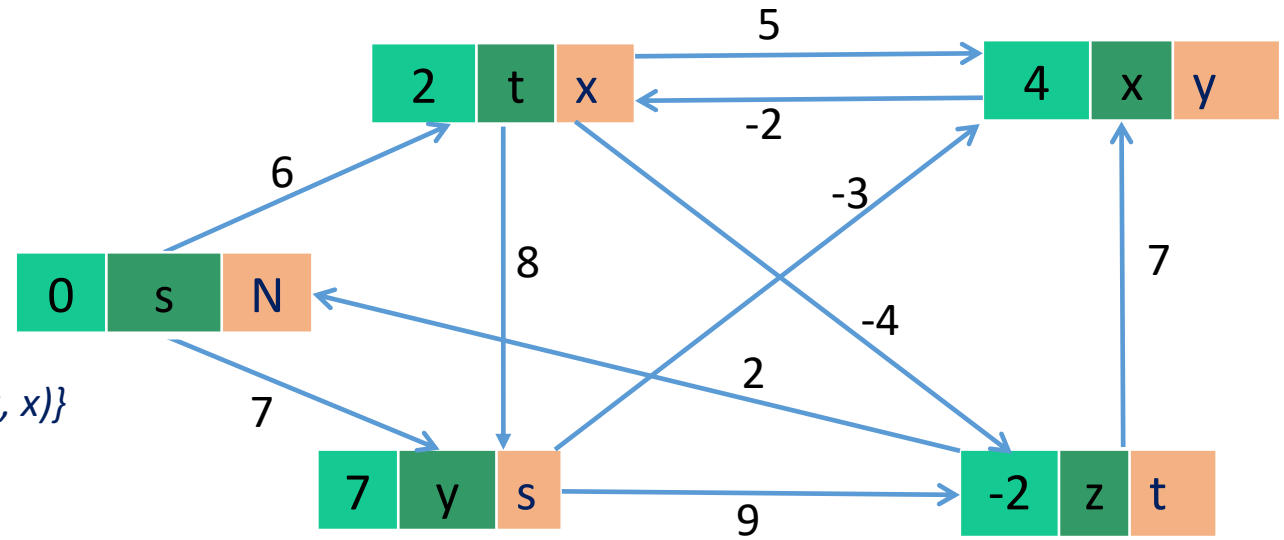
Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $y, x, w$ );     RELAX ( $y, z, w$ );

$4 > 7 + -3$ ;

$-2 > 7 + 9$ ;



# SSSP: Bellman-Ford (9)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

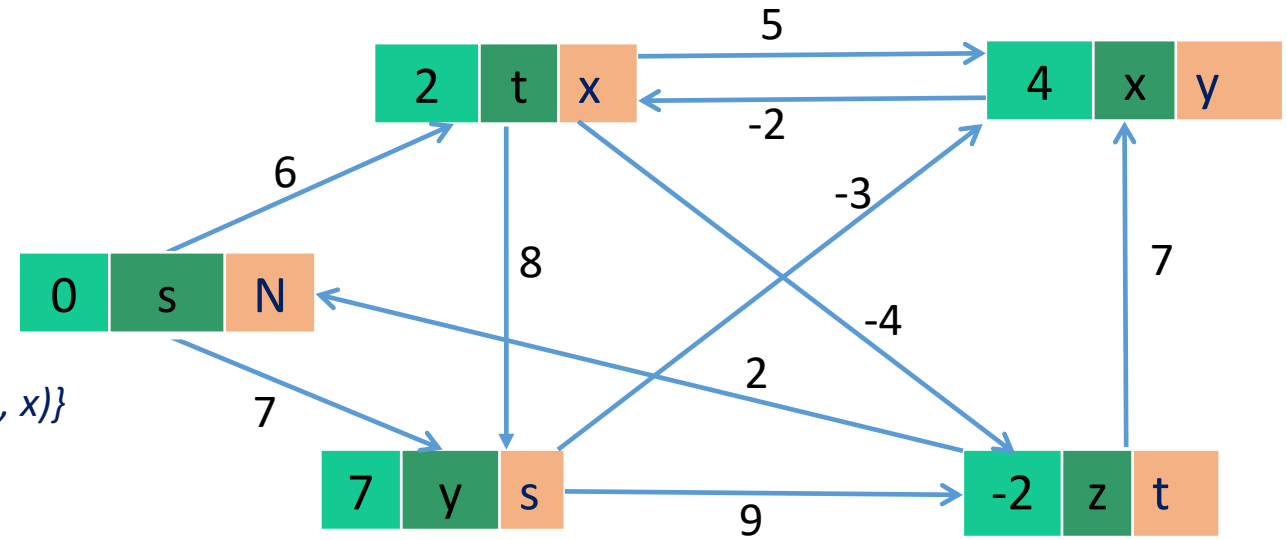
Step 2 : for  $i = 2$

Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (\mathbf{x, t}), (z, s), (z, x)\}$

Step 4 : RELAX ( $x, t, w$ );

$2 > 4 + -2;$



# SSSP: Bellman-Ford (10)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1   if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 2$

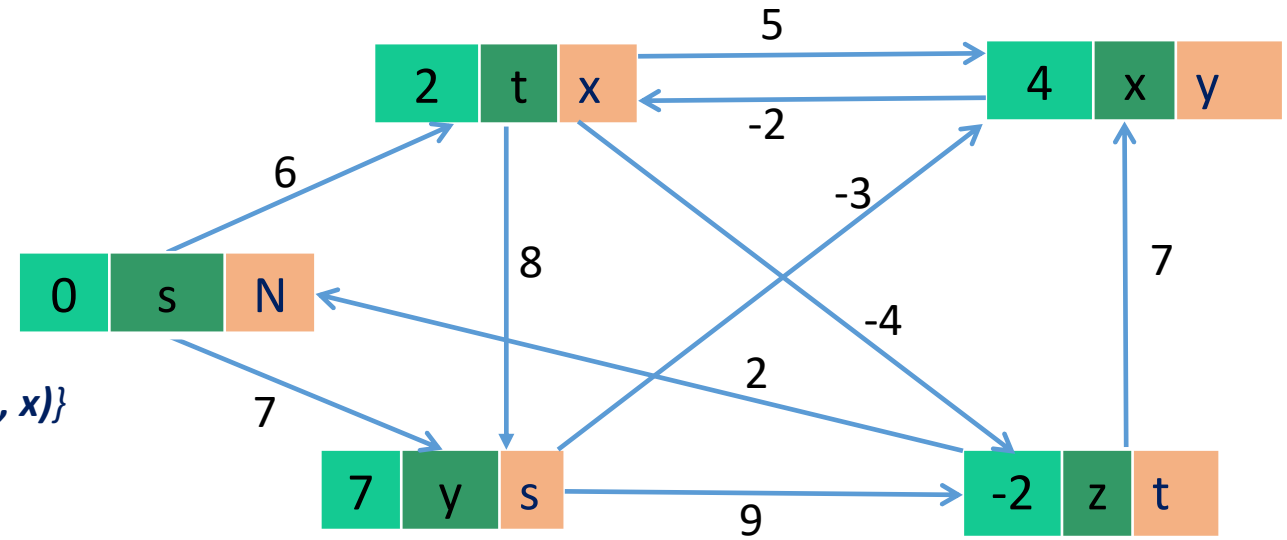
Step 3 : # of edges =

$\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $z, s, w$ );     RELAX ( $z, x, w$ );

$0 > -2 + 2;$

$4 > -2 + 7;$





# SSSP: Bellman-Ford (11)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

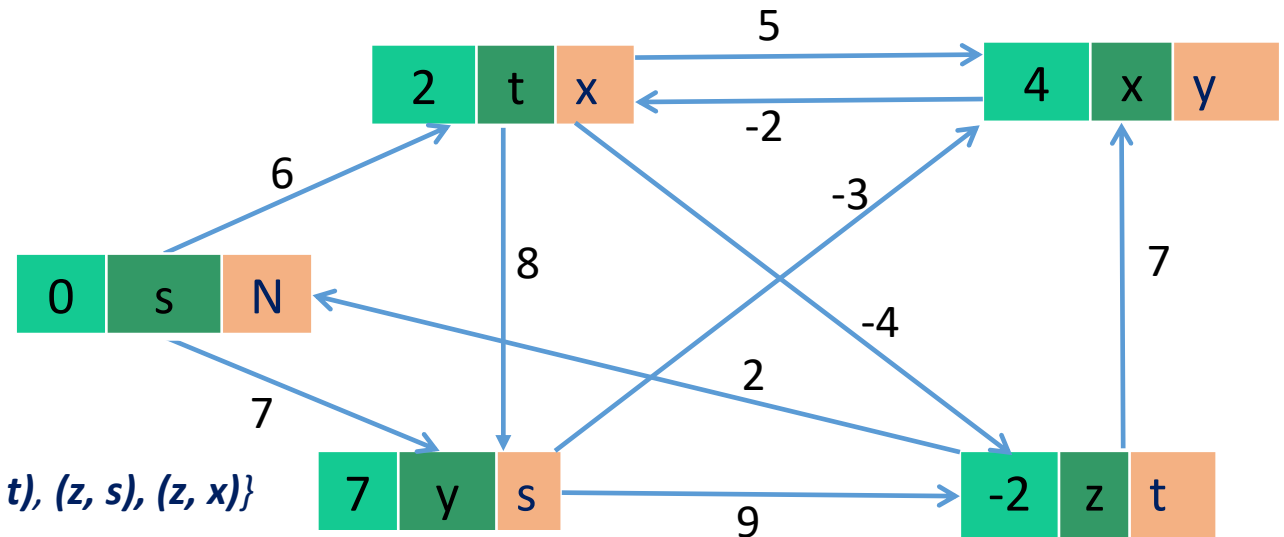
4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$



Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for  $i = 3$

Step 3 : # of edges =  $\{(s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x)\}$

Step 4 : RELAX ( $s, t, w$ ); RELAX ( $s, y, w$ ); RELAX ( $t, x, w$ ); RELAX ( $t, y, w$ ); RELAX ( $t, z, w$ ); RELAX ( $y, x, w$ ); RELAX ( $y, z, w$ ); RELAX ( $x, t, w$ );

$2 > 0 + 6;$

$7 > 0 + 7;$

$4 > 2 + 5;$

$7 > 2 + 8;$

$-2 > 2 + -4;$

$4 > 7 + -3;$

$-2 > 7 + 9;$

$2 > 4 + -2$

RELAX ( $z, s, w$ ); RELAX ( $z, x, w$ );

$0 > -2 + 2;$

$4 > -2 + 7;$

# SSSP: Bellman-Ford (12)

```

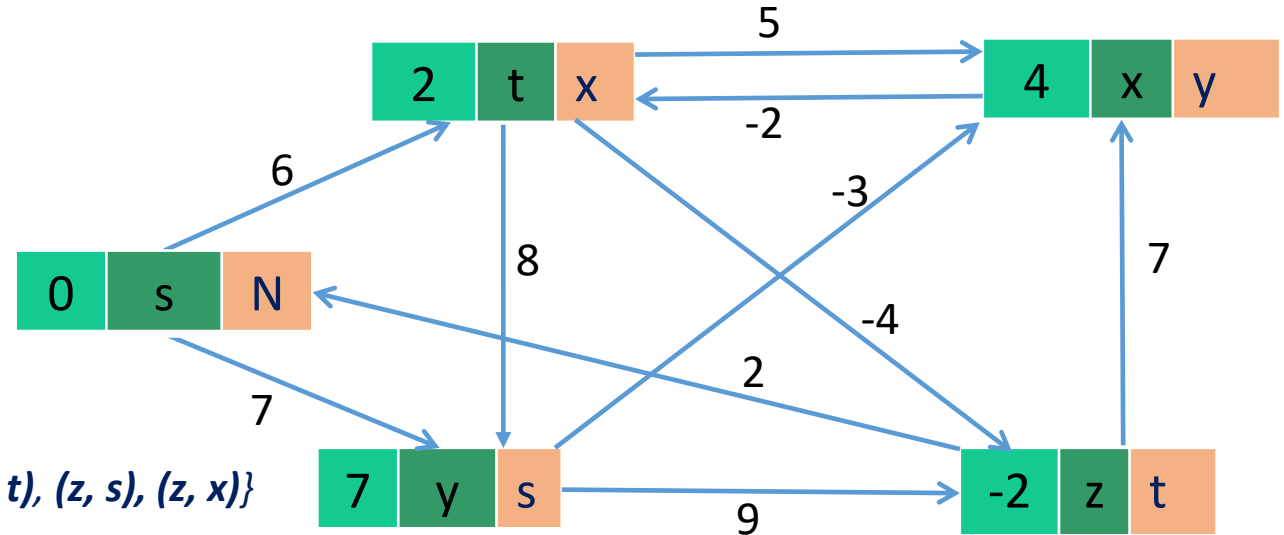
BELLMAN-FORD (G, w, s) {
1 INITIALIZE-SINGLE-SOURCE (G, s)
2 for i = 1 to |G.V| - 1
3   for each edge (u, v) ∈ G.E
4     RELAX (u, v, w)
5 for each edge (u, v) ∈ G.E
6   if v. d > u. d + w(u, v)
7     return False
8 return True
    
```

```

INITIALIZE-SINGLE-SOURCE (G, s) {
1 for each v ∈ G.V
2   v. d = ∞
3   v. π = NIL
4 s. d = 0
    
```

```

RELAX (u, v, w) {
1 if v. d > u. d + w(u, v)
2   v. d = u. d + w(u, v)
3   v. π = u
    
```



Step 1 : INITIALIZE-SINGLE-SOURCE

Step 2 : for i = 4

Step 3 : # of edges = { (s, t), (s, y), (t, x), (t, y), (t, z), (y, x), (y, z), (x, t), (z, s), (z, x) }

Step 4 : RELAX (s, t, w); RELAX (s, y, w); RELAX (t, x, w); RELAX (t, y, w); RELAX (t, z, w); RELAX (y, x, w); RELAX (y, z, w); RELAX (x, t, w);

2 > 0 + 6;      7 > 0 + 7;      4 > 2 + 5;      7 > 2 + 8;      -2 > 2 + -4;      4 > 7 + -3;      -2 > 7 + 9;      2 > 4 + -2

RELAX (z, s, w); RELAX (z, x, w);

0 > -2 + 2;      4 > -2 + 7;

# SSSP: Bellman-Ford (13)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

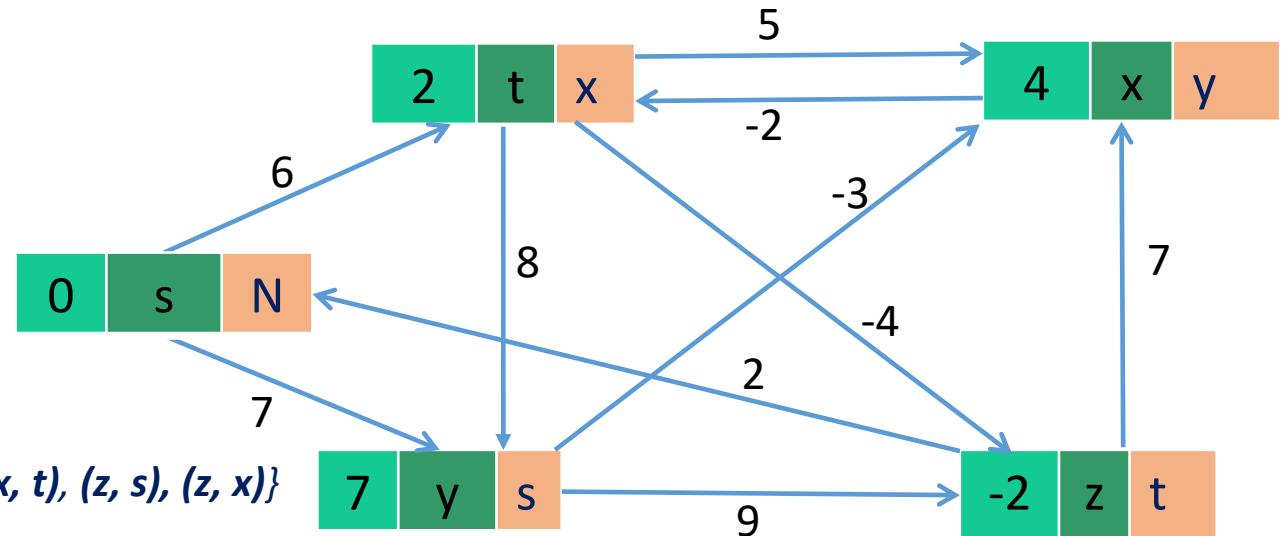
4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$



Step 5 : # of edges = { ( $s, t$ ), ( $s, y$ ), ( $t, x$ ), ( $t, y$ ), ( $t, z$ ), ( $y, x$ ), ( $y, z$ ), ( $x, t$ ), ( $z, s$ ), ( $z, x$ ) }

Step 6-7: RELAX ( $s, t, w$ ); RELAX ( $s, y, w$ ); RELAX ( $t, x, w$ ); RELAX ( $t, y, w$ ); RELAX ( $t, z, w$ ); RELAX ( $y, x, w$ ); RELAX ( $y, z, w$ ); RELAX ( $x, t, w$ );

$2 > 0 + 6;$

$7 > 0 + 7;$

$4 > 2 + 5;$

$7 > 2 + 8;$

$-2 > 2 + -4;$

$4 > 7 + -3;$

$-2 > 7 + 9;$

$2 > 4 + -2$

RELAX ( $z, s, w$ ); RELAX ( $z, x, w$ );

$0 > -2 + 2;$

$4 > -2 + 7;$

Step 8: return

# SSSP: Bellman-Ford (14)

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3   for each edge  $(u, v) \in G.E$

4       RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6   if  $v.d > u.d + w(u, v)$

7       return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2    $v.d = \infty$

3    $v.\pi = \text{NIL}$

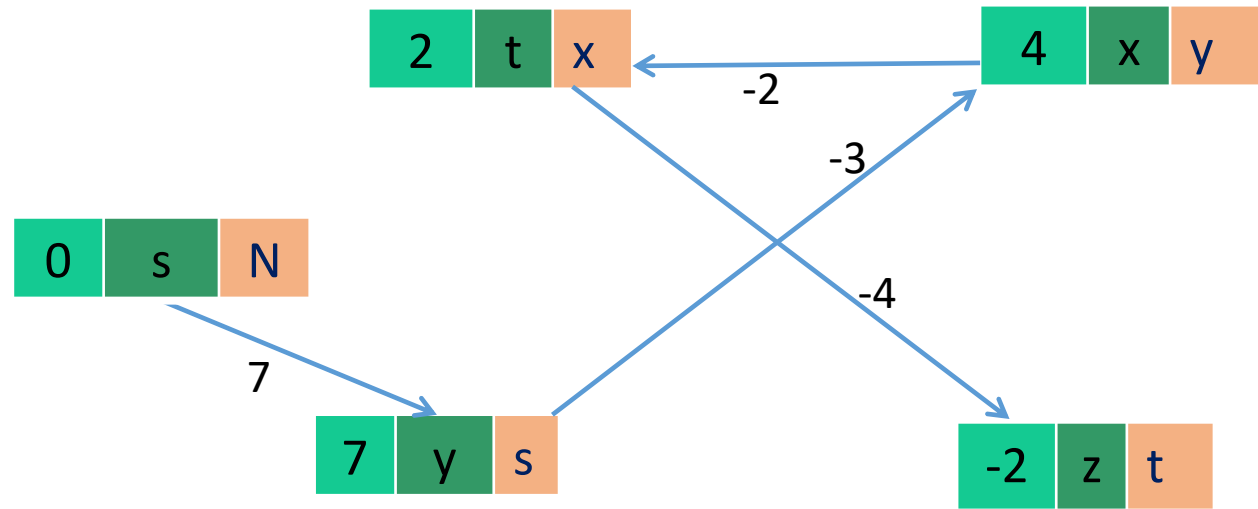
4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

2    $v.d = u.d + w(u, v)$

3    $v.\pi = u$



# SSSP: Bellman-Ford $\rightarrow$ Time complexity analysis

*BELLMAN-FORD* ( $G, w, s$ ) {

1 INITIALIZE-SINGLE-SOURCE ( $G, s$ )

2 for  $i = 1$  to  $|G.V| - 1$

3     for each edge  $(u, v) \in G.E$

4         RELAX ( $u, v, w$ )

5 for each edge  $(u, v) \in G.E$

6     if  $v.d > u.d + w(u, v)$

7         return False

8 return True

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

$\leftarrow$  # of vertices  $\rightarrow O(V)$

$\leftarrow$  Outer loop  $\rightarrow O(V)$ ; Inner loop  $\rightarrow O(E)$   
 $\rightarrow O(VE)$

$\leftarrow O(E)$

**Total time complexity:  $O(V) + O(VE) + O(E)$**

RELAX ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

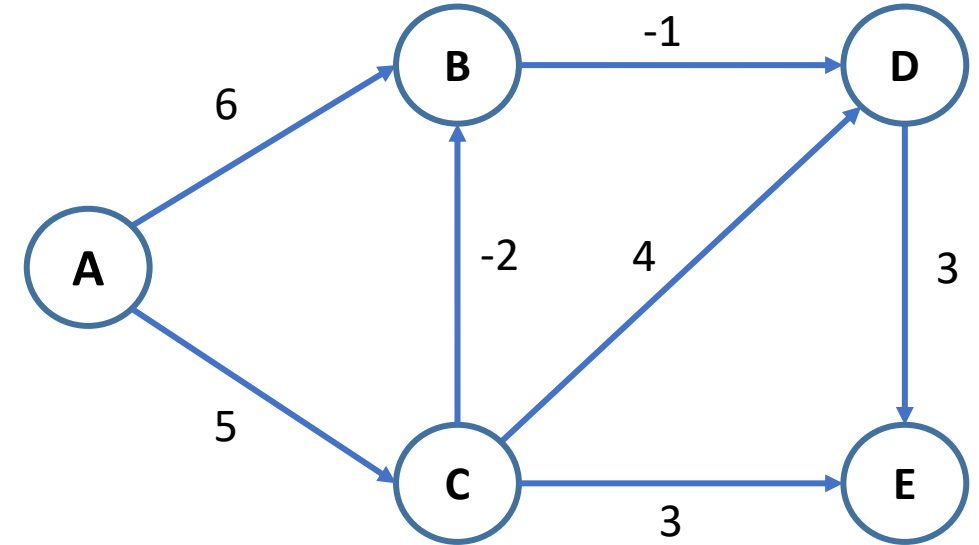
2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

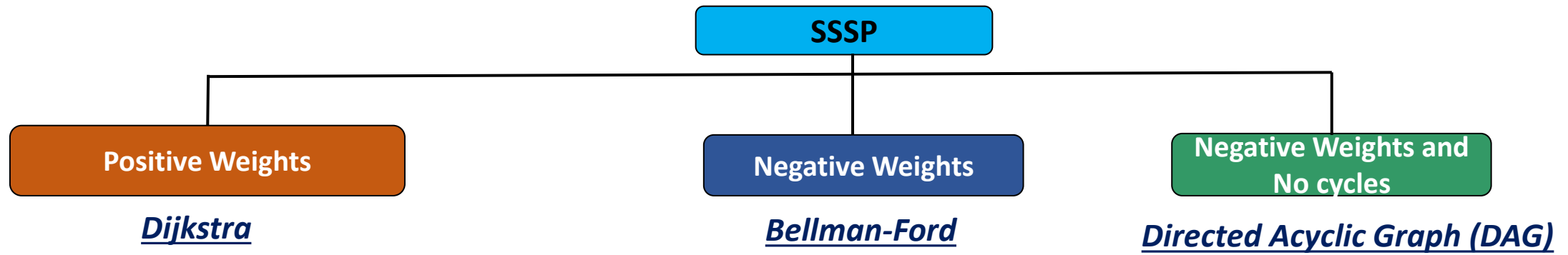
# Exercise: Bellman-Ford

Consider the following digraph starting at vertex A and apply Bellman-Ford single source shortest path algorithm on it.

- What is the minimum distance between vertex A and E
  - 3
  - 2
  - 1
  - 5



# Recap: SSSP



```
DIJKSTRA (G, w, s) {  
1 INITIALIZE-SINGLE-SOURCE (G, s)  
2 S =  $\emptyset$   
3 Q = G.V  
4 while Q  $\neq \emptyset$  ;  
5   u = EXTRACT-MIN(Q)  
6   S = S  $\cup$  {u}  
7   for each vertex v  $\in$  Q. Adj[u]  
8     RELAX (u, v, w)  
}
```

```
INITIALIZE-SINGLE-SOURCE (G, s) {  
1 for each v  $\in$  G.V  
2   v. d =  $\infty$   
3   v.  $\pi$  = NIL  
4 s. d = 0  
  
RELAX (u, v, w) {  
1 if v.d > u.d + w(u, v)  
2   v. d = u.d + w(u, v)  
3   v.  $\pi$  = u  
}
```

```
BELLMAN-FORD (G, w, s) {  
1 INITIALIZE-SINGLE-SOURCE (G, s)  
2 for i = 1 to |G.V| - 1  
3   for each edge (u, v)  $\in$  G.E  
4     RELAX (u, v, w)  
5 for each edge (u, v)  $\in$  G.E  
6   if v. d > u. d + w(u, v)  
7     return False  
8 return True
```

Total time complexity:  $O(V) + O(V \log V) + O(V) + O(E \log V)$   
:  $O(E \log V)$

Total time complexity:  $O(V) + O(VE) + O(E)$   
:  $O(VE)$

# SSSP: Directed Acyclic Graph (DAG)

- ❖ Directed Graph
- ❖ No cycles
- ❖ Topological sort can be applied on any DAG  $\rightarrow O(V+E)$
- ❖ Works for negative edges

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

1 Topological sort the vertices of  $G$

2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )

3 for each vertex  $u$ , taken in topologically sorted order

4     for each edge  $v \in G.Adj[u]$

5         *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

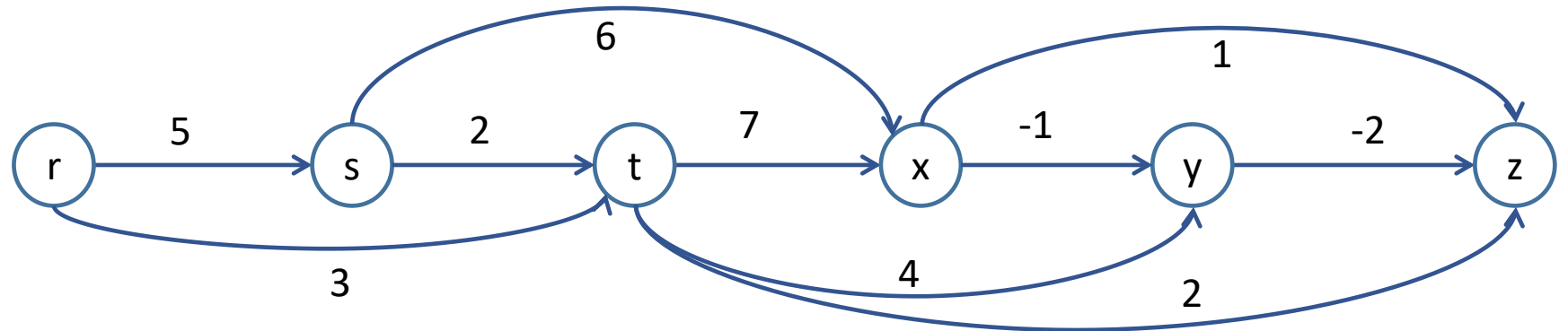
4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$





# SSSP: DAG (1)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

1 Topological sort the vertices of  $G$

2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )

3 **for** each vertex  $u$ , taken in topologically sorted order

4     **for** each edge  $v \in G.Adj[u]$

5         *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

1 **for** each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

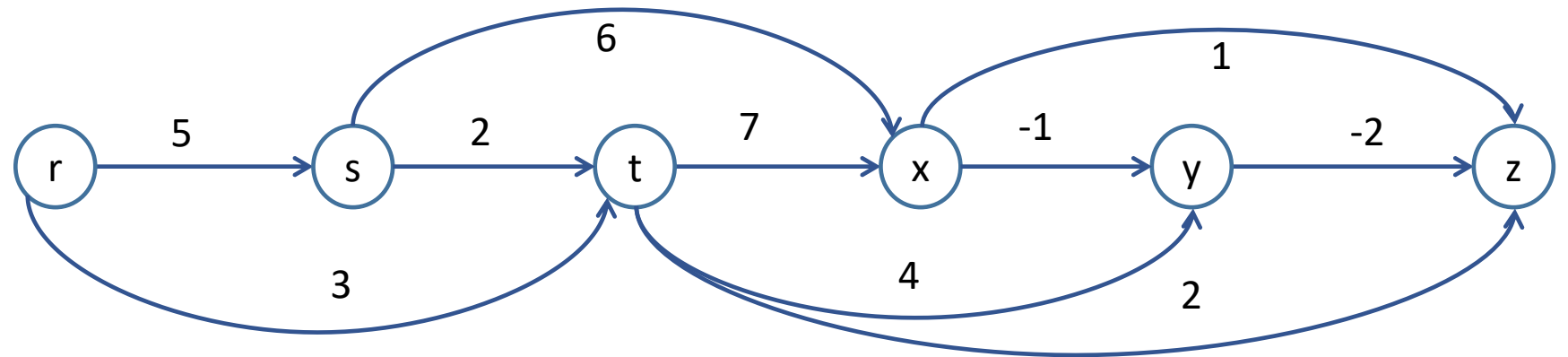
*RELAX* ( $u, v, w$ ) {

1     **if**  $v.d > u.d + w(u, v)$

2          $v.d = u.d + w(u, v)$

3          $v.\pi = u$

*Step 1: Topological order*



# SSSP: DAG (1)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

1 Topological sort the vertices of  $G$

2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )

3 for each vertex  $u$ , taken in topologically sorted order

4   for each edge  $v \in G.Adj[u]$

5       *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

1 for each  $v \in G.V$

2    $v.d = \infty$

3    $v.\pi = \text{NIL}$

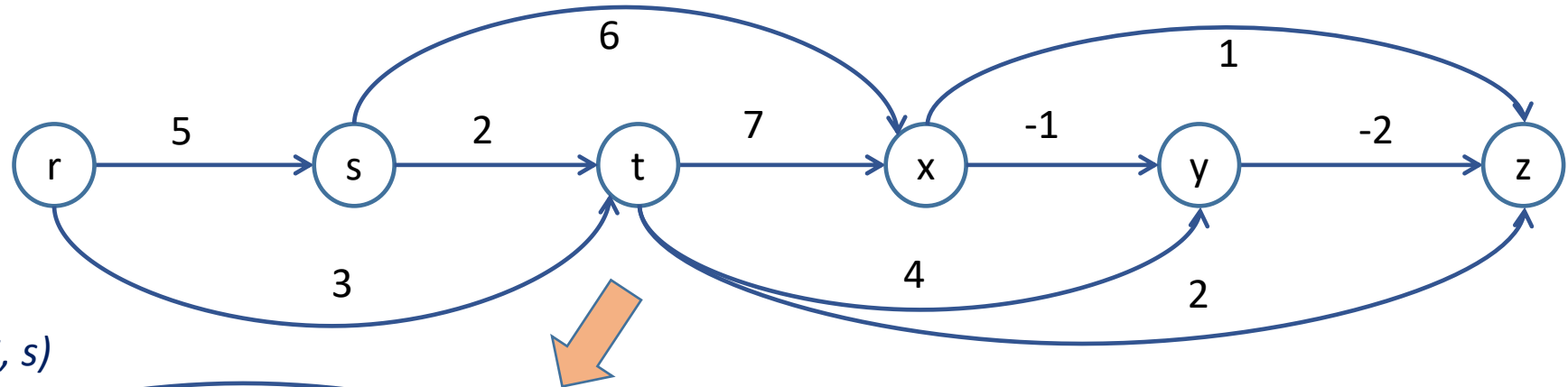
4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

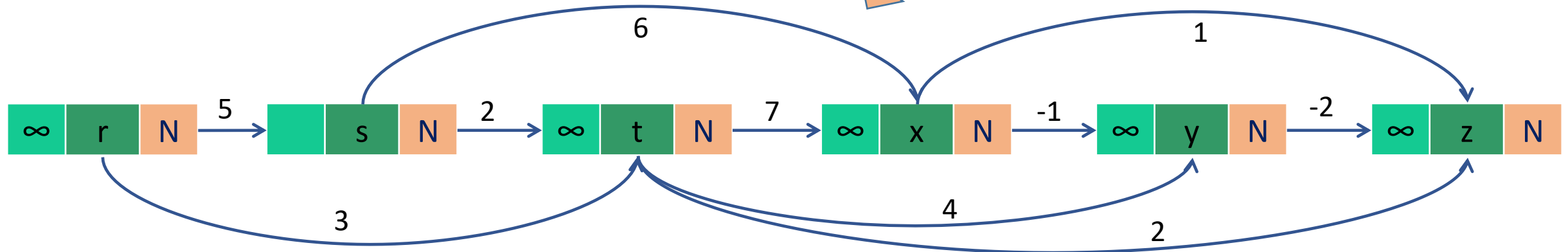
2    $v.d = u.d + w(u, v)$

3    $v.\pi = u$



Step 1: Topological order

Step 2: *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )



# SSSP: DAG (2)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

- 1 Topological sort the vertices of  $G$
- 2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )
- 3 for each vertex  $u$ , taken in topologically sorted order
- 4   for each edge  $v \in G. \text{Adj}[u]$
- 5     *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

- 1 for each  $v \in G.V$
- 2    $v.d = \infty$
- 3    $v.\pi = \text{NIL}$
- 4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

- 1 if  $v.d > u.d + w(u, v)$
- 2    $v.d = u.d + w(u, v)$
- 3    $v.\pi = u$

*Step 1: Topological order*

*Step 2: INITIALIZE-SINGLE-SOURCE* ( $G, s$ )

*Step 3:  $u = \{r, s, t, x, y, z\}$*

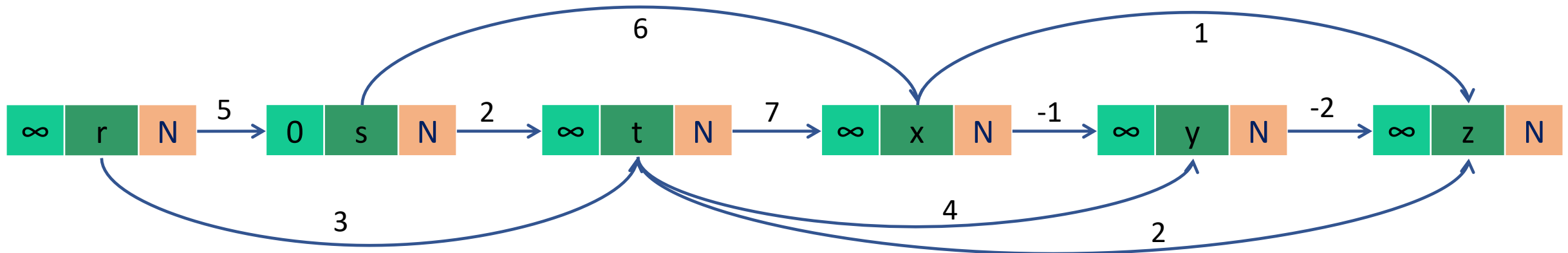
$u = r$

*Step 5: RELAX* ( $r, s, w$ ); *RELAX* ( $r, t, w$ )

$0 > \infty + 5;$      $\infty > \infty + 3;$

*Step 4:  $v = G. \text{Adj}[r]$*

$= \{s, t\}$



# SSSP: DAG (3)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

- 1 Topological sort the vertices of  $G$
- 2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )
- 3 **for** each vertex  $u$ , taken in topologically sorted order
- 4     **for** each edge  $v \in G. \text{Adj}[u]$
- 5         *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

- 1 **for** each  $v \in G.V$
- 2      $v.d = \infty$
- 3      $v.\pi = \text{NIL}$
- 4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

- 1     **if**  $v.d > u.d + w(u, v)$
- 2          $v.d = u.d + w(u, v)$
- 3          $v.\pi = u$

Step 3:  $u = \{r, s, t, x, y, z\}$

$u = s$

Step 4:  $v = G. \text{Adj}[s]$

$= \{t, x\}$

Step 5: *RELAX* ( $s, t, w$ ); *RELAX* ( $s, x, w$ )

$\infty > 0 + 2;$

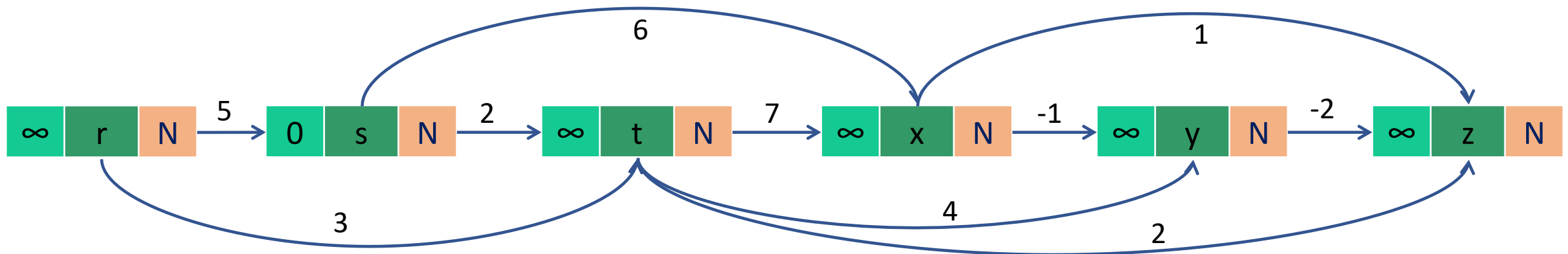
$t.d = 2$

$t.\pi = s$

$\infty > 0 + 6;$

$x.d = 6$

$x.\pi = s$



# SSSP: DAG (4)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

- 1 Topological sort the vertices of  $G$
- 2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )
- 3 for each vertex  $u$ , taken in topologically sorted order
- 4   for each edge  $v \in G. \text{Adj}[u]$
- 5     *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

- 1 for each  $v \in G.V$
- 2    $v.d = \infty$
- 3    $v.\pi = \text{NIL}$
- 4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

- 1 if  $v.d > u.d + w(u, v)$
- 2    $v.d = u.d + w(u, v)$
- 3    $v.\pi = u$

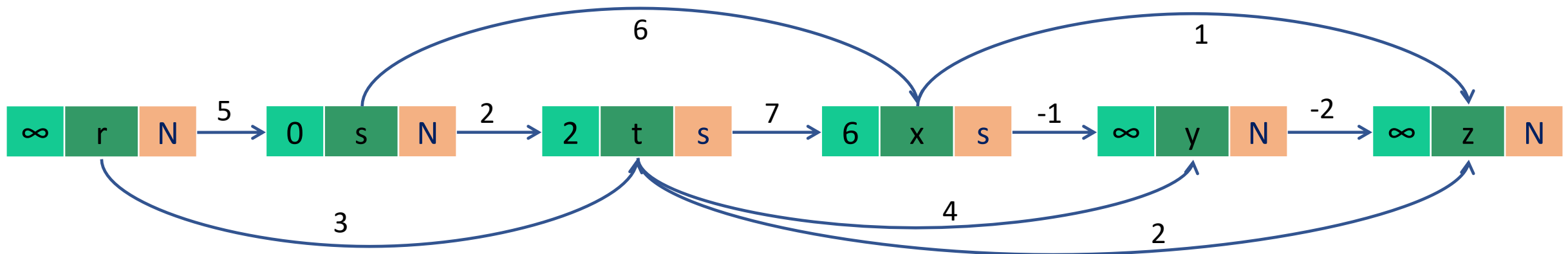
Step 3:  $u = \{r, s, t, x, y, z\}$   
 $u = t$

Step 4:  $v = G. \text{Adj}[t]$   
 $= \{x, y, z\}$

Step 5: *RELAX* ( $t, x, w$ );  
 $6 > 2 + 7$

*RELAX* ( $t, y, w$ );  
 $\infty > 2 + 4$   
 $y.d = 6$   
 $y.\pi = t$

*RELAX* ( $t, z, w$ );  
 $\infty > 2 + 2$   
 $z.d = 4$   
 $z.\pi = t$



# SSSP: DAG (5)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

- 1 Topological sort the vertices of  $G$
- 2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )
- 3 **for** each vertex  $u$ , taken in topologically sorted order
- 4     **for** each edge  $v \in G. \text{Adj}[u]$
- 5         *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

- 1 **for** each  $v \in G.V$
- 2      $v.d = \infty$
- 3      $v.\pi = \text{NIL}$
- 4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

- 1     **if**  $v.d > u.d + w(u, v)$
- 2          $v.d = u.d + w(u, v)$
- 3          $v.\pi = u$

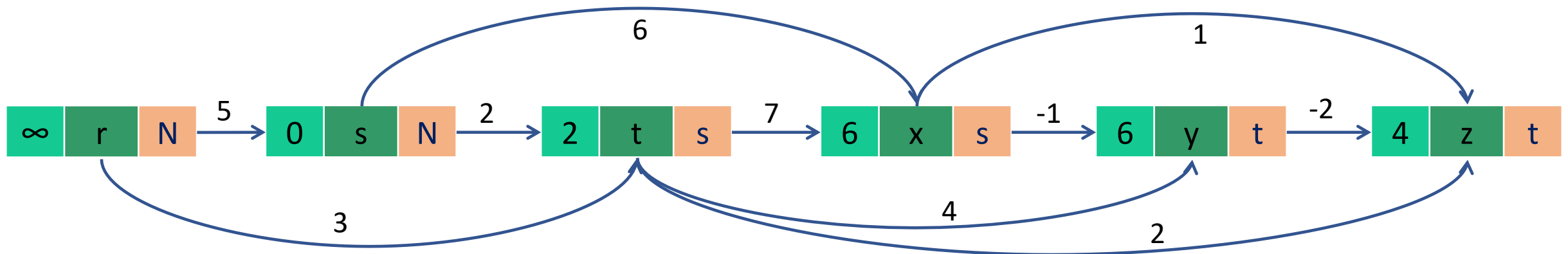
Step 3:  $u = \{r, s, t, x, y, z\}$   
 $u = x$

Step 4:  $v = G. \text{Adj}[x]$   
 $= \{y, z\}$

Step 5: *RELAX* ( $x, y, w$ ); *RELAX* ( $x, z, w$ )

$6 > 6 + -1$   
 $y.d = 5$   
 $y.\pi = x$

$4 > 6 + 1$



# SSSP: DAG (6)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

- 1 Topological sort the vertices of  $G$
- 2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )
- 3 for each vertex  $u$ , taken in topologically sorted order
- 4   for each edge  $v \in G. \text{Adj}[u]$
- 5     *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

- 1 for each  $v \in G.V$
- 2    $v.d = \infty$
- 3    $v.\pi = \text{NIL}$
- 4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

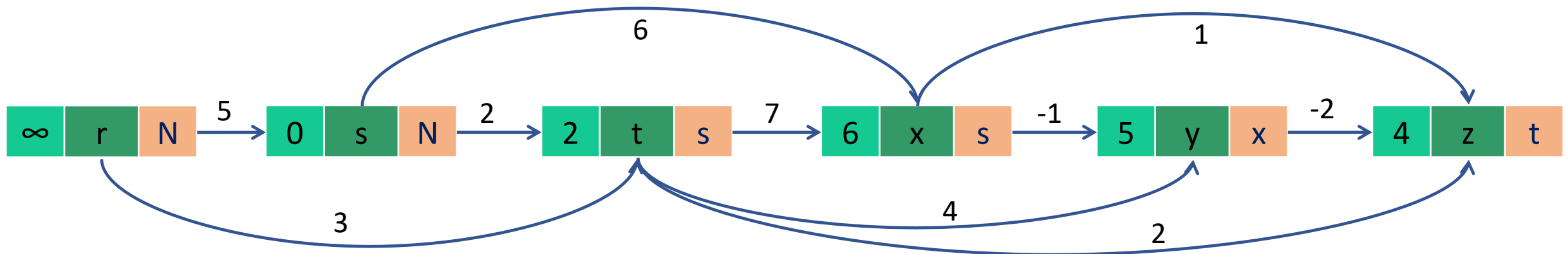
- 1 if  $v.d > u.d + w(u, v)$
- 2    $v.d = u.d + w(u, v)$
- 3    $v.\pi = u$

Step 3:  $u = \{r, s, t, x, y, z\}$   
 $u = y$

Step 4:  $v = G. \text{Adj}[y]$   
 $= \{z\}$

Step 5: *RELAX* ( $y, z, w$ );

$4 > 5 + -2$   
 $z.d = 3$   
 $z.\pi = y$



# SSSP: DAG (7)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

- 1 Topological sort the vertices of  $G$
- 2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )
- 3 for each vertex  $u$ , taken in topologically sorted order
- 4   for each edge  $v \in G. \text{Adj}[u]$
- 5     *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

- 1 for each  $v \in G.V$
- 2    $v.d = \infty$
- 3    $v.\pi = \text{NIL}$
- 4  $s.d = 0$

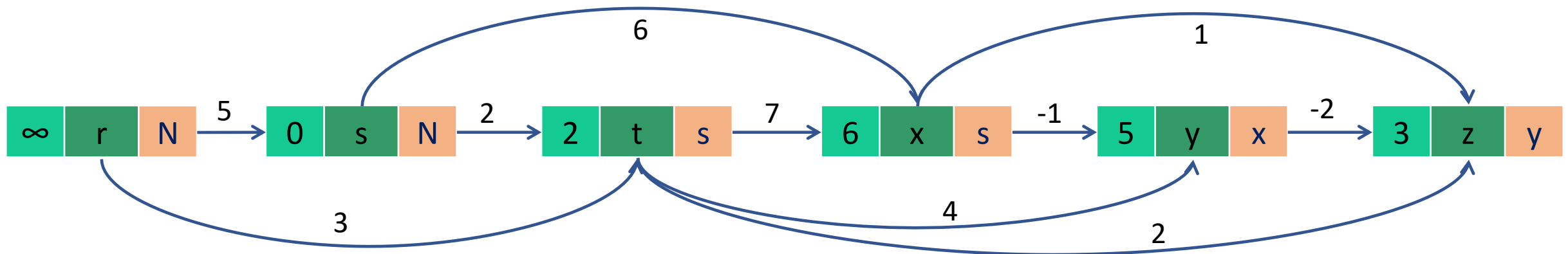
*RELAX* ( $u, v, w$ ) {

- 1 if  $v.d > u.d + w(u, v)$
- 2    $v.d = u.d + w(u, v)$
- 3    $v.\pi = u$

Step 3:  $u = \{r, s, t, x, y, z\}$

$u = z$

Step 4:  $v = G. \text{Adj}[z]$   
 $= \{\}$





# SSSP: DAG (8)

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

1 Topological sort the vertices of  $G$

2 *INITIALIZE-SINGLE-SOURCE* ( $G, s$ )

3 **for** each vertex  $u$ , taken in topologically sorted order

4     **for** each edge  $v \in G. Adj[u]$

5         *RELAX* ( $u, v, w$ )

*INITIALIZE-SINGLE-SOURCE* ( $G, s$ ) {

1 **for** each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

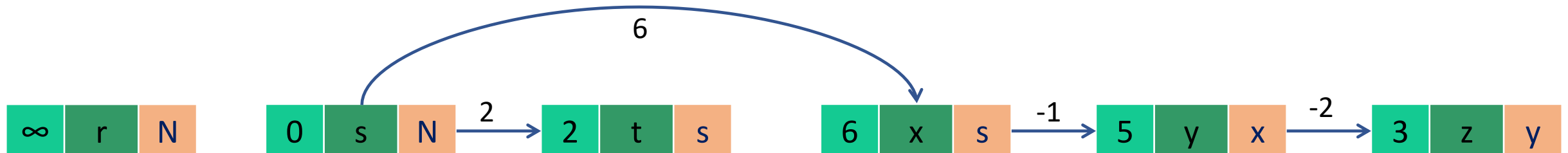
4  $s.d = 0$

*RELAX* ( $u, v, w$ ) {

1     **if**  $v.d > u.d + w(u, v)$

2          $v.d = u.d + w(u, v)$

3          $v.\pi = u$



# SSSP: DAG $\rightarrow$ Time Complexity Analysis

*DAG-SHORTEST-PATH* ( $G, w, s$ ) {

1 Topological sort the vertices of  $G$  }  $\leftarrow O(V + E)$

2 INITIALIZE-SINGLE-SOURCE ( $G, s$ ) }  $\leftarrow O(V)$

3 for each vertex  $u$ , taken in topologically sorted order

4     for each edge  $v \in G.Adj[u]$

5         RELAX ( $u, v, w$ )

$\leftarrow$  Outer loop  $\rightarrow O(V)$ ; Inner loop  $\rightarrow O(E)$

INITIALIZE-SINGLE-SOURCE ( $G, s$ ) {

1 for each  $v \in G.V$

2      $v.d = \infty$

3      $v.\pi = \text{NIL}$

4  $s.d = 0$

RELAX ( $u, v, w$ ) {

1 if  $v.d > u.d + w(u, v)$

2      $v.d = u.d + w(u, v)$

3      $v.\pi = u$

Total time complexity:  $O(V+E)$

# thank you!

email:

[k.kondepu@iitdh.ac.in](mailto:k.kondepu@iitdh.ac.in)