

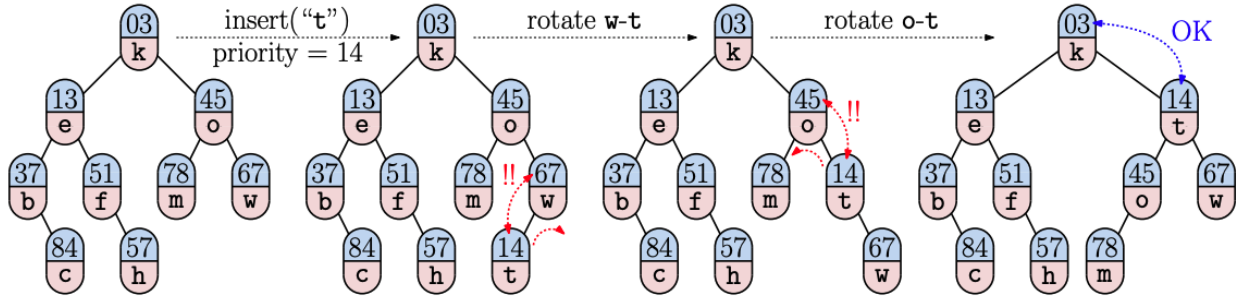
CS204L : Data Structures & Algorithms
Mid - Semester Examination

Select the problem based on the number assigned to you. [100 Marks]

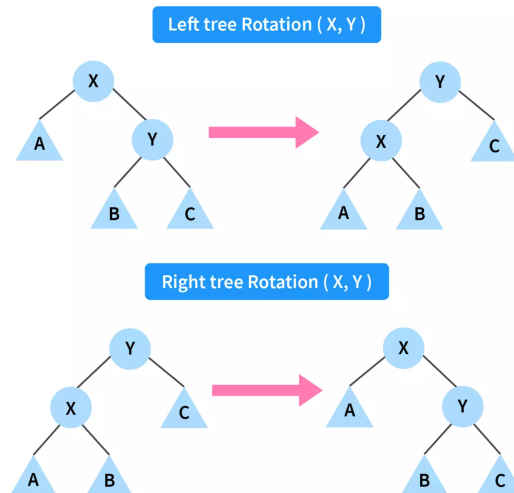
1. Implement a Binary Search Tree with each node has a pair (**key, priority**), where we maintain the BST property with respect to the *key*, i.e., for a node q , all its left children have *key* values less than the *key* value of the node and all the right children have their *key* values greater than or equal to the *key* value of the node, i.e., $x.key < q.key \leq y.key, x \in L_q, y \in R_q$. This being recursively true for all nodes implies this condition must be satisfied by all nodes of BST.

Also, for each node q , a MAX-HEAP property is maintained with respect to the *priority* values, i.e., the *priority* value of the node q is greatest among him and all its children's *priority* values, i.e., (**$q.left.priority < q.priority > q.right.priority$**) which is recursively true for all nodes, implies that this condition must be satisfied by all nodes of BST.

Implement: **insert(Node(key,priority))** where we insert the node into the tree. Do a normal BST insert. Then you need to use left rotation or right rotation to maintain heap property. Example:



Here the *key* values are in pink and *priority* values are in blue. Whenever the *priority* value property is broken, do a left or right rotation; **in the example, it follows a min-heap property but implements it as a max-heap property in the actual problem.** [Safely assume all *key* values and *priority* values are distinct]

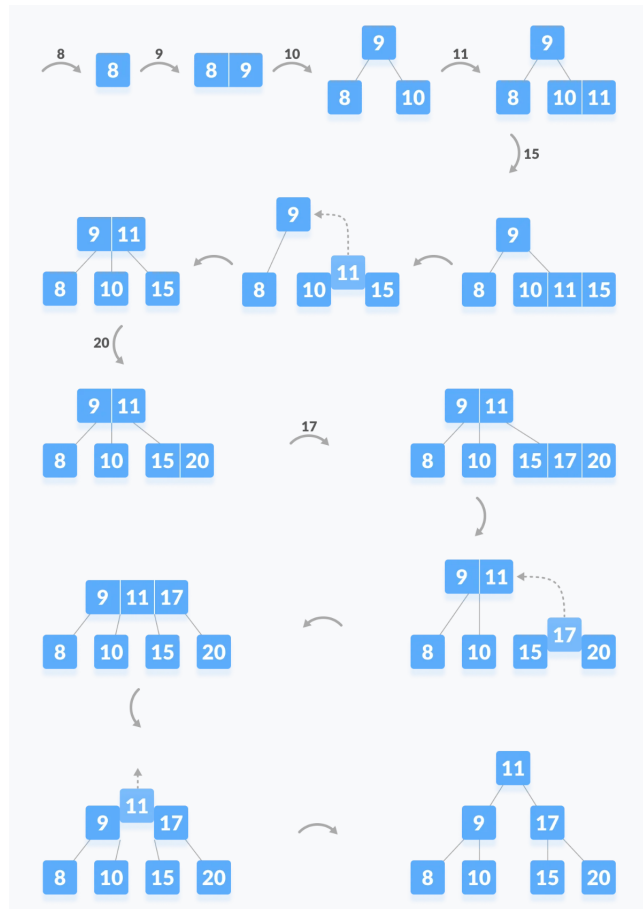


Implement: **extractMaxpriority(Node *root)** which is similar to a heap-extract-max except here you need to delete the node with max priority and rearrange the tree such that all other properties are satisfied.

Hint: Do a normal BST delete using the successor node. Now remove the successor node and do rotation on the node (downwards if required) where the successor values are installed.

2. Implement a generalised Search Tree that can have more than one key element for each node. So a generalised search tree of order “**m**” must satisfy:
 - (a) Every node has at most m children and $m - 1$ keys.
 - (b) Every item inside the node is kept in non decreasing order.

For example, if we need to insert the following keys: 8, 9, 10, 11, 15, 20, 17 in a tree of order, $m = 3$.



For selecting the element in a node that is to be shifted in the upper level you can use **median()**. [Assume all keys are distinct]

You can choose to solve the problem below if you find yourself unable to solve the problem above [60 marks]

Implement a K-ary Max Heap with insertion and Extract Max operations.