

Grokking the Coding Interview: Patterns for Coding Questions

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Pattern: Two Heaps

Pattern: Subsets

Pattern: Modified Binary Search

Pattern: Bitwise XOR

Pattern: Top 'K' Elements

- Introduction
- Top 'K' Numbers (easy)
- Kth Smallest Number (easy)
- 'K' Closest Points to the Origin (easy)
- Connect Ropes (easy)
- Top 'K' Frequent Numbers (medium)
- Frequency Sort (medium)
- Kth Largest Number in a Stream (medium)
- 'K' Closest Numbers (medium)
- Maximum Distinct Elements (medium)
- Sum of Elements (medium)
- Rearrange String (hard)
- Problem Challenge 1
- Solution Review: Problem Challenge 1
- Problem Challenge 2
- Solution Review: Problem Challenge 2
- Problem Challenge 3
- Solution Review: Problem Challenge 3**

Pattern: K-way merge

- Introduction
- Merge K Sorted Lists (medium)
- Kth Smallest Number in M Sorted Lists (Medium)
- Kth Smallest Number in a Sorted Matrix (Hard)
- Smallest Number Range (Hard)
- Problem Challenge 1
- Solution Review: Problem Challenge 1

Pattern : 0/1 Knapsack (Dynamic Programming)

- Introduction
- 0/1 Knapsack (medium)
- Equal Subset Sum Partition (medium)
- Subset Sum (medium)
- Minimum Subset Sum Difference

Solution Review: Problem Challenge 3

We'll cover the following

- Frequency Stack (hard)
- Solution
- Code
 - Time complexity
 - Space complexity

Frequency Stack (hard)

Design a class that simulates a Stack data structure, implementing the following two operations:

- `push(int num)`: Pushes the number 'num' on the stack.
- `pop()`: Returns the most frequent number in the stack. If there is a tie, return the number which was pushed later.

Example:

```
After following push operations: push(1), push(2), push(3), push(2), push(1), push(2), push(5)

1. pop() should return 2, as it is the most frequent number
2. Next pop() should return 1
3. Next pop() should return 2
```

Solution

This problem follows the [Top 'K' Elements](#) pattern, and shares similarities with [Top 'K' Frequent Numbers](#).

We can use a **Max Heap** to store the numbers. Instead of comparing the numbers we will compare their frequencies so that the root of the heap is always the most frequently occurring number. There are two issues that need to be resolved though:

- How can we keep track of the frequencies of numbers in the heap? When we are pushing a new number to the **Max Heap**, we don't know how many times the number has already appeared in the **Max Heap**. To resolve this, we will maintain a **HashMap** to store the current frequency of each number. Thus whenever we push a new number in the heap, we will increment its frequency in the **HashMap** and when we pop, we will decrement its frequency.
- If two numbers have the same frequency, we will need to return the number which was pushed later while popping. To resolve this, we need to attach a sequence number to every number to know which number came first.

In short, we will keep three things with every number that we push to the heap:

- ```
1. number // value of the number
2. frequency // current frequency of the number when it was pushed to the heap
3. sequenceNumber // a sequence number, to know what number came first
```

### Code

Here is what our algorithm will look like:

```
Java Python3 C++ JS

1 const Heap = require('./collections/heap'); //http://www.collectionsjs.com
2
3
4 class Element {
5 constructor(number, frequency, sequenceNumber) {
6 this.number = number;
7 this.frequency = frequency;
8 this.sequenceNumber = sequenceNumber;
9 }
10
11 compare(other) {
12 // higher frequency wins
13 if (this.frequency !== other.frequency) {
14 return this.frequency - other.frequency;
15 }
16 // if both elements have same frequency, return the element that was pushed later
17 return this.sequenceNumber - other.sequenceNumber;
18 }
19 }
20
21 class FrequencyStack {
22 constructor() {
23 this.sequenceNumber = 0;
24 this.frequencyMap = {};
25 this.maxHeap = new Heap([], null, ((a, b) => a.compare(b)));
26 }
27
28 }
```

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Create

Minimum Subset Sum Difference (hard)

Problem Challenge 1

Solution Review: Problem Challenge 1

Problem Challenge 2

Solution Review: Problem Challenge 2

Pattern: Topological Sort (Graph)

Introduction

Topological Sort (medium)

Tasks Scheduling (medium)

Tasks Scheduling Order (medium)

All Tasks Scheduling Orders (hard)

Alien Dictionary (hard)

Problem Challenge 1

Solution Review: Problem Challenge 1

Problem Challenge 2

Solution Review: Problem Challenge 2

Miscellaneous

Kth Smallest Number (hard)

Conclusions

Where to Go from Here

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Time complexity

The time complexity of `push()` and `pop()` is  $O(\log N)$  where 'N' is the current number of elements in the heap.

Space complexity

We will need  $O(N)$  space for the heap and the map, so the overall space complexity of the algorithm is  $O(N)$ .

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Problem Challenge 3Introduction

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