

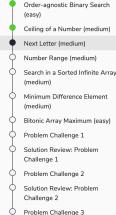


Solution Review: Problem

Pattern: Modified **Binary Search**

Introduction

Challenge 3



Pattern: Bitwise XOR

Introduction

Solution Review: Problem Challenge 3

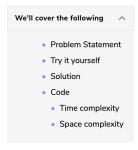


Elements



'K' Closest Numbers (medium)

Next Letter (medium)



Problem Statement

Given an array of lowercase letters sorted in ascending order, find the smallest letter in the given array greater than a given 'key'.

Assume the given array is a circular list, which means that the last letter is assumed to be connected with the first letter. This also means that the smallest letter in the given array is greater than the last letter of the array and is also the first letter of the array.

Write a function to return the next letter of the given 'key'.

Example 1:

```
Input: ['a', 'c', 'f', 'h'], key = 'f'
Output: 'h
Explanation: The smallest letter greater than 'f' is 'h' in the given array.
```

Example 2:

```
Explanation: The smallest letter greater than 'b' is 'c'.
```

Example 3:

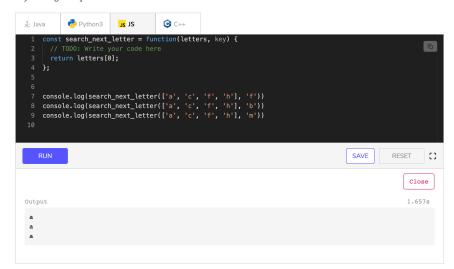
```
Input: ['a', 'c', 'f', 'h'], key = 'm'
Explanation: As the array is assumed to be circular, the smallest letter greater than 'm' i
```

Example 4:

```
Input: ['a', 'c', 'f', 'h'], key = 'h'
Explanation: As the array is assumed to be circular, the smallest letter greater than 'h' i
```

Try it yourself

Try solving this question here:



Solution

The problem follows the Binary Search pattern. Since Binary Search helps us find an element in a sorted array efficiently we can use a modified version of it to find the next letter



 \oplus

앟

Tracks

Ö

გ+

Refer a

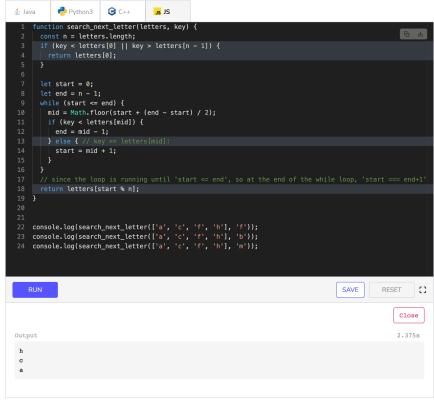
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We can use a similar approach as discussed in Ceiling of a Number. There are a couple of differences though:

- 1. The array is considered circular, which means if the 'key' is bigger than the last letter of the array or if it is smaller than the first letter of the array, the key's next letter will be the first letter of the array.
- 2. The other difference is that we have to find the next biggest letter which can't be equal to the 'key'. This means that we will ignore the case where key == arr[middle] . To handle this case, we can update our start range to start = middle +1.

In the end, instead of returning the element pointed out by start, we have to return the letter pointed out by start % array_length. This is needed because of point 2 discussed above. Imagine that the last letter of the array is equal to the 'key'. In that case, we have to return the first letter of the input array.

Here is what our algorithm will look like; the most important changes are in the highlighted lines:



Since, we are reducing the search range by half at every step, this means that the time complexity of our algorithm will be O(logN) where 'N' is the total elements in the given array.

Space complexity

The algorithm runs in constant space O(1).

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