New report

O. Denas

June 5, 2017

Contents

1	Current performance – not working	1
2	Input properties 2.1 Node properties	
3	wl() optimizations 3.1 sandbox tests 3.2 full tests	7
4	parent optimizations 4.1 sandbox tests 4.2 full tests	19 19 22
5	parallelization	22

1 Current performance – not working

2 Input properties

All experiments are performed on the following inputs (the mutation period refers to the index sequence):

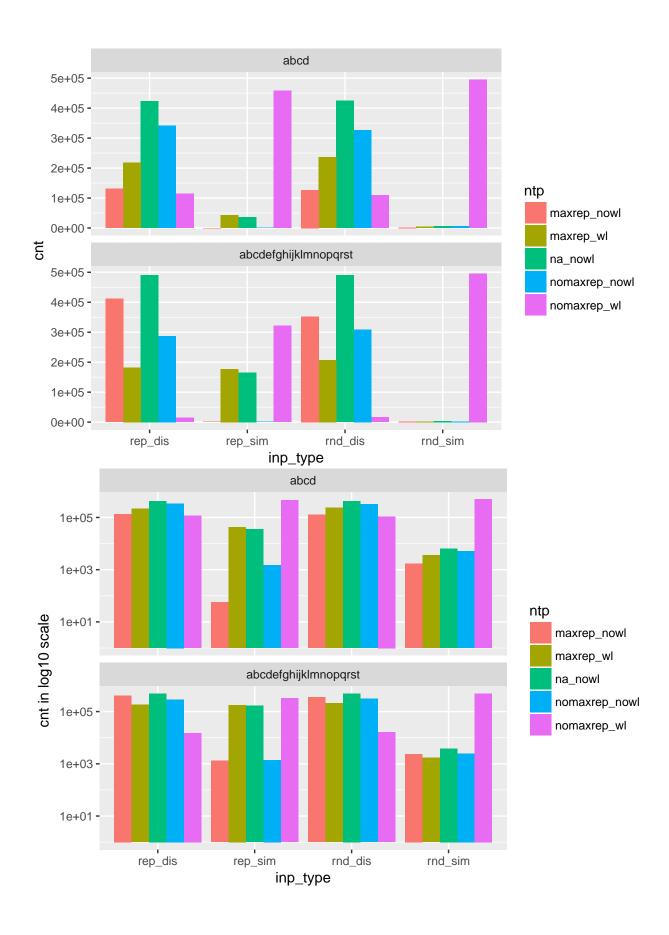
Table 1: Input datasets

s_type	t_type	slen	tlen	alp
rep	dis	100'000'000	500'000	abcdefghijklmnopqrst
rep	dis	100'000'000	500'000	abcd
rep	$_{ m sim}$	100'000'000	500'000	abcdefghijklmnopqrst
rep	\sin	100'000'000	500'000	abcd
rnd	dis	100'000'000	500'000	abcdefghijklmnopqrst
rnd	dis	100'000'000	500'000	abcd
rnd	\sin	100'000'000	500'000	abcdefghijklmnopqrst
rnd	$_{ m sim}$	100'000'000	500'000	abcd

In this section, we run the matching statistics algorithm and gather information on the way.

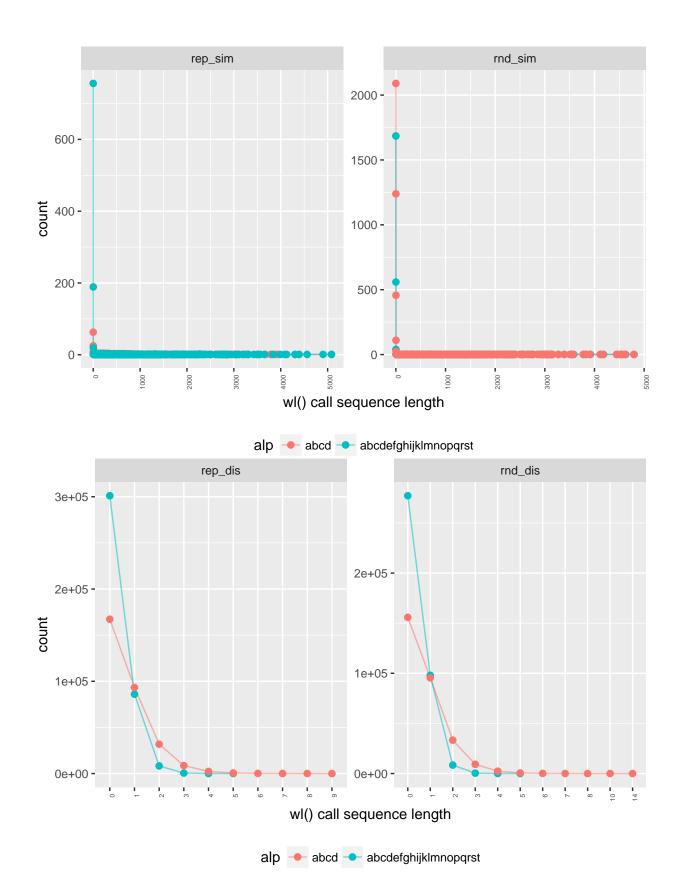
2.1 Node properties

Before making a $\mathtt{wl}(\mathtt{v}, \mathtt{c})$ call register whether \mathtt{v} is a maximal repeat, and whether the Weiner link exists or not.



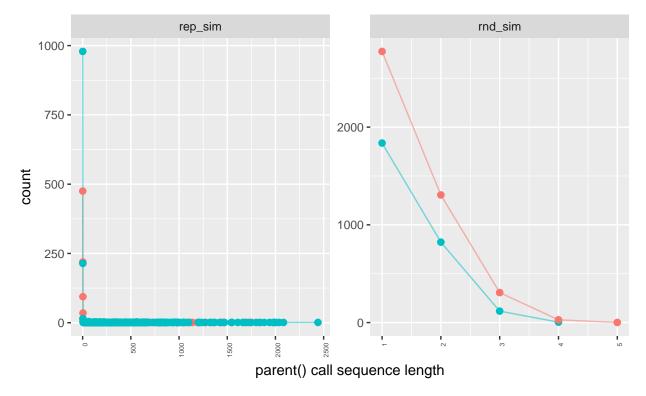
2.2 Consecutive wl() calls

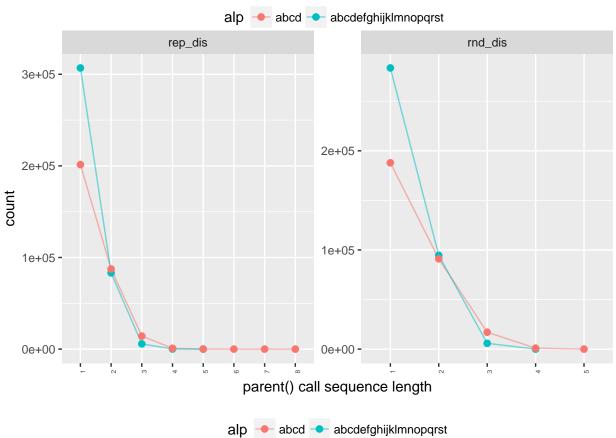
During the MS vector construction, count the number of consecutive wl() calls due to matches between reversed indexed string and the query. In other words count the k-length iterations of the while cycle.



2.3 Consecutive parent() calls

During the MS vector construction, count the number of consecutive parent() calls after a failed wl() call.





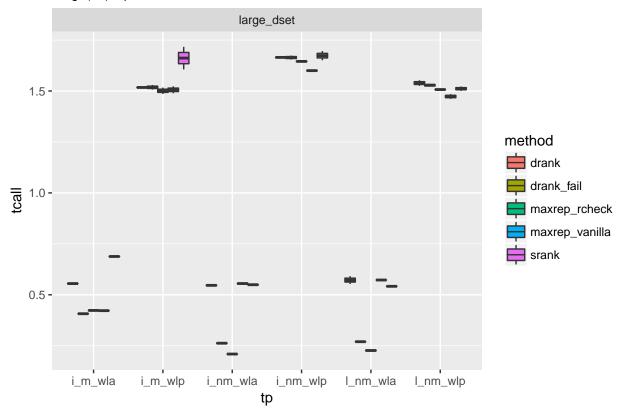
3 wl() optimizations

3.1 sandbox tests

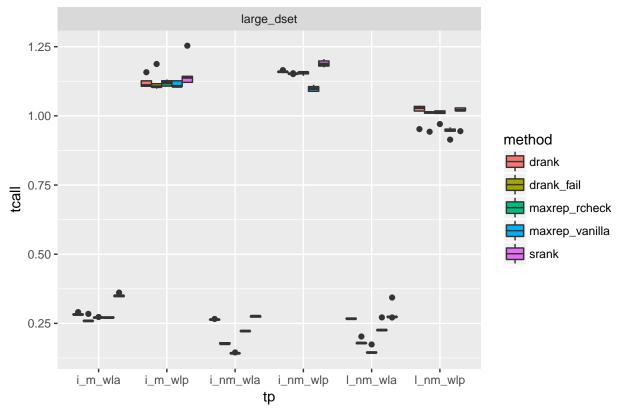
The 4 plots correspond to combinations of index size (one 100Mb and one 1Mb) and alphabet size (one 4 and one 20). Each plot shows 6 measurements based on

- node type: leaf (l) or internal (i)
- maximality: maximal (m) or non-maximal (nm)
- WL presence: present (wlp) or absent (wla)

Large(20) alphabet



Small(4) alphabet



3.2 full tests

I run the program several times with and without the optimization. The pointrange plots report (median, with quartile ranges) the relative difference of each optimized time from the average non-optimized time in the construction time of the MS vector which is

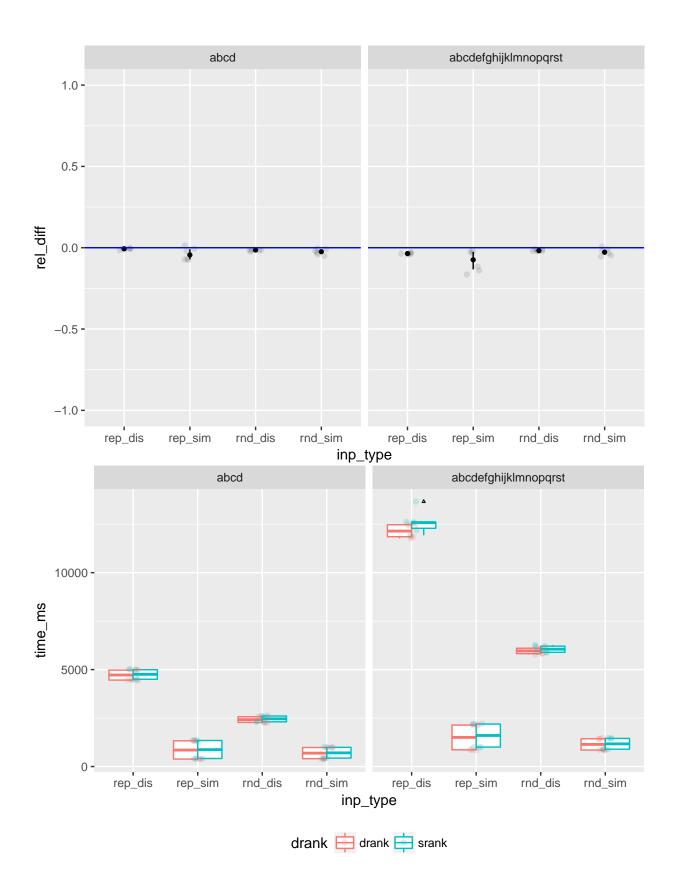
$$d^{(i)} = \frac{t_{\text{opt}}^{(i)} - \bar{t}_{\text{non_opt}}}{\max\{t_{\text{opt}}^{(i)}, \bar{t}_{\text{non_opt}}\}}$$

with $\bar{t}_{\mathtt{non_opt}} = 1/n \sum t_{\mathtt{non_opt}}^{(i)}$. Hence negative values indicate a speedup by the optimization over the average non-optimized time: -0.5 is a 2x speedup etc, in general the speed up is $-1/d^{(i)}$.

The boxplots report the raw times for the MS construction.

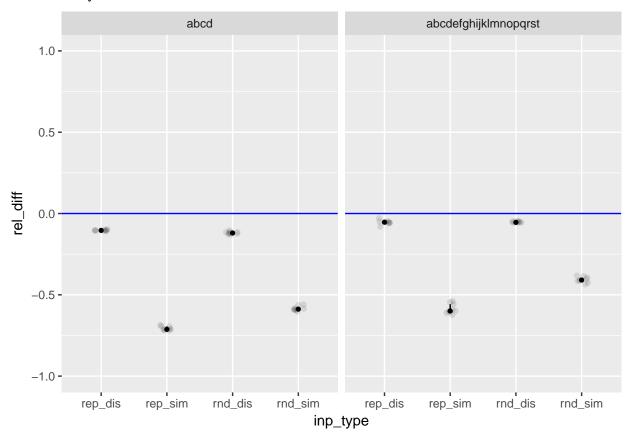
3.2.1 drank

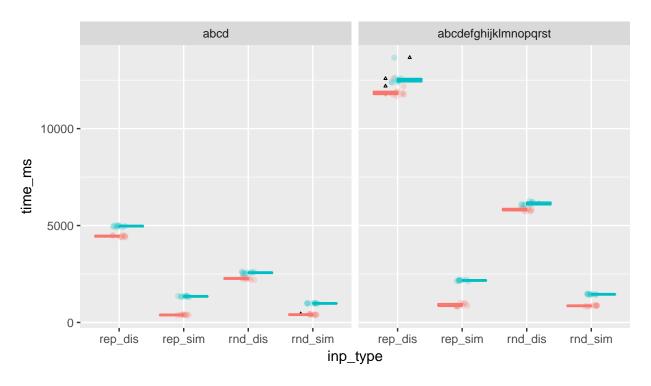
A tibble: 2 x 2
drank a
<chr> <int>
1 drank 48
2 srank 48



3.2.2 Lazy

```
## # A tibble: 2 x 2
## lazy a
## <chr> <int>
## 1 lazy 72
## 2 nonlazy 72
```

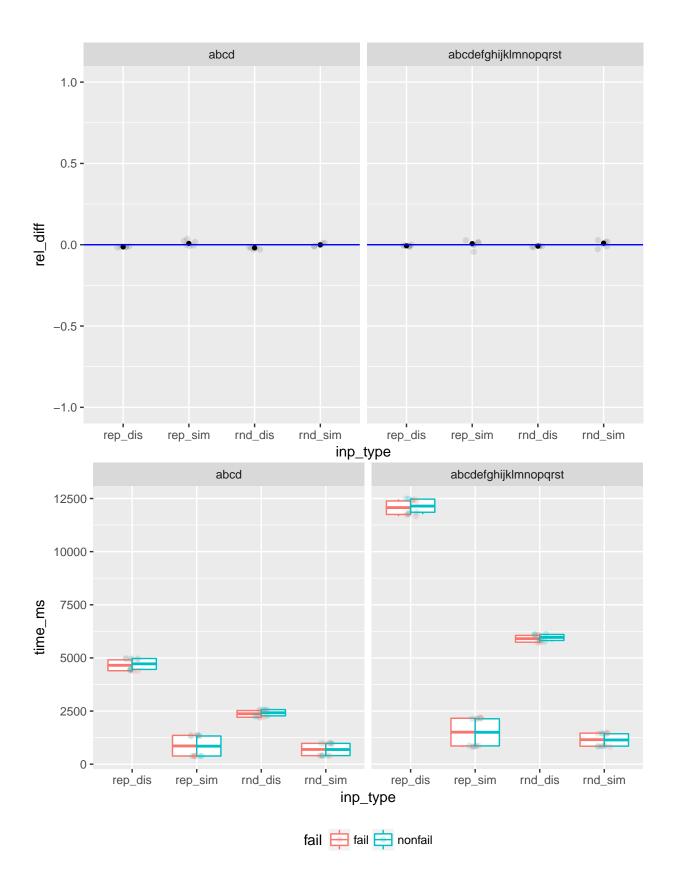




lazy 🖨 lazy 🖨 nonlazy

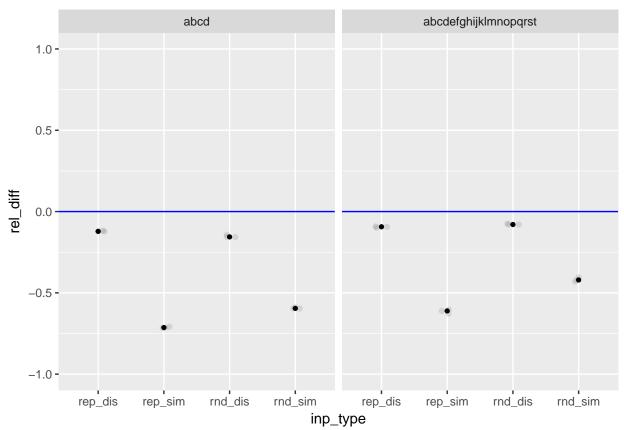
3.2.3 Fail

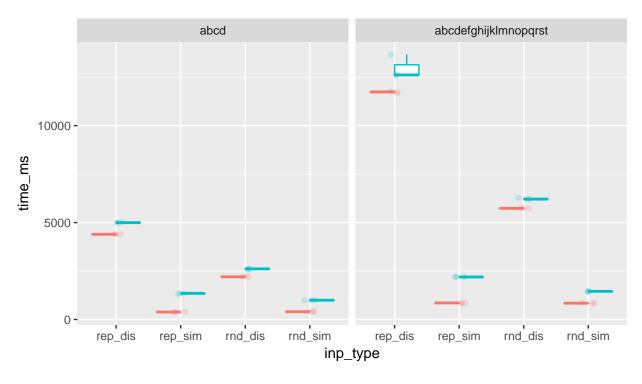
A tibble: 2 x 2
fail a
<chr> <int>
1 fail 48
2 nonfail 48



3.2.4 all flags (lazy, fail, double_rank) vs. single rank

A tibble: 2 x 2
drank a
<chr> <int>
1 drank 24
2 srank 24

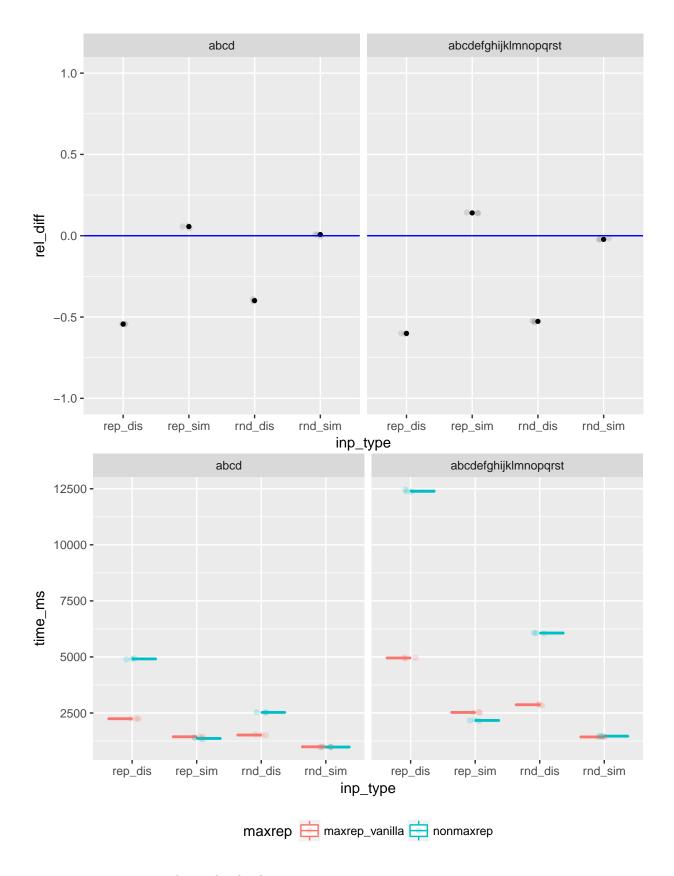




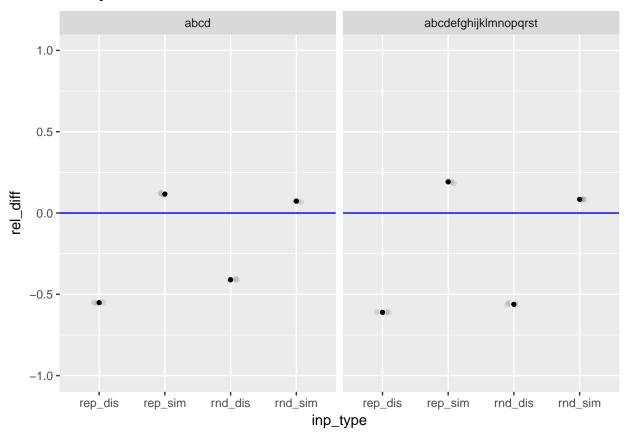
fail ⊨ fail 🖶 nonfail

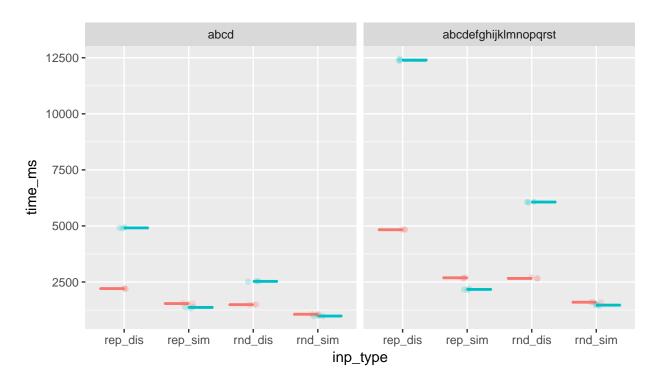
3.2.5 maxrep

3.2.5.1 maxrep_vanilla vs. non-maxrep



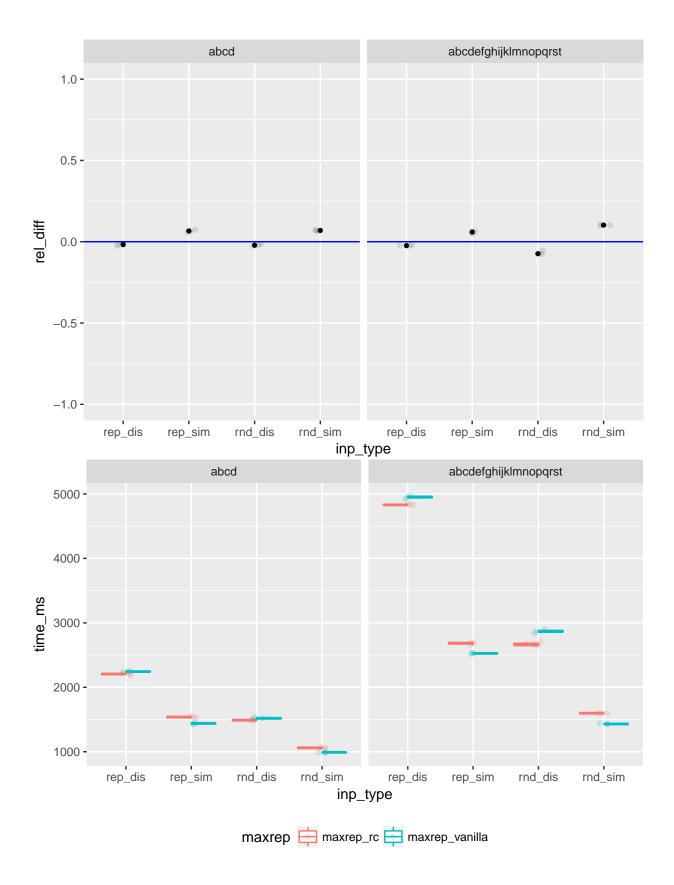
3.2.5.2 maxrep+rank_and_check vs. non-maxrep





maxrep \rightleftharpoons maxrep_rc \rightleftharpoons nonmaxrep

$3.2.5.3 \quad \text{maxrep+rank_and_check vs. maxrep_vanilla}$

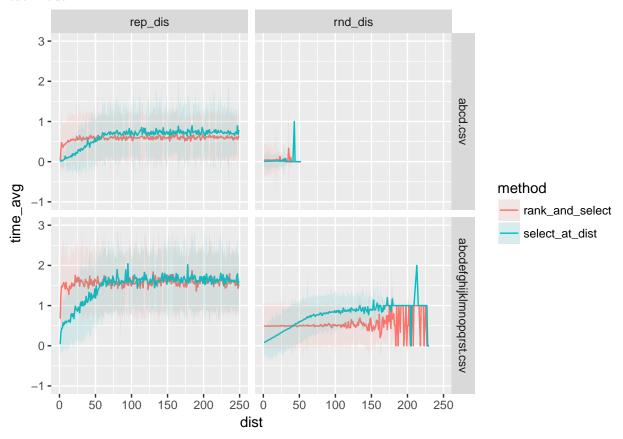


4 parent optimizations

4.1 sandbox tests

4.1.1 select at dist

Traverse the tree and time the call select_at_dist (labeled f) and the call select(rank()) (labeled s) on each node.



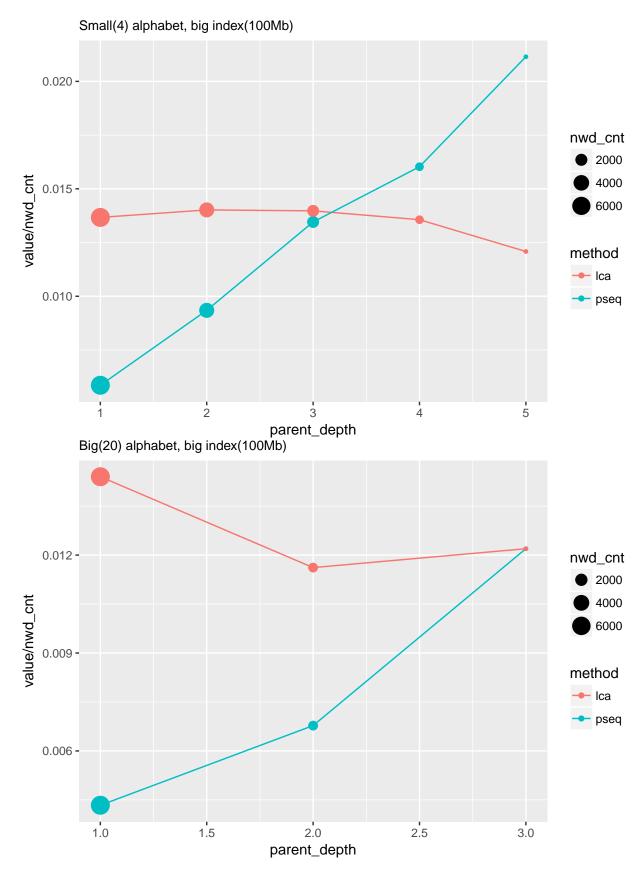
4.1.2 lowest maximal ancestor

Generate all tuples (v, c, d) from a tree where, n is an internal node, c a character, and d = depth(v) - depth(u) with u the lowest maximal ancestor.

Then shuffle and time the two ways of finding u: using the maxrep or with a sequence of parent calls.

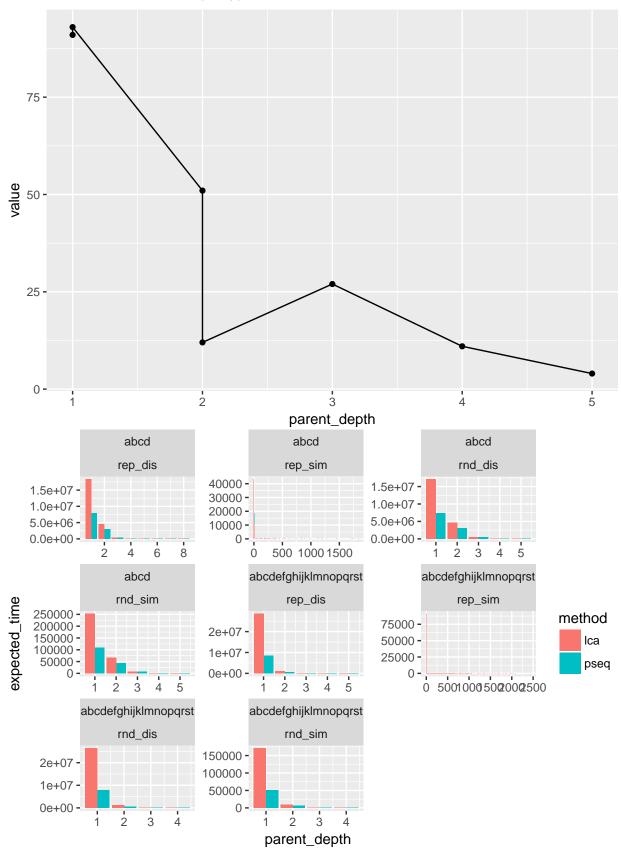
The plots show the total time / nr of calls for a given depth with the size of the point representing the nr of calls for a given depth value.

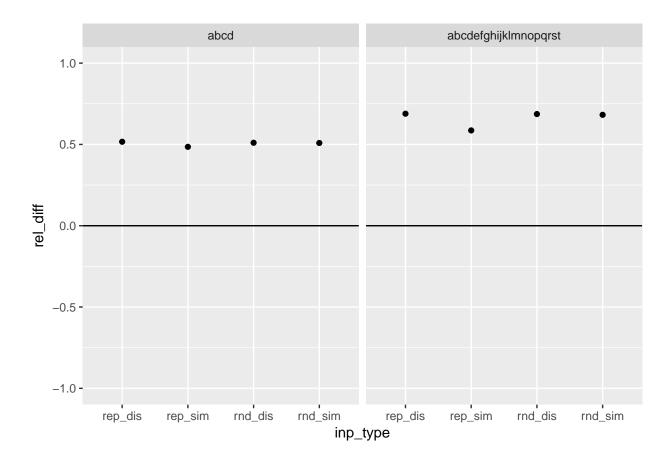
When the nr of calls is small (this happens for large depth lengths) the time resolution is not big enough and the numbers are not to be trusted – hence I have removed the data from the plots.



By counting the number of consecutive parent calls (for each number of consecutive calls) we project the

runtime of each method for an input type as below.





4.2 full tests

5 parallelization

Run the program on 1, 2, 4, 8 and 16 threads and measure the time it takes to build the RUNS and MS vectors.

TODO: there is currently a bug on the parallel version of the program for particular inputs. At the moment fixing it is not a priority.