University of Jyväskylä - Course TIEJ6003 INTRODUCTION TO QUANTUM COMPUTING

Day-02

Prof. Ofer Shir oshir@alumni.Princeton.EDU



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BUILDING BLOCKS OF COMPUTATION: QUBITS AND Q-GATES

$$|0\rangle := \begin{pmatrix} 1 \\ 0 \end{pmatrix} \qquad |1\rangle := \begin{pmatrix} 0 \\ 1 \end{pmatrix}.$$

the single qubit

The lowest computational unit is a single spin- $\frac{1}{2}$ particle whose Hamiltonian is under control — and is called a **qubit**.

It is a complex superposition of the 'up' and 'down' eigenkets (denoted as either $\{|\uparrow\rangle,|\downarrow\rangle\}$, $\{|+\rangle,|-\rangle\}$, or $\{|0\rangle,|1\rangle\}$):

$$|\psi\rangle := \alpha |\uparrow\rangle + \beta |\downarrow\rangle = \alpha |0\rangle + \beta |1\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix},$$
 (1)

with $|\alpha|^2 + |\beta|^2 = 1$.

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producing any state?

One can produce any state by setting the SGE apparatus to a certain rotation angle. For a rotation over an angle θ about the \hat{z} -axis, the rotated spin operator becomes (using the fact that Pauli's matrices are unitary, $\sigma_z^2 = 1$)

$$S_{\theta|\hat{z}} = \exp\left(i\theta \cdot \frac{\sigma_z}{2}\right) = \cos\left(\frac{\theta}{2}\right)\mathbf{1} + i\sin\left(\frac{\theta}{2}\right)\sigma_z.$$
 (2)

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generalization

Let the vector $\hat{r} := (\cos \varphi \cdot \sin \theta, \cos \theta \cdot \sin \theta, \cos \theta)^T$ represent a target direction in the 3D space of the SGE, then the following operator, which possesses eigenvalues of $\frac{\hbar}{2}$,

$$S_{\hat{r}} := \begin{bmatrix} \cos \theta & \sin \theta \cdot \exp(-i\varphi) \\ \sin \theta \cdot \exp(i\varphi) & -\cos \theta \end{bmatrix}, \tag{3}$$

becomes the measurement of this generalized SGE in the \hat{r} -direction.

Given this setup, we are able to construct any desired qubit,

$$|\psi\rangle := \cos\left(\frac{\theta}{2}\right)|0\rangle + \exp\left(i\varphi\right)\sin\left(\frac{\theta}{2}\right)|1\rangle$$

up to an undetectable global phase.*

global phase and the Bloch vector

*A generic quantum state is described by

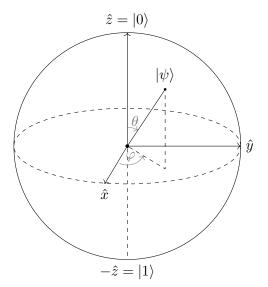
$$|\psi\rangle := \exp\left(i\gamma\right) \left[\cos\left(\frac{\theta}{2}\right)|0\rangle + \exp\left(i\varphi\right)\sin\left(\frac{\theta}{2}\right)|1\rangle\right],$$

with γ entitled the *global phase*, which is undetectable.

Conventionally, \hat{r} is called the Bloch vector, and the corresponding Bloch Sphere (next chart) provides a spatial description and a visualization of the qubit's state.

A so-called *pure* state always resides on the sphere's surface; *ensembles*, which are not discussed in this course, reside within.

the Bloch sphere



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quantum gates

Operators that are applied to qubits must be unitary:

$$\mathcal{U}\mathcal{U}^{\dagger}=\mathbf{1}.$$

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Why?

- Probabilities remain consistent across quantum operations.
- Quantum computations are reversible.

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A basic example is the so-called phase gate (Pauli's Z!):

$$\phi := \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}. \tag{4}$$

the Hadamard gate

The Hadamard gate allows to obtain an equal superposition between the states, which makes it especially useful in the starting point of quantum computations:

$$H := \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix} \tag{5}$$

$$\frac{\alpha|0\rangle + \beta|1\rangle}{\text{initial state}} \qquad \qquad \underline{H} \qquad \qquad \alpha\left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) + \beta\left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)$$

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other 1-qubit gates

The phase gate (denoted as S), and the T gate (a.k.a. the $\frac{\pi}{8}$, for historical reasons):

$$S := \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \qquad T := \begin{bmatrix} 1 & 0 \\ 0 & \exp(i\pi/4) \end{bmatrix} \tag{6}$$

Recall the Pauli operators:

$$\sigma_x = X := \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \qquad \sigma_y = Y := \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \qquad \sigma_z = Z := \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$
 (7)

JYU: intro2QC (COM3)

2 qubits

The Hilbert space of a 2-qubit system holds four states:

$$\{\left|00\right\rangle,\left|01\right\rangle,\left|10\right\rangle,\left|11\right\rangle\}$$

Notation may differ in the literature, e.g., when using commas to distinguish the qubits $(|0,0\rangle)$, or when aggregating kets $(|0\rangle|0\rangle)$ — but we will use herein $|00\rangle$.

Every state may be represented as

$$|\psi\rangle := \alpha_{00} |00\rangle + \alpha_{01} |01\rangle + \alpha_{10} |10\rangle + \alpha_{11} |11\rangle$$

subject to $\sum_{i,j=0}^{1} |\alpha_{ij}|^2 = 1$.

2-qubit representation

The explicit representation is scaled-up by applying a tensor multiplication:

$$|00\rangle := |0\rangle \otimes |0\rangle = \begin{pmatrix} 1 & \times & \begin{pmatrix} 1 \\ 0 \\ 0 & \times & \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}. \tag{8}$$

JYU: intro2QC (COM3)

2-qubit measurement

given a 2-qubit state, a measurement may target the pair (leading to a 2-qubit state collapse according to (P-5)), or alternatively, target only a single qubit, e.g., measuring the first qubit:

$$|\psi\rangle \xrightarrow{\text{1st-qubit}} \begin{cases} 0 & \text{with probability } |\alpha_{00}|^2 + |\alpha_{01}|^2 \\ 1 & \text{with probability } |\alpha_{10}|^2 + |\alpha_{11}|^2 \end{cases}$$
 (9)

and leading to a partial (1-qubit) state collapse: $|\psi_0\rangle = \frac{\alpha_{00}|0\rangle + \alpha_{01}|1\rangle}{\sqrt{|\alpha_{00}|^2 + |\alpha_{01}|^2}}$ if $|0\rangle$ is measured, or $|\psi_1\rangle = \frac{\alpha_{10}|0\rangle + \alpha_{11}|1\rangle}{\sqrt{|\alpha_{10}|^2 + |\alpha_{11}|^2}}$ if otherwise $|1\rangle$ is measured.

entanglement

Certain states, e.g., $|\psi\rangle = \frac{1}{\sqrt{2}} \left(|00\rangle + |11\rangle\right)$ clearly possess high correlation between the two qubits — and therefore, measuring only a single qubit provides information on the other qubit without leading to its collapse. Such states feature the so-called **entanglement** property, which is highly beneficial for computation and communication.

Bell states

The 4 states in which the 2 qubits are entangled:

$$|\phi_{+}\rangle = |\beta_{00}\rangle := \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

$$|\psi_{+}\rangle = |\beta_{01}\rangle := \frac{1}{\sqrt{2}} (|01\rangle + |10\rangle)$$

$$|\phi_{-}\rangle = |\beta_{10}\rangle := \frac{1}{\sqrt{2}} (|00\rangle - |11\rangle)$$

$$|\psi_{-}\rangle = |\beta_{11}\rangle := \frac{1}{\sqrt{2}} (|01\rangle - |10\rangle)$$

Notably, these states form a basis.

2-qubit gates

2-qubit gates are 4×4 unitary matrices, which may be scaled-up using a tensor product, as in the Hadamard case:

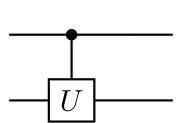
An application example:

$$H^{\otimes 2} |00\rangle = \frac{1}{2} (|00\rangle + |10\rangle + |01\rangle + |11\rangle)$$

controlled operations

"if A is **true** then do B" is a common type of controlled operation that is very useful in computation. In QC, a controlled-U operation is a 2-qubit gate, with a *control* qubit and a *target* qubit, and an arbitrary 1-qubit unitary operation U.

If the control qubit is true then U is applied to the target qubit, otherwise it is unchanged:



 $|c\rangle |t\rangle \rightarrow |c\rangle U^c |t\rangle$

controlled-NOT

The controlled-NOT operation is the most widely-used controlled gate featuring a XOR:

$$|c\rangle\,|t\rangle \to |c\rangle \oplus |t\rangle$$

It is called the CNOT gate, whose matrix representation and circuit are as follows:

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$|a\rangle \xrightarrow{\qquad} |a\rangle$$

$$|b\rangle \xrightarrow{\qquad} |b \oplus a\rangle$$

2-qubit entanglement: aftermath

The unitary gates operating on 2-qubit systems may either act on each qubit independently (and thus be represented as a tensor product of 1-qubit operators, as in the Hadamard case above), or act in a correlated manner (featuring entanglement, and thus having no tensor product form — as in the CNOT gate).

Obtaining entanglement:

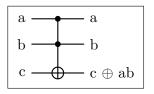
$$\begin{array}{c|c} \alpha \left| 0 \right\rangle + \beta \left| 1 \right\rangle & & \\ & \left| 0 \right\rangle & & \alpha \left| 00 \right\rangle + \beta \left| 11 \right\rangle \end{array}$$

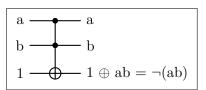
classical computational universality

Simulating classical circuits on a quantum computer is feasible using quantum gates.

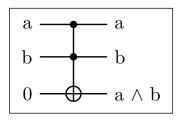
Obtaining such a computational universality is realizable simply by simulating the NAND logical gate.

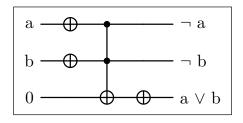
The **Toffoli** gate is a 3-qubit gate that flips the state of the third qubit if the first two qubits are both in state $|1\rangle$:





Toffoli as AND & OR





TODO: verify truth tables!

no-cloning theorem for qubits

The no-cloning theorem is a fundamental result in QM that states it is impossible to create an identical copy of an arbitrary unknown quantum state. Formally, the theorem can be stated as follows—

no-cloning theorem for qubits

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Theorem: There does not exist a unitary operation \mathcal{U} such that for any arbitrary quantum state $|\psi\rangle$ and a fixed state $|0\rangle$, the following transformation holds:

$$\mathcal{U}(|\psi\rangle|0\rangle) = |\psi\rangle|\psi\rangle$$

no-cloning theorem - proof

Proof (by contradiction):

Consider two arbitrary quantum states $|\psi\rangle$ and $|\phi\rangle$.

Assume the opposite that there exists a unitary cloning operator \mathcal{U} ,

$$\mathcal{U}\left(\left|\psi\right\rangle\left|0\right\rangle\right)=\left|\psi\right\rangle\left|\psi\right\rangle$$

$$\mathcal{U}(|\phi\rangle|0\rangle) = |\phi\rangle|\phi\rangle$$

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$$\mathcal{U}(|\psi\rangle|0\rangle) = |\psi\rangle|\psi\rangle$$

$$\mathcal{U}(|\phi\rangle|0\rangle) = |\phi\rangle|\phi\rangle$$

Since \mathcal{U} is a linear operator, it must preserve the inner product:

$$(\langle \psi | \langle 0 |) \mathcal{U}^{\dagger} \mathcal{U} (|\phi\rangle | 0 \rangle) = (\langle \psi | \langle \psi |) (|\phi\rangle | \phi \rangle)$$

Given that $\mathcal{U}^{\dagger}\mathcal{U} = \mathbf{1}$, we obtain the following result:

$$\langle \psi | \phi \rangle = \langle \psi | \phi \rangle \langle \psi | \phi \rangle$$

no-cloning theorem - proof cont'd

But $|\psi\rangle$ and $|\phi\rangle$ are arbitrary states that generally satisfy

$$0 \leq \langle \psi | \phi \rangle \leq 1,$$

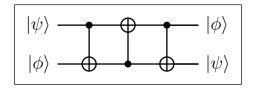
and thus the contradiction is achieved when $|\psi\rangle$ and $|\phi\rangle$ are chosen not to be orthogonal (nor identical).

Hence, no such unitary operator \mathcal{U} exists, and the theorem holds.



swapping rather than cloning

Cloning is not possible, but swapping is:



Information-wise, think of the swap as rearranging the existing information between the two qubits (rather than information exchange). The CNOT sequence might make the information accessible in a different way, but the total information content remains the same.

Tomorrow: teleporting rather than cloning