Arrow's & May's Theorems*

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1 Motivation & Preliminaries

In this set of notes, we first examine Arrow's Theorem, a famous impossibility result in social choice theory. Originally stated in [1] by Kenneth J. Arrow in 1951, we have the unfortunate result that all "reasonable" voting rules, involving 3 or more alternatives, must be dictatorial. After our melancholic endeavor of proving Arrow's result, we will consider May's Theorem, which will provide slight relief, for the case of 2 alternatives.

We start with some definitions, following closely the exposition of [2].

Let $N = \{1, ..., n\}$ be a finite set of voters, and let A be a finite set of alternatives—or candidates. Consider the following definitions.

Definition 1 (Weak and Linear Orders). A binary relation on a finite set A is a weak order if it is both complete and transitive. A linear order is a weak order that is additionally antisymmetric. Denote the set of weak orders \succeq on A by $\mathcal{R}(A)$ and the set of linear orders \succeq on A by $\mathcal{L}(A)$. Note that \succ denotes the strict part of \succeq .

Weak orders are used to model preferences permitting ties, and linear orders are used to model strict preferences. The preference of $i \in N$ is denoted by \succ_i . We now define social welfare functions, the central object of Arrow's Theorem.

Definition 2 (Social Welfare Functions (SWFs)). A social welfare function f is a map of the form $f: \mathcal{L}(A)^n \to \mathcal{R}(A)$.

Think of Definition 2 as a type of voting rule: taking in a profile of voters' preferences $P = (\succeq_1, \ldots, \succeq_n)$, f returns aggregates the preferences into a weak order. Arrow suggested two natural axioms that SWFs should satisfy. We state them here.

Axiom 1 (Weakly Paretian). An SWF f is weakly Paretian if for all $x_1, x_2 \in A$, $x_1 \succ_f x_2$ whenever $x_1 \succ_i x_2$ for all $i \in N$.

Axiom 2 (Independence of Irrelevant Alternatives (IIA)). An SWF f is independent of irrelevant alternatives if for all $x_1, x_2 \in A$, the relative ranking of x_1 and x_2 by f depends only on the relative rankings of x_1 and x_2 provided by the individuals, and not on the individuals' rankings of some irrelevant alternative x_3 .

Intuitively, Axiom 1 specifies that if all voters rank one candidate over another, the SWF must reflect this preference. Axiom 2 ensures that SWF's ranking of two alternatives shouldn't depend on voters' preferences involving a third "irrelevant alternative." Axiom 2 can be thought of as a guard against voters' having an incentive to strategize and misreport their true preferences.

We now define the notion of a dictatorial SWF.

^{*}Originally written to better understand basic ideas in social choice theory.

Definition 3 (Dictatorial SWFs). An SWF f is dictatorial if there exists $i^* \in N$ such that for all $x_1, x_2 \in A$, $x_1 \succ_{i^*} x_2$ implies $x_1 \succ_f x_2$. We refer to i^* as the dictator under f.

Now, we define the central object of May's Theorem: social choice functions.

Definition 4 (Social Choice Functions (SCFs)). A social choice function f is a map of the form $f: \mathcal{L}(A)^n \to 2^A \setminus \emptyset$, where 2^A is the power set of A.

Think of Definition 4 as a type of voting rule: taking in a profile of voters' preferences $P = (\succ_1, \ldots, \succ_n)$, f returns a set of "winners." If |f(P)| = 1, we say f is single-valued on P. In this case, we may use the semantics of $f : \mathcal{L}(A)^n \to A$. We say f is resolute if it is single-valued for all profiles. We now state some natural axioms, just as we did for SWFs.

Axiom 3 (Anonymity). An SCF f is anonymous if each pair of voters are interchangeable. That is, $f(P) = f(P^*)$ for profiles P and P^* , whenever P^* is obtained from P by swapping the ballots cast by two voters i and j. Moreover, we say that f is dictatorial with dictator i^* if f(P) corresponds to the top-ranked alternative of i^* for all profiles P.

Axiom 4 (Neutrality). An SCF f is neutral if each pair of alternatives are interchangeable. That is, when P^* is obtained from P by swapping the positions of alternatives x_1 and x_2 in every ballot, $f(P^*)$ is obtained from P by a similar swap. Moreover, we say that f is imposed if there exists an unelectable candidate x; i.e. for no profile P does $f(P) = \{x\}$.

Axiom 5 (Monotonicity & Positive Responsiveness). An SCF f is monotone if for a preference P, $x \in f(P)$ and for P^* obtained from P by just having one voter rank x higher in their ballot, $x \in f(P^*)$. We say f is positive responsive if $x \in f(P)$ and for P^* obtained from P by just having one voter rank x higher in their ballot, $\{x\} = f(P^*)$.

Intuitively, Axiom 3 specifies that the SCF treats all voters equally: a ballot cast by one voter yields the same preference as the same ballot cast by another voter. Note that nondictatoriality is a very weak form of anonymity. Axiom 4 ensures that permuting alternatives' identities on the ballots yields an analogous permutation in the results. Note that nonimposition is a very weak form of neutrality. Finally, Axiom 5 requires that alternatives are not negatively affected by voters ranking them higher.

Remark. As we will explore soon, Axiom 5 is especially insightful when there are only two alternatives; positive responsiveness, in particular, helps us grapple with ties.

2 The Heart of Arrow's Theorem

We build up to Arrow's Theorem with some lemmas and definitions. Here, we follow the general argument of [2] with some reorganization and notational differences. For the remainder of this section, $|A| \geq 3$ unless otherwise stated.

Lemma 1 (Dictatorial SWF \implies Weakly Paretian and IIA). Any dictatorial SWF f is both weakly Paretian and IIA.

Proof. Let f be a dictatorial SWF with dictator i^* . Consider arbitrary $x_1, x_2 \in A$.

If for all $i \in N$ $x_1 \succ_i x_2$, then it must be the case that $x_1 \succ_{i^*} x_2$, so $x_1 \succ_f x_2$. Therefore, f is weakly Paretian.

Since the ordering \succ_f is equivalent to the ordering \succ_{i^*} , it is also immediate that the ranking of x_1 and x_2 under f is equivalent to that of that of i^* and doesn't depend on the preferences that i^* has on a third alternative.

Definition 5 (Coalitions). A subset $C \subseteq N$ is called a coalition. We say C is decisive over (x_1, x_2) if $x_1 \succ_f x_2$ whenever $x_1 \succ_i x_2$ for all $i \in C$. Additionally, we say C is weakly decisive over (x_1, x_2) if $x_1 \succ_f x_2$ whenever $x_1 \succ_i x_2$ for all $i \in C$ and $x_2 \succ_j x_1$ for all $j \notin C$.

Lemma 2 (Field Expansion, Weakly Decisive \implies Decisive). A weakly Paretian and IIA SWF f, with a weakly decisive coalition C over (x_1, x_2) , is a decisive for all alternatives.

Proof. Consider mutually distinct $x_1, x_2, x'_1, x'_2 \in A$, and let C be weakly decisive over (x_1, x_2) . We will show that C is decisive over the alternatives (x'_1, x'_2) .

Let $x'_1 \succ_i x_1 \succ_i x_2 \succ_i x'_2$ for all $i \in C$, and for all $j \notin C$, let $x'_1 \succ_j x_1$, $x_2 \succ_j x'_2$, and $x_2 \succ_j x_1$, Since C is weakly decisive over (x_1, x_2) , we must have $x_1 \succ_f x_2$. Since f is weakly Paretian, we also have $x'_1 \succ_f x_1$ and $x_2 \succ_f x'_2$. By transitivity of \succ_f , we get $x'_1 \succ_f x'_2$. So, C is decisive over (x'_1, x'_2) .

Because the choice of (x'_1, x'_2) was arbitrary, C is decisive over all pairs of alternatives.

Remark. To be very explicit, the construction of preferences in Lemma 2 is to exploit two assumptions. Since $x_1 \succ_i x_2$ for $i \in C$ and $x_2 \succ_j x_1$ for $j \notin C$, we can use the weakly decisiveness of C to conclude $x_1 \succ_f x_2$. Since $x_1' \succ_k x_1$ and $x_2 \succ_k x_2'$ for all $k \in N$, we can use the weakly Paretian nature of f to conclude $x_1' \succ_f x_1$ and $x_2 \succ_f x_2'$.

Importantly, in our proof, we also did not need to consider how voters outside C rank x_1' versus x_2' . We obtained $x_1' \succ_f x_2'$ only from exploiting the weak decisiveness of C, the weakly Paretian SWF f, and the transitivity of \succ_f . Also, note that by IIA, we only needed to make sure that all $i \in C$ had $x_1' \succ_i x_2'$. If instead $x_2 \succ_i x_1$, a similar argument would hold.

Lemma 3 (Splitting). Let $C \subseteq N$ be a decisive coalition, with respect to some pair of alternatives. Additionally, let $|C| \ge 2$. Then, we can write $C = C_1 \cup C_2$ with $C_1 \ne \emptyset$, $C_2 \ne \emptyset$, and $C_1 \cap C_2 = \emptyset$, where either C_1 or C_2 is decisive over all pairs of alternatives.

Proof. Recalling that $|A| \geq 3$, suppose $x_1 \succ_i x_2 \succ_i x_3$ for all $i \in C_1$, $x_2 \succ_j x_3 \succ_j x_1$ for all $j \in C_2$, and $x_3 \succ_k x_1 \succ_k x_2$ for all $k \notin C_1 \cup C_2$. Because C is decisive, $x_2 \succ_f x_3$. Then, either $x_1 \succ_f x_3$ or $x_3 \succeq_f x_1$.

- Case 1 $(x_1 \succ_f x_3)$: We see that the preferences in C_1 match those aggregated by f. Since f is IIA, whenever voters in C_1 rank x_1 above x_3 , the SWF does the same. So, C_1 is weakly decisive over (x_1, x_3) . But by Lemma 2, C_1 is decisive for all pairs of alternatives.
- Case 2 $(x_3 \succeq_f x_1)$: By transitivity, and $x_2 \succ_f x_3$, we have that $x_2 \succ_f x_1$. So, the preferences in C_2 match those aggregated by f. Since f is IIA, whenever voters in C_2 rank x_2 above x_1 , the SWF does the same. So, C_2 is weakly decisive over (x_1, x_2) . But by Lemma 2, C_2 is decisive for all pairs of alternatives.

We have shown the desired result.

Theorem 1 (Arrow, Weakly Paretian and IIA \iff Dictatorial SWF). When $|A| \ge 3$, an SWF is weakly Paretian and IIA if and only if it is dictatorial.

Proof. By Lemma 1, we need only show that an arbitrary weakly Paretian and IIA SWF f is dictatorial. Note that N is a decisive coalition since f is weakly Paretian. Then, we can apply Lemma 3 repeatedly to obtain smaller and smaller decisive coalitions. Once we obtain a singleton decisive coalition, we are done. The element of the singleton decisive coalition is the dictator under f. Note that the inductive argument is valid since N is finite.

3 May's Theorem

We state and prove May's Theorem, with relative ease compared to the previous section. We again follow [2].

Theorem 2 (May, Majority Rule is Best for 2 Alternatives). For two alternatives and an odd number of voters, majority rule is the unique resolute, anonymous, neutral, and monotone SCF. For two alternatives and any number of voters, it is the unique anonymous, neutral, and positively responsive SCF.

Proof. Let $A = \{x, y\}$. Trivially, majority rule satisfies all the above properties.

For uniqueness, with any other SCF, we choose a profile where x wins, but with fewer votes than y. Suppose we switch enough ballots to reverse the number of votes x and y each have. Monotonicity implies that x still wins; however, neutrality and anonymity implies that y wins. If x and y tie, meaning $\{x,y\} \in f(P)$, but x has fewer votes than y, positive responsiveness similarly contradicts neutrality and anonymity.

References

- [1] Arrow, K.J.: Social Choice and Individual Values. Yale University Press (1951)
- [2] Brandt, F., Conitzer, V., Endriss, U., Jérôme, L., Procaccia, A.D.: Handbook of Computational Social Choice. Cambridge University Press (2016), https://procaccia.info/wp-content/uploads/2020/03/comsoc.pdf