Apache Performance Tuning

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Warning

This document is partially out of date and might be inaccurate.

- Hardware and Operating System Issues
- Run-Time Configuration Issues
- Compile-Time Configuration Issues
- Appendix: Detailed Analysis of a Trace

See also

- Comments

Apache 2.4 is a general-purpose webserver, designed to provide a balance of flexibility, portability, and performance. Although it has not been designed specifically to set benchmark records, Apache 2.4 is capable of high performance in many real-world situations.

This document describes the options that a server administrator can configure to tune the performance of an Apache 2.4 installation. Some of these configuration options enable the httpd to better take advantage of the capabilities of the hardware and OS, while others allow the administrator to trade functionality for speed.

Hardware and Operating System Issues

The single biggest hardware issue affecting webserver performance is RAM. A webserver should never ever have to swap, as swapping increases the latency of each request beyond a point that users consider "fast enough". This causes users to hit stop and reload, further increasing the load. You can, and should, control the `MaxRequestWorkers` setting so that your server does not spawn so many children that it starts swapping. The procedure for doing this is simple: determine the size of your average Apache process, by looking at your process list via a tool such as `top`, and divide this into your total available memory, leaving some room for other processes.

Beyond that the rest is mundane: get a fast enough CPU, a fast enough network card, and fast enough disks, where "fast enough" is something that needs to be determined by experimentation.

Operating system choice is largely a matter of local concerns. But some guidelines that have proven generally useful are:

- Run the latest stable release and patch level of the operating system that you choose. Many OS suppliers have introduced significant performance improvements to their TCP stacks and thread libraries in recent years.
- If your OS supports a `sendfile(2)` system call, make sure you install the release and/or patches needed to enable it. (With Linux, for example, this means using Linux 2.4 or later. For early releases of Solaris 8, you may need to apply a patch.) On systems where it is available, `sendfile` enables Apache to deliver static content faster and with lower CPU utilization.

Run-Time Configuration Issues

Related Modules	Related Directives
mod_dir mpm_common mod_status	AllowOverride DirectoryIndex HostnameLookups EnableMMAP EnableSendfile KeepAliveTimeout MaxSpareServers MinSpareServers Options StartServers

HostnameLookups and other DNS considerations

Prior to Apache 1.3, `HostnameLookups` defaulted to `On`. causing an extra latency penalty for every request due to a DNS lookup to complete before the request was finished. In Apache 2.4 this setting defaults to `Off`. If you need to have addresses in your log files resolved to hostnames, please consider post-processing rather than forcing Apache to do it in the first place. It is recommended that you do this sort of post-processing of your log files on some machine other than the production web server machine, in order that this activity not adversely affect server performance.

If you use any `Allow from domain` or `Deny from domain` directives (i.e., using a hostname, or a domain name, rather than an IP address) then you will pay for two DNS lookups (a reverse, followed by a forward lookup to make sure that the reverse is not being spoofed). For best performance, whenever it is possible, use IP addresses rather than domain names.

> **Warning:**
>
> Please use the `Require` directive with Apache 2.4; more info in the related upgrading guide.

Note that it's possible to scope the directives, such as within a `<Location "/server-status">` section. In this case the DNS lookups are only performed on requests matching the criteria. Here's an example which disables lookups except for `.html` and `.cgi` files:

```
<Files ~ "\.(html|cgi)$">
  HostnameLookups on
</Files>
```

But even still, if you just need DNS names in some CGIs you could consider doing the `gethostbyname` call in the specific CGIs that need it.

FollowSymLinks and SymLinksIfOwnerMatch

Wherever in your URL-space you do not have an `Options FollowSymLinks`, or you do have an `Options SymLinksIfOwnerMatch`, Apache will need to issue extra system calls to check up on symlinks. (One extra call per filename component.) For example, if you had:

```
DocumentRoot "/www/htdocs"
<Directory "/">
  Options SymLinksIfOwnerMatch
</Directory>
```

and a request is made for the URI `/index.html`, then Apache will perform `lstat(2)` on `/www`, `/www/htdocs`, and `/www/htdocs/index.html`. The results of these `lstats` are never cached, so they will occur on every single request. If you really desire the symlinks security checking, you can do something like this:

```
DocumentRoot "/www/htdocs"
<Directory "/">
  Options FollowSymLinks
</Directory>

<Directory "/www/htdocs">
  Options -FollowSymLinks +SymLinksIfOwnerMatch
</Directory>
```

This at least avoids the extra checks for the `DocumentRoot` path. Note that you'll need to add similar sections if you have any `Alias` or `RewriteRule` paths outside of your document root. For highest performance, and no symlink protection, set `FollowSymLinks` everywhere, and never set `SymLinksIfOwnerMatch`.

AllowOverride

Wherever in your URL-space you allow overrides (typically `.htaccess` files), Apache will attempt to open `.htaccess` for each filename component. For example,

```
DocumentRoot "/www/htdocs"
<Directory "/">
  AllowOverride all
</Directory>
```

and a request is made for the URI `/index.html`. Then Apache will attempt to open `/.htaccess`, `/www/.htaccess`, and `/www/htdocs/.htaccess`. The solutions are similar to the previous case of `Options FollowSymLinks`. For highest performance use `AllowOverride None` everywhere in your filesystem.

Negotiation

If at all possible, avoid content negotiation if you're really interested in every last ounce of performance. In practice the benefits of negotiation outweigh the performance penalties. There's one case where you can speed up the server. Instead of using a wildcard such as:

```
DirectoryIndex index
```

Use a complete list of options:

```
DirectoryIndex index.cgi index.pl index.shtml index.htm
```

where you list the most common choice first.

Also note that explicitly creating a `type-map` file provides better performance than using `MultiViews`, as the necessary information can be determined by reading this single file, rather than having to scan the directory for files.

If your site needs content negotiation, consider using `type-map` files, rather than the `Options MultiViews` directive to accomplish the negotiation. See the [Content Negotiation] documentation for a full discussion of the methods of negotiation, and instructions for creating `type-map` files.

Memory-mapping

In situations where Apache 2.x needs to look at the contents of a file being delivered--for example, when doing server-side-include processing--it normally memory-maps the file if the OS supports some form of `mmap(2)`.

On some platforms, this memory-mapping improves performance. However, there are cases where memory-mapping can hurt the performance or even the stability of the httpd:

- On some operating systems, `mmap` does not scale as well as `read(2)` when the number of CPUs increases. On multiprocessor Solaris servers, for example, Apache 2.x sometimes delivers server-parsed files faster when `mmap` is disabled.
- If you memory-map a file located on an NFS-mounted filesystem and a process on another NFS client machine deletes or truncates the file, your process may get a bus error the next time it tries to access the mapped file content.

For installations where either of these factors applies, you should use `EnableMMAP off` to disable the memory-mapping of delivered files. (Note: This directive can be overridden on a per-directory basis.)

Sendfile

In situations where Apache 2.x can ignore the contents of the file to be delivered -- for example, when serving static file content -- it normally uses the kernel sendfile support for the file if the OS supports the `sendfile(2)` operation.

On most platforms, using sendfile improves performance by eliminating separate read and send mechanics. However, there are cases where using sendfile can harm the stability of the httpd:

- Some platforms may have broken sendfile support that the build system did not detect, especially if the binaries were built on another box and moved to such a machine with broken sendfile support.
- With an NFS-mounted filesystem, the kernel may be unable to reliably serve the network file through its own cache.

For installations where either of these factors applies, you should use `EnableSendfile off` to disable sendfile delivery of file contents. (Note: This directive can be overridden on a per-directory basis.)

Recycle child processes

`MaxConnectionsPerChild` limits the numbers of connections that a child process can handle during its lifetime (by default set to `0` - unlimited). This affects all the MPMs, even the ones using threads. For example, each process created by the `worker` MPM spawns multiple threads that will handle connections, but this does not influence the overall count. It only means that the sum of requests handled by all the threads spawned by a single process will be counted against the `MaxConnectionsPerChild` value.

`MaxConnectionsPerChild` should not have any limit in the optimal use case, since there should not be any reason to force a process kill other than software bugs causing memory leaks or excessive CPU usage.

When keep-alives are in use, a process (or a thread spawned by a process) will be kept busy doing nothing but waiting for more requests on the already open connection. The default `KeepAliveTimeout` of `5` seconds attempts to minimize this effect. The tradeoff here is between network bandwidth and server resources.

Compile-Time Configuration Issues

Choosing an MPM

Apache 2.x supports pluggable concurrency models, called Multi-Processing Modules (MPMs). When building Apache, you must choose an MPM to use. There are platform-specific MPMs for some platforms: `mpm_netware`, `mpmt_os2`, and `mpm_winnt`. For general Unix-type systems, there are several MPMs from which to choose. The choice of MPM can affect the speed and scalability of the httpd:

- The `worker` MPM uses multiple child processes with many threads each. Each thread handles one connection at a time. Worker generally is a good choice for high-traffic servers because it has a smaller memory footprint than the prefork MPM.
- The `event` MPM is threaded like the Worker MPM, but is designed to allow more requests to be served simultaneously by passing off some processing work to supporting threads, freeing up the main threads to work on new requests.

- The `prefork` MPM uses multiple child processes with one thread each. Each process handles one connection at a time. On many systems, prefork is comparable in speed to worker, but it uses more memory. Prefork's threadless design has advantages over worker in some situations: it can be used with non-thread-safe third-party modules, and it is easier to debug on platforms with poor thread debugging support.

For more information on these and other MPMs, please see the MPM documentation.

Modules

Since memory usage is such an important consideration in performance, you should attempt to eliminate modules that you are not actually using. If you have built the modules as DSOs, eliminating modules is a simple matter of commenting out the associated `LoadModule` directive for that module. This allows you to experiment with removing modules and seeing if your site still functions in their absence.

If, on the other hand, you have modules statically linked into your Apache binary, you will need to recompile Apache in order to remove unwanted modules.

An associated question that arises here is, of course, what modules you need, and which ones you don't. The answer here will, of course, vary from one web site to another. However, the *minimal* list of modules which you can get by with tends to include `mod_mime`, `mod_dir`, and `mod_log_config`. `mod_log_config` is, of course, optional, as you can run a web site without log files. This is, however, not recommended.

Atomic Operations

Some modules, such as `mod_cache` and recent development builds of the worker MPM, use APR's atomic API. This API provides atomic operations that can be used for lightweight thread synchronization.

By default, APR implements these operations using the most efficient mechanism available on each target OS/CPU platform. Many modern CPUs, for example, have an instruction that does an atomic compare-and-swap (CAS) operation in hardware. On some platforms, however, APR defaults to a slower, mutex-based implementation of the atomic API in order to ensure compatibility with older CPU models that lack such instructions. If you are building Apache for one of these platforms, and you plan to run only on newer CPUs, you can select a faster atomic implementation at build time by configuring Apache with the `--enable-nonportable-atomics` option:

```
./buildconf
./configure --with-mpm=worker --enable-nonportable-
atomics=yes
```

The `--enable-nonportable-atomics` option is relevant for the following platforms:

- Solaris on SPARC
  By default, APR uses mutex-based atomics on Solaris/SPARC. If you configure with `--enable-nonportable-atomics`, however, APR generates code that uses a SPARC v8plus opcode for fast hardware compare-and-swap. If you configure Apache with this option, the atomic operations will be more efficient (allowing for lower CPU utilization and higher concurrency), but the resulting executable will run only on UltraSPARC chips.
- Linux on x86
  By default, APR uses mutex-based atomics on Linux. If you configure with `--enable-nonportable-atomics`, however, APR generates code that

uses a 486 opcode for fast hardware compare-and-swap. This will result in more efficient atomic operations, but the resulting executable will run only on 486 and later chips (and not on 386).

mod_status and ExtendedStatus On

If you include `mod_status` and you also set `ExtendedStatus On` when building and running Apache, then on every request Apache will perform two calls to `gettimeofday(2)` (or `times(2)` depending on your operating system), and (pre-1.3) several extra calls to `time(2)`. This is all done so that the status report contains timing indications. For highest performance, set `ExtendedStatus off` (which is the default).

accept Serialization - Multiple Sockets

> **Warning:**
>
> This section has not been fully updated to take into account changes made in the 2.x version of the Apache HTTP Server. Some of the information may still be relevant, but please use it with care.

This discusses a shortcoming in the Unix socket API. Suppose your web server uses multiple `Listen` statements to listen on either multiple ports or multiple addresses. In order to test each socket to see if a connection is ready, Apache uses `select(2)`. `select(2)` indicates that a socket has *zero* or *at least one* connection waiting on it. Apache's model includes multiple children, and all the idle ones test for new connections at the same time. A naive implementation looks something like this (these examples do not match the code, they're contrived for pedagogical purposes):

```
for (;;) {
          for (;;) {
            fd_set accept_fds;

            FD_ZERO (&accept_fds);
            for (i = first_socket; i <= last_socket; +
              FD_SET (i, &accept_fds);
            }
            rc = select (last_socket+1, &accept_fds, N
            if (rc < 1) continue;
            new_connection = -1;
            for (i = first_socket; i <= last_socket; +
              if (FD_ISSET (i, &accept_fds)) {
                new_connection = accept (i, NULL, NULL
                if (new_connection != -1) break;
              }
            }
            if (new_connection != -1) break;
          }
          process_the(new_connection);
        }
```

But this naive implementation has a serious starvation problem. Recall that multiple children execute this loop at the same time, and so multiple children will block at `select` when they are in between requests. All those blocked children will awaken and return from `select` when a single request appears on any socket. (The number of children which awaken varies depending on the operating system and timing issues.) They will all then fall down into the loop and try to `accept` the connection. But only one will succeed (assuming there's still only one connection ready). The rest will be *blocked* in `accept`. This effectively locks those children into serving requests from that one socket and no other sockets, and they'll be stuck there until enough new requests appear on that socket to wake them all up. This starvation problem was first documented in PR#467. There are at least two solutions.

One solution is to make the sockets non-blocking. In this case the `accept` won't block the children, and they will be allowed to continue immediately. But this wastes CPU time. Suppose you have ten idle children in `select`, and one connection arrives. Then nine of those children will wake up, try to `accept` the connection, fail, and loop back into `select`, accomplishing nothing. Meanwhile none of those children are servicing requests that occurred on other sockets until they get back up to the `select` again. Overall this solution does not seem very fruitful unless you have as many idle CPUs (in a multiprocessor box) as you have idle children (not a very likely situation).

Another solution, the one used by Apache, is to serialize entry into the inner loop. The loop looks like this (differences highlighted):

```
for (;;) {accept_mutex_on ();for (;;) {
            fd_set accept_fds;

            FD_ZERO (&accept_fds);
            for (i = first_socket; i <= last_socket; +
              FD_SET (i, &accept_fds);
            }
            rc = select (last_socket+1, &accept_fds, N
            if (rc < 1) continue;
            new_connection = -1;
            for (i = first_socket; i <= last_socket; +
              if (FD_ISSET (i, &accept_fds)) {
                new_connection = accept (i, NULL, NULL
                if (new_connection != -1) break;
              }
            }
            if (new_connection != -1) break;
          }accept_mutex_off ();process the new_connect
        }
```

The functions `accept_mutex_on` and `accept_mutex_off` implement a mutual exclusion semaphore. Only one child can have the mutex at any time. There are several choices for implementing these mutexes. The choice is defined in `src/conf.h` (pre-1.3) or `src/include/ap_config.h` (1.3 or later). Some architectures do not have any locking choice made, on these architectures it is unsafe to use multiple [`Listen`] directives.

The [`Mutex`] directive can be used to change the mutex implementation of the `mpm-accept` mutex at run-time. Special considerations for different mutex implementations are documented with that directive.

Another solution that has been considered but never implemented is to partially serialize the loop -- that is, let in a certain number of processes. This would only be of interest on multiprocessor boxes where it's possible that multiple children could run simultaneously, and the serialization actually doesn't take advantage of the full bandwidth. This is a possible area of future investigation, but priority remains low because highly parallel web servers are not the norm.

Ideally you should run servers without multiple [`Listen`] statements if you want the highest performance. But read on.

accept Serialization - Single Socket

The above is fine and dandy for multiple socket servers, but what about single socket servers? In theory they shouldn't experience any of these same problems because all children can just block in `accept(2)` until a connection arrives, and no starvation results. In practice this hides almost the same "spinning" behavior discussed above in the non-blocking solution. The way that most TCP stacks are implemented, the kernel actually wakes up all processes blocked in `accept` when a single connection arrives. One of those processes gets the connection and returns to user-space. The rest spin in the kernel and go back to sleep when they discover there's no connection for them. This spinning is hidden from the user-land code, but it's there nonetheless. This can result in the same load-spiking wasteful behavior that a non-blocking solution to the multiple sockets case can.

For this reason we have found that many architectures behave more "nicely" if we serialize even the single socket case. So this is actually the default in almost all cases. Crude experiments under Linux (2.0.30 on a dual Pentium pro 166 w/128Mb RAM) have shown that the serialization of the single socket case causes less than a 3% decrease in requests per second over unserialized single-socket. But unserialized single-socket showed an extra 100ms latency on each request. This latency is probably a wash on long haul lines, and only an issue on LANs. If you want to override the single socket serialization, you can define `SINGLE_LISTEN_UNSERIALIZED_ACCEPT`, and then single-socket servers will not serialize at all.

Lingering Close

As discussed in draft-ietf-http-connection-00.txt section 8, in order for an HTTP server to **reliably** implement the protocol, it needs to shut down each direction of the communication independently. (Recall that a TCP connection is bi-directional. Each half is independent of the other.)

When this feature was added to Apache, it caused a flurry of problems on various versions of Unix because of shortsightedness. The TCP specification does not state that the `FIN_WAIT_2` state has a timeout, but it doesn't prohibit it. On systems without the timeout, Apache 1.2 induces many sockets stuck forever in the `FIN_WAIT_2` state. In many cases this can be avoided by simply upgrading to the latest TCP/IP patches supplied by the vendor. In cases where the vendor has never released patches (*i.e.*, SunOS4 -- although folks with a source license can patch it themselves), we have decided to disable this feature.

There are two ways to accomplish this. One is the socket option `SO_LINGER`. But as fate would have it, this has never been implemented properly in most TCP/IP stacks. Even on those stacks with a proper implementation (*i.e.*, Linux 2.0.31), this method proves to be more expensive (cputime) than the next solution.

For the most part, Apache implements this in a function called `lingering_close` (in `http_main.c`). The function looks roughly like this:

```
void lingering_close (int s)
        {
          char junk_buffer[2048];

          /* shutdown the sending side */
          shutdown (s, 1);

          signal (SIGALRM, lingering_death);
          alarm (30);

          for (;;) {
            select (s for reading, 2 second timeout);
            if (error) break;
            if (s is ready for reading) {
              if (read (s, junk_buffer, sizeof (junk_bu
                break;
              }
              /* just toss away whatever is here */
            }
          }

          close (s);
        }
```

This naturally adds some expense at the end of a connection, but it is required for a reliable implementation. As HTTP/1.1 becomes more prevalent, and all connections are persistent, this expense will be amortized over more requests. If you want to play with fire and disable this feature, you can define `NO_LINGCLOSE`, but this is not recommended at all. In particular, as HTTP/1.1 pipelined persistent connections come into use, `lingering_close` is an absolute necessity (and pipelined connections are faster, so you want to support them).

Scoreboard File

Apache's parent and children communicate with each other through something called the scoreboard. Ideally this should be implemented in shared memory. For those operating systems that we either have access to, or have been given detailed ports for, it typically is implemented using shared memory. The rest default to using an on-disk file. The on-disk file is not only slow, but it is unreliable (and less featured). Peruse the `src/main/conf.h` file for your architecture, and look for either `USE_MMAP_SCOREBOARD` or `USE_SHMGET_SCOREBOARD`. Defining one of those two (as well as their companions `HAVE_MMAP` and `HAVE_SHMGET` respectively) enables the supplied shared memory code. If your system has another type of shared memory, edit the file `src/main/http_main.c` and add the hooks necessary to use it in Apache. (Send us back a patch too, please.)

> Historical note: The Linux port of Apache didn't start to use shared memory until version 1.2 of Apache. This oversight resulted in really poor and unreliable behavior of earlier versions of Apache on Linux.

DYNAMIC_MODULE_LIMIT

If you have no intention of using dynamically loaded modules (you probably don't if you're reading this and tuning your server for every last ounce of performance), then you should add `-DDYNAMIC_MODULE_LIMIT=0` when building your server. This will save RAM that's allocated only for supporting dynamically loaded modules.

Appendix: Detailed Analysis of a Trace

Here is a system call trace of Apache 2.0.38 with the worker MPM on Solaris 8. This trace was collected using:

```
truss -l -p httpd_child_pid.
```

The `-l` option tells truss to log the ID of the LWP (lightweight process--Solaris' form of kernel-level thread) that invokes each system call.

Other systems may have different system call tracing utilities such as `strace`, `ktrace`, or `par`. They all produce similar output.

In this trace, a client has requested a 10KB static file from the httpd. Traces of non-static requests or requests with content negotiation look wildly different (and quite ugly in some cases).

```
/67:    accept(3, 0x00200BEC, 0x00200C0C, 1) (sleeping...)
/67:    accept(3, 0x00200BEC, 0x00200C0C, 1)            = 9
```

In this trace, the listener thread is running within LWP #67.

> Note the lack of `accept(2)` serialization. On this particular platform, the worker MPM uses an unserialized accept by default unless it is listening on multiple ports.

```
/65:    lwp_park(0x00000000, 0)                          = 0
/67:    lwp_unpark(65, 1)                                = 0
```

Upon accepting the connection, the listener thread wakes up a worker thread to do the request processing. In this trace, the worker thread that handles the request is mapped to LWP #65.

```
/65:    getsockname(9, 0x00200BA4, 0x00200BC4, 1)        = 0
```

In order to implement virtual hosts, Apache needs to know the local socket address used to accept the connection. It is possible to eliminate this call in many situations (such as when there are no virtual hosts, or when `Listen` directives are used which do not have wildcard addresses). But no effort has yet been made to do these optimizations.

```
/65:    brk(0x002170E8)                                  = 0
/65:    brk(0x002190E8)                                  = 0
```

The `brk(2)` calls allocate memory from the heap. It is rare to see these in a system call trace, because the httpd uses custom memory allocators (`apr_pool` and `apr_bucket_alloc`) for most request processing. In this trace, the httpd has just been started, so it must call `malloc(3)` to get the blocks of raw memory with which to create the custom memory allocators.

```
/65:    fcntl(9, F_GETFL, 0x00000000)                    = 2
/65:    fstat64(9, 0xFAF7B818)                           = 0
/65:    getsockopt(9, 65535, 8192, 0xFAF7B918, 0xFAF7B910, 2190656) = 0
/65:    fstat64(9, 0xFAF7B818)                           = 0
/65:    getsockopt(9, 65535, 8192, 0xFAF7B918, 0xFAF7B914, 2190656) = 0
/65:    setsockopt(9, 65535, 8192, 0xFAF7B918, 4, 2190656) = 0
/65:    fcntl(9, F_SETFL, 0x00000082)                    = 0
```

Next, the worker thread puts the connection to the client (file descriptor 9) in non-blocking mode. The `setsockopt(2)` and `getsockopt(2)` calls are a side-effect of how Solaris' libc handles `fcntl(2)` on sockets.

```
/65:    read(9, " G E T   / 1 0 k . h t m".., 8000)      = 97
```

The worker thread reads the request from the client.

```
/65:    stat("/var/httpd/apache/httpd-8999/htdocs/10k.html", 0xFAF7B978) = 0
/65:    open("/var/httpd/apache/httpd-8999/htdocs/10k.html", O_RDONLY) = 10
```

This httpd has been configured with `Options FollowSymLinks` and `AllowOverride None`. Thus it doesn't need to `lstat(2)` each directory in the path leading up to the requested file, nor check for `.htaccess` files. It simply calls `stat(2)` to verify that the file: 1) exists, and 2) is a regular file, not a directory.

```
/65:    sendfilev(0, 9, 0x00200F90, 2, 0xFAF7B53C)       = 10269
```

In this example, the httpd is able to send the HTTP response header and the requested file with a single `sendfilev(2)` system call. Sendfile semantics vary among operating systems. On some other systems, it is necessary to do a `write(2)` or `writev(2)` call to send the headers before calling `sendfile(2)`.

```
/65:    write(4, " 1 2 7 . 0 . 0 . 1   -   ".., 78)      = 78
```

This `write(2)` call records the request in the access log. Note that one thing missing from this trace is a `time(2)` call. Unlike Apache 1.3, Apache 2.x uses `gettimeofday(3)` to look up the time. On some operating systems, like Linux or Solaris, `gettimeofday` has an optimized implementation that doesn't require as much overhead as a typical system call.

```
/65:    shutdown(9, 1, 1)                                      = 0
/65:    poll(0xFAF7B980, 1, 2000)                              = 1
/65:    read(9, 0xFAF7BC20, 512)                               = 0
/65:    close(9)                                               = 0
```

The worker thread does a lingering close of the connection.

```
/65:    close(10)                                              = 0
/65:    lwp_park(0x00000000, 0)         (sleeping...)
```

Finally the worker thread closes the file that it has just delivered and blocks until the listener assigns it another connection.

```
/67:    accept(3, 0x001FEB74, 0x001FEB94, 1) (sleeping...)
```

Meanwhile, the listener thread is able to accept another connection as soon as it has dispatched this connection to a worker thread (subject to some flow-control logic in the worker MPM that throttles the listener if all the available workers are busy). Though it isn't apparent from this trace, the next `accept(2)` can (and usually does, under high load conditions) occur in parallel with the worker thread's handling of the just-accepted connection.